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Representation and Partial Automation of the Principia Logico-Metaphysica in Isabelle/HOL

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Abstract

We present an embedding of the second-order fragment of the Theory of Abstract Objects as described in Edward Zalta's upcoming work Principia Logico-Metaphysica (PLM[12]) in the automated reasoning framework Isabelle/HOL. The Theory of Abstract Objects is a metaphysical theory that reifies property patterns, as they for example occur in the abstract reasoning of mathematics, as *abstract objects* and provides an axiomatic framework that allows to reason about these objects. It thereby serves as a fundamental metaphysical theory that can be used to axiomatize and describe a wide range of philosophical objects, such as Platonic forms or Leibniz' concepts, and has the ambition to function as a foundational theory of mathematics. The target theory of our embedding as described in chapters 7-9 of PLM[12] employs a modal relational type theory as logical foundation for which a representation in functional type theory is known to be challenging[8].

Nevertheless we arrive at a functioning representation of the theory in the functional logic of Isabelle/HOL based on a semantical representation of an Aczel-model of the theory. Based on this representation we construct an implementation of the deductive system of PLM ([12, Chap. 9]) which allows to automatically and interactively find and verify theorems of PLM.

Our work thereby supports the concept of shallow semantical embeddings of logical systems in HOL as a universal tool for logical reasoning as promoted by Christoph Benzmüller[1].

The most notable result of the presented work is the discovery of a previously unknown paradox in the formulation of the Theory of Abstract Objects. The embedding of the theory in Isabelle/HOL played a vital part in this discovery. Furthermore it was possible to immediately offer several options to modify the theory to guarantee its consistency. Thereby our work could provide a significant contribution to the development of a proper grounding for object theory.

Contents

1. Introduction

Calculemus!

Leibniz

1.1. Universal Logical Reasoning[1]

The concept of understanding rational argumentation and reasoning using formal logical systems has a long tradition and can already be found in the study of syllogistic arguments by Aristotle. Since then a large variety of formal systems has evolved, each using different syntactical and semantical structures to capture specific aspects of logical reasoning (e.g. propositional logic, first-order/higher-order logic, modal logic, free logic, etc.). This diversity of formal systems gives rise to the question, whether a *universal* logic can be devised, that would be capable of expressing statements of all existing specialized logical systems and provide a basis for meta-logical considerations like the equivalence of or relations between those systems.

The idea of a universal logical framework is very prominent in the works of Gottfried Wilhelm Leibniz (1646-1716) with his concept of a *characteristica universalis*, i.e. a universal formal language able to express metaphysical, scientific and mathematical concepts. Based thereupon he envisioned the *calculus ratiocinator*, a universal logical calculus with which the truth of statements formulated in the characteristica universalis could be decided purely by formal calculation and thereby in an automated fashion, an idea that became famous under the slogan: *Calculemus!*

Nowadays with the rise of powerful computer systems such a universal logical framework could have repercussions throughout the sciences and may be a vital part of human-machine interaction in the future. Leibniz' ideas have inspired recent efforts to use functional higher-order logic (HOL) as such a universal logical language and to represent various logical systems by the use of *shallow semantical embeddings*[1].

Notably this approach received attention due to the formalization, validation and analysis of Gödel's ontological proof of the existence of God by Christoph Benzmüller[5], for which higher-order modal logic was embedded in the computerized logic framework Isabelle/HOL.

[1] This introductory section is based on the description of the topic in [1].

1.2. Shallow Semantical Embeddings in HOL

A semantic embedding of a target logical system defines the syntactic elements of the target language in a background logic (e.g. in a framework like Isabelle/HOL) based on their semantics. This way the background logic can be used as meta-logic to argue about the semantic truth of syntactic statements in the embedded logic.

A *deep* embedding represents the complete syntactic structure of the target language separately from the background logic, i.e. every term, variable symbol, connective, etc. of the target language is represented as a syntactic object and then the background logic is used to evaluate a syntactic expression by quantifying over all models that can be associated with the syntax. Variable symbols of the target logic for instance would be represented as constants in the background logic and a proposition would be considered semantically valid if it holds for all possible denotations an interpretation function can assign to them.

While this approach will work for most target logics, it has several drawbacks. It is likely that there are principles that are shared between the target logic and the background logic, such as α-conversion for λ-expressions or the equivalence of terms with renamed variables in general. In a deep embedding these principles usually have to be explicitly shown to hold for the syntactic representation of the target logic, which is usually connected with significant complexity. Furthermore if the framework used for the background logic allows automated reasoning, the degree of automation that can be achieved in the embedded logic is limited, as any reasoning in the target logic will have to consider the meta-logical evaluation process in the background logic which will usually be complex.

A *shallow* embedding uses a different approach based on the idea that most contemporary logical systems are semantically characterized by the means of set theory. A shallow embedding defines primitive syntactic objects of the target language such as variables or propositions using a set theoretic representation. For example propositions in a modal logic can be represented as functions from possible worlds to truth values in a non-modal logic.

