## Linear programming

## Optimization Problems

- General optimization problem

$$
\max \left\{z(x) \mid f_{j}(x) \leq 0, x \in D\right\} \text { or } \min \left\{z(x) \mid f_{j}(x) \leq 0, x \in D\right\}
$$

where $D \subseteq \mathbb{R}^{n}, f_{j}: D \rightarrow \mathbb{R}$, for $j=1, \ldots, m, z: D \rightarrow \mathbb{R}$.

- Linear optimization problem

$$
\max \left\{c^{T} x \mid A x \sum b, x \in \mathbb{R}^{n}\right\}, \text { with } c \in \mathbb{R}^{n}, A \in \mathbb{R}^{m \times n}, b \in \mathbb{R}^{m}
$$

- Integer optimization problem: $x \in \mathbb{Z}^{n}$
- 0-1 optimization problem: $x \in\{0,1\}^{n}$


## Feasible and optimal solutions

- Consider the optimization problem

$$
\max \left\{z(x) \mid f_{j}(x) \leq 0, x \in D, j=1, \ldots, m\right\}
$$

- A feasible solution is a vector $x^{*} \in D \subseteq \mathbb{R}^{n}$ such that $f_{j}\left(x^{*}\right) \leq 0$, for all $j=1, \ldots, m$.
- The set of all feasible solutions is called the feasible region.
- An optimal solution is a feasible solution such that

$$
z\left(x^{*}\right)=\max \left\{z(x) \mid f_{j}(x) \leq 0, x \in D, j=1, \ldots, m\right\}
$$

- Feasible or optimal solutions
- need not exist,
- need not be unique.


## Transformations

- $\min \{z(x) \mid x \in S\}=-\max \{-z(x) \mid x \in S\}$.
- $f(x) \geq a$ if and only if $-f(x) \leq-a$.
- $f(x)=a$ if and only if $f(x) \leq a \wedge-f(x) \leq-a$.


## Lemma

Any linear programming problem can be brought to the form

$$
\max \left\{c^{\top} x \mid A x \leq b\right\} \text { or } \max \left\{c^{T} x \mid A x=b, x \geq 0\right\}
$$

Proof: a) $a^{T} x \leq \beta \rightsquigarrow a^{T} x+x^{\prime}=\beta, x^{\prime} \geq 0$ (slack variable)
b) $x$ free $\rightsquigarrow x=x^{+}-x^{-}, x^{+}, x^{-} \geq 0$.

## Practical problem solving

1. Model building
2. Model solving
3. Model analysis

## Example: Production problem

A firm produces $n$ different goods using $m$ different raw materials.

- $b_{i}$ : available amount of the $i$-th raw material
- $a_{i j}$ : number of units of the $i$-th material needed to produce one unit of the $j$-th good
- $c_{j}$ : revenue for one unit of the $j$-th good.

Decide how much of each good to produce in order to maximize the total revenue $\rightsquigarrow$ decision variables $x_{j}$.

## Linear programming formulation

$$
\begin{aligned}
& \max c_{1} x_{1}+\cdots+c_{n} x_{n} \\
& \text { w.r.t. } a_{11} x_{1}+\cdots+a_{1 n} x_{n} \leq b_{1} \text {, } \\
& \begin{array}{c}
\vdots \\
a_{m 1} x_{1}+\cdots+a_{m n} x_{n} \leq b_{m},
\end{array} \\
& x_{1}, \quad \ldots \quad, x_{n} \geq 0 .
\end{aligned}
$$

In matrix notation:

$$
\max \left\{c^{\top} x \mid A x \leq b, x \geq 0\right\}
$$

where $A \in \mathbb{R}^{m \times n}, b \in \mathbb{R}^{m}, c \in \mathbb{R}^{n}, x \in \mathbb{R}^{n}$.

## Geometric illustration



$$
\begin{array}{cr}
\max & x_{1}+x_{2} \\
\text { w.r.t. } & x_{1}+2 x_{2} \leq 3 \\
& 2 x_{1}+x_{2} \leq 3 \\
& x_{1}, x_{2} \geq 0
\end{array}
$$

