

## Metaheuristics and Local Search

### Discrete optimization problems

- Variables  $x_1, \dots, x_n$ .
- Variable domains  $D_1, \dots, D_n$ , with  $D_j \subseteq \mathbb{Z}$ .
- Constraints  $C_1, \dots, C_m$ , with  $C_i \subseteq D_1 \times \dots \times D_n$ .
- Objective function  $f : D_1 \times \dots \times D_n \rightarrow \mathbb{R}$ , to be minimized.

### Solution approaches

- Complete (exact) algorithms  $\rightsquigarrow$  *systematic search*
  - Integer linear programming
  - Finite domain constraint programming
- Approximate algorithms
  - Heuristic approaches  $\rightsquigarrow$  *heuristic search*
    - \* Constructive methods: construct solutions from partial solutions
    - \* **Local search**: improve solutions through neighborhood search
    - \* **Metaheuristics**: Combine basic heuristics in higher-level frameworks
  - Polynomial-time approximation algorithms for NP-hard problems

### Metaheuristics

- Heuriskein (ευρισκειν): to find
- Meta: beyond, in an upper level
- *Survey paper*: C. Blum, A. Roli: Metaheuristics in Combinatorial Optimization, ACM Computing Surveys, Vol. 35, 2003. <http://iridia.ulb.ac.be/meta/newsite/downloads/ACSUR-blum-rol.pdf>

### Characteristics

- Metaheuristics are strategies that “guide” the search process.
- The goal is to efficiently explore the search space in order to find (near-) optimal solutions.
- Metaheuristics range from simple local search to complex learning procedures.
- Metaheuristic algorithms are approximate and usually non-deterministic.
- They may incorporate mechanisms to avoid getting trapped in confined areas of the search space.

### Characteristics <sup>(2)</sup>

- The basic concepts of metaheuristics permit an abstract level description.
- Metaheuristics are not problem-specific.
- Metaheuristics may make use of domain-specific knowledge in the form of heuristics that are controlled by the upper level strategy.

- Today more advanced metaheuristics use search experience (embodied in some form of memory) to guide the search.

## Intensification and diversification

*Glover and Laguna 1997*

The main difference between intensification and diversification is that during an intensification stage the search focuses on examining neighbors of elite solutions. . . . The diversification stage on the other hand encourages the search process to examine unvisited regions and to generate solutions that differ in various significant ways from those seen before.

## Classification of metaheuristics

- Single point search (trajectory methods) vs. population-based search
- Nature-inspired vs. non-nature inspired
- Dynamic vs. static objective function
- One vs. various neighborhood structures
- Memory usage vs. memory-less methods

### I. Trajectory methods

- Basic local search: iterative improvement
- Simulated annealing
- Tabu search
- Explorative search methods
  - Greedy Randomized Adaptive Search Procedure (GRASP)
  - Variable Neighborhood Search (VNS)
  - Guided Local Search (GLS)
  - Iterated Local Search (ILS)

### Local search

- Find an initial solution  $s$
- Define a neighborhood  $\mathcal{N}(s)$
- Explore the neighborhood
- Proceed with selected neighbor

### Simple descent

```

procedure SimpleDescent(solution s)
  repeat
    choose  $s' \in \mathcal{N}(s)$ 
    if  $f(s') < f(s)$  then
       $s \leftarrow s'$ 
    end if
  until  $f(s') \geq f(s), \forall s' \in \mathcal{N}(s)$ 
end