The shallow embedding aims to equationally define only the syntactic elements of the target logic that are not already present in the background logic or whose semantics behaves differently than in the background logic, while preserving as much of the logical structure of the background logic as possible. The modal box operator for example can be represented as a quantification over all possible worlds, satisfying an accessibility relation, while negation and quantification can be directly represented using the negation and quantification of the background logic (preserving the dependency on possible worlds).

This way basic principles of the background logic (such as alpha conversion) can often be directly applied to the embedded logic and the equational, definitional nature of the representation preserves a larger degree of automation. Furthermore, axioms in the embedded logic can often be equivalently stated in the background logic, which makes the

construction of models for the system easier and again increases the degree of automation that can be retained.

The shallow semantical embedding of modal logic was the basis for the analysis of Gödel's ontological argument[5] and the general concept has shown great potential as a universal tool for logical embeddings while retaining the existing infrastructure for automation as for example present in a framework like Isabelle/HOL[2].

1.3. Relational Type Theory vs. Functional Type Theory

The universality of this approach has since been challenged by Paul Oppenheimer and Edward Zalta who argue in the paper *Relations Versus Functions at the Foundations of Logic: Type-Theoretic Considerations*[8] that relational type theory is more general than functional type theory. In particular they argue that the Theory of Abstract Objects, which is founded in relational type theory, cannot be properly characterized in functional type theory.

This has led to the question whether a shallow semantical embedding of the Theory of Abstract Objects in a functional logic framework like Isabelle/HOL is at all possible, which is the core question the work presented here attempts to examine and partially answer.

One of their main arguments is that unrestricted λ-expressions as present in functional type theory lead to an inconsistency when combined with one of the axioms of the theory and indeed it has been shown for early attempts on embedding the theory that despite significant efforts to avoid the aforementioned inconsistency by excluding problematic λ-expressions in the embedded logic, it could still be reproduced using an appropriate construction in the background logic[3].

The solution presented here circumvents this problem by identifying λ-expressions as one element of the target language that behaves differently than their counterparts in the background logic and consequently by representing λ-expressions of the target logic using a new *defined* kind of λ-expressions. This forces λ-expressions in the embedded logic to have a particular semantics that is inspired by the *Aczel-model* of the target theory (see 2.6) and avoids prior inconsistencies. The mentioned issue and the employed solution is discussed in more detail in sections 3.2 and 3.4.7.

[2] See [1] for an overview and an description of the ambitions of the approach.

[3] Early attempts of an embedding by Christoph Benzmüller (see https://github.com/cbenzmueller/PrincipiaMetaphysica) were discussed in his university lecture *Computational Metaphysics* (FU Berlin, SS2016) and the proof of their inconsistency in the author's final project for the course inspired the continued research in this master's thesis.

1.4. Overview of the following Chapters

The following chapters are structured as follows:

- The second chapter gives an overview of the motivation and structure of the target theory of the embedding, the Theory of Abstract Objects. It also introduces the *Aczel-model* of the theory, that was adapted as the basis for the embedding.
- The third chapter is a detailed documentation of the concepts and technical structure of the embedding. This chapter references the Isabelle theory that can be found in the appendix.
- The fourth chapter consists of a technical discussion about some of the issues encountered during the construction of the embedding due to limitations of the logic framework Isabelle/HOL and the solutions that were employed.
- The last chapter discusses the relation between the embedding and the target theory of PLM and describes some of the results achieved using the embedding. Furthermore it states some open questions for future research.

This entire document is generated from an Isabelle theory file and thereby in particular all formal statements in the third chapter are well-formed terms, resp. verified valid theorems in the constructed embedding unless the contrary is stated explicitly.

2. The Theory of Abstract Objects

> It is widely supposed that every entity falls into one of two categories: Some are concrete; the rest abstract. The distinction is supposed to be of fundamental significance for metaphysics and epistemology.
>
> *Stanford Encyclopedia of Philosophy[9]*

2.1. Motivation

As the name suggests the Theory of Abstract Objects revolves around *abstract objects* and is thereby a metaphysical theory. As Zalta puts it: "Whereas physics attempts a systematic description of fundamental and complex concrete objects, metaphysics attempts a systematic description of fundamental and complex abstract objects. [...] The theory of abstract objects attempts to organize these objects within a systematic and axiomatic framework. [...] [We can] think of abstract objects as possible and actual property-patterns. [...] Our theory of abstract objects will *objectify* or *reify* the group of properties satisfying [such a] pattern."[13][1]

So what is the fundamental distinction between abstract and concrete objects? The analysis in the Theory of Abstract Objects is based on a distinction between two fundamental modes of predication that is based on the ideas of Ernst Mally. Whereas objects that are concrete (the Theory of Abstract Objects calls them *ordinary objects*) are characterized by the classical mode of predication, i.e. *exemplification*, a second mode of predication is introduced that is reserved for abstract objects. This new mode of predication is called *encoding* and formally written as xF (x *encodes* F) in contrast to Fx (x *exemplifies* F).

