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Algorithms

WS 2012/13

Exercises 3

1. Modulo Arithmetic (Niveau II) Prove the following theorem:

For any positive intergers a and n, if d = gcd(a, n) (the greatest common divisor of a and n), then

$$\langle a \rangle = \langle d \rangle = \{0, d, 2d, \dots, n-d\}$$

and thus

 $|\langle a \rangle| = n/d$

 $(\langle a \rangle := \{ a \cdot i \mod n \mid i \in \mathbb{N} \}).$

Hint: Use Bezout's lemma. It states that if a and b are nonzero integers with greatest common divisor d, then there exist integers x and y such that ax + by = d

- 2. Hashing (Niveau II) Consider a version of the division method in which $h(k) = k \mod m$, where $m=2^p-1$ and k is a character string interpreted in radix 2^p . Show that if string x can be derived from string y by permuting its characters, then x and y hash to the same value.
- 3. Expected value (Niveau I) Proof: Define pi = Pr (exactly i probes access occupied slots) for i = 0, 1, 2, ... (Note that for $i > n, p_i = 0$). The expected number of probes is then $\sum_{i=0}^{\infty} i \cdot p_i$. Now define $q_i = Pr$ (at least i probes access occupied slots). Show that $\sum_{i=0}^{\infty} i \cdot p_i = \sum_{i=1}^{\infty} q_i$.