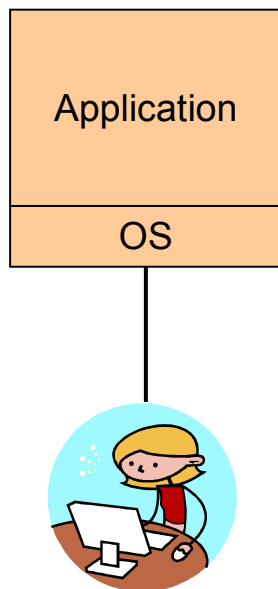


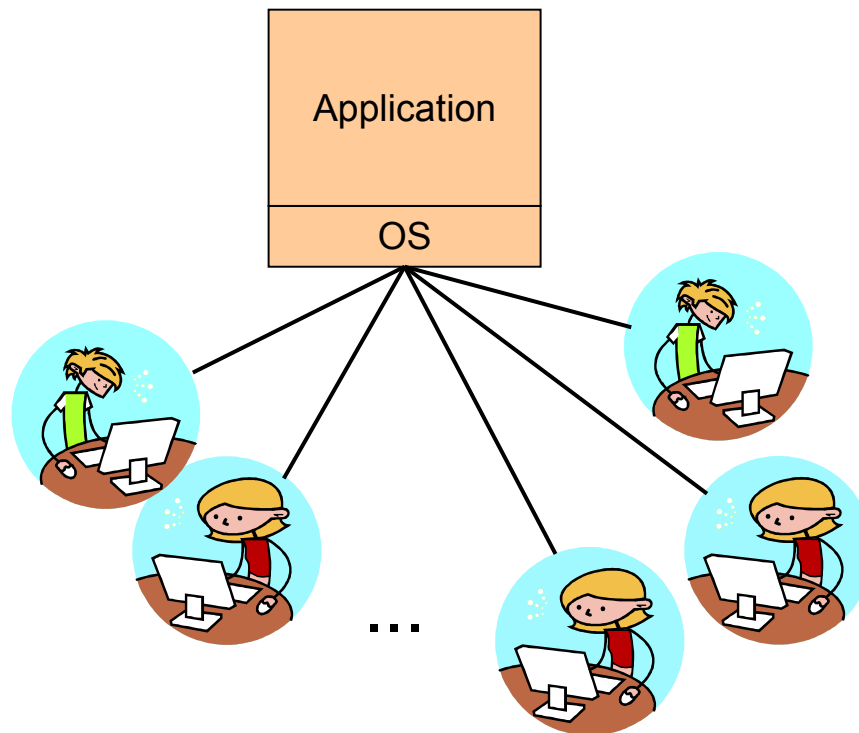
What is Work Management?

More User

Single User Environment

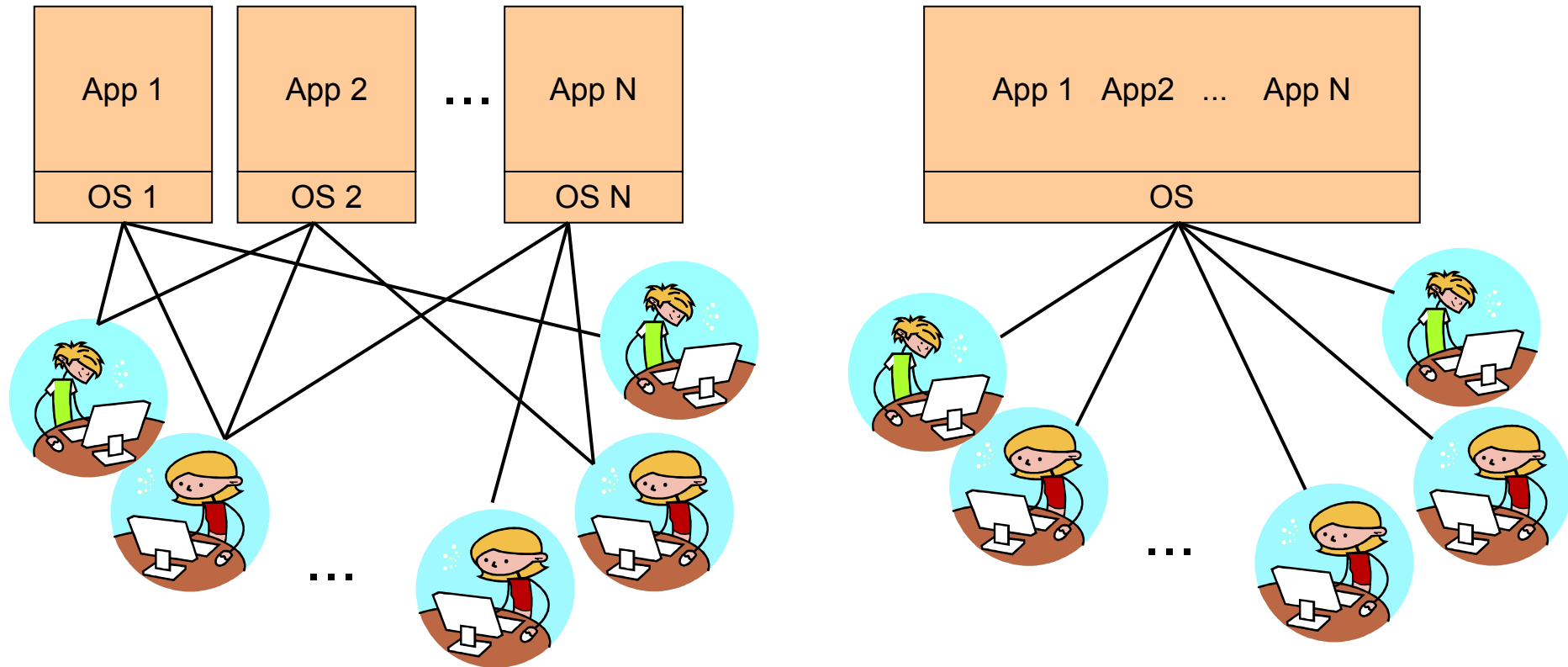


Multi User Environment



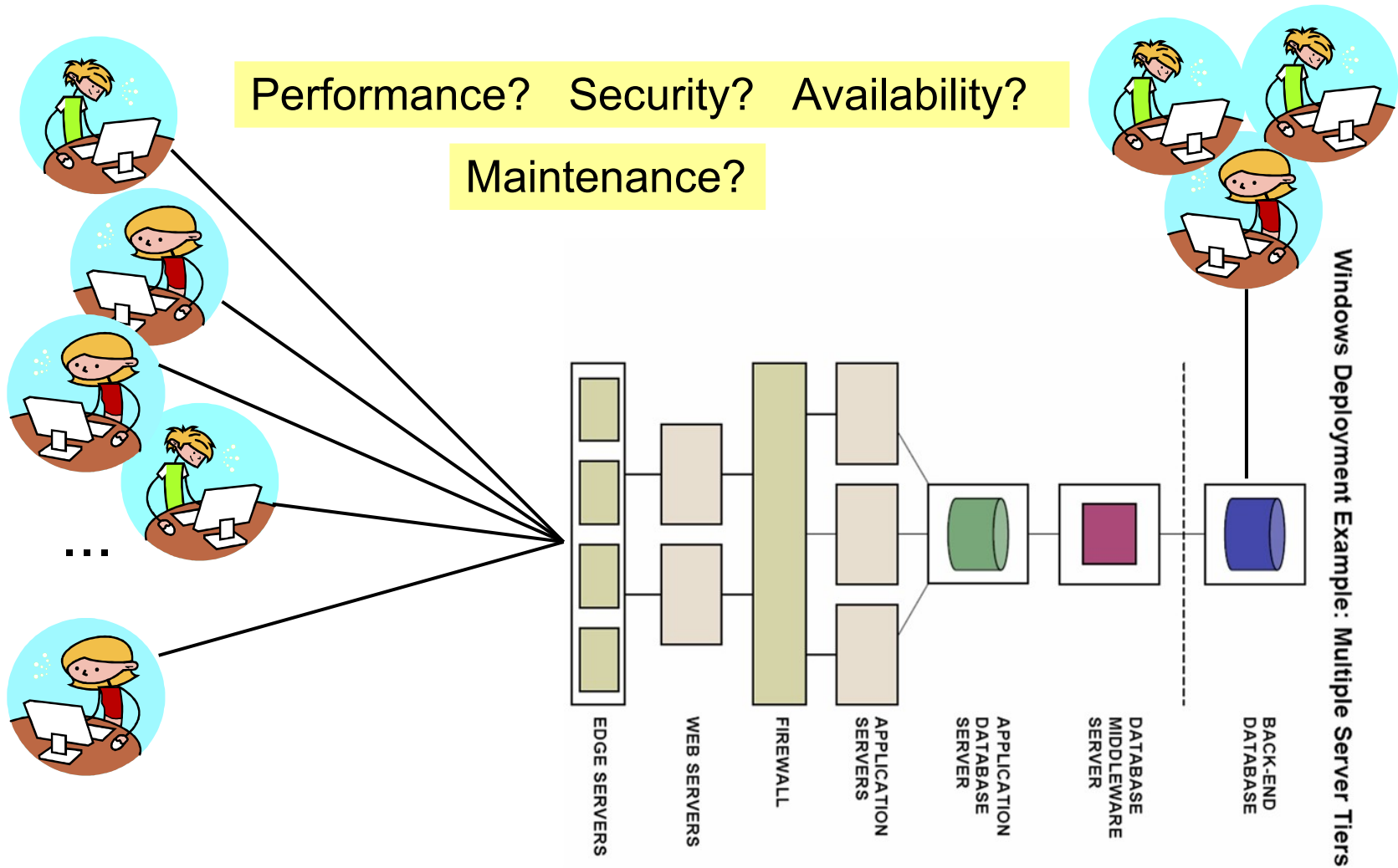
Performance? Security? Availability?

More Applications



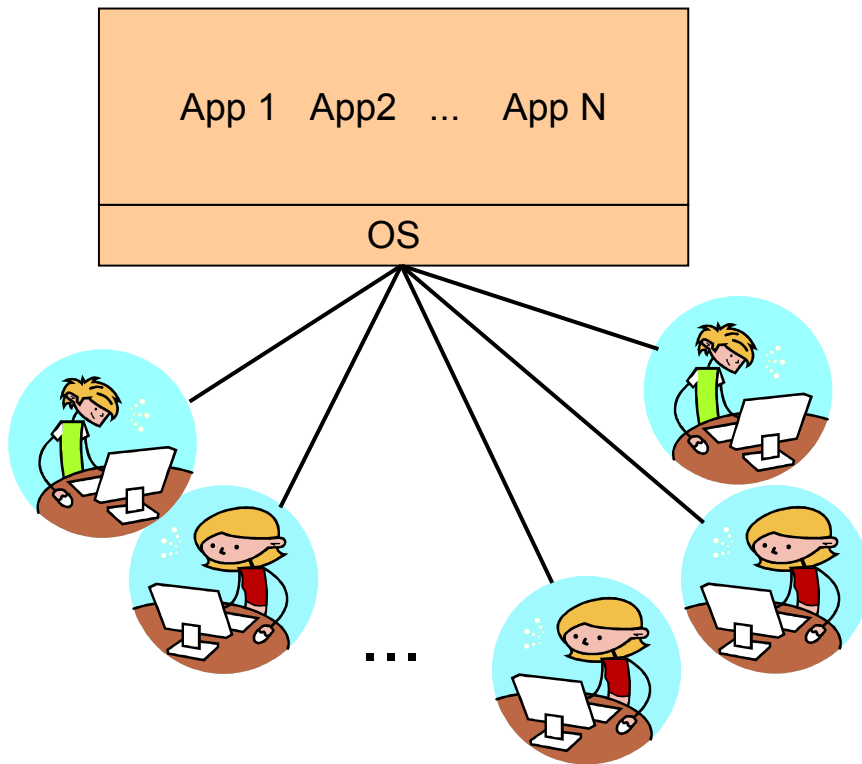
Performance? Security? Availability?
Maintenance?

Reality is more complex!



Work Management

Running Multiple Applications in a single OS Instance



Performance
(or how to share resources)

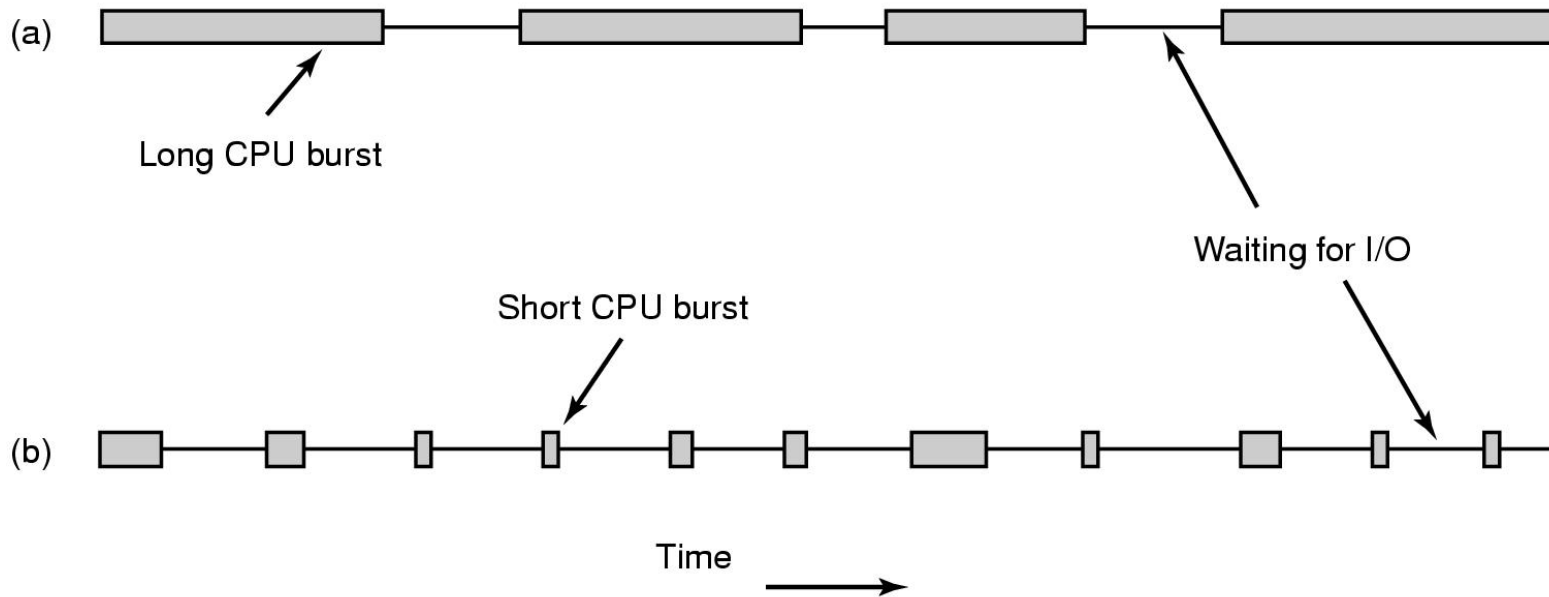
Security → Lecture 4

Availability

Do not confuse work management and LPAR!

Process Scheduling

Introduction to Scheduling



- Bursts of CPU usage alternate with periods of I/O wait
 - a CPU-bound process
 - an I/O bound process

Scheduling? Time slice?
Memory management?