Mally informally introduces this second mode of predication in order to represent sentences about fictional objects. In his thinking, concrete objects, that for example have a fixed spatiotemporal location, a body and shape, etc., only *exemplify* their properties and are characterized by the properties they exemplify. Sentences about fictional objects such as "Sherlock Holmes is a detective" have a different meaning. Stating that "Sherlock Holmes is a detective" does not imply that there is some concrete object that is

[1] The introduction to the theory in this and the next section is based on the documentation of the theory in [13] and [14], which is paraphrased and summarized throughout the sections. Further references about the topic include [12], [11], [10].

Sherlock Holmes and this object exemplifies the property of being a detective - it rather states that the concept we have of the fictional character Sherlock Holmes includes the property of being a detective. Sherlock Holmes is not concrete, but an abstract object that is *determined* by the properties Sherlock Holmes is given by the fictional works involving him as character. This is expressed using the second mode of predication *Sherlock Holmes encodes the property of being a detective.*

To clarify the difference between the two concepts note that any object either exemplifies a property or its negation. The same is not true for encoding. For example it is not determinate whether Sherlock Holmes has a mole on his left foot. Therefore the abstract object Sherlock Holmes neither encodes the property of having a mole on his left foot, nor the property of not having a mole on his left foot[2].

The theory even allows for an abstract object to encode properties that no object could possibly exemplify and reason about them, for example the quadratic circle. In classical logic meaningful reasoning about a quadratic circle is impossible - as soon as I suppose that an object *exemplifies* the properties of being a circle and of being quadratic, this will lead to a contradiction and every statement becomes derivable.

In the Theory of Abstract Objects on the other hand there is an abstract object that encodes exactly these two properties and it is possible to reason about it. For example we can state that this object *exemplifies* the property of *being thought about by the reader of this paragraph.* This shows that the Theory of Abstract Objects provides the means to reason about processes of human thought in a much broader sense than classical logic would allow.

It turns out that by the means of abstract objects and encoding the Theory of Abstract Objects can be used to represent and reason about a large variety of concepts that regularly occur in philosophy, mathematics or linguistics.

In [13] the principal objectives of the theory are summarized as follows:

- To describe the logic underlying (scientific) thought and reasoning by extending classical propositional, predicate, and modal logic.
- To describe the laws governing universal entities such as properties, relations, and propositions (i.e., states of affairs).
- To identify *theoretical* mathematical objects and relations as well as the *natural* mathematical objects such as natural numbers and natural sets.
- To analyze the distinction between fact and fiction and systematize the various relationships between stories, characters, and other fictional objects.
- To systematize our modal thoughts about possible (actual, necessary) objects, states of affairs, situations and worlds.
- To account for the deviant logic of propositional attitude reports, explain the informativeness of identity statements, and give a general account of the objective and cognitive content of natural language.

[2] see [14]

- To axiomatize philosophical objects postulated by other philosophers, such as Forms (Plato), concepts (Leibniz), monads (Leibniz), possible worlds (Leibniz), nonexistent objects (Meinong), senses (Frege), extensions of concepts (Frege), noematic senses (Husserl), the world as a state of affairs (early Wittgenstein), moments of time, etc.

The Theory of Abstract Objects has therefore the ambition and the potential to serve as a foundational theory of metaphysics as well as mathematics and can provide a simple unified axiomatic framework that allows reasoning about a huge variety of concepts throughout the sciences. This makes the attempt to represent the theory using the universal reasoning approach of shallow semantical embeddings outlined in the previous chapter particularly challenging and at the same time rewarding, if successful.

A successful implementation of the theory which allows to utilize the existing sophisticated infrastructure for automated reasoning present in a framework like Isabelle/HOL would not only strongly support the applicability of shallow semantical embeddings as a universal reasoning tool, but could also aid in spreading the utilization of the theory itself as a foundational theory for various scientific fields by enabling convenient interactive and automated reasoning in a verified framework.

2.2. Basic Principles

Although the formal language of the theory is introduced in the next section, some of the basic concepts of the theory are presented in advance to provide further motivation for the formalism.

The following are the two most important principles of the theory (see [13]):

- $\exists\, x(A!x \;\&\; \forall\, F(xF \equiv \varphi))$
- $x = y \equiv \Box\forall\, F(xF \equiv yF)$

The first statement asserts that for every condition on properties φ there exists an abstract object that encodes exactly those properties satisfying φ, whereas the second statement holds for two abstract objects x and y and states that they are equal, if and only if they necessarily encode the same properties.

Together these two principles clarify the notion of abstract objects as the reification of property patterns: Any set of properties is objectified as a distinct abstract object.

Using these principles it is already possible to postulate interesting abstract objects.

