

Graph Algorithms

I. Shortest paths

- $D = (V, A)$ directed graph, $s, t \in V$.
- A *walk* is a sequence $P = (v_0, a_1, v_1, \dots, a_k, v_k), k \geq 0$, where a_i is an arc from v_{i-1} to v_i , for $i = 1, \dots, k$.
- P is a *path*, if v_0, \dots, v_k are all different.
- If $s = v_0$ and $t = v_k$, P is a *s-t walk* resp. *s-t path of length k* (i.e., each arc has length 1).
- The *distance* from s to t is the minimum length of any *s-t path* (and $+\infty$ if no *s-t path* exists).

Shortest paths with unit lengths

Algorithm (Breadth-first search)

Initialization: $V_0 = \{s\}$

Iteration: $V_{i+1} = \{v \in V \setminus (V_0 \cup V_1 \cup \dots \cup V_i) \mid (u, v) \in A, \text{ for some } u \in V_i\}$,
until $V_{i+1} = \emptyset$.

Running time: $O(|A|)$

- V_i is the set of nodes with distance i from s .
- The algorithm computes shortest paths from s to all reachable nodes.
- Can be described by a directed tree $T = (V', A')$ with root s such that each $u-v$ path in T is a shortest $u-v$ path in D .

Shortest paths with non-negative lengths

- Length function $l : A \rightarrow \mathbb{Q}_+ = \{x \in \mathbb{Q} \mid x \geq 0\}$
- For a walk $P = (v_0, a_1, v_1, \dots, a_k, v_k)$ define $l(P) = \sum_{i=1}^k l(a_i)$.

Algorithm (Dijkstra 1959)

Initialization: $U = V, f(s) = 0, f(v) = \infty$, for $v \in V \setminus \{s\}$

Iteration: Find $u \in U$ with $f(u) = \min\{f(v) \mid v \in U\}$.

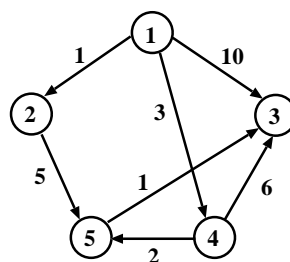
For all $a = (u, v) \in A$ with $f(v) > f(u) + l(a)$ let $f(v) = f(u) + l(a)$.

Let $U \leftarrow U \setminus \{u\}$, until $U = \emptyset$.

Upon termination, $f(v)$ gives the length of a shortest path from s to v .

Running time: $O(|V|^2)$ (can be improved to $O(|A| + |V| \log |V|)$.)

Example



Iteration	u	U	$f[1]$	$f[2]$	$f[3]$	$f[4]$	$f[5]$
0	–	{1, 2, 4, 3, 5}	0	∞	∞	∞	∞
1	1	{2, 3, 4, 5}	0	1	10	3	∞
2	2	{3, 4, 5}	0	1	10	3	6
3	4	{3, 5}	0	1	9	3	5
4	5	{3}	0	1	6	3	5
5	3	{}	0	1	6	3	5

Application: Longest common subsequence

- Sequences $a = a_1, \dots, a_m$ and $b = b_1, \dots, b_n$
- Find the longest common subsequence of a and b (obtained by removing symbols in a or b).

Modeling as a shortest path problem

- Grid graph with nodes $(i, j), 0 \leq i \leq m, 0 \leq j \leq n$.
- Horizontal and vertical arcs of length 1.
- Diagonal arcs $((i-1, j-1), (i, j))$ of length 0, if $a_i = b_j$.

The diagonal arcs on a shortest path from $(0, 0)$ to (m, n) define a longest common subsequence.

Circuits of negative length

- Consider arbitrary length functions $l : A \rightarrow \mathbb{Q}$.
- A *directed circuit* is a walk $P = (v_0, a_1, v_1, \dots, a_k, v_k)$ with $k \geq 1$ and $v_0 = v_k$ such that v_1, \dots, v_k and a_1, \dots, a_k are all different.
- If $D = (V, A)$ contains a directed circuit of negative length, there exist s - t walks of arbitrary small negative length.