Scheduling Algorithm Goals

All systems

Fairness - giving each process a fair share of the CPU

Policy enforcement - seeing that stated policy is carried out

Balance - keeping all parts of the system busy

Batch systems

Throughput - maximize jobs per hour

Turnaround time - minimize time between submission and termination

CPU utilization - keep the CPU busy all the time

Interactive systems

Response time - respond to requests quickly

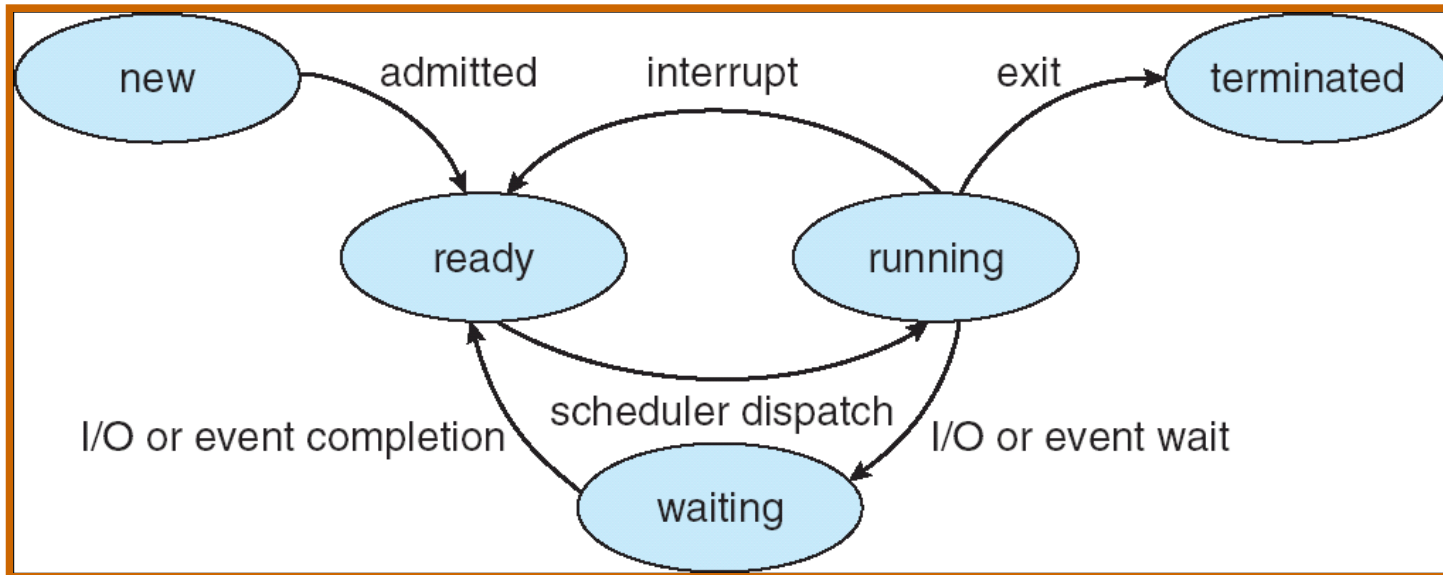
Proportionality - meet users' expectations

Real-time systems

Meeting deadlines - avoid losing data

Predictability - avoid quality degradation in multimedia systems

Diagram of Process State



- As a process executes, it changes *state*
 - new: The process is being created
 - running: Instructions are being executed
 - waiting: The process is waiting for some event to occur
 - ready: The process is waiting to be assigned to a processor
 - terminated: The process has finished execution

Preemptive and Non-preemptive Scheduling

- CPU scheduling decisions may take place, if a process
 - (1) switches from running to waiting state
 - (2) switches from running to ready state
 - (3) switches from waiting to ready state
 - (4) terminates
- When scheduling take place only under (1) or (4) the scheduling scheme is non-preemptive or preemptive.
- Examples:
 - non-preemptive: Windows 3.1
 - preemptive: Windows 95 or later, Mac OS X

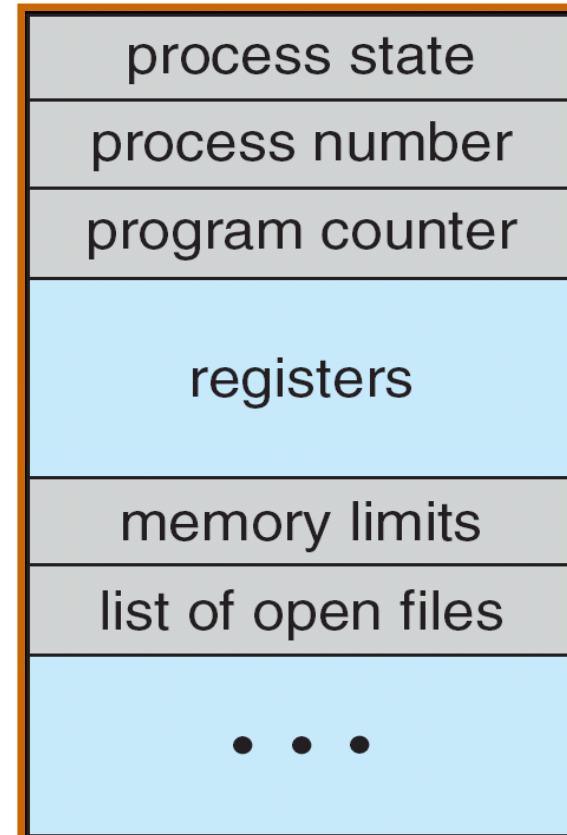
Context Switch

- When CPU switches to another process, the system must save the state of the old process and load the saved state for the new process
- Context-switch time is overhead; the system does no useful work while switching
- Time dependent on hardware support

Process Control Block (PCB)

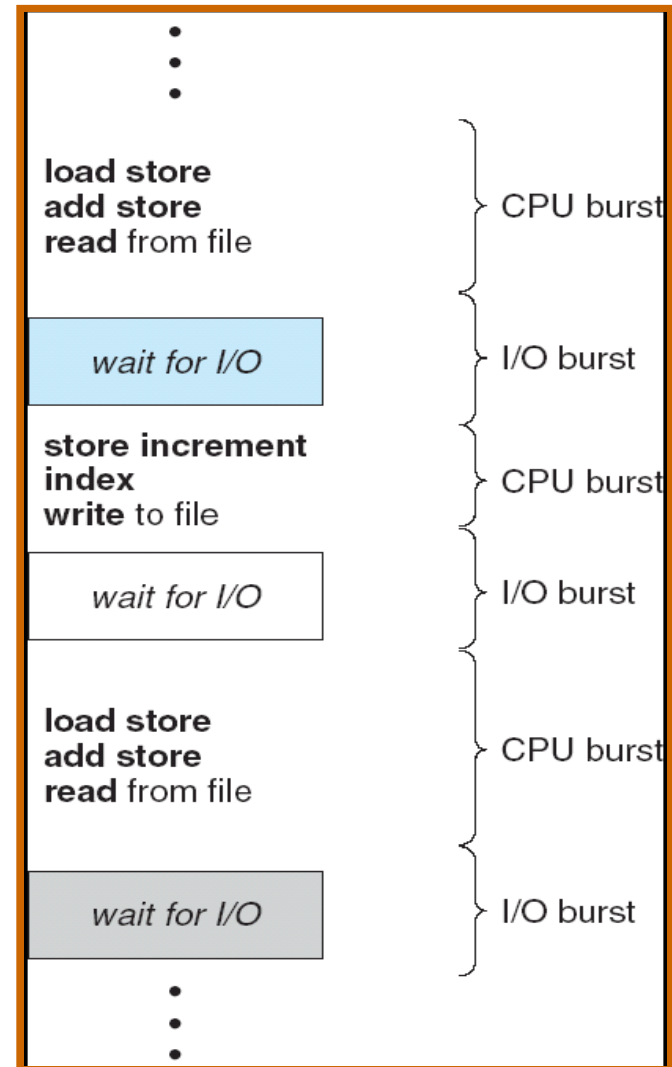
• Information associated with each process:

- Process state
- Program counter
- CPU registers
- CPU scheduling information
- Memory-management information
- Accounting information
- I/O status information

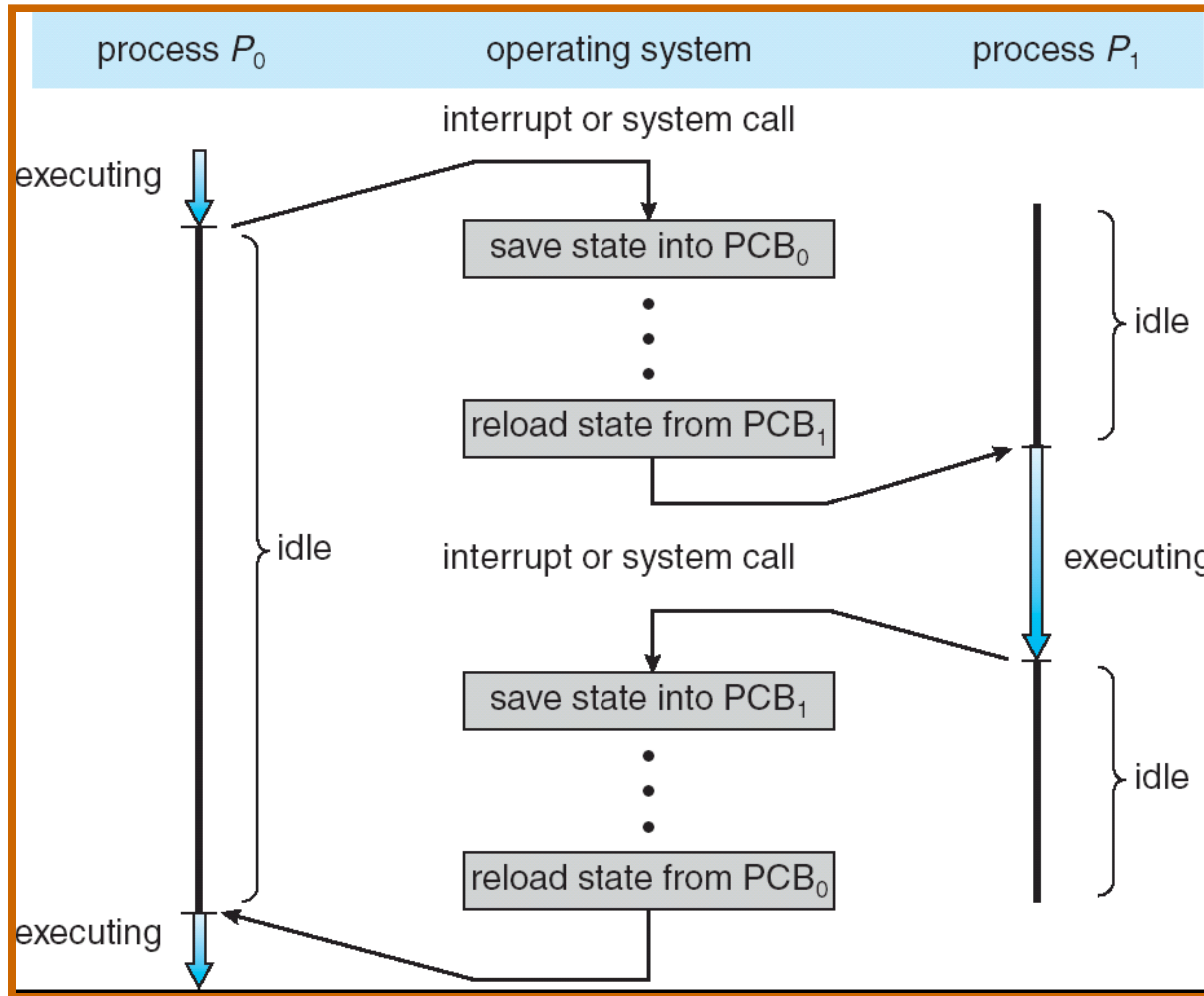


Single Process Execution

- CPU-I/O Burst Cycle – Process execution consists of a *cycle* of CPU execution and I/O wait
- CPU burst distribution



CPU Switch From Process to Process



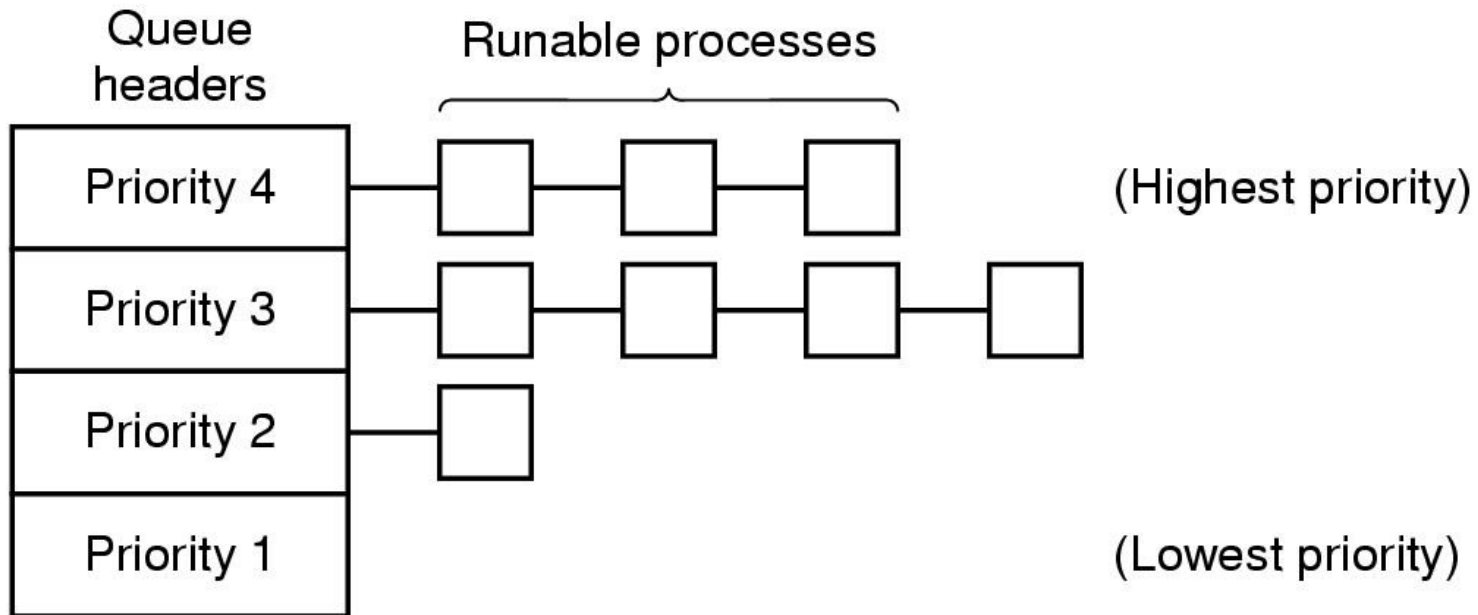
Source: Silberschatz, Galvin and Gagne: Operating System Concepts - 7th Edition

Efficient process scheduling?

Scheduling Criteria

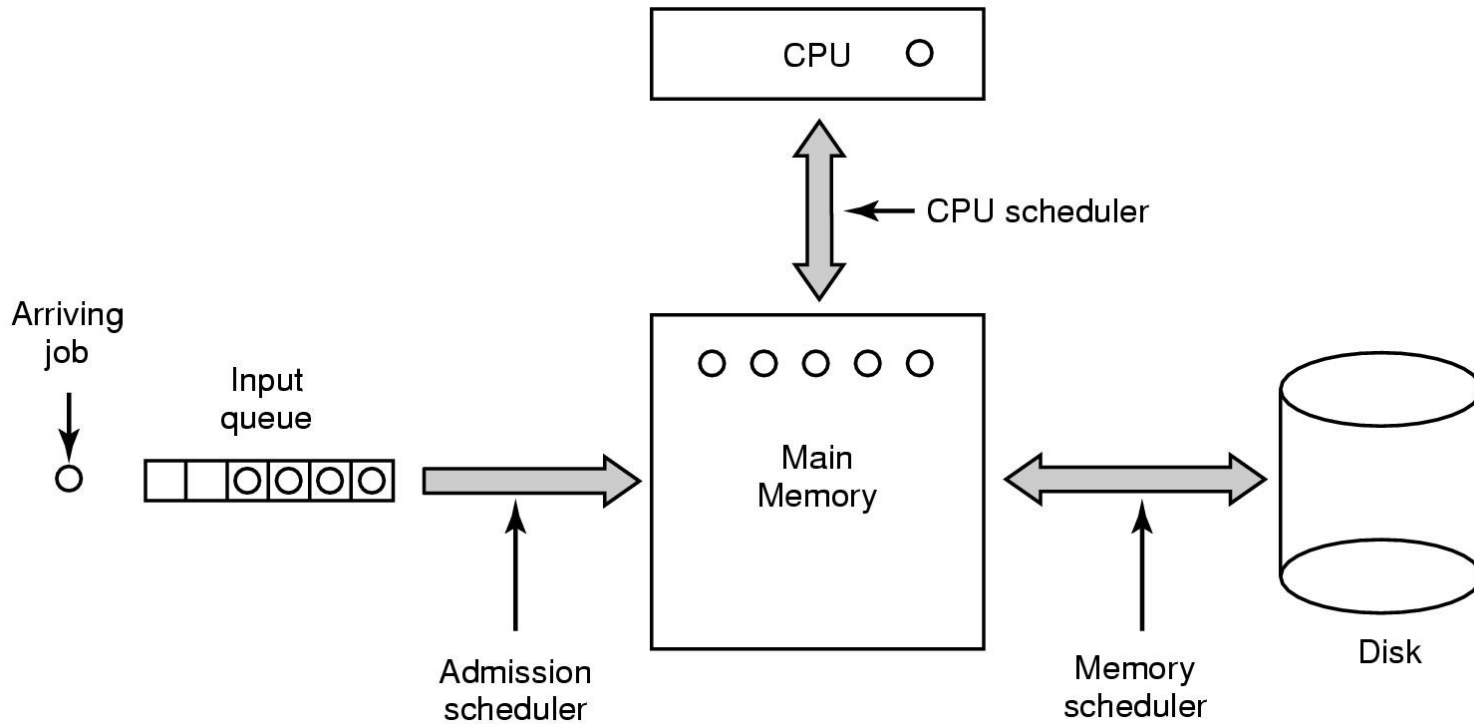
- **CPU utilization** – keep the CPU as busy as possible
- **Throughput** – # of processes that complete their execution per time unit
- **Turnaround time** – amount of time to execute a particular process
- **Waiting time** – amount of time a process has been waiting in the ready queue
- **Response time** – amount of time it takes from when a request was submitted until the first response is produced, **not** output (for time-sharing environment)