For example the Leibnizian concept of an (ordinary) individual u can be defined as *the (unique) abstract object that encodes all properties that u exemplifies*, formally: $\iota x\ A!x \;\&\; \forall\, F\ (xF \equiv Fu)$

Other interesting examples include possible worlds, Platonic Forms or even basic logical objects like truth values. The theory allows to formulate purely *syntactic* definitions of objects like possible worlds and truth values and from these definitions it can be *derived* that there are two truth values or that the application of the modal box operator to a

proposition is equivalent to the proposition being true in all possible worlds (where *being true in a possible world* is again defined syntactically).

This is an impressive property of the Theory of Abstract Objects: it can *syntactically* define objects that are usually only considered semantically.

2.3. The Language of PLM

The target of the embedding is the second-order fragment of object theory as described in chapter 7 of Edward Zalta's upcoming *Principia Logico-Metaphysica* (PLM)[12]. The logical foundation of the theory uses a second-order modal logic (without primitive identity) formulated using relational type theory that is modified to admit *encoding* as a second mode of predication besides the traditional *exemplification*. In the following an informal description of the important aspects of the language is provided; for a detailed and fully formal description and the type-theoretic background refer to the respective chapters of PLM[12].

A compact description of the language can be given in Backus-Naur Form (BNF)[12, Definition (6)], as shown in figure 2.1, in which the following grammatical categories are used:

δ	individual constants
ν	individual variables
Σ^n	n-place relation constants ($n \geq 0$)
Ω^n	n-place relation variables ($n \geq 0$)
α	variables
κ	individual terms
Π^n	n-place relation terms ($n \geq 0$)
Φ^*	propositional formulas
Φ	formulas
τ	terms

The language distinguishes between two types of basic formulas, namely (non-propositional) *formulas* that *may* contain encoding subformulas and *propositional formulas* that *may not* contain encoding subformulas. Only propositional formulas may be used in λ-expressions. The main reason for this distinction will be explained in section 3.2.

Note that there is a case in which propositional formulas *can* contain encoding expressions. This is due to the fact that *subformula* is defined in such a way that xQ is *not* a subformula of $\iota x(xQ)$[3]. Thereby $F\iota x(xQ)$ is a propositional formula and $[\lambda y\ F\iota x(xQ)]$ a well-formed λ-expression. On the other hand xF is not a propositional formula and therefore $[\lambda x\ xF]$ not a well-formed λ-expression. This fact will become relevant in the discussion in section 5.2, that describes a paradox in the formulation of the theory in the draft of PLM at the time of writing[4].

[3] For a formal definition of subformula refer to definition (8) in [12].

[4] At the time of writing several options are being considered that can restore the consistency of the theory while retaining all theorems of PLM.

Figure 2.1.: BNF grammar of the language of PLM[12, p. 170]

	δ	::=	$a_1, a_2, \ldots$
	ν	::=	$x_1, x_2, \ldots$
$(n \geq 0)$	Σ^n	::=	$P^n_1, P^n_2, \ldots$
$(n \geq 0)$	Ω^n	::=	$F^n_1, F^n_2, \ldots$
	α	::=	$\nu \mid \Omega^n \ (n \geq 0)$
	κ	::=	$\delta \mid \nu \mid \iota\nu\varphi$
$(n \geq 1)$	Π^n	::=	$\Sigma^n \mid \Omega^n \mid [\lambda \nu_1 \ldots \nu_n \ \varphi^*]$
	Π^0	::=	$\Sigma^0 \mid \Omega^0 \mid [\lambda \ \varphi^*] \mid \varphi^*$
	φ^*	::=	$\Pi^n \kappa_1 \ldots \kappa_n \ (n \geq 1) \mid \Pi^0 \mid (\neg\varphi^*) \mid (\varphi^* \rightarrow \varphi^*) \mid \forall\alpha\varphi^* \mid (\Box\varphi^*) \mid (\mathcal{A}\varphi^*)$
	φ	::=	$\kappa_1\Pi^1 \mid \varphi^* \mid (\neg\varphi) \mid (\varphi \rightarrow \varphi) \mid \forall\alpha\varphi \mid (\Box\varphi) \mid (\mathcal{A}\varphi)$
	τ	::=	$\kappa \mid \Pi^n \ (n \geq 0)$

Furthermore the theory contains a designated relation constant $E!$ to be read as *being concrete*. Using this constant the distinction between ordinary and abstract objects is defined as follows:

- $O! =_{df} [\lambda x \ \Diamond E!x]$
- $A! =_{df} [\lambda x \ \neg\Diamond E!x]$

So ordinary objects are possibly concrete, whereas abstract objects cannot possibly be concrete.

The language does not contain a primitive identity, but *defines* an identity for each type of term as follows:

ordinary objects	$x =_E y =_{df} O!x \ \& \ O!y \ \& \ \Box(\forall F \ Fx \equiv Fy)$
individuals	$x = y =_{df} x =_E y \vee (A!x \ \& \ A!y \ \& \ \Box(\forall F \ xF \equiv yF))$
one-place relations	$F^1 = G^1 =_{df} \Box(\forall x \ xF^1 \equiv xG^1)$
zero-place relations	$F^0 = G^0 =_{df} [\lambda y \ F^0] = [\lambda y \ G^0]$

The identity for n-place relations for $n \geq 2$ is defined in terms of the identity of one-place relations, see (16)[12] for the full details.