```

## Deepest descent

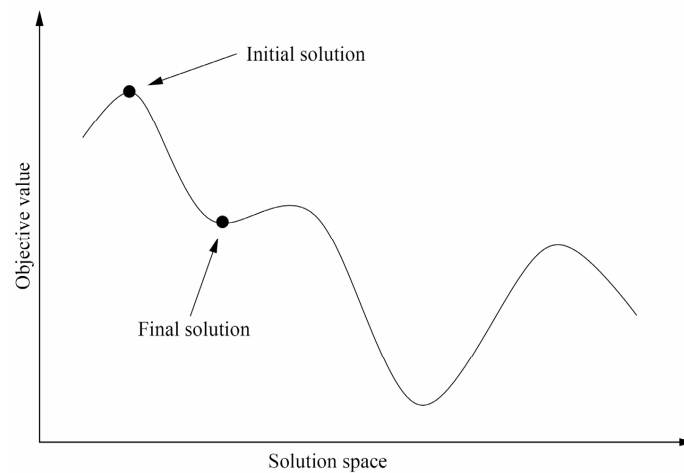
```

procedure DeepestDescent(solution s)
  repeat
    choose  $s' \in \mathcal{N}(s)$  with  $f(s') \leq f(s''), \forall s'' \in \mathcal{N}(s)$ 
    if  $f(s') < f(s)$  then
       $s \leftarrow s'$ 
    end if
  until  $f(s') \geq f(s), \forall s' \in \mathcal{N}(s)$ 
end

```

**Problem:** Local minima

## Local and global minima



## Multistart and deepest descent

```

procedure Multistart
  iter ← 1
  f(Best) ← ∞
  repeat
    choose a starting solution  $s_0$  at random
     $s \leftarrow \text{DeepestDescent}(s_0)$ 
    if  $f(s) < f(\text{Best})$  then
      Best ← s
    end if
    iter ← iter + 1
  until iter = IterMax
end

```

## Simulated annealing

*Kirkpatrick 83*

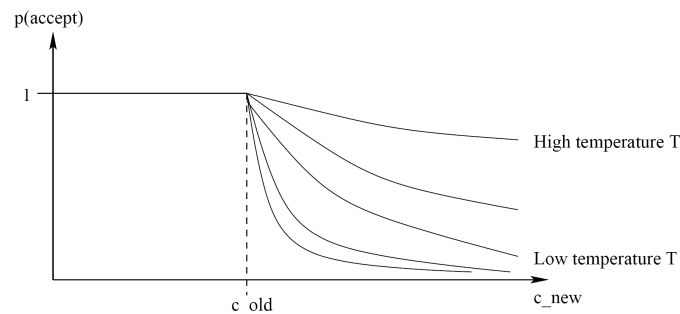
- *Anneal*: to heat and then slowly cool (esp. glass or metal) to reach minimal energy state
- Like standard local search, but sometimes accept worse solution.

- Select random solution from the neighborhood and accept it with probability  $\rightsquigarrow$  Boltzmann distribution

$$p = \begin{cases} 1, & \text{if } f(\text{new}) < f(\text{old}), \\ \exp(-(f(\text{new}) - f(\text{old}))/T), & \text{else.} \end{cases}$$

- Start with high temperature  $T$ , and gradually lower it  $\rightsquigarrow$  cooling schedule

### Acceptance probability



### Algorithm

```

s ← GenerateInitialSolution()
T ← T0
while termination conditions not met do
  s' ← PickAtRandom(N(s))
  if (f(s') < f(s)) then
    s ← s'
  else
    Accept s' as new solution with probability p(T, s', s)
  endif
  Update(T)
endwhile

```

### Tabu search

Glover 86

- Local search with short term memory, to escape local minima and to avoid cycles.
- *Tabu list*: Keep track of the last  $r$  moves, and don't allow going back to these.
- *Allowed set*: Solutions that do not belong to the tabu list.
- Select solution from allowed set, add to tabu list, and update tabu list.

### Basic algorithm

```

s ← GenerateInitialSolution()
TabuList ← ∅
while termination conditions not met do
  s ← ChooseBestOf( $\mathcal{N}(s) \setminus \text{TabuList}$ )
  Update(TabuList)
endwhile

```

### Choices in tabu search

- Neighborhood
- Size of tabu list  $\rightsquigarrow$  *tabu tenure* (static/dynamic)
- Kind of tabu to use (complete solutions vs. attributes)  $\rightsquigarrow$  *tabu conditions*
- *Aspiration criteria* (exceptions to tabu conditions)
- Termination condition
- Long-term memory: recency, frequency, quality, influence

### Refined algorithm

```

s ← GenerateInitialSolution()
Initialize TabuLists ( $TL_1, \dots, TL_r$ )
k ← 0
while termination conditions not met do
  AllowedSet(s, k) ← { s' ∈  $\mathcal{N}(s)$  |
    s' does not violate a tabu condition
    or satisfies at least one aspiration condition }
  s ← ChooseBestOf(AllowedSet(s, k))
  UpdateTabuListsAndAspirationConditions()
  k ← k + 1
endwhile