Scheduling in Interactive Systems



A scheduling algorithm with four priority classes

Scheduling in Batch Systems



Three level scheduling

Source: Andrew S. Tanenbaum: Modern Operating Systems 2th Edition

Example: IBM i

aka i5/OS, OS/400

5.1 Operator's View

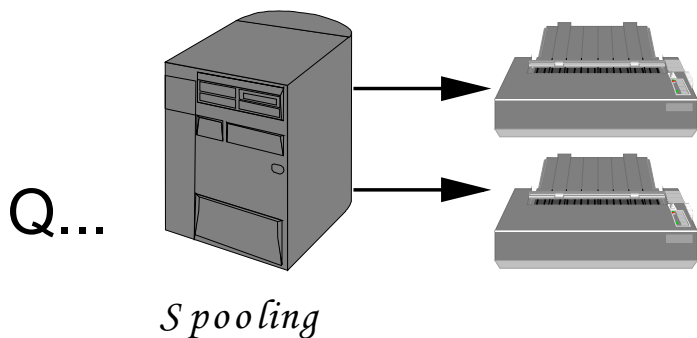
Job

Qualified job name

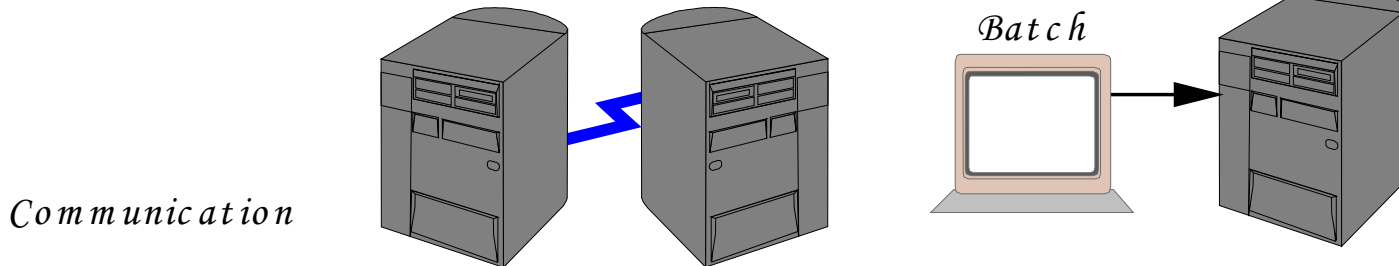
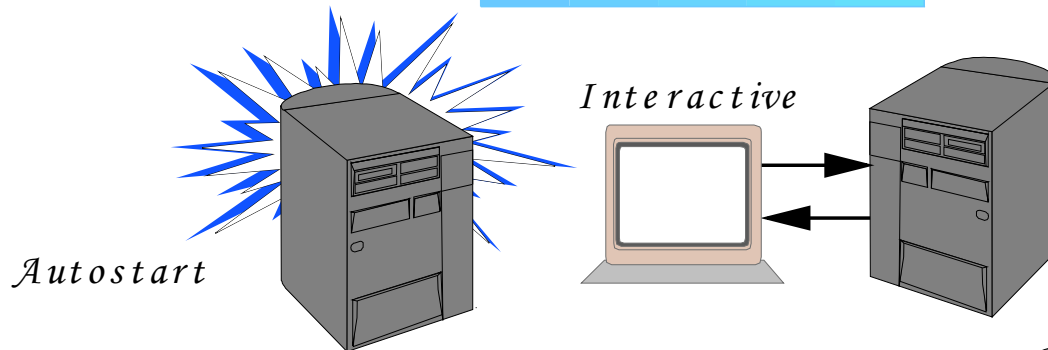


Job Types

System jobs



User jobs



Job Properties

010671/As2401/Qpadev0002 Properties - 10.31.226.222

Server | Security | Date/Time | International | Threads | Resources | Other

General | Performance

Job name: Qpadev0002
 User: As2401
 Number: 010671
 Type: Interactive
 Job user identity: As2401
 Detailed status: Waiting for workst
 Function: Major menu
 Entered system: 10/16/04 9:44:32 P
 Started: 10/16/04 9:44:32 P
 Subsystem: Qinter
 Library: Qsys
 Job description name: Qdftjobd
 Library: Qgpl

010671/As2401/Qpadev0002 Properties - 10.31.226.222

Server | Security | Date/Time | International | Threads | Resources | Other

General | Performance | Printer

Run priority (1-99): 20
 Time slice (1-9999999): 2000
 Default wait time (1-9999999): 30

Performance statistics calculated since job start

CPU time:
 Total: 50 milliseconds
 Total database: 0 milliseconds
 Maximum: No maximum

Total disk I/O count: 55
 Total lock wait time: 0 milliseconds
 Temporary storage:
 Used: 1 MB
 Maximum: No maximum

Elapsed Perform

010671/As2401/Qpadev0002 Properties - 10.31.226.222

Server | Security | Date/Time | International | Threads | Resources | Other

General | Performance | Printer Output | Messages | Job Log

Detach printer output after job has ended

Printer: Prt01 Browse...

Output queue: Oproutq Browse...

Library: As24

Priority on output queue: 5

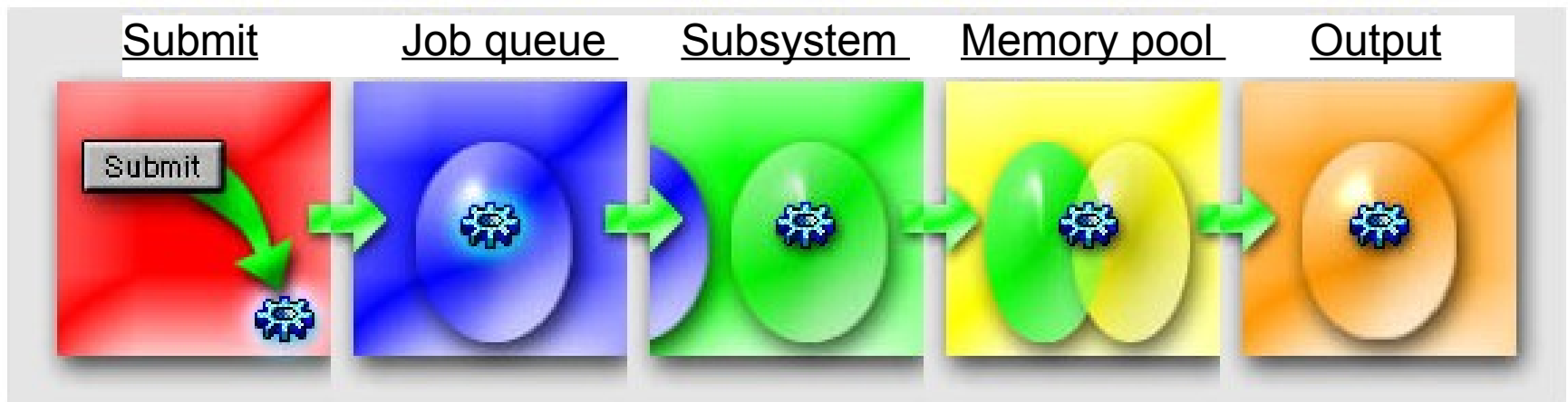
Printed page footer:

Format when using Print key: Print header

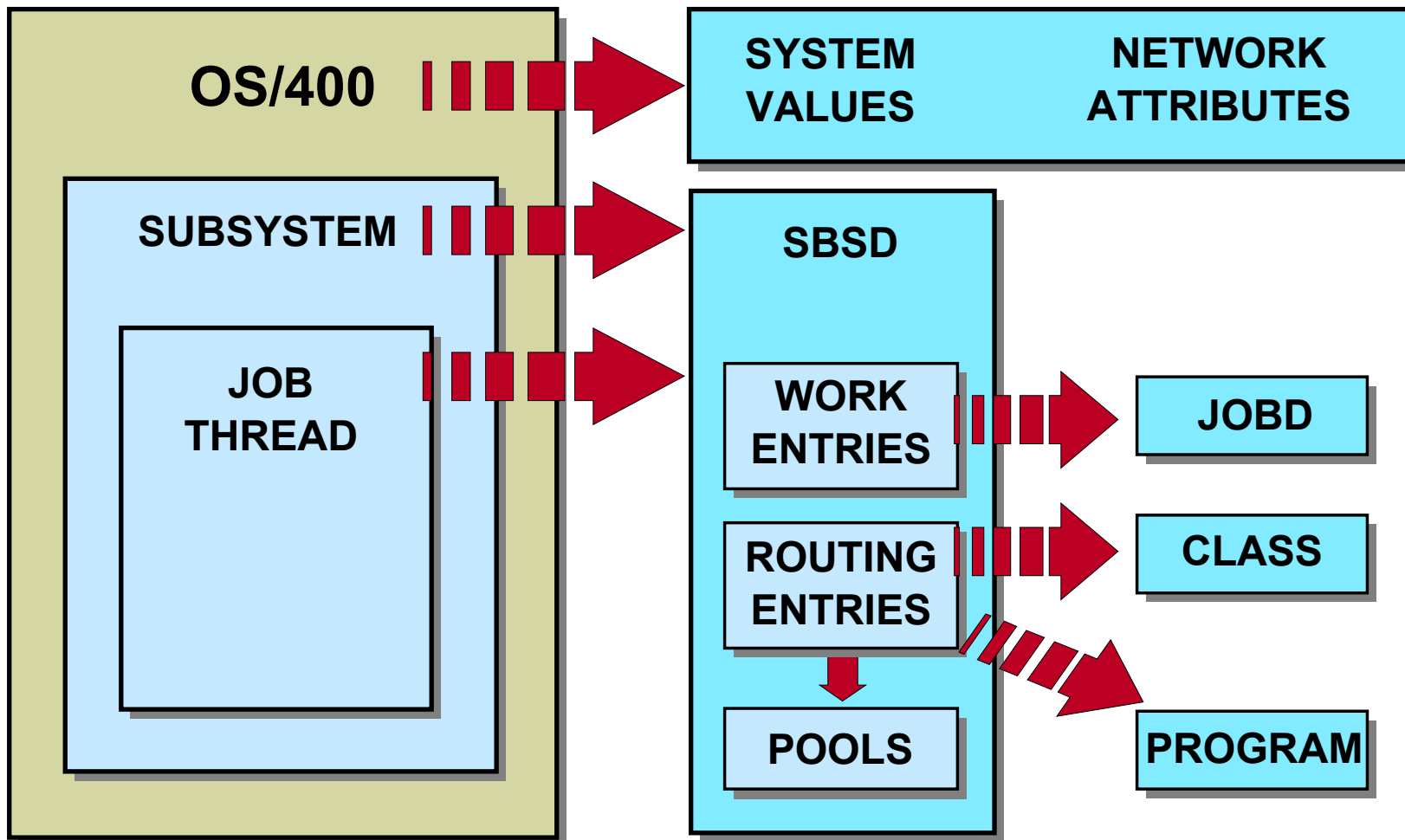
Printer Output

OK Cancel Help ?

The Life Cycle of a Job



Work Management Objects



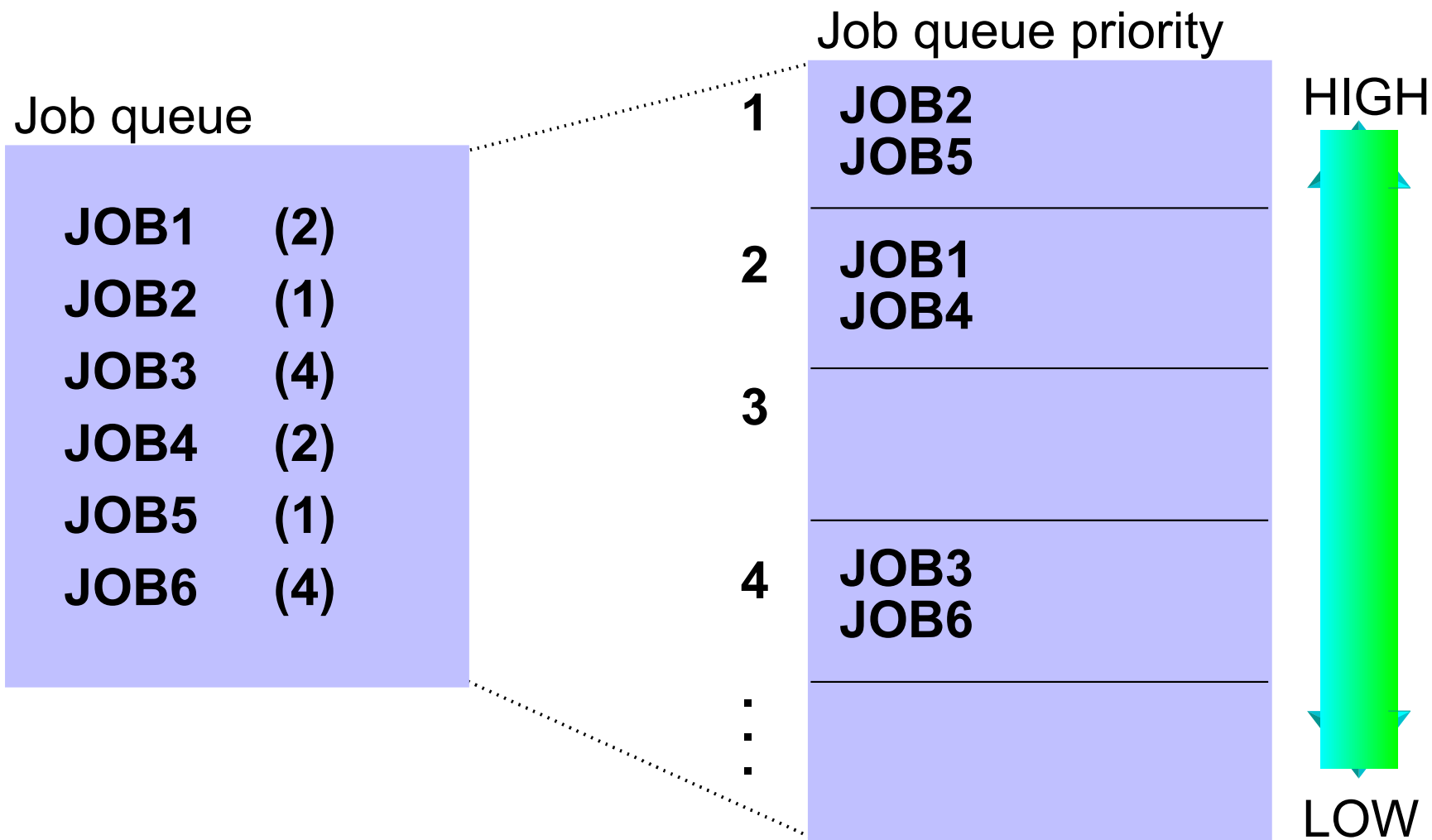
System Values

Name in iSeries Navigator	Description of system value	Name in character-based interface
Dynamically adjust job priorities of interactive jobs	Sets the job priorities of interactive jobs	QDYNPTYADJ
Dynamically adjust job priorities within priority bands	Sets job priorities within priority bands to be dynamically adjusted.	QDYNPTYSCD
Automatically adjust memory pools and activity levels	Specifies when to automatically adjust memory pools and activity levels.	QPFRADJ
Maximum eligible threads	Specifies the maximum number of eligible threads.	QMAXACTLVL
Machine memory pool size	Specifies the size of the machine memory pool.	QMCHPOOL
Base memory pool minimum size	Specifies the minimum base memory pool.	QBASPOOL
Base memory pool maximum eligible threads	Specifies maximum number of eligible threads.	QBASACTLVL
Move interactive jobs to base pool at end of time slice	Specifies whether to move interactive jobs to base pool at the end of the time slice or not.	QTSEPOOL

Network Attribute

Current System Name	ABC
Local Network ID	APPN
APPN Node type	*ENDNODE
Alert Status	*OFF
Alert Primary Focal Point	*NO
Message Queue/Obj. Dist. Output Queue/Obj. Dist.	QSYS/QSYSOPR QGPL/QPRINT
Action for DDM Request	*OBJAUT
PC Support Request Access	*OBJAUT

Job Queue



Job Description

```

                                Display Job Description
                                System:  HYDRA

Job description:  QDFTJOB      Library:  QGPL

User profile . . . . . : *RQD
CL syntax check . . . . . : *NOCHK
Hold on job queue . . . . . : *NO
End severity . . . . . : 30
Job date . . . . . : *SYSVAL
Job switches . . . . . : 00000000
Inquiry message reply . . . . . : *RQD
Job priority (on job queue) . . . . . : 5
Job queue . . . . . : QBATCH
  Library . . . . . : QGPL
Output priority (on output queue) . . . . . : 5
Printer device . . . . . : *USRPRF
Output queue . . . . . : *USRPRF
  Library . . . . . :

More...

Press Enter to continue.

F3=Exit  F12=Cancel

Message logging:
Level . . . . . : 4
Severity . . . . . : 0
Text . . . . . : *NOLIST
    
```

User Profile

Display User Profile - Basic

User profile : AS2401

Storage information:

Maximum storage allowed : *NOMAX

Storage used : 312

Storage used on independent ASP : *NO

Highest scheduling priority : 3

Job description : QDFTJOB

Library : QGPL

Accounting code :

Message queue : AS2401

Library : QUSRSYS

Message queue delivery : *NOTIFY

Message queue severity : 00

Output queue : AS24OUTQ

Library : AS2401

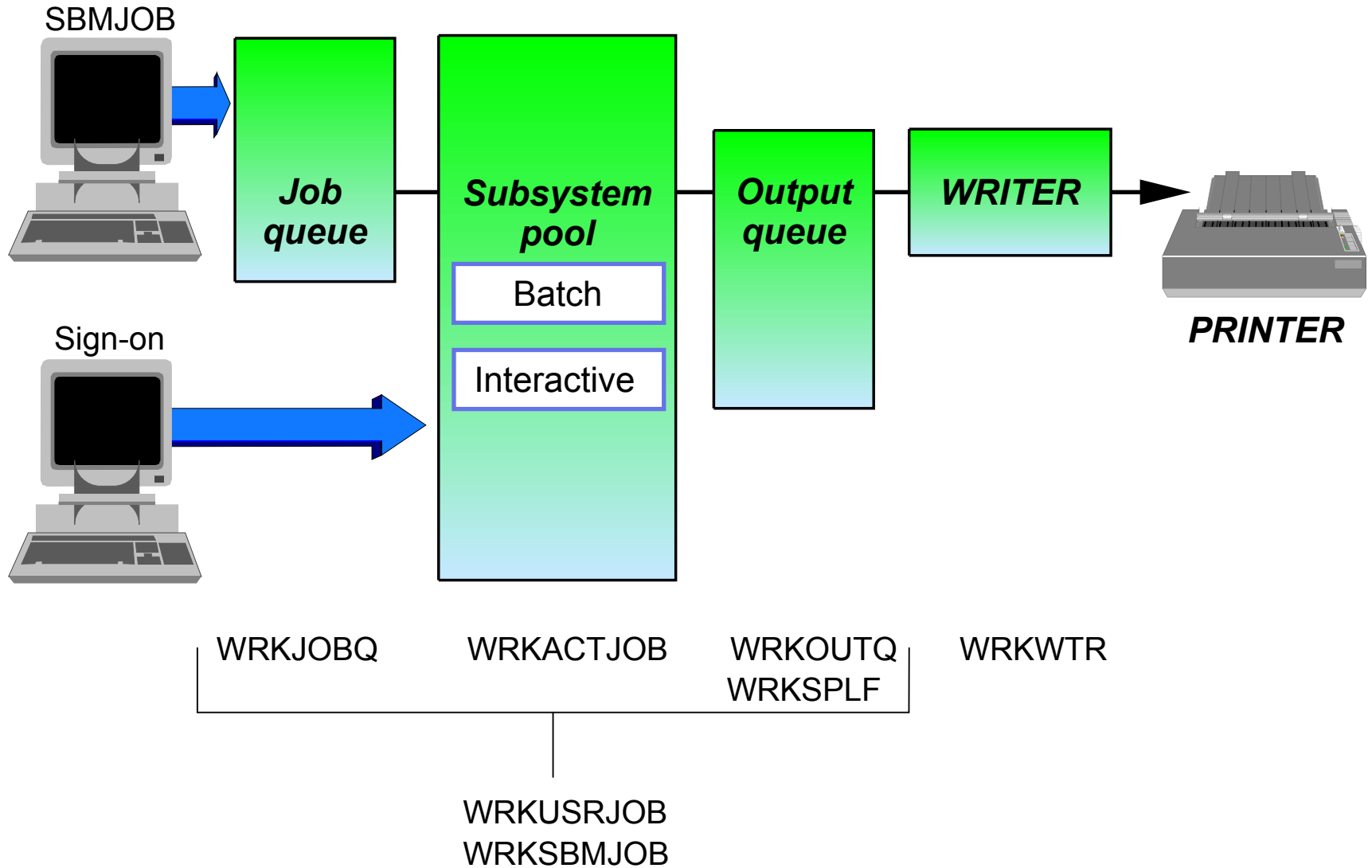
Printer device : *WRKSTN

More...