The identity for ordinary objects follows Leibniz' law of the identity of indiscernibles: Two ordinary objects that necessarily exemplify the same properties are identical. Abstract objects, however, are only identical if they necessarily *encode* the same properties. As mentioned in the previous section this goes along with the concept of abstract objects as the reification of property patterns.

Notably the identity for properties has a different definition than one would expect from classical logic. Classically two properties are considered identical if and only if they necessarily are *exemplified* by the same objects. The Theory of Abstract Objects, however, defines two properties to be identical if and only if they are necessarily *encoded* by the same (abstract) objects. This has some interesting consequences that will be

described in more detail in section 2.5 which describes the *hyperintensionality* of relations in the theory.

2.4. The Axioms

Based on the language above, an axiom system is defined that constructs a S5 modal logic with an actuality operator, axioms for definite descriptions that go along with Russell's analysis of descriptions, the substitution of identicals as per the defined identity, α-, β-, η- and a special ι-conversion for λ-expressions, as well as dedicated axioms for encoding. A full accounting of the axioms in their representation in the embedding is found in section 3.10. For the original axioms refer to [12, Chap. 8]. At this point the axioms of encoding are the most relevant, namely:

- $xF \rightarrow \Box xF$
- $O!x \rightarrow \neg\exists\, F\, xF$
- $\exists\, x\ (A!x\ \&\ \forall\, F\ (xF \equiv \varphi))$,
  provided x doesn't occur free in φ

So encoding is modally rigid, ordinary objects do not encode properties and most importantly the comprehension axiom for abstract objects that was already mentioned above:

For every condition on properties φ there exists an abstract object, that encodes exactly those properties, that satisfy φ.

2.5. Hyperintensionality of Relations

An interesting property of the Theory of Abstract Objects results from the definition of identity for one-place relations. Recall that two properties are defined to be identical if and only if they are *encoded* by the same (abstract) objects. The theory imposes no restrictions whatsoever on which properties an abstract object encodes. Let for example F be the property *being the morning star* and G be the property *being the evening star*. Since the morning star and the evening star are actually both the planet Venus, every object that *exemplifies* F will also *exemplify* G and vice-versa: $\Box\forall\, x\ Fx \equiv Gx$. However the concept of being the morning star is different from the concept of being the evening star. The Theory of Abstract Objects therefore does not prohibit the existence of an abstract object that *encodes* F, but does *not* encode G. Therefore by the definition of identity for properties it does *not* hold that $F = G$. As a matter of fact the Theory of Abstract Objects does not force $F = G$ for any F and G. It rather stipulates what needs to be proven, if $F = G$ is to be established, namely that they are necessarily encoded by the same objects. Therefore if two properties *should* be equal in some context an axiom has to be added to the theory that allows to prove that both properties are encoded by the same abstract objects.

The fact that the following relation terms do *not* necessarily denote the same relations illustrates the extent of this *hyperintensionality*:

$$[\lambda y\ p \vee \neg p] \text{ and } [\lambda y\ q \vee \neg q]$$
$$[\lambda y\ p\ \&\ q] \text{ and } [\lambda y\ q\ \&\ p]$$

Of course the theory can be extended in such a way that these properties are equal. However, without additional axioms their equality is not derivable.

Although the relations of object theory are hyperintensional entities, propositional reasoning is still governed by classical extensionality. For example properties that are necessarily exemplified by the same objects can be substituted for each other in an exemplification formula, the law of the excluded middle can be used in propositional reasoning, etc.

The Theory of Abstract Objects is an *extensional* theory of *intensional* entities[12, (130)].

