

Prof. Dr. Knut Reinert,  
Dr. Yaron Goldstein,  
Sandro Andreotti

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# Algorithms

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## Exercises 6

### 1. Tree decomposition (Niveau I)

How large is the largest piece of a any tree decomposition for a graph  $G$  of  $n$  nodes if  $G$  is a clique? Prove your answer.

### 2. Tree decomposition (Niveau I)

Prove the following theorem:

Let  $G = (V, E)$  be a graph,  $T$  be a tree decomposition of  $G$ , and  $(x, y)$  an edge in  $T$ . The deletion of  $(x, y)$  divides  $T$  into two components  $X$  and  $Y$ . Let  $V_x$  and  $V_y$  be the ‘pieces’ of  $x$  and  $y$ , respectively. Then deleting the set  $V_x \cap V_y$  from  $V$  disconnects  $G$  into the two subgraphs  $G_X - (V_x \cap V_y)$  and  $G_Y - (V_x \cap V_y)$ .

( $G_M$  for  $M = X, Y$  is the subgraph of  $G$  that consists of all nodes in the ‘pieces’ of  $M$ .)

### 3. Tree decomposition (Niveau II)

Prove the following theorem:

If graph  $G$  contains a  $(w + 1)$ -linked set of size at least  $3w$ , then  $G$  has tree-width at least  $w$ .

Suppose, by way of contradiction, that  $G$  has a  $(w + 1)$ -linked set  $X$  of size at least  $3w$ , and it also has a nonredundant TD  $(T; \{V_i\})$  of width less than  $w$ . The idea of the proof is to find a piece  $V_i$  that is “centered” with respect to  $X$ , so that when some part of  $V_i$  is deleted from  $G$ , one small subset of  $X$  is separated from another.

#### 4. Tree decomposition (Niveau I)

Use the algorithm presented in the lecture to compute a tree decomposition of the graph below. Use  $w = 4$ :