Press Enter to continue.

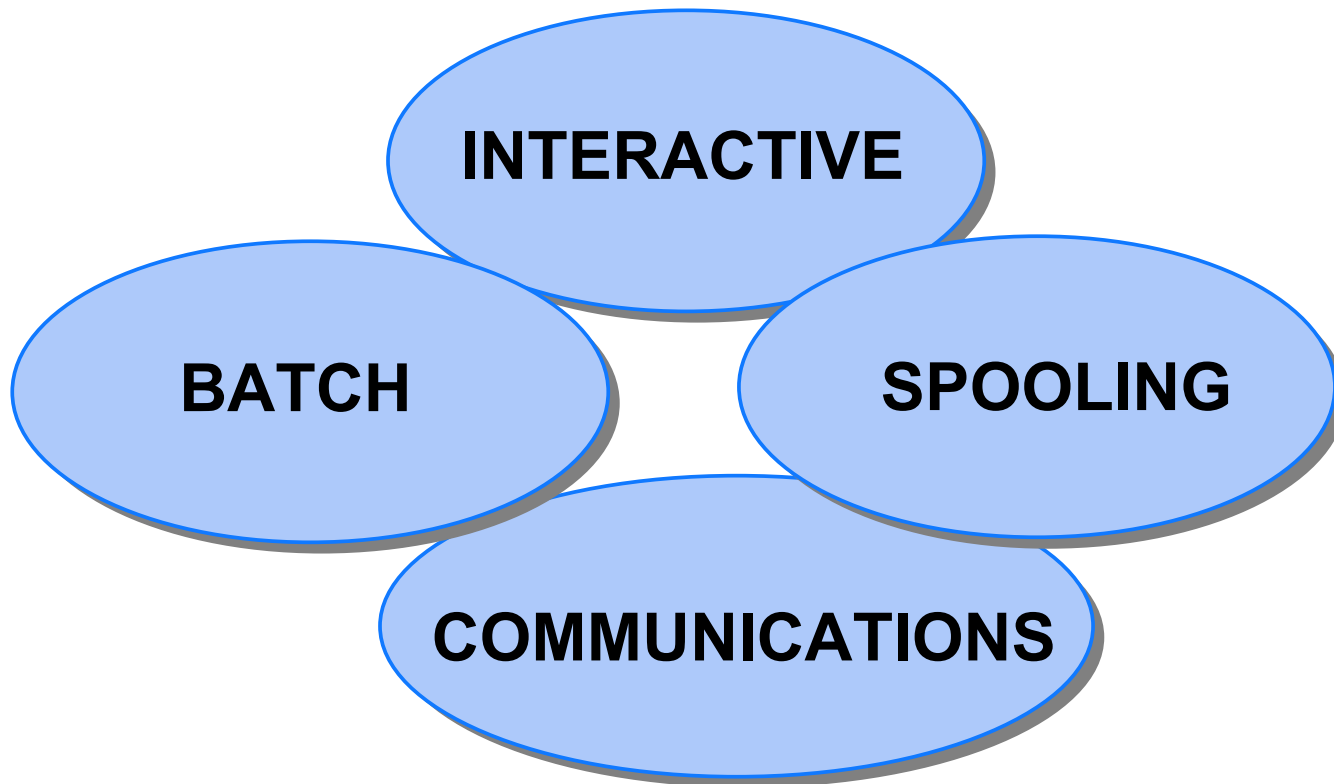
F3=Exit F12=Cancel

Job Processing



Subsystems

A specialized environment for handling a certain type of work or function, such as:

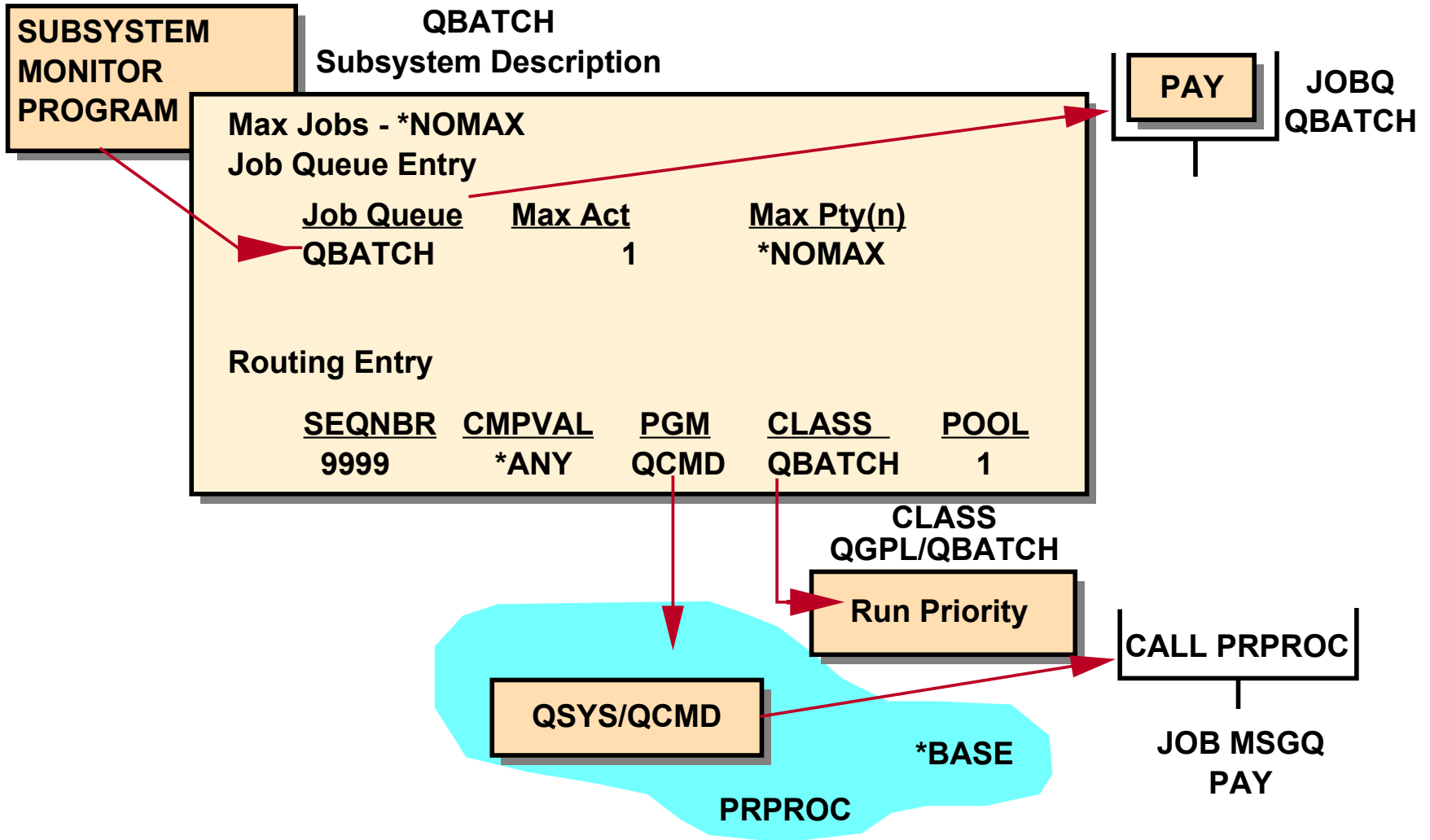


Storage Pools

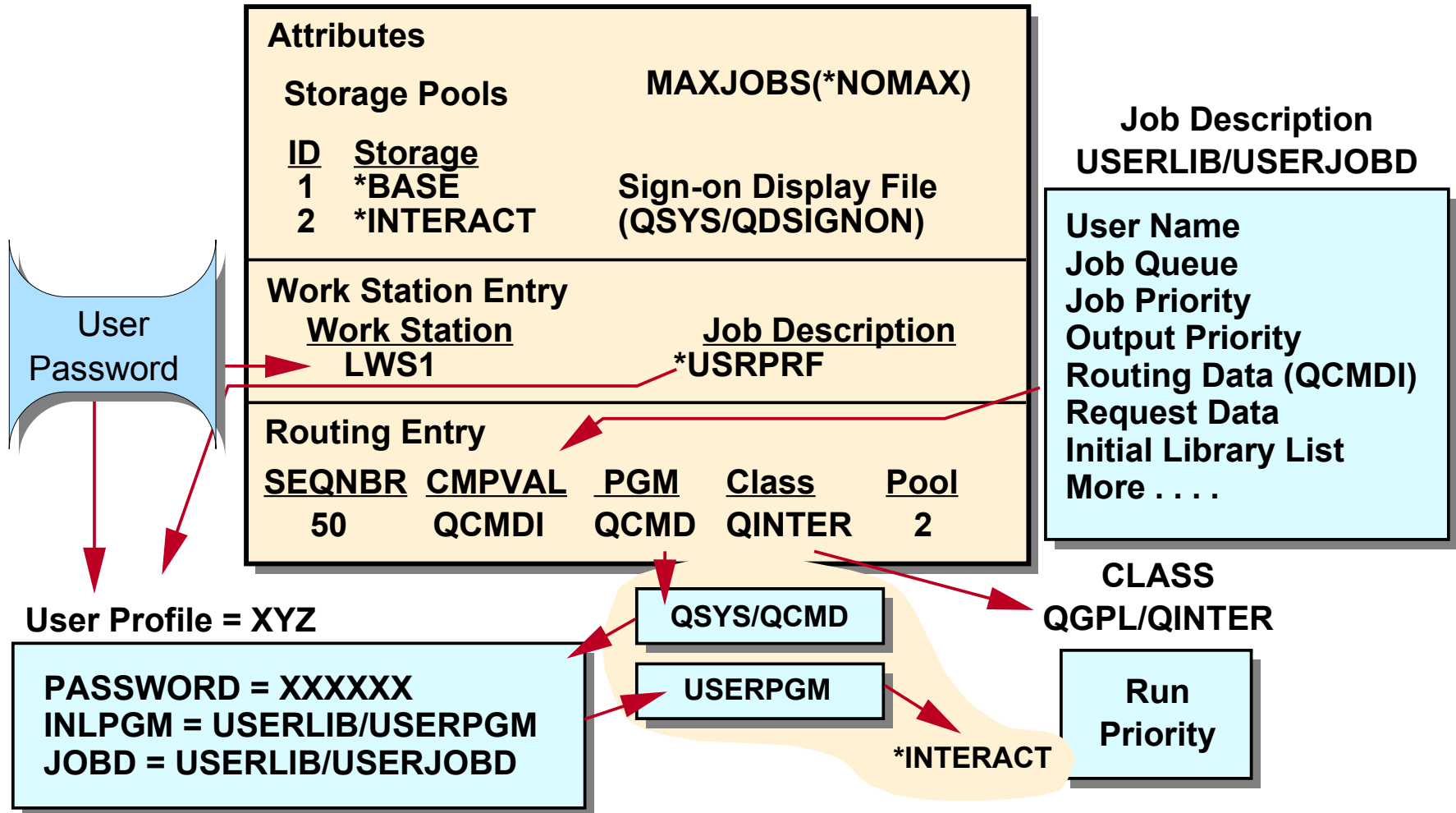
- Logical Areas of Main Storage for Jobs:
 - POOL SIZE sets amounts of storage
 - Noncontiguous pages
 - Deallocated when set to *NOSTG (No storage)
 - PAGING is restricted to pool
 - ACTIVITY LEVEL determines how many job threads may have main storage
 - Two types of pools: SHARED AND PRIVATE

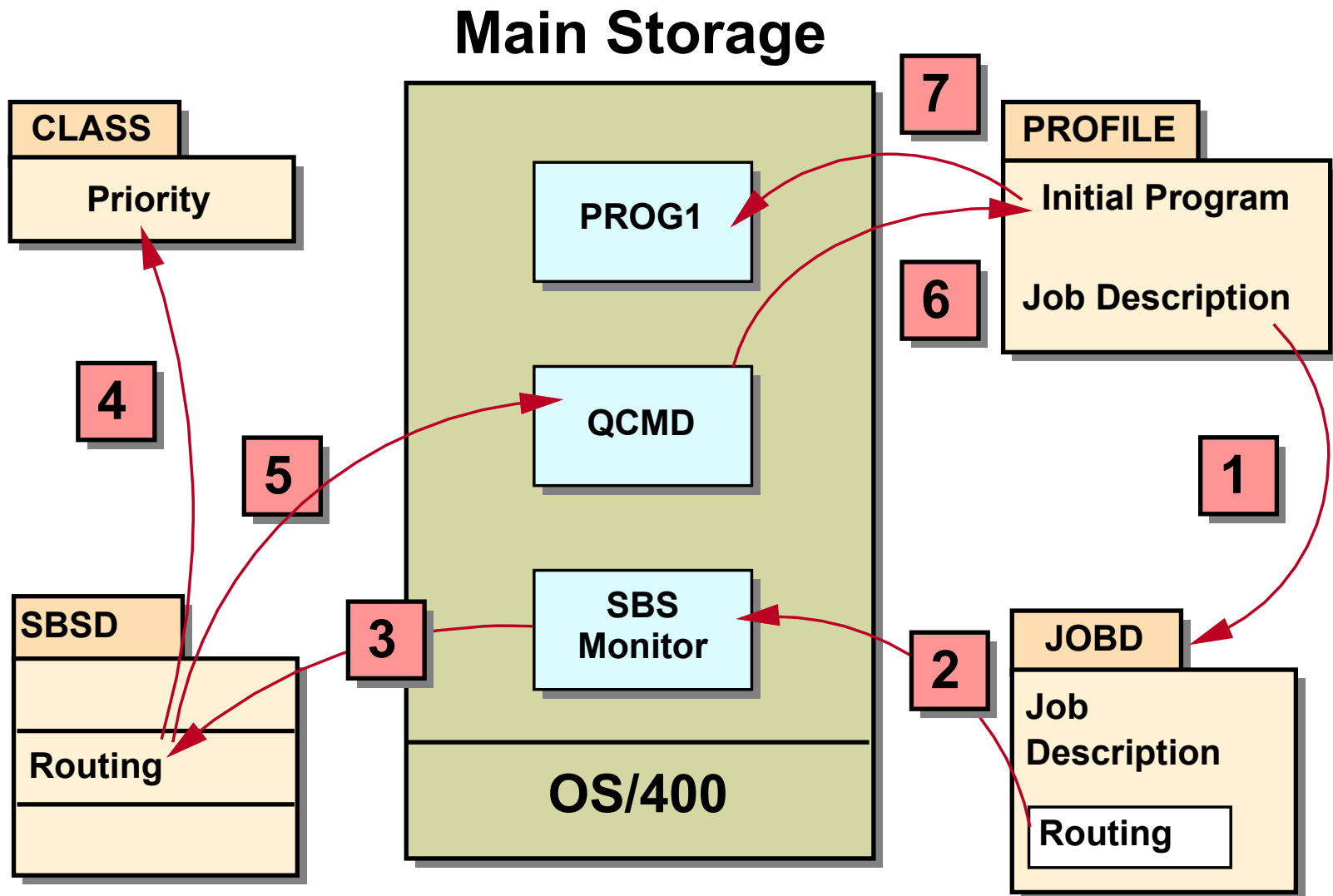
*SPOOL POOL	(4)
*INTERACT POOL	(3)
*BASE POOL	(2)
*MACHINE POOL	(1)

SBMJOB CMD(CALL PRPROC) RTGDTA(QCMDB) JOB(PAY) JOBQ(QBATCH)