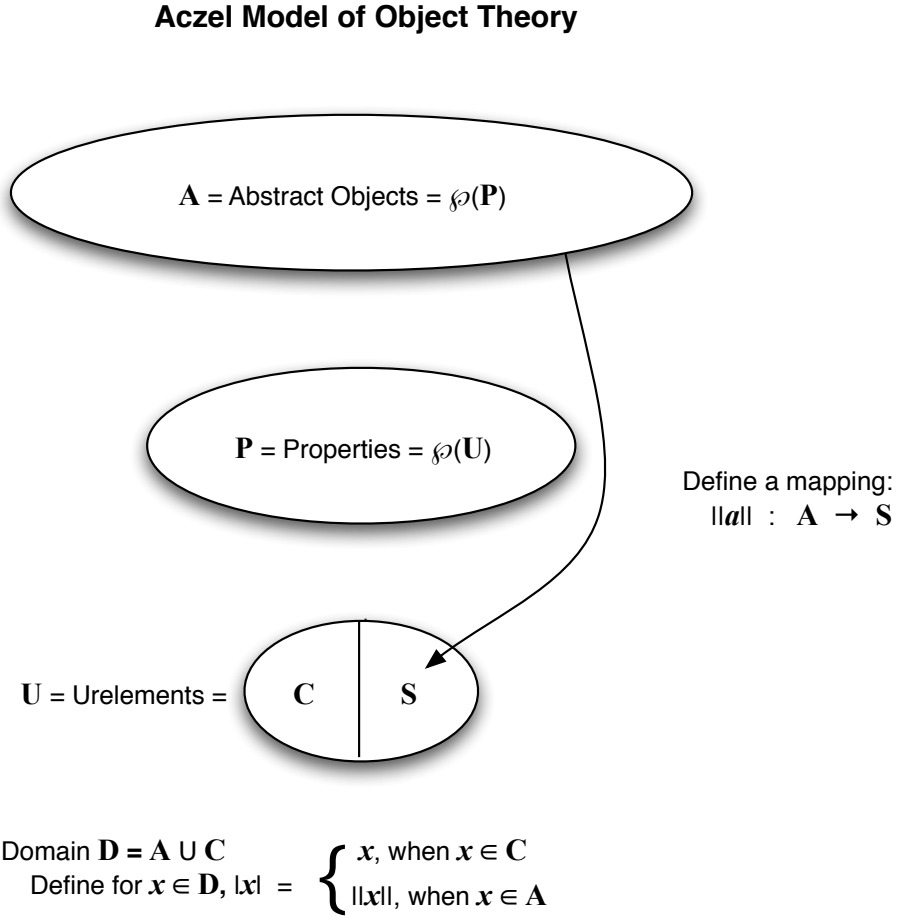
2.6. The Aczel-Model

When thinking about a model for the theory one will quickly notice the following problem: The comprehension axiom for abstract objects implies that for each set of properties there exists an abstract object encoding exactly those properties. Considering the definition of identity there therefore exists an injective map from the power set of properties to the set of abstract objects. On the other hand for an object y the term $[\lambda x\ Rxy]$ constitutes a property. If for distinct abstract objects these properties were distinct, this would result in a violation of Cantor's theorem, since this would mean that there is an injective map from the power set of properties to the set of properties. So does the Theory of Abstract Objects as constructed above have a model? An answer to this question was provided by Peter Aczel[5] who proposed the model structure illustrated in figure 2.2.

In the Aczel-model abstract objects are represented by sets of properties. This of course validates the comprehension axiom of abstract objects. Properties on the other hand are not naively represented by sets of objects, which would lead to a violation of Cantor's theorem, but rather as the sets of *urelements*. Urelements are partitioned into two groups, ordinary urelements (C in the illustration) and special urelements (S in the illustration). Ordinary urelements can serve as the denotations of ordinary objects. Every abstract object on the other hand has a special urelement as its proxy. Which properties an abstract object exemplifies depends solely on its proxy. However, the map from abstract objects to special urelements is not injective; more than one abstract object can share the same proxy. This way a violation of Cantor's theorem is avoided. As a consequence there are abstract objects, that cannot be distinguished by the properties

[5] In fact to our knowledge Dana Scott proposed a first model for the theory before Peter Aczel that we believe is a special case of an Aczel-model with only one *special urelement*.

Figure 2.2.: Illustration of the Aczel-Model, courtesy of Edward Zalta



they exemplify. Interestingly the existence of abstract objects that are exemplification-indistinguishable is a theorem of PLM, see (197)[12].

Although the Aczel-model illustrated in figure 2.2 is non-modal, the extension to a modal version is straightforward by introducing primitive possible worlds as in the Kripke semantics of modal logic.

Further note that relations in the Aczel-model are *extensional*. Since properties are represented as the power set of urelements, two properties are in fact equal if they are

exemplified by the same objects. Consequently statements like $[\lambda\ p \vee \neg p] = [\lambda\ q \vee \neg q]$ are true in the model, although they are not derivable from the axioms of object theory as explained in the previous section.

For this reason an *intensional* variant of the Aczel-model is developed and used as the basis of the embedding. The technicalities of this model are described in the next chapter (see 3.3.1).

3. The Embedding

3.1. The Framework Isabelle/HOL

The embedding is implemented in Isabelle/HOL, that provides a functional higher-order logic that serves as meta-logic. An introduction to Isabelle/HOL can be found in [7][1]. For a general introduction to HOL and its automation refer to [2].

The Isabelle theory containing the embedding is included in the appendix and documented in this chapter. Throughout the chapter references to the various sections of the appendix can be found.

This document itself is generated from a separate Isabelle theory that imports the complete embedding. The terms and theorems discussed throughout this chapter (starting from 3.4) are well-formed terms or valid theorems in the embedding, unless the contrary is stated explicitly. Furthermore the *pretty printing* facility of Isabelle's document generation has been utilized to make it easier to distinguish between the embedded logic and the meta-logic: all expressions that belong to the embedded logic are printed in blue color throughout the chapter.

