On the relation between interaction and state transition graphs

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- T. Lorenz's Algorithms
- ASTG checking

2 Current work and results

- Current work
- Application on 3-node MAPK cascade model

3 Future work and so on







Asynchronous State Transition Graph (ASTG)









u: component*x*: state

$$Res(u,x) = \left\{ v \in Pre(u) | \begin{array}{c} \varepsilon(v,u) = + & \wedge & x_v \ge \vartheta(v,u) \\ \varepsilon(v,u) = - & \wedge & x_v < \vartheta(v,u) \end{array} \right\}$$

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Model: M = (I, K)



Preliminaries



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state	$Res(X_1, \cdot)$	$K(X_1, Res(X_1, \cdot))$
00	$\{X_1\}$	0
01	$\{X_1, X_2\}$	2
10	$\{X_1\}$	0
11	$\{X_1, X_2\}$	2
20	ϕ	0
21	$\{X_2\}$	0





Preliminaries



$$Res(u,x) = \left\{ v \in Pre(u) | \begin{array}{c} \varepsilon(v,u) = + & \wedge & x_v \ge \vartheta(v,u) \\ \varepsilon(v,u) = - & \wedge & x_v < \vartheta(v,u) \end{array} \right\}$$





ASTG T = (X, S)



Extreme state a state $x, \forall u \in V, x_u \in \{0, max_u\}.$



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a row of states
$$\tau(x^0, u) = (x^0, \dots, x^k),$$

 $x_u^i = x^0 + l\mathbf{e}_u, \forall l \in \{0, \dots, k\}.$

isomorphic *u*-rows

$$(x^0, \dots, x^k), (y^0, \dots, y^k)$$

if $(x_i, x_j) \in S$ iff $(y_i, y_j) \in S, \forall i, j \in \{0, \dots, k\}.$



ASTG T = (X, S)



Extreme state a state $x, \forall u \in V, x_u \in \{0, max_u\}.$

 $u - \mathbf{row}$

a row of states $\tau(x^0, u) = (x^0, ..., x^k),$ $x_u^i = x^0 + l\mathbf{e}_u, \forall l \in \{0, ..., k\}.$

Extreme row. an u-row $\tau(x^0, u)$ with extreme state x^0 and $k = max_u$.



Theorem

For any model M = (I, K), the state transition graph T_M is uniquely determined by IG I and the extremal rows of T_M .[1]



Proposition





ASTG description using extremal rows



Corollary

Let $u, v \in V$, $u \neq v$, and $x \in X$ with $x_v = x_u = 0$. For every $i \leq max_v \tau(x + i\mathbf{e}_v, u)$ with start state $(x + i\mathbf{e}_v)$. Then there are:

- If $(v, u) \notin E$, all rows are isomorphic.
- If $(v, u) \in E$, all rows in $\tau^0, \ldots, \tau^{\vartheta(v, u)-1}$ and $\tau^{\vartheta(v, u)}, \ldots, \tau^{\max_v}$ are isomorphic respectively. [2]



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Model conditions



 $(u, v) \in E$ is **visible**, if $\exists \omega \subseteq Pre\{v\} \setminus u$, $K(v, \omega) \neq K(v, \omega \cup \{u\})$. Eg. (X_2, X_1) $\exists K(X_1, \{X_1\}) \neq K(X_1, \{X_1, X_2\})$ **Visibility Model**

> A model M = (I, K)where $\forall (u, v) \in E$ is visible.





 $(u, v) \in E$ is **observable**, if $\exists \omega \subseteq Pre\{v\} \setminus u$, $K(v, \omega) < K(v, \omega \cup \{u\})$. Eg. (X_2, X_1) $\exists K(X_1, \{X_1\}) < K(X_1, \{X_1, X_2\})$ **Observibility Model** A model M = (I, K)where $\forall (u, v) \in E$ is observable.







A model M satisfies Snoussi-condition $\forall u \in V, \, \forall \zeta \subseteq \omega \subseteq Pre(u),$ $K(u,\zeta) \le K(u,\omega).$

This model is

Snoussi model.

Eg. (X_2, X_1) $K(X_1, \phi) = K(X_1, \{X_1\})$ $K(X_1, \{X_2\}) < K(X_1, \{X_1, X_2\})$ Satisfies Snoussi condition.



Model conditions





Work Review T. Lor

T. Lorenz's Algorithms

T. Lorenz's Algorithms: Visibility Model

ASTG





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T. Lorenz's Algorithms: Observability-Snoussi Model

Observability-Snoussi Model M = (I, K)

A model M = (I, K), with minimal number

ASTG





Inverse engineering: From ASTG \rightarrow Model M(I, K)

- Sometimes, the output *M* can not reproduce the input. Why?
- Multiple edges in the IG corresponding to the input ASTG is not allowed.
- Note: in the corresponding IG, thresholds from one component to others can be the same.(*Different with Thomas' model*)



Multiple Edges X_1, X_2, X are components.











ASTG Checking

Aim The input ASTG should be checked before using T. Lorenz's algorithms.

• Idea: the ASTG which will be the input of T. Lorenz's algorithms, the IG of this ASTG should have no multiple edges.





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- Method: on small scale, enumerate all possible ASTGs and get corresponding models using T. Lorenz's algorithms.
- Analysis.
 - How many IGs share the same attractors?
 - $\bullet~$ The changes on ASTGs $\longleftrightarrow~$ the changes on IGs.



Fig. Experiment process map



Exploring relations between IG and ASTG

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Fig. Experiment process map



Application on 3-node MAPK cascade model

- A reduced MAPK model from Kirsten. ;)
- Set input RTK to 1, the corresponding ASTG is in below on the right.



ASTG: Orange path is the cyclic attractor. Dashed arrows mean in the back of the cube. All directed arrows are transitions.



Application on 3-node MAPK cascade model

- A reduced MAPK model from Kirsten. ;)
- Set input RTK to 1, the corresponding ASTG is in below on the right.
- Experiment: how many IGs share the same cyclic attractors?
 - Enumerate all possible transitions on $X := \{0, 1\}^3$ which keeps the cyclic cycle. ($2^6 = 64$)
 - Use T. Lorenz's Algorithm: Visibility Model, get all visible models.
 - See the different IGs.



ASTG: Orange path is the cyclic attractor. Dashed arrows mean in the back of the cube. All directed arrows are transitions.



Current work and results

Results

• Different ASTG & corresponding IG.





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Current work and results App

Results

• Different ASTG & corresponding IG.





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• Different ASTG & corresponding IG.





• Eight kinds of interaction graphs.









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Analysis

• Eg. To understand why $Merk \rightarrow Mek$ and Erk - - - |Mek.





Analysis



• Eg. To understand why *Merk* \rightarrow *Mek* and *Erk* - - - |Mek.



Analysis



• Eg. To understand why $Merk \rightarrow Mek$ and Erk - - - |Mek.



- Other small network options.
 - Firdevs' Signaling models (4 components).
 - GI/S cell cycle regulation(2 components).
 - Bacteriophage λ infection (4 components).
 - Neural development in Rat (CNS, 4 components).
 - Circadian clock (from 2 to 5 components). The last four, from Adam.
- Open question. What kind of graph is a real ASTG? What kind of ASTG doesn't carriy any IG?



Thanks for your attention!

Questions are warmly welcome!

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