

Linear programming

Optimization Problems

- *General optimization problem*

$$\max\{z(x) \mid f_j(x) \leq 0, x \in D\} \text{ or } \min\{z(x) \mid f_j(x) \leq 0, x \in D\}$$

where $D \subseteq \mathbb{R}^n$, $f_j : D \rightarrow \mathbb{R}$, for $j = 1, \dots, m$, $z : D \rightarrow \mathbb{R}$.

- *Linear optimization problem*

$$\max\{c^T x \mid Ax \leq b, x \in \mathbb{R}^n\}, \text{ with } c \in \mathbb{R}^n, A \in \mathbb{R}^{m \times n}, b \in \mathbb{R}^m$$

- *Integer optimization problem*: $x \in \mathbb{Z}^n$
- *0-1 optimization problem*: $x \in \{0, 1\}^n$

Feasible and optimal solutions

- Consider the optimization problem

$$\max\{z(x) \mid f_j(x) \leq 0, x \in D, j = 1, \dots, m\}$$

- A *feasible solution* is a vector $x^* \in D \subseteq \mathbb{R}^n$ such that $f_j(x^*) \leq 0$, for all $j = 1, \dots, m$.
- The set of all feasible solutions is called the *feasible region*.
- An *optimal solution* is a feasible solution such that

$$z(x^*) = \max\{z(x) \mid f_j(x) \leq 0, x \in D, j = 1, \dots, m\}.$$

- Feasible or optimal solutions
 - need not exist,
 - need not be unique.

Transformations

- $\min\{z(x) \mid x \in S\} = \max\{-z(x) \mid x \in S\}$.
- $f(x) \geq a$ if and only if $-f(x) \leq -a$.
- $f(x) = a$ if and only if $f(x) \leq a \wedge -f(x) \leq -a$.

Lemma

Any linear programming problem can be brought to the form

$$\max\{c^T x \mid Ax \leq b\} \text{ or } \max\{c^T x \mid Ax = b, x \geq 0\}.$$

Proof: a) $a^T x \leq \beta \rightsquigarrow a^T x + x' = \beta, x' \geq 0$ (*slack variable*)

b) x free $\rightsquigarrow x = x^+ - x^-, x^+, x^- \geq 0$.

Practical problem solving

1. Model building
2. Model solving
3. Model analysis

Example: Production problem

A firm produces n different goods using m different raw materials.

- b_i : available amount of the i -th raw material
- a_{ij} : number of units of the i -th material needed to produce one unit of the j -th good
- c_j : revenue for one unit of the j -th good.

Decide how much of each good to produce in order to maximize the total revenue \rightsquigarrow *decision variables* x_j .

Linear programming formulation

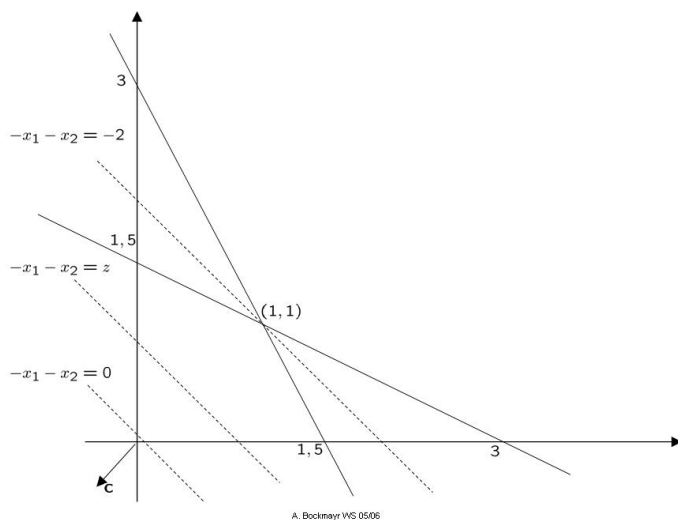
$$\begin{aligned} \max \quad & c_1 x_1 + \cdots + c_n x_n \\ \text{w.r.t.} \quad & a_{11} x_1 + \cdots + a_{1n} x_n \leq b_1, \\ & \vdots \\ & a_{m1} x_1 + \cdots + a_{mn} x_n \leq b_m, \\ & x_1, \quad \dots, \quad x_n \geq 0. \end{aligned}$$

In matrix notation:

$$\max \{c^T x \mid Ax \leq b, x \geq 0\},$$

where $A \in \mathbb{R}^{m \times n}$, $b \in \mathbb{R}^m$, $c \in \mathbb{R}^n$, $x \in \mathbb{R}^n$.

Geometric illustration



$$\begin{aligned} \max \quad & x_1 + x_2 \\ \text{w.r.t.} \quad & x_1 + 2x_2 \leq 3 \\ & 2x_1 + x_2 \leq 3 \\ & x_1, x_2 \geq 0 \end{aligned}$$