For technical reasons this color coding could not be used for the raw Isabelle theory in the appendix. Still note the use of bold print for the quantifiers and connectives of the embedded logic.

3.2. A Russell-style Paradox

One of the major challenges of an implementation of the Theory of Abstract Objects in functional logic is the fact that a naive representation of the λ-expressions of the theory using the unrestricted, β-convertible λ-expressions of functional logic results in the following paradox (see [8, pp. 24-25]):

Assume $[\lambda x\ \exists F\ (xF\ \&\ \neg Fx)]$ were a valid λ-expression denoting a relation. Now the comprehension axiom of abstract objects requires the following:

$$\exists x\ (A!x\ \&\ \forall F\ (xF \equiv F = [\lambda x\ \exists F\ (xF\ \&\ \neg Fx)]))$$

So there is an abstract object that encodes only the property $[\lambda x\ \exists F\ (xF\ \&\ \neg Fx)]$. Let b be such an object. Now first assume b exemplifies $[\lambda x\ \exists F\ (xF\ \&\ \neg Fx)]$. By β-reduction this implies that there exists a property, that b encodes, but does not exemplify. Since

[1] An updated version is available at http://isabelle.in.tum.de/doc/tutorial.pdf or in the documentation of the current Isabelle release, see http://isabelle.in.tum.de/.

b only encodes $[\lambda x\ \exists F\ (xF\ \&\ \neg Fx)]$, but does also exemplify it by assumption this is a contradiction.

Now assume b does not exemplify $[\lambda x\ \exists F\ (xF\ \&\ \neg Fx)]$. By β-reduction it follows that there does not exist a property that b encodes, but does not exemplify. Since b encodes $[\lambda x\ \exists F\ (xF\ \&\ \neg Fx)]$ by construction and does not exemplify it by assumption this is again a contradiction.

This paradox is prevented in the formulation of object theory by disallowing encoding subformulas in λ-expressions, so in particular $[\lambda x\ \exists F\ (xF\ \&\ \neg Fx)]$ is not part of the language. However during the construction of the embedding it was discovered that this restriction is not sufficient to prevent paradoxes in general. This is discussed in section 5.2. The solution used in the embedding is described in section 3.4.7.

3.3. Basic Concepts

The introduction mentioned that shallow semantical embeddings were used to successfully represent different varieties of modal logic by implementing them using Kripke semantics. The advantage here is that Kripke semantics is well understood and there are extensive results about its soundness and completeness that can be utilized in the analysis of semantical embeddings (see [3]).

For the Theory of Abstract Objects the situation is different. Section 2.6 already established that even a modal version of the traditional Aczel-model is extensional and therefore theorems are true in it, that are not derivable from the axioms of object theory. On the other hand the last section showed that care has to be taken to ensure the consistency of an embedding of the theory in functional logic.

For this reason the embedding first constructs a hyperintensional version of the Aczel-model that serves as a provably consistent basis for the theory. Then several abstraction layers are implemented on top of the model structure in order to enable reasoning that is independent of the particular representation. These concepts are described in more detail in the following sections.

3.3.1. Hyperintensional Aczel-model

As mentioned in section 2.6 it is straightforward to extend the traditional (non-modal) Aczel-model to a modal version by introducing primitive possible worlds following the Kripke semantics for a modal S5 logic.

Relations in the resulting Aczel-model are, however, still *extensional*. Two relations that are necessarily exemplified by the same objects are equal. The Aczel-model that is used as the basis for the embedding therefore introduces *states* as another primitive besides possible worlds. Truth values are represented as ternary functions from states and possible worlds to booleans; relations as functions from urelements, states and possible worlds to booleans.

Abstract objects are still defined as sets of one-place relations and the division of urelements into ordinary urelements and special urelements, that serve as proxies for abstract objects, is retained as well. Consequently encoding can still be defined as set membership of a relation in an abstract object. Exemplification is defined as function application of a relation to the urelement corresponding to an individual, a state and a possible world.

The semantic truth evaluation of a proposition in a given possible world is defined as its evaluation for a designated *actual state* and the possible world.

Logical connectives are defined to behave classically in the *actual state*, but have undefined behavior in other states.

The reason for this construction becomes apparent if one considers the definition of the identity of relations: relations are considered identical if they are *encoded* by the same abstract objects. In the constructed model encoding depends on the behavior of a relation in all states. Two relations can necessarily be *exemplified* by the same objects in the actual state, but still not be identical, since they can differ in other states. Therefore hyperintensionality of relations is achieved.

The dependency on states is not limited to relations, but introduced to propositions, connectives and quantifiers as well, although the semantic truth conditions of formulas only depend on the evaluation for the actual state. The reason for this is to be able to define λ-expressions (see section 3.4.7) and to extend the hyperintensionality of relations to them. Since the behavior of logical connectives is undefined in states other than the actual state, the behavior of λ-expressions - although classical in the actual state - remains undefined for different states.