QINTER Subsystem Description

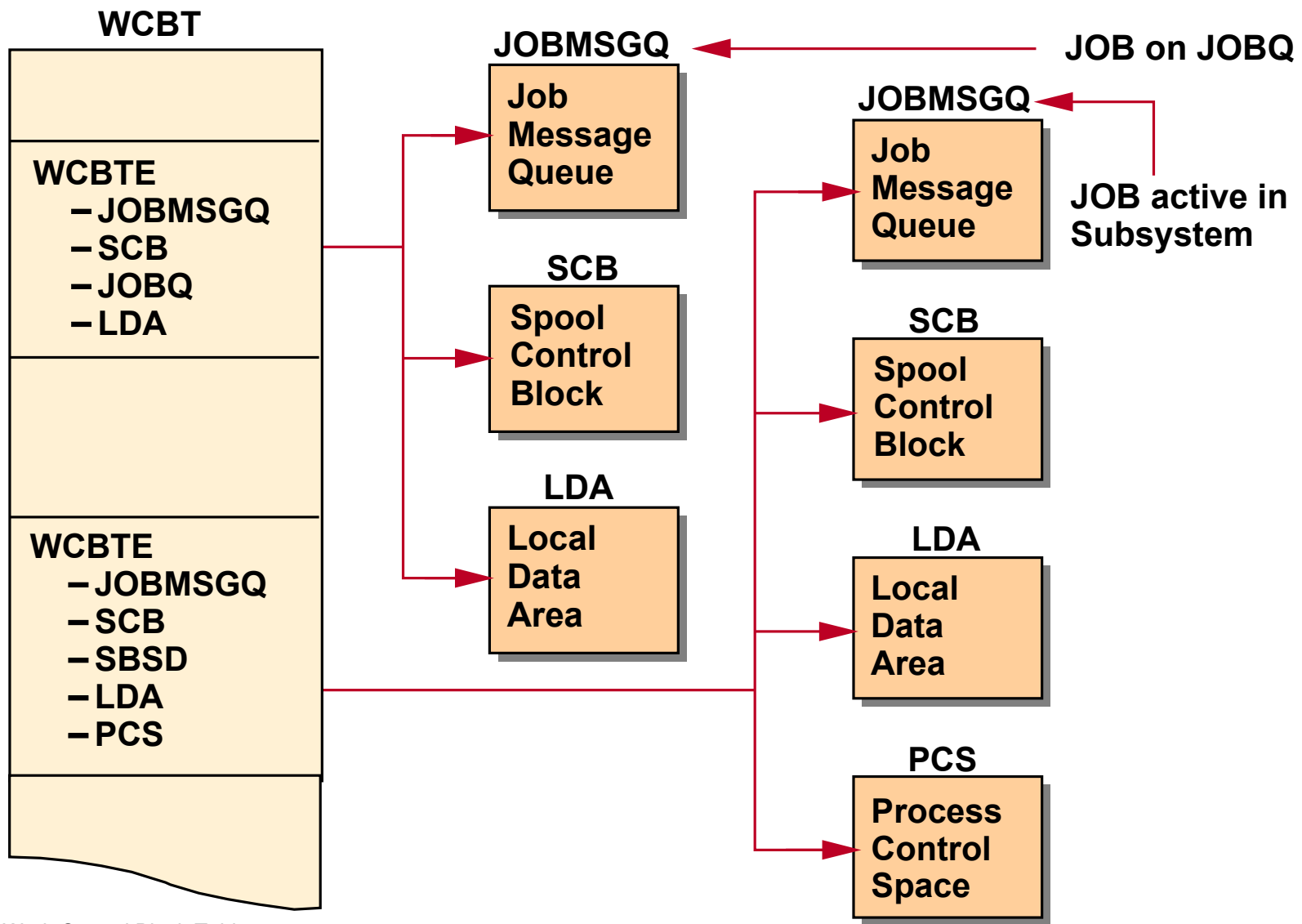




5.2 System's View

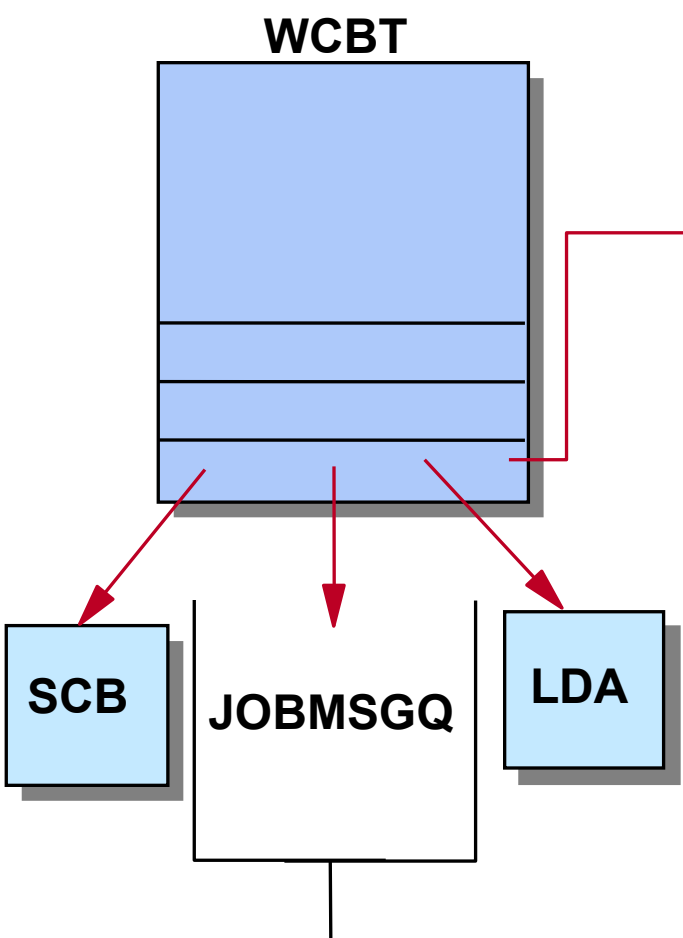
Job Structure and Task Dispatching

- **Process**
 - MI System object called process-control space (PCS)
 - Tie together the resources for program invocation
 - Consists of
 - A executable program (can be shared)
 - Program's data
 - State informations
 - Could also be defined as a program in execution
- **Job**
 - OS/400 object that contains a process structure to manage system resources
- **Thread**
 - A portion of a process
 - A process may have more than one thread
 - Each thread consists of
 - Set of registers representing the state of the processor
 - A control stack
 - Automatic program storage (local variables)
 - System resources allocated to the process are shared with other threads within the process

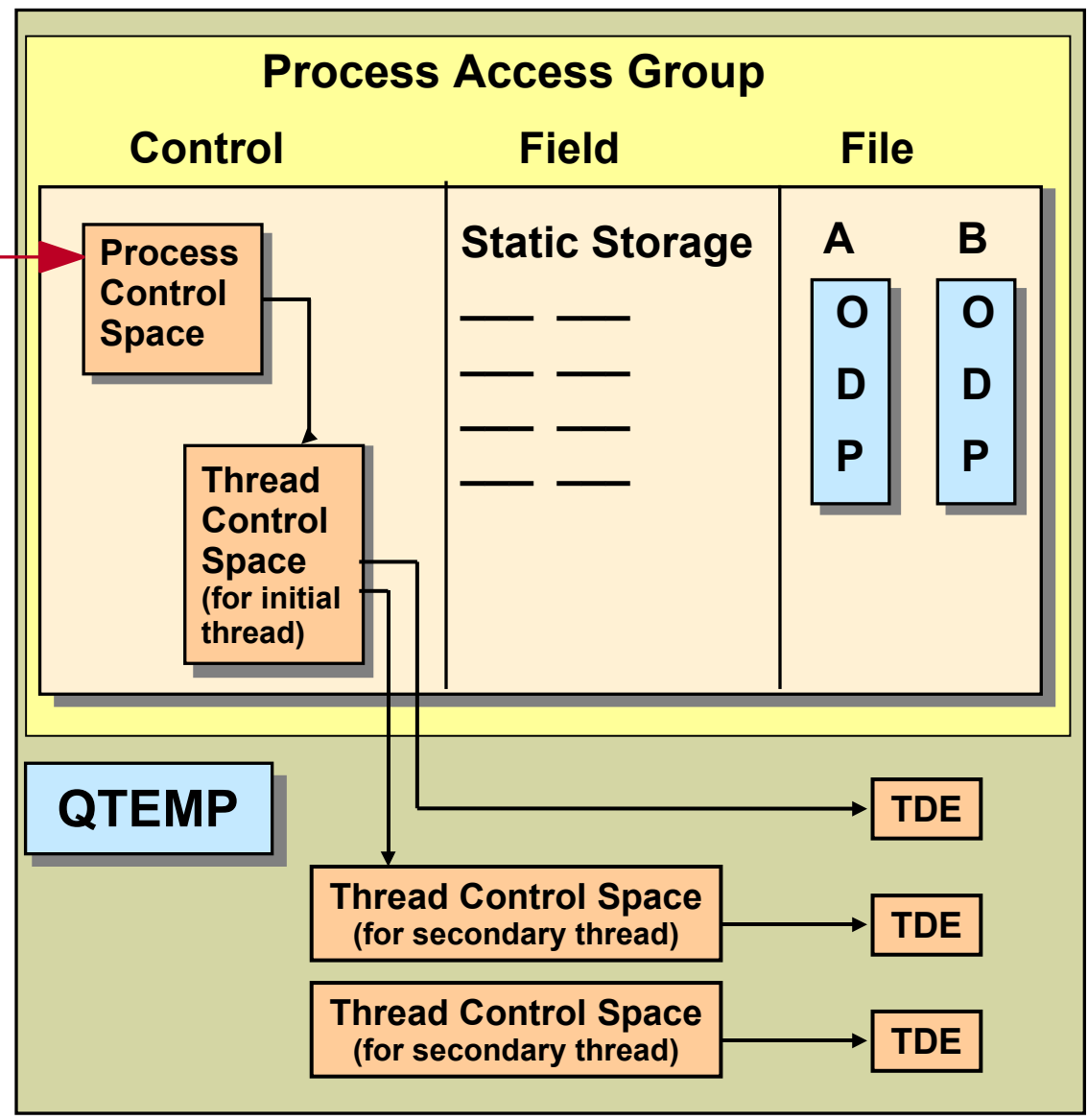


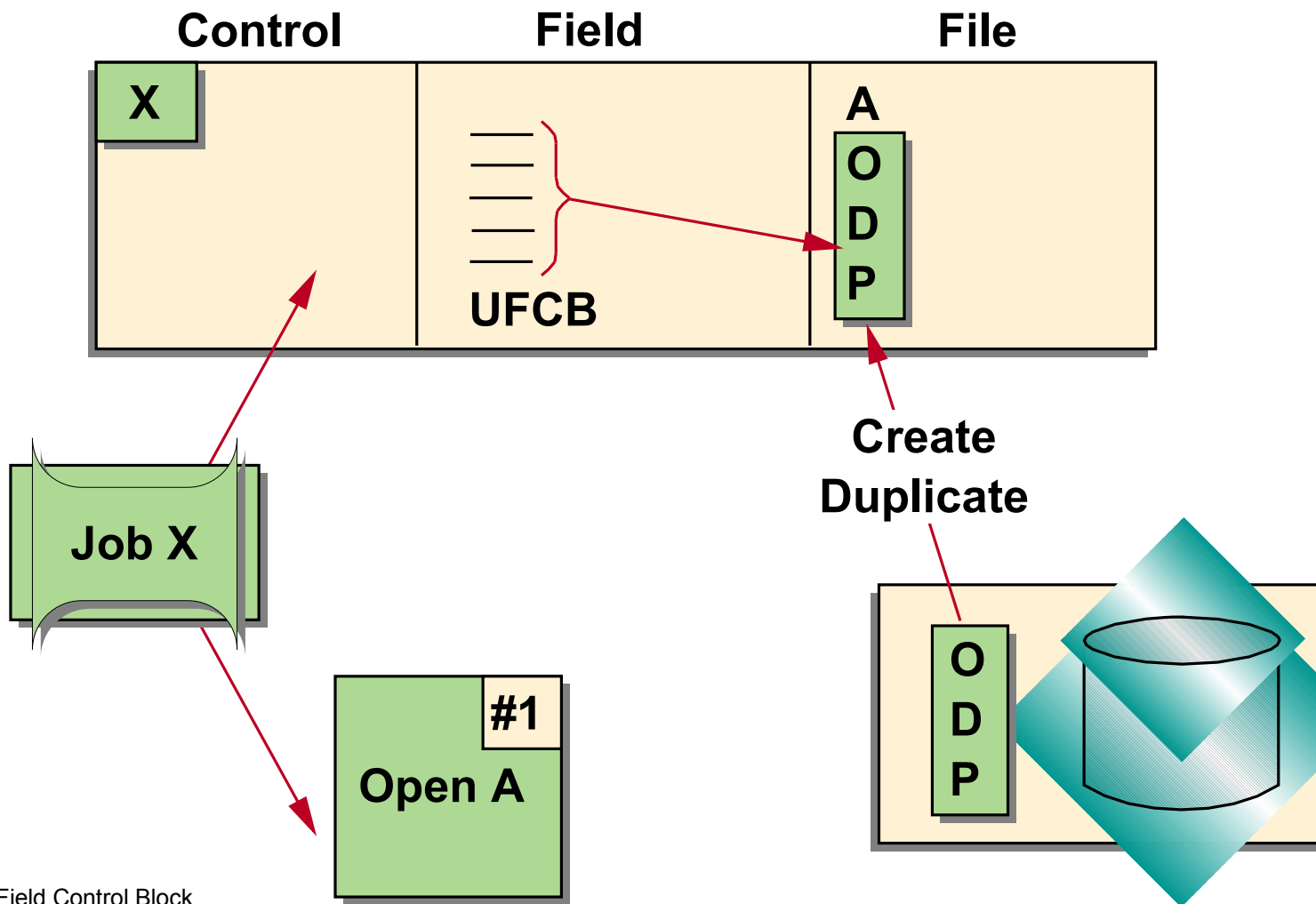
WCBT – Work Control Block Table
 WCBTE – Work Control block Table Entry