```

## II. Population-based search

Use a set (i.e. a population) of solutions rather than a single solutions

- Evolutionary computation
- Ant colony optimization

### Evolutionary computation

- Idea: Mimic evolution - obtain better solutions by combining current ones.
- Keep several current solutions, called *population* or *generation*.
- Create new generation:
  - select a pool of promising solutions, based on a *fitness function*.
  - create new solutions by combining solutions in the pool in various ways  $\rightsquigarrow$  *recombination, crossover*.
  - add random *mutations*.

- *Variants*: Evolutionary programming, evolutionary strategies, genetic algorithms

## Algorithm

```

P ← GenerateInitialPopulation()
Evaluate(P)
while termination conditions not met do
  P' ← Recombine(P)
  P'' ← Mutate(P')
  Evaluate(P'')
  P ← Select(P'' ∪ P)
endwhile

```

## Crossover and mutations

- Individuals (solutions) often coded as bit or integer vectors
- *Crossover* operations provide new individuals, e.g.
 

101101		0110	↔	101101		1011
000110		1011	↔	000110		0110
- *Mutations* often helpful, e.g., swap random bit.

## Further issues

- Individuals (“genotypes”) vs. solutions (“phenotypes”): individuals are not necessarily solutions
- Evolution process: generational replacement vs. steady state fixed vs. variable population size
- Use of neighborhood structure to define recombination partners
- Two-parent vs. multi-parent crossover
- Infeasible individuals: reject/penalize/repair
- Intensification by local search
- Diversification by mutations

## Ant colony optimization

*Dorigo 92*

- Observation: Ants are able to find quickly the shortest path from their nest to a food source ↔ how ?
- Each ant leaves a *pheromone* trail.
- When presented with a path choice, they are more likely to choose the trail with higher pheromone concentration.
- The shortest path gets high concentrations because ants choosing it can return more often.

## Ant colony optimization <sup>(2)</sup>

- Ants are simulated by individual (ant) agents ↔ *swarm intelligence*

- Artificial ants incrementally construct solutions by adding components to a partial solution.
- By dispatching a number of ants, the pheromone levels associated with the components are adjusted according to how useful they are.
- Pheromone levels may also *evaporate* to discourage suboptimal solutions.

## Construction graph

- Complete graph  $G = (C, L)$ 
  - $C$  solution components
  - $L$  connections
- Pheromone trail values  $\tau_i, \tau_{ij} \in \mathcal{T}$ , for  $c_i \in C, l_{ij} \in L$
- Heuristic values  $\eta_i, \eta_{ij} \in \mathcal{H}$
- Moves in the graph depend on transition probabilities (use only  $\tau_i, \eta_i$ )

$$p(c_r | s_a[c_i]) = \begin{cases} \frac{[\eta_r]^\alpha [\tau_r]^\beta}{\sum_{c_u \in J(s_a[c_i])} [\eta_u]^\alpha [\tau_u]^\beta}, & \text{if } c_r \in J(s_a[c_i]), \\ 0, & \text{otherwise.} \end{cases}$$

$s_a$  denotes the solution constructed by ant  $a$ ,  $c_i$  its last component.

$J(s_a[c_i])$  denotes the set of components allowed to be added.

## Algorithm: Ant System (AS)

InitializePheromoneValues

**while** termination conditions not met **do**

**for** all ants  $a \in \mathcal{A}$  **do**

$s_a \leftarrow \text{ConstructSolution}(\mathcal{T}, \mathcal{H})$

**endfor**

  ApplyOnlineDelayedPheromoneUpdate( $\mathcal{T}, \{s_a \mid a \in \mathcal{A}\}$ )

**endwhile**

- ApplyOnlineDelayedPheromoneUpdate  $\rightsquigarrow$  increase of pheromone on solutions components found in high-quality solutions
- The Ant System (AS) algorithm may be generalized to the Ant Colony Optimization (ACO) Metaheuristics.