Prof. Dr. Alexander Bockmayr, Prof. Dr. Knut Reinert, Sandro Andreotti

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# Discrete Mathematics for Bioinformatics (P1) WS 2010/11

#### Exercises 5

#### 1. Bellman-Ford (Niveau I)

- (a) Use the Bellman-Ford algorithm to determine the shortest path from source A to any other node in the graph.
- (b) Let D = (V, A), n = |V| be a directed graph. Prove that D contains a circuit of negative length reachable from s if and only if  $f_n(v) \neq f_{n-1}(v)$ , for some  $v \in V$ , where  $f_k(v) = \min\{l(P)|P \text{ is an } s v \text{ walk traversing at most } k \text{ arcs}\}$



2. Network Flow (Niveau II) Prove the Theorem: For a network (V, E, s, t) with capacities cap :  $E \to \mathbb{R}_+$  the maximum value of a flow is equal to the minimum capacity of an (s, t)-cut:

 $\max\{\operatorname{val}(f) \mid f \text{ is a flow}\} = \min\{\operatorname{cap}(S,T) \mid (S,T) \text{ is an } (s,t)\text{-cut}\}$ 

Hint: Show that the following conditions are equivalent:

- (a) f is a maximum flow.
- (b) The residual network  $G_f$  contains no augmenting path.
- (c) val(f) = cap(S, T) for some cut (S, T) of G

## 3. Ford-Fulkerson (Niveau I)

(a) Use the Ford-Fulkerson algorithm to find a maximum flow in the network



Start with the initial flow f. An edge label f/c means initial flow f and capacity c.

(b) Find a minimum cut proving the maximality of the flow.

### 4. Matching and Bipartite Graphs (Niveau I)

(a) Apply the matching augmenting algorithm for bipartite graphs to the graph below and compute a maximum cardinality matching from the initial matching.