Temporary Job Structure Objects

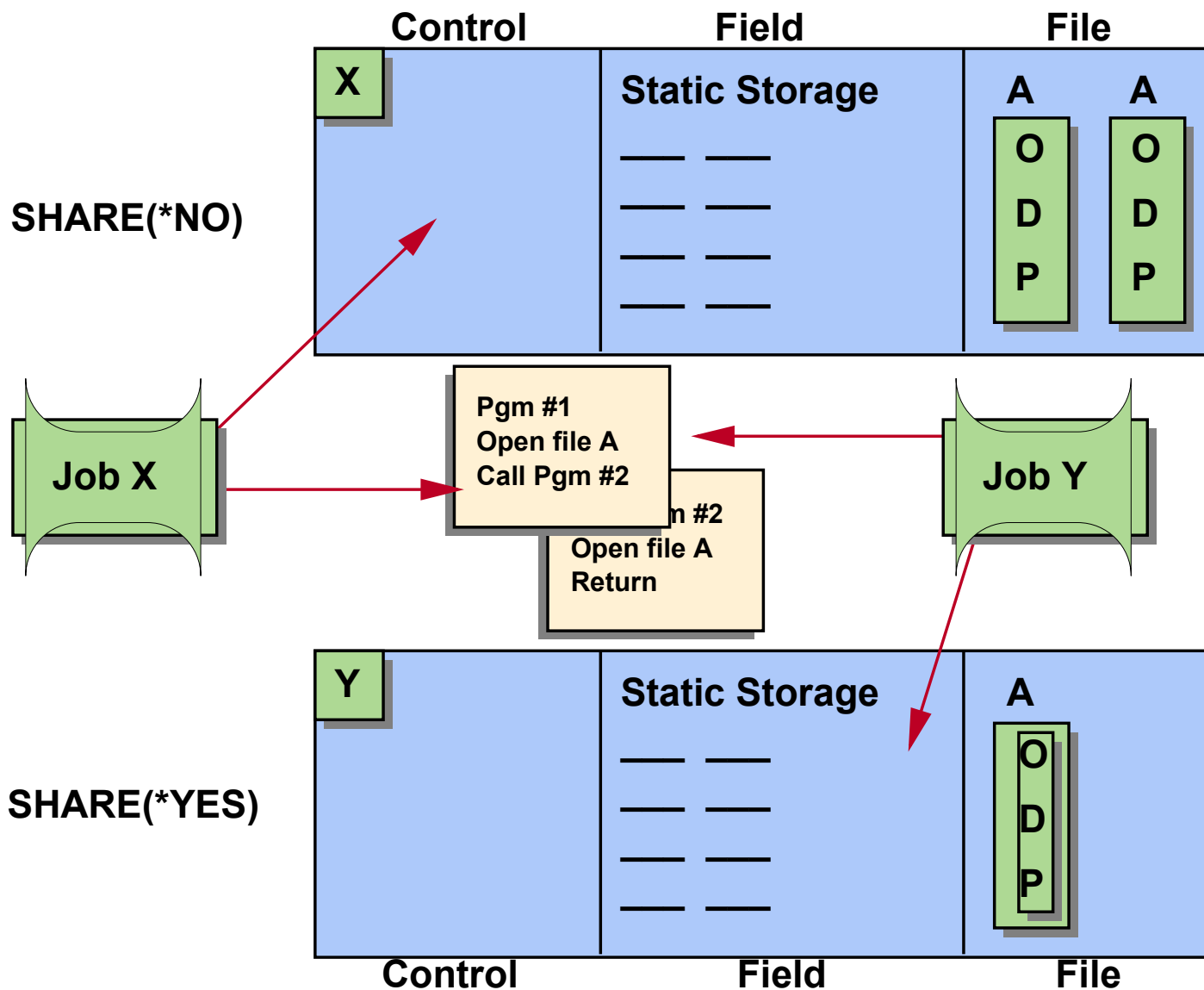


ODP – Open Data Path
TDE – Task Dispatching Element

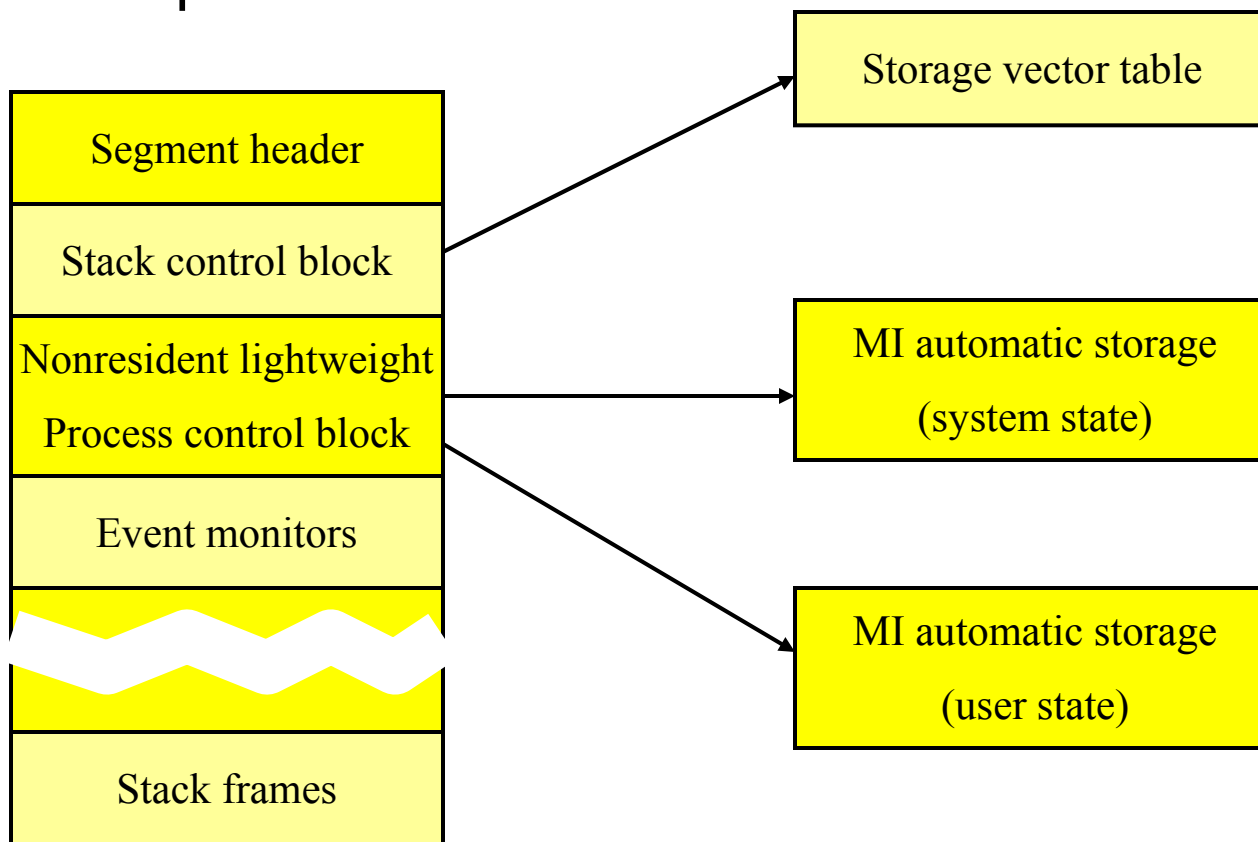




UFCB – User Field Control Block

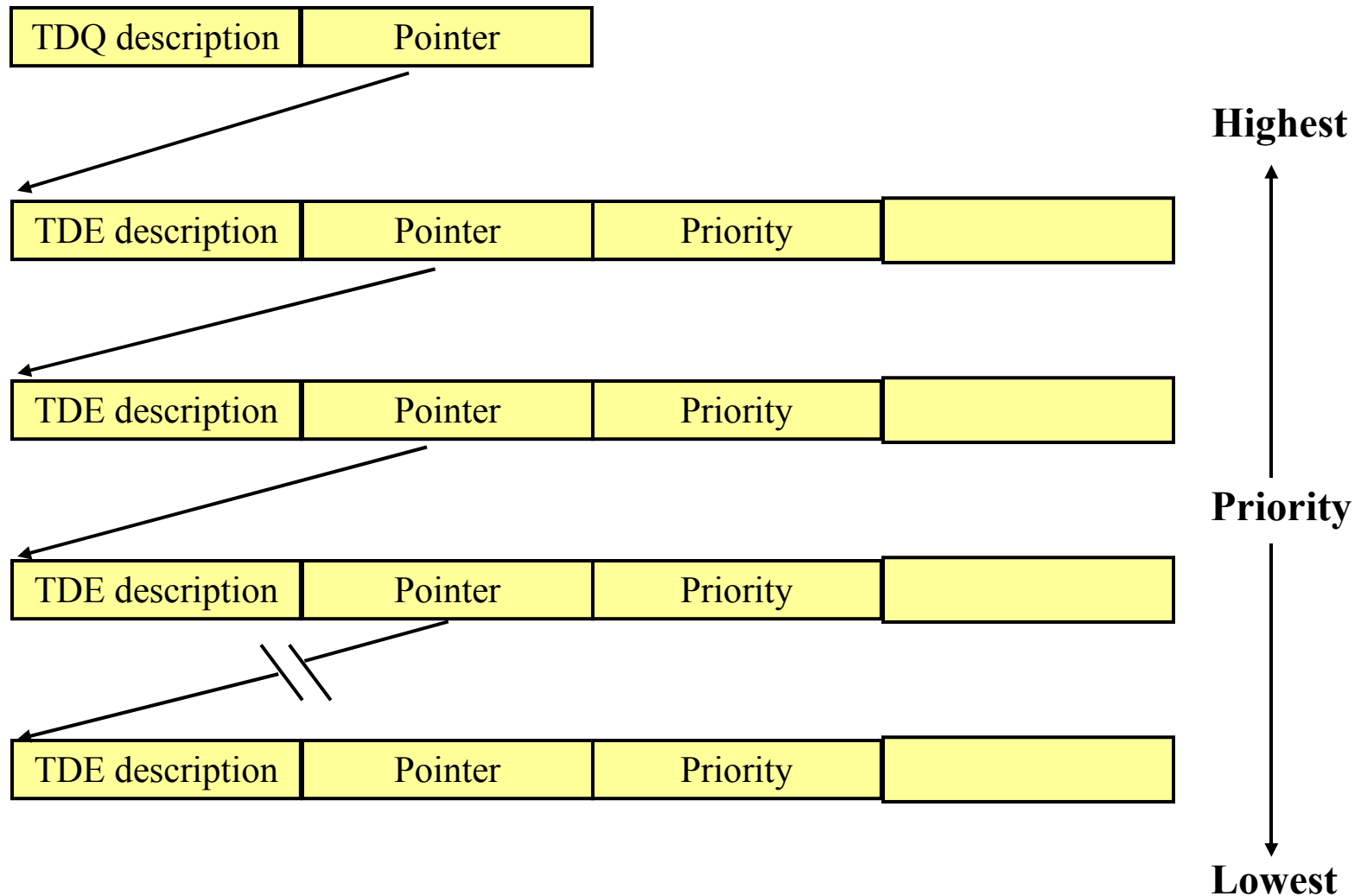


- are temporary SLIC objects
- will be recycled at process level or at systems level when process ends



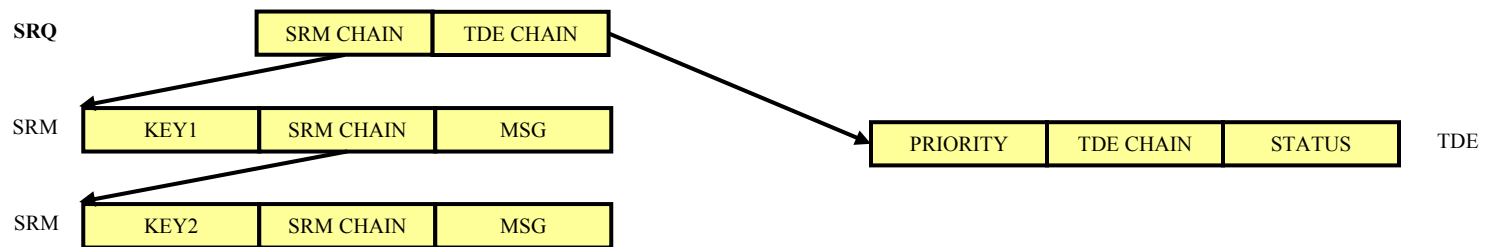
Source: Frank G. Soltis, Fortress Rochester, Fig. 13.5

Task Dispatching Queues

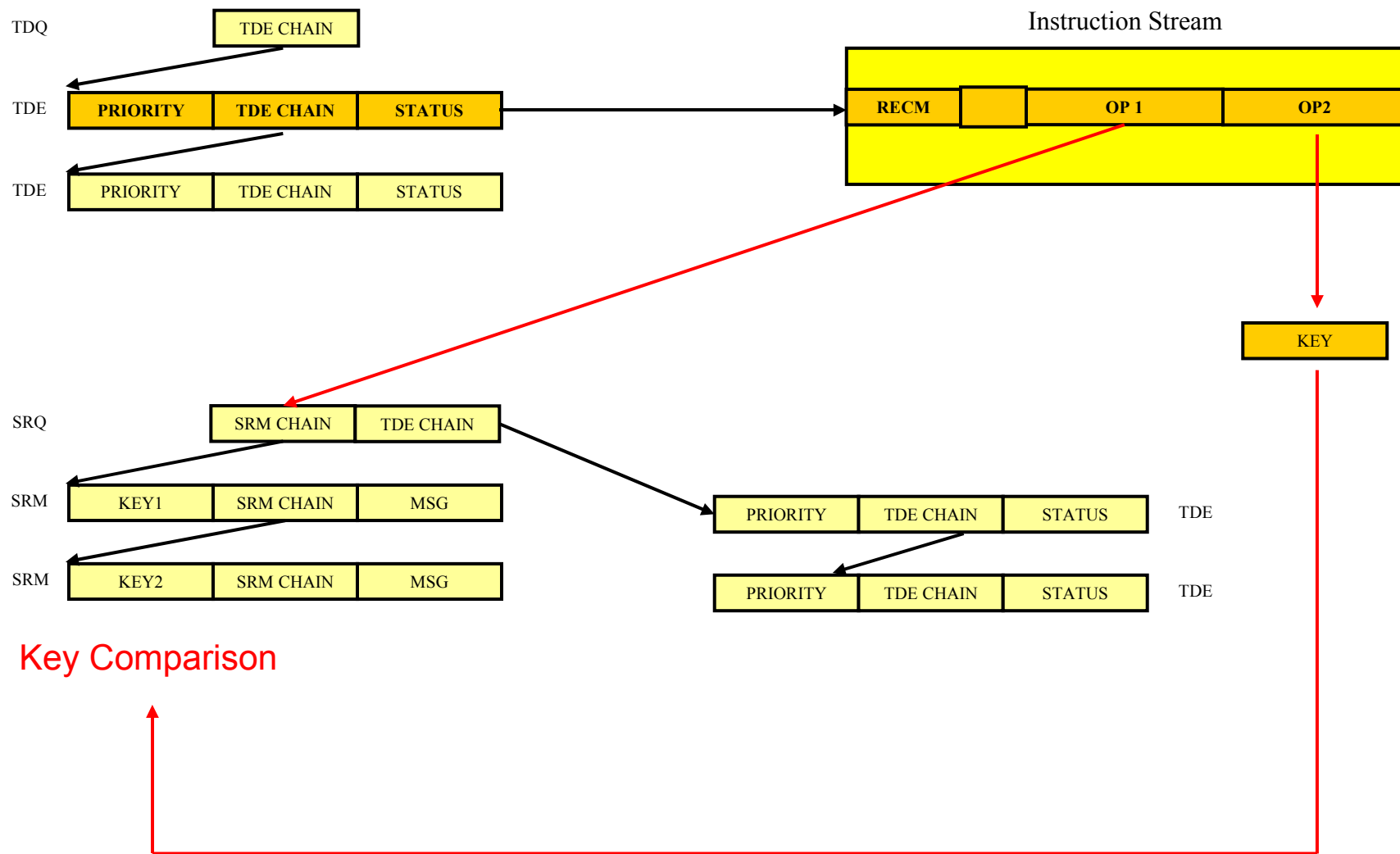


Send/Receive Queues

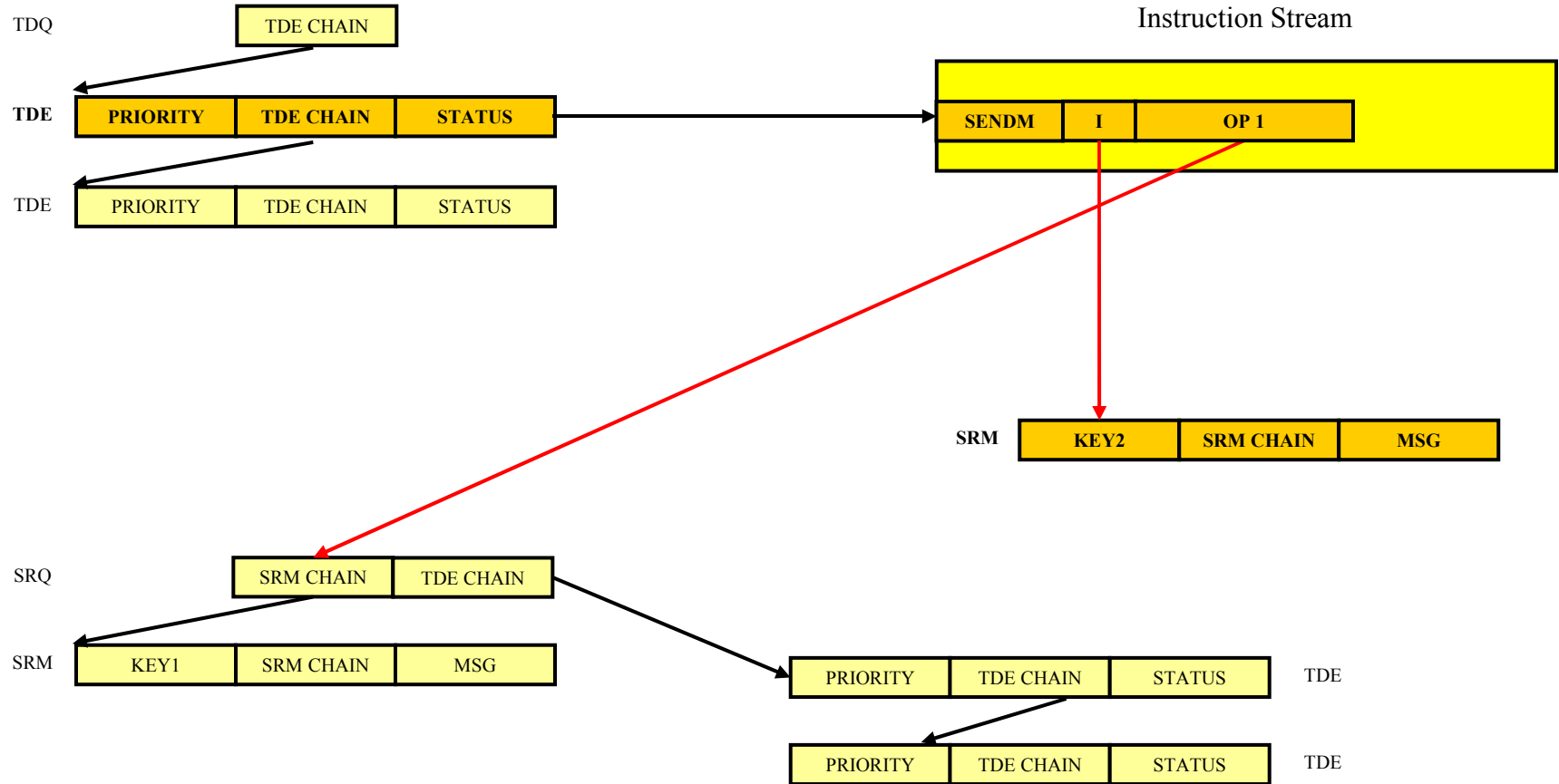
- Data structure used as “mailbox” for tasks
- SEND MESSAGE enqueued a SRM
- Execute a RECEIVE MESSAGE operations to obtain a message (dequeing)
 - If no message is available the task has to wait