Proposition

Let $D = (V, A)$ be a directed graph without circuits of negative length.

For any $s, t \in V$ for which there exists at least one s - t walk, there exists a shortest s - t walk, which is a path.

Shortest paths with arbitrary lengths

$D = (V, A), n = |V|, l : A \rightarrow \mathbb{Q}$.

Algorithm (Bellman-Ford 1956/58)

Compute $f_0, \dots, f_n : V \rightarrow \mathbb{R} \cup \{\infty\}$ in the following way:

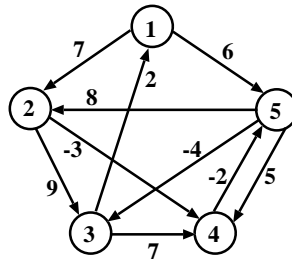
Initialization: $f_0(s) = 0, f_0(v) = \infty$, for $v \in V \setminus \{s\}$

Iteration: For $k = 1, \dots, n$ and all $v \in V$:

$$f_k(v) = \min\{f_{k-1}(v), \min_{(u,v) \in A} (f_{k-1}(u) + l(u, v))\}$$

Running time: $O(|V||A|)$

Example



Iteration k	$f_k[1]$	$f_k[2]$	$f_k[3]$	$f_k[4]$	$f_k[5]$
0	0	∞	∞	∞	∞
1	0	7	∞	∞	6
2	0	7	2	4	6
3	0	7	2	4	2
4	0	7	-2	4	2

Properties

- For each $k = 0, \dots, n$ and each $v \in V$:

$$f_k(v) = \min\{l(P) \mid P \text{ is an } s\text{-}v \text{ walk traversing at most } k \text{ arcs}\}$$

(by induction)

- If D contains no circuits of negative length, $f_{n-1}(v)$ is the length of a shortest path from s to v .

Finding an explicit shortest path

- When computing f_0, \dots, f_n determine a predecessor function $p : V \rightarrow V$ by setting $p(v) = u$ whenever $f_{k+1}(v) = f_k(u) + l(u, v)$.
- At termination, $v, p(v), p(p(v)), \dots, s$ gives the reverse of a shortest s - v path.

Theorem

Given $D = (V, A)$, $s, t \in V$ and $l : A \rightarrow \mathbb{Q}$ such that D contains no circuit of negative length, a shortest s - t path can be found in time $O(|V||A|)$.

Remark

D contains a circuit of negative length reachable from s if and only if $f_n(v) \neq f_{n-1}(v)$, for some $v \in V$.

NP-completeness

For directed graphs containing circuits of negative length, the problem becomes NP-complete:

Theorem

The decision problem

Input: Directed graph $D = (V, A)$, $s, t \in V$, $l : A \rightarrow \mathbb{Z}$, $L \in \mathbb{Z}$

Question: Does there exist an s - t path P with $l(P) \leq L$?

is NP-complete.

Corollary

The shortest path problem with arbitrary lengths is NP-complete.

The longest path problem with non-negative lengths is NP-complete.

Application: Knapsack problem

- Knapsack, volume 8, 5 articles

Article i	Volume a_i	Value c_i
1	5	4
2	3	7
3	2	3
4	2	5
5	1	4

- Objective: Select articles fitting into the knapsack and maximizing the total value.

Possible models

- *Shortest path model*

- Directed graph with nodes $(i, x), 0 \leq i \leq 6, 0 \leq x \leq 8$.
- Arcs from $(i-1, x)$ to (i, x) resp. $(i, x+a_i)$ of length 0 resp. $-c_i$, for $0 \leq i \leq 5$.
- Arcs from $(5, x)$ to $(6, 8)$ of length 0, for $0 \leq x \leq 6$.
- A shortest path from $(0, 0)$ to $(6, 8)$ gives an optimal solution.

↪ *pseudo-polynomial algorithm*

- *Linear 0-1 model*

$$\max\{4x_1 + 7x_2 + 3x_3 + 5x_4 + 4x_5 \mid 5x_1 + 3x_2 + 2x_3 + 2x_4 + x_5 \leq 8, x_1, \dots, x_5 \in \{0, 1\}\}$$