In summary, since the semantic truth of a proposition solely depends on its evaluation for the designated actual state, in which the logical connectives are defined to behave classically, the reasoning about propositions remains classical, as desired. On the other hand the additional dependency on states allows a representation of the hyperintensionality of relations.

The technical details of the implementation are described in section 3.4.

3.3.2. Layered Structure

Although the constructed variant of the Aczel-model preserves the hyperintensionality of relations in the theory, it is still known that there are true theorems in this model that are not derivable from the axioms of object theory (see 3.12).

Given this lack of a model with a well-understood degree of soundness and completeness, the embedding uses a different approach than other semantical embeddings, namely the embedding is divided into several *layers* as follows:

- The first layer represents the primitives of PLM using the described hyperintensional and modal variant of the Aczel-model.
- In a second layer the objects of the embedded logic constructed in the first layer are considered as primitives and some of their semantic properties are derived using the background logic as meta-logic.

- The third layer derives the axiom system of PLM mostly using the semantics of the second layer and partly using the model structure directly.
- Based on the third layer the deductive system PLM as described in [12, Chap. 9] is derived solely using the axiom system of the third layer and the fundamental meta-rules stated in PLM. The model structure and the constructed semantics are explicitly not used in any proofs. Thereby the reasoning in this last layer is independent of the first two layers.

The rationale behind this approach is the following: The first layer provides a representation of the embedded logic that is provably consistent. Only minimal axiomatization is necessary, whereas the main construction is purely definitional. Since the subsequent layers don't contain any additional axiomatization (the axiom system in the third layer is *derived*) their consistency is thereby guaranteed as well.

The second layer tries to abstract away from the details of the representation by implementing an approximation of the formal semantics of PLM[2]. The long time goal would be to arrive at the representation of a complete semantics in this layer, that would be sufficient to derive the axiom system in the next layer and which any specific model structure would have to satisfy. Unfortunately this could not be achieved so far, but it was possible to lay some foundations for future work.

At the moment full abstraction from the representation layer is only achieved after deriving the axiom system in the third layer. Still it can be reasoned that in any model of object theory the axiom system has to be derivable and therefore by disallowing all further proofs to rely on the representation layer and model structure directly the derivation of the deductive system PLM is universal. The only exceptions are the primitive meta-rules of PLM: modus ponens, RN (necessitation) and GEN (universal generalization), as well as the deduction rule. These rules do not follow from the axiom system itself, but are derived from the semantics in the second layer (see 3.11.2). Still as the corresponding semantical rules will again have to be derivable for *any* model, this does not have an impact on the universality of the subsequent reasoning.

The technical details of the constructed embedding are described in the following sections.

[2]Our thanks to Edward Zalta for supplying us with a preliminary version of the corresponding unpublished chapter of PLM.

3.4. The Representation Layer

The first layer of the embedding (see A.1) implements the variant of the Aczel-model described in section 3.3.1 and builds a representation of the language of PLM in the logic of Isabelle/HOL. This process is outlined step by step throughout this section.

3.4.1. Primitives

The following primitive types are the basis of the embedding (see A.1.1):

- Type i represents possible worlds in the Kripke semantics.
- Type j represents *states* as described in section 3.3.1.
- Type *bool* represents meta-logical truth values (*True* or *False*) and is inherited from Isabelle/HOL.
- Type ω represents ordinary urelements.
- Type σ represents special urelements.

Two constants are introduced:

- The constant *dw* of type i represents the designated actual world.
- The constant *dj* of type j represents the designated actual state.

Based on the primitive types above the following types are defined (see A.1.2):

- Type o is defined as the set of all functions of type $j \Rightarrow i \Rightarrow bool$ and represents propositions in the embedded logic.
- Type υ is defined as **datatype** $\upsilon = \omega\upsilon\ \omega \mid \sigma\upsilon\ \sigma$. This type represents urelements and an object of this type can be either an ordinary or a special urelement (with the respective type constructors $\omega\upsilon$ and $\sigma\upsilon$).
- Type Π_0 is defined as a synonym for type o and represents zero-place relations.
- Type Π_1 is defined as the set of all functions of type $\upsilon \Rightarrow j \Rightarrow i \Rightarrow bool$ and represents one-place relations (for an urelement a one-place relation evaluates to a truth value in the embedded logic; for an urelement, a state and a possible world it evaluates to a meta-logical truth value).
- Type Π_2 is defined as the set of all functions of type $\upsilon \Rightarrow \upsilon \Rightarrow j \Rightarrow i \Rightarrow bool$ and represents two-place relations.
- Type Π_3 is defined as the set of all functions of type $\upsilon \Rightarrow \upsilon \Rightarrow \upsilon \Rightarrow j \Rightarrow i \Rightarrow bool$ and represents three-place relations.
- Type α is defined as a synonym of the type of sets of one-place relations Π_1 *set*, i.e