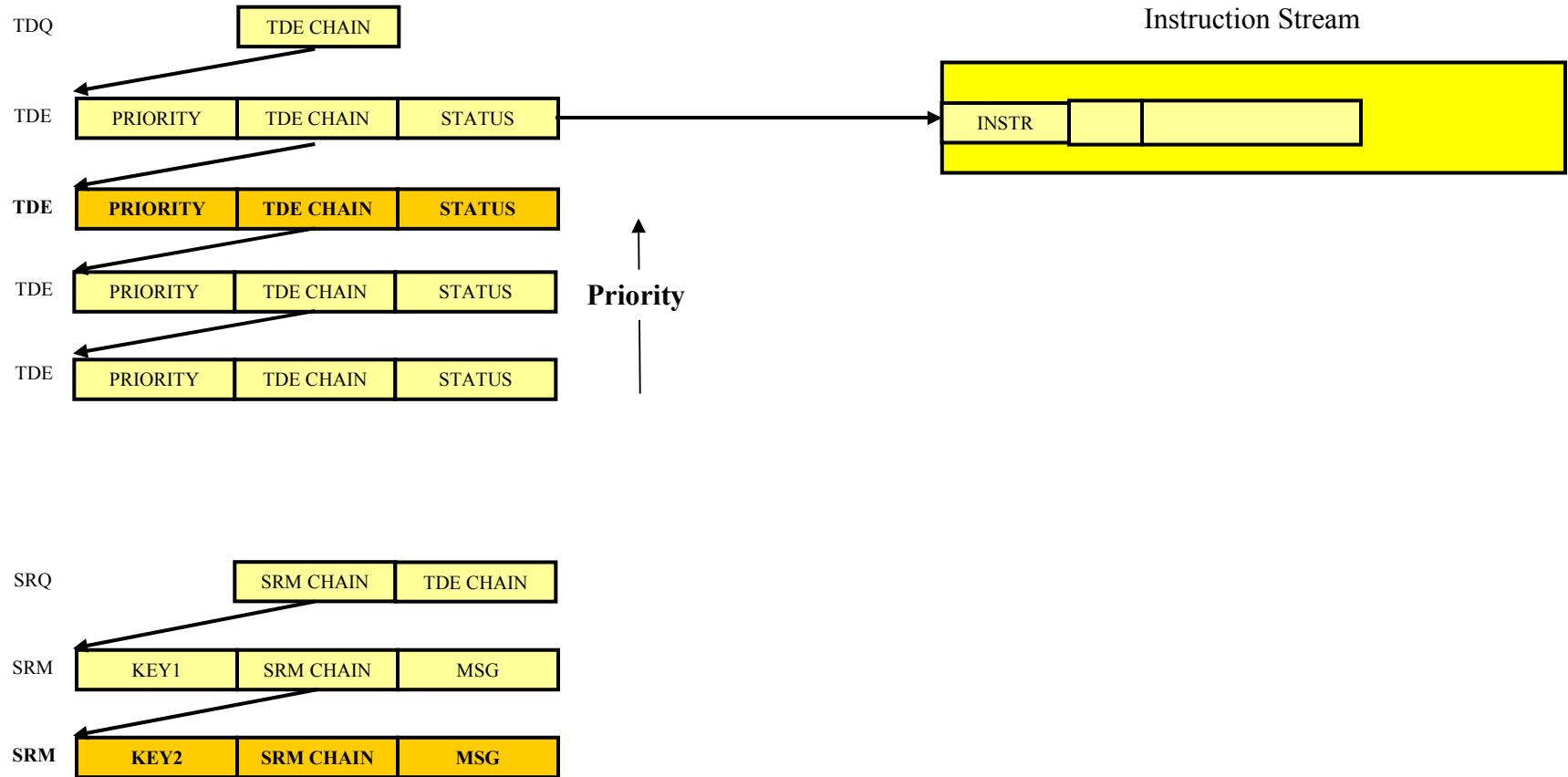


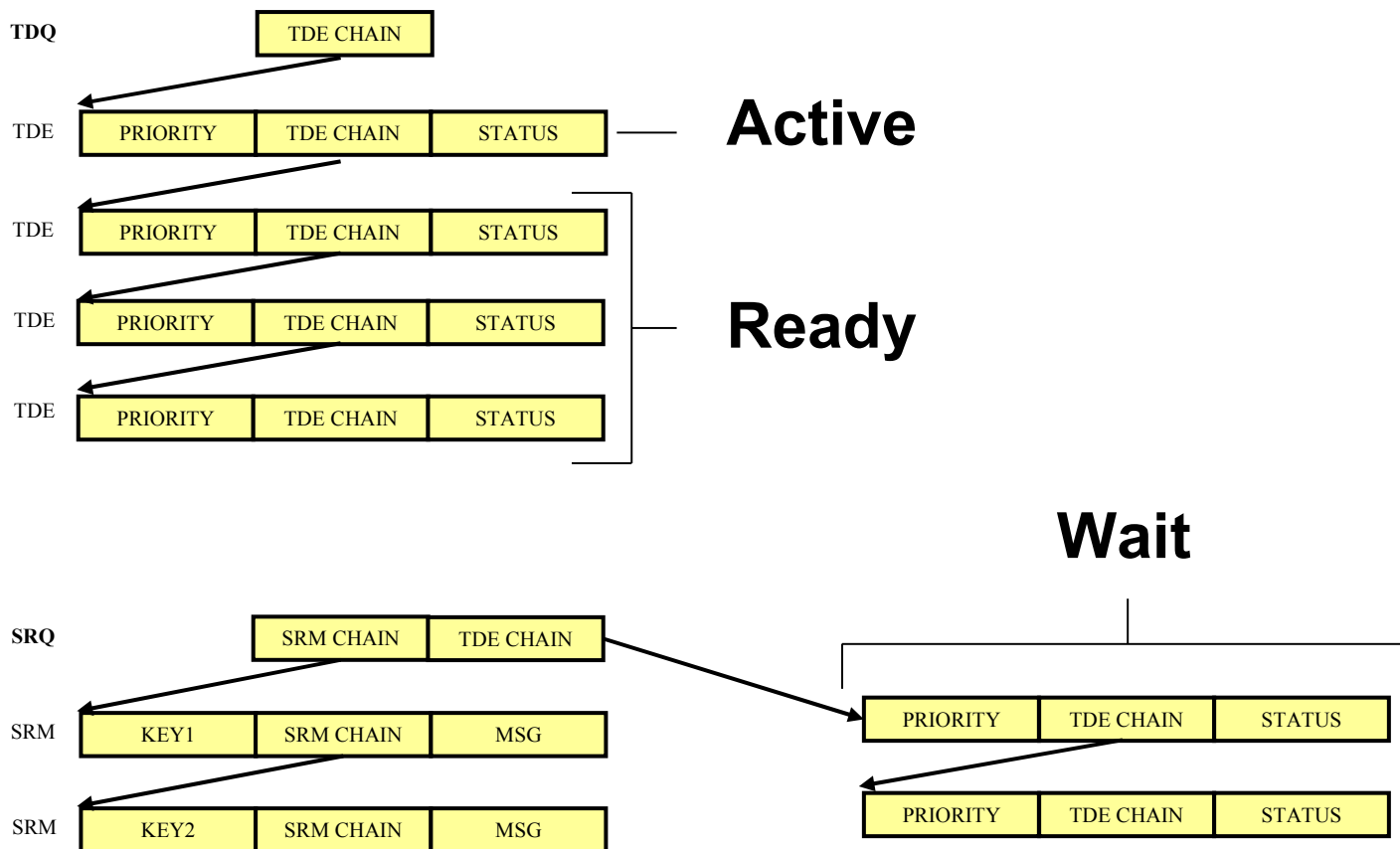
SRQ – Send/Receive Queue
 SRM – Send/Receive Message
 RECM – Recieve Message
 SENDM – Send Message



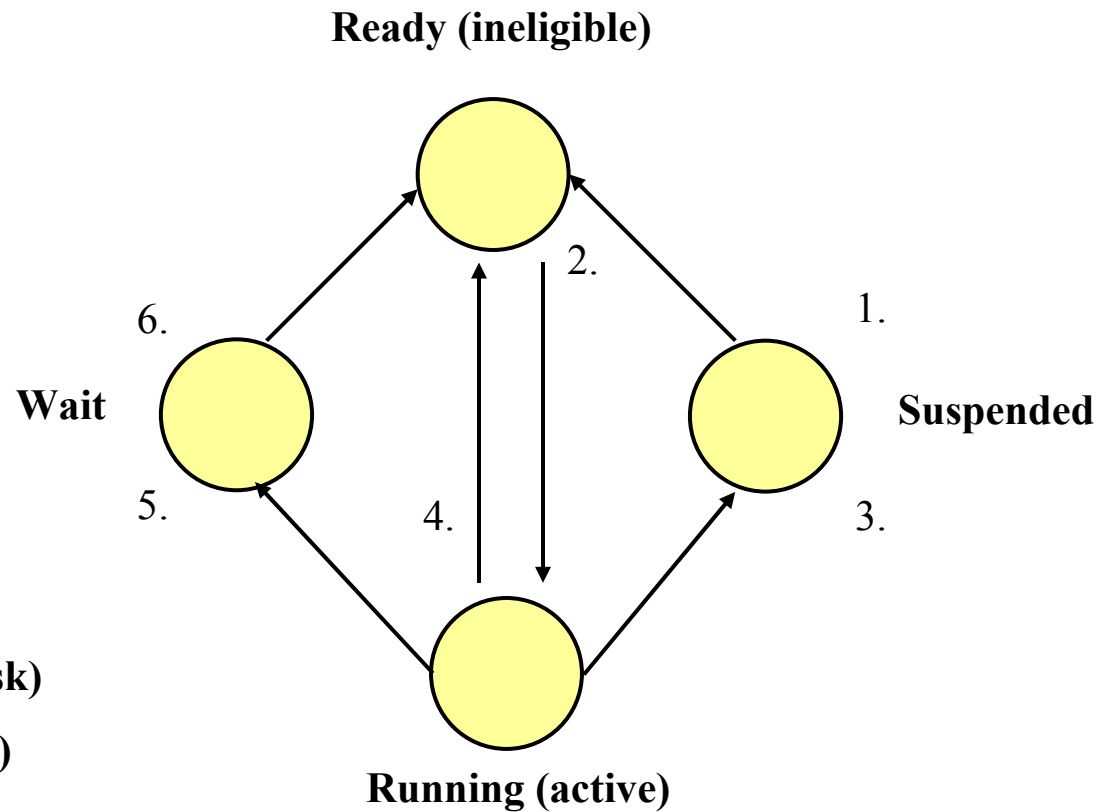
SENDM (1/2)





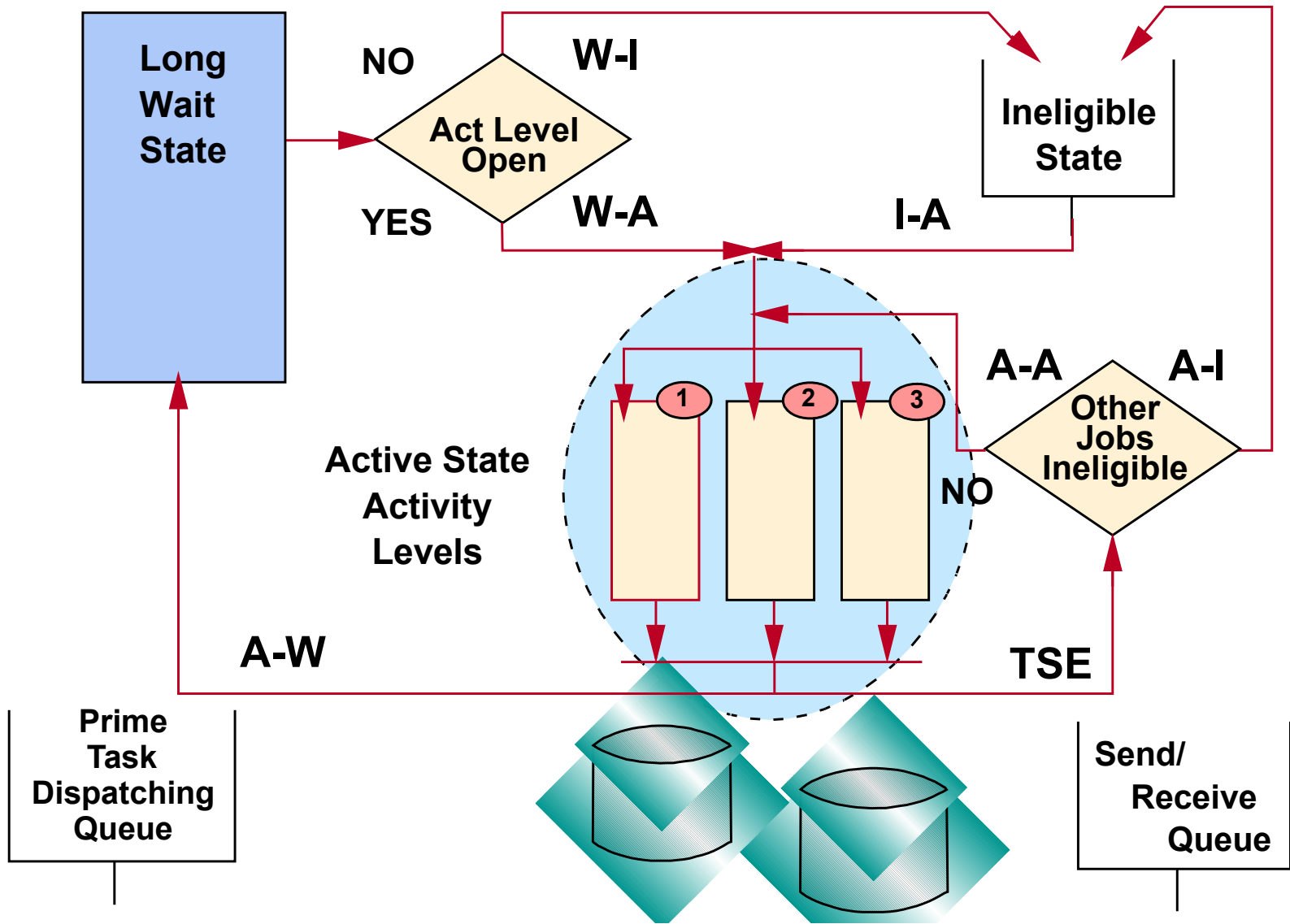


1. **Initiate task (start work)**
2. **Run task (dispatch task)**
3. **Suspend task (work done)**



1. **Preempt task (by another task)**
2. **Wait(for I/O or another task)**
3. **Signal (task is done waiting)**

Thread States and Transitions



Observing Thread State Transitions

```

Work with System Status
                                06/21/XX
                                10:11:12
% CPU used . . . . . :          61.7
% DB capability . . . . :
Elapsed time . . . . . :      00:06:30
Jobs in system . . . . . :       494
% perm addresses . . . . :
% temp addresses . . . . :
    Permanent . . . . . :       3.063
    Temporary . . . . . :       3.962
                                Auxiliary storage:
                                System ASP . . . . . :      3258 M
                                % system ASP . . . . . :     18.6657
                                Total . . . . . :           3258 M
                                Current unprotect used . . . . :    437 M
                                Maximum unprotect . . . . . :    437 M

```

Type changes (if allowed), press Enter.

System	Pool	Reserved	Max
Pool	Size (M)	Size (M)	Active
1	3.61	2.26	+++
2	3.97	0	10
3	.23	0	3
4	8.19	0	25

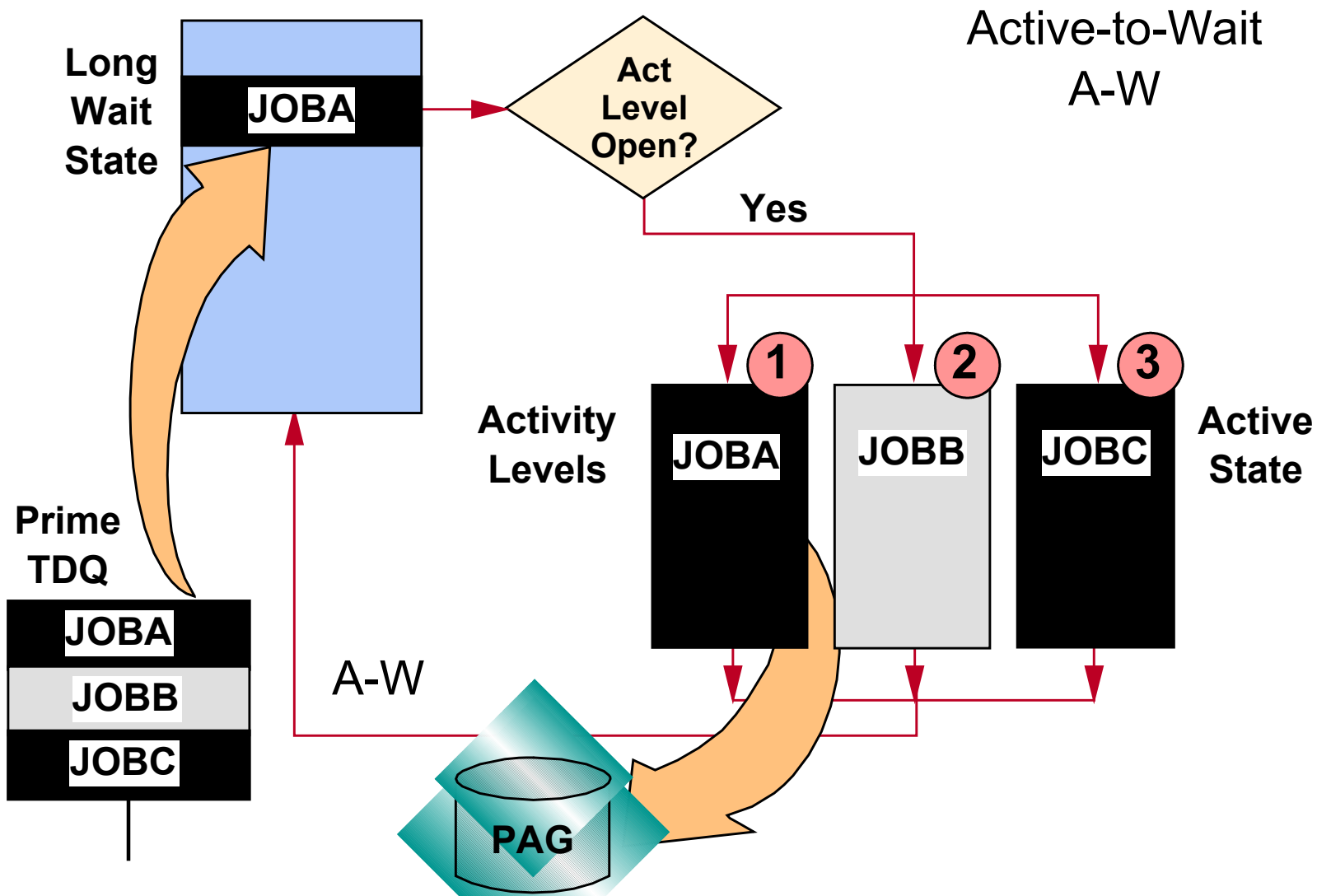
Active -->	Wait -->	Active -->
Wait	Inel	Inel
7.5	.0	.0
.4	.0	.0
.0	.0	.0
88.4	.0	.0

Bottom

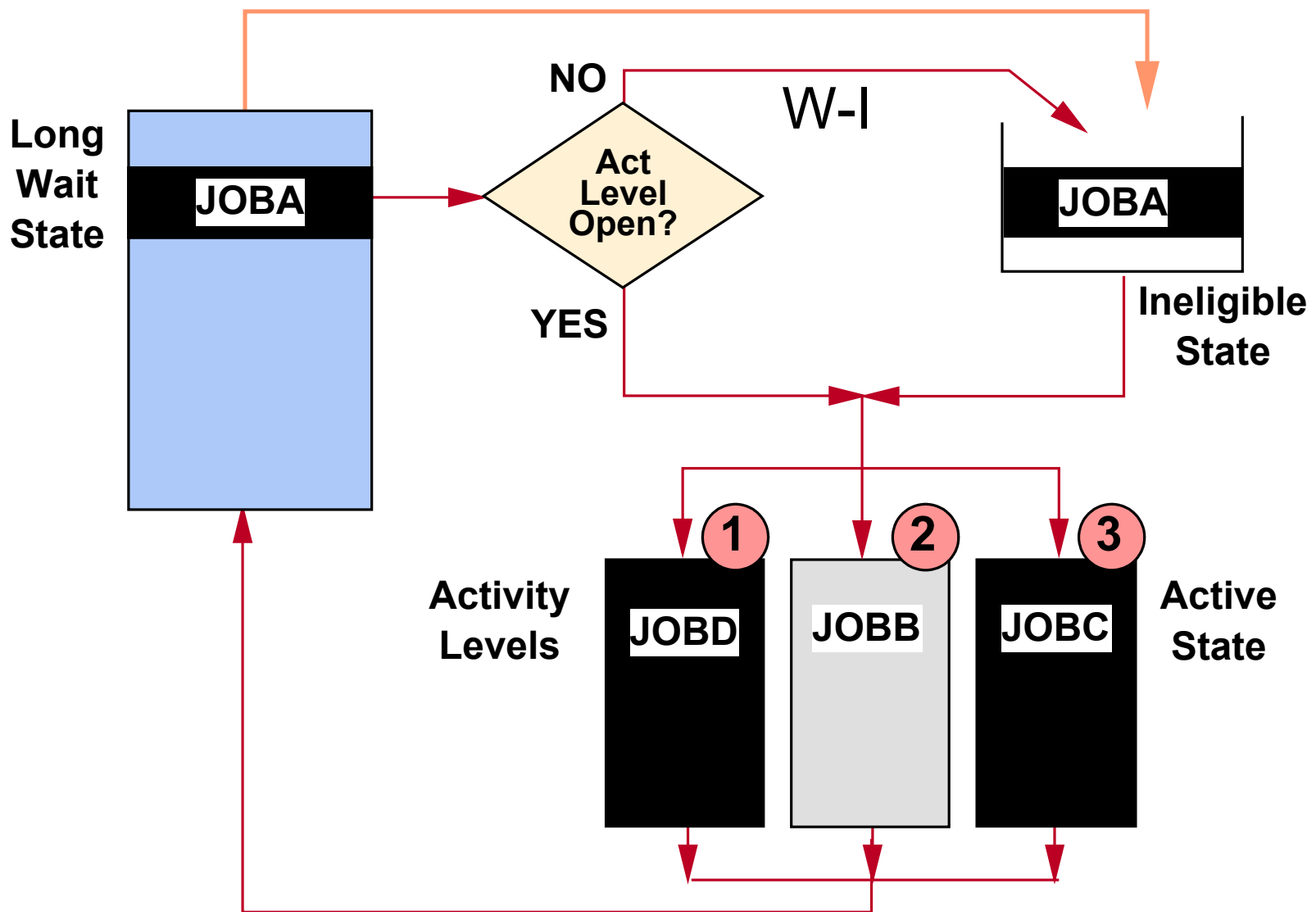
Command

===>

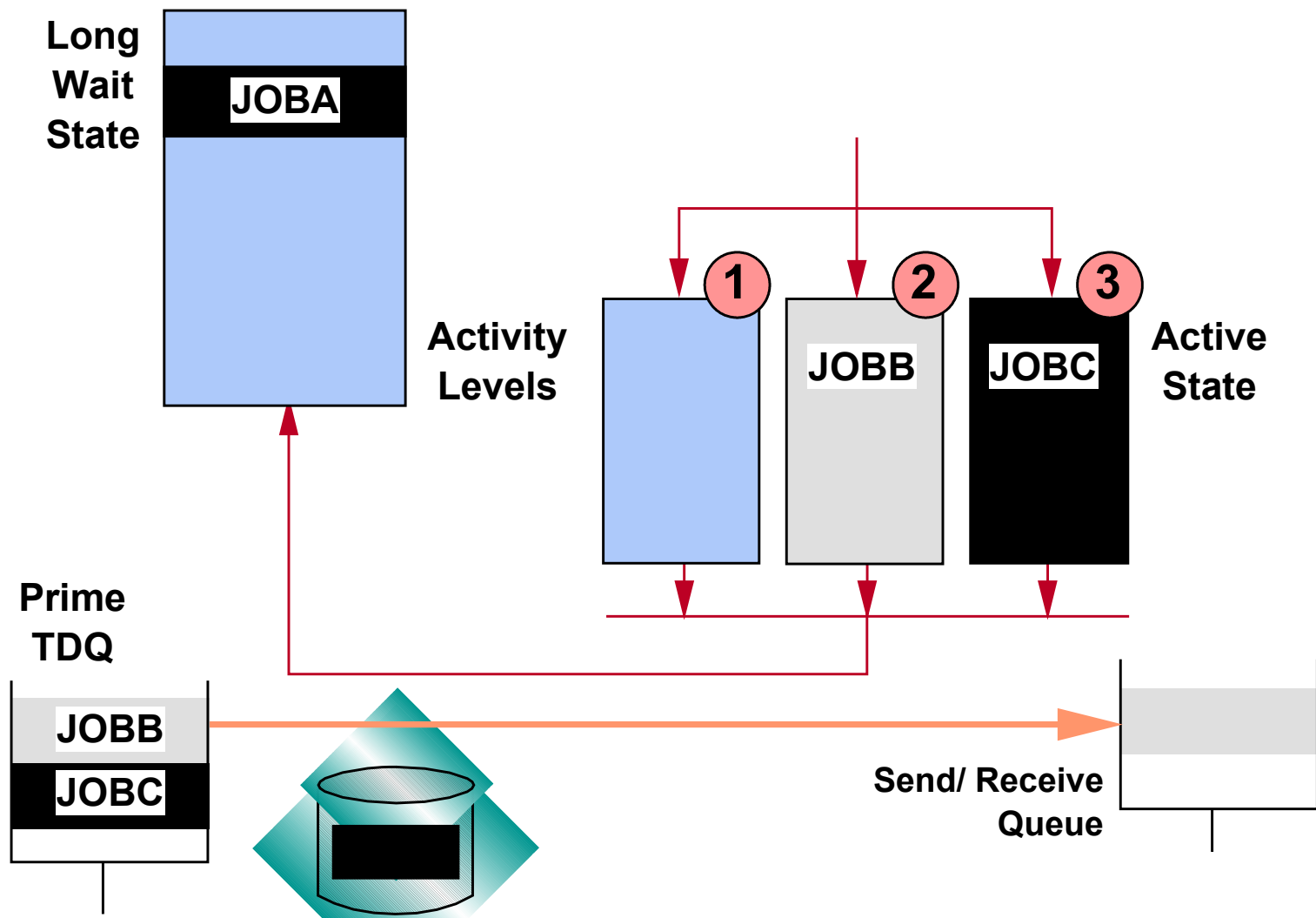
4=Prompt F5=Refresh F9=Retrieve F10=Restart
F11=Display pool data F12=Cancel F14= Work with subsystems F24=More Keys



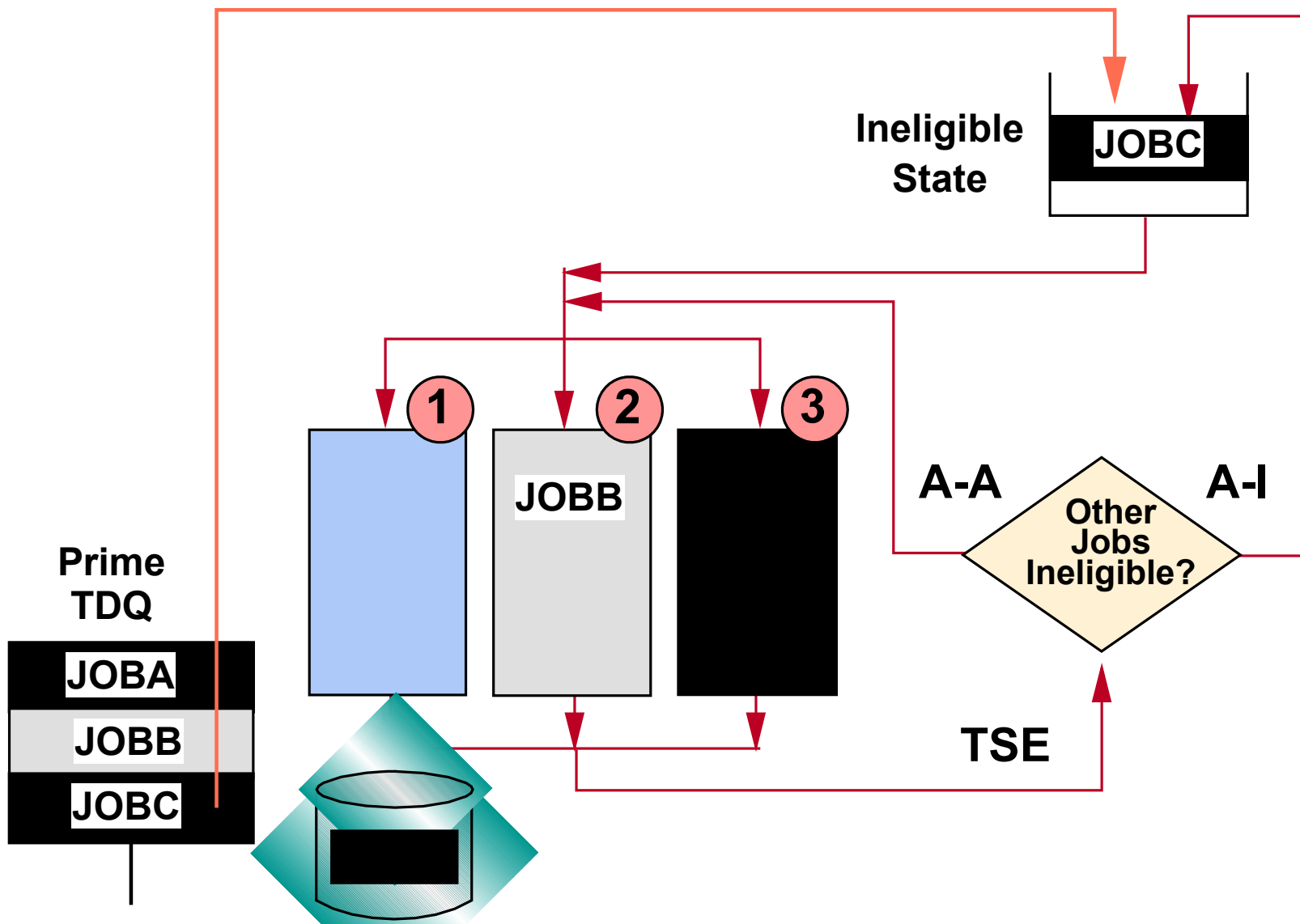
Wait-to-Ineligible



Wait in Activity Level



Active-to-Ineligible



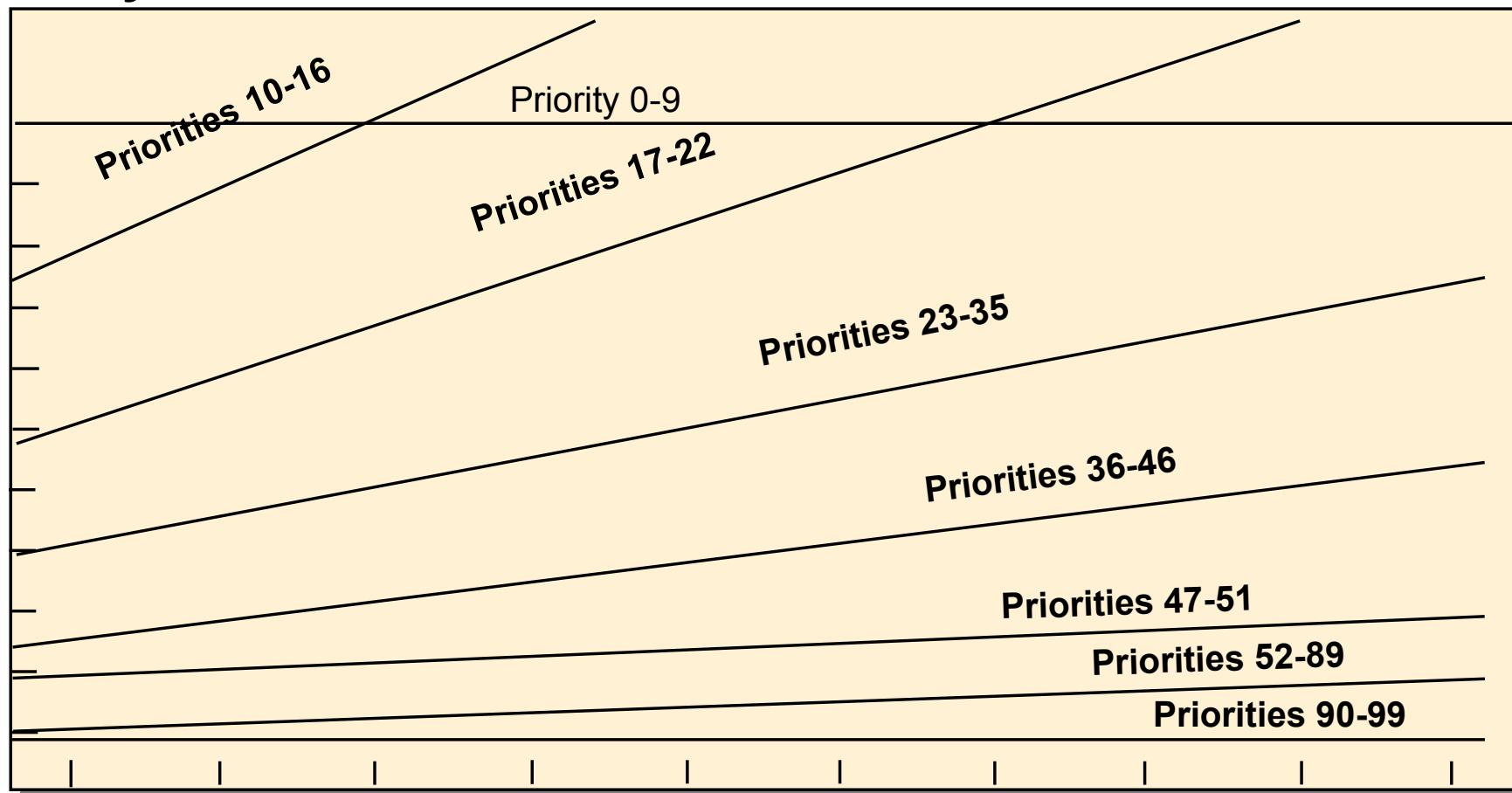
New algorithms on RISC systems

- No job can monopolize the CPU
 - Jobs not using CPU have their DELAY COST increased
- Low-priority jobs have a chance to progress
- iSeries server optimization for batch
- Interactive job priority lowered when batch demand is high

Can be turned on/off

- QDYNPTYSCD and QDYNPTYADJ

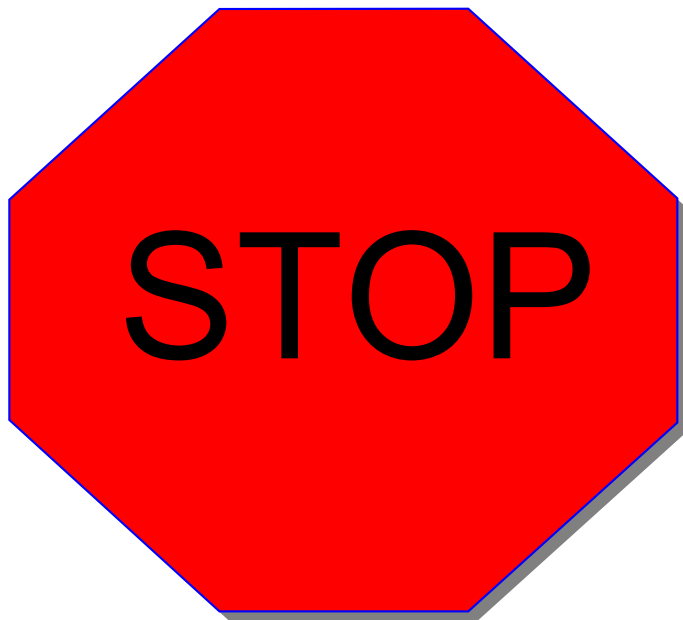
Delay Cost



Time (Millions of cycles)

Multi-Processor Considerations (1/2)

Dispatching the first n top TDEs on the TDQ?



Cache Affinity
and
Performance?!

Multi-Processor Considerations (2/2)

Task Dispatching is based on priority, cache affinity and eligibility

Three bits in a TDE

- Eligibility bit
 - one bit for every processor
 - indicates that a task is eligible to run on on corresponding processor
- Active field
 - one bit for every processor
 - indicates on which processor the task is currently running
- Affinity field
 - one bit for every processor
 - Indicates on which processor the task was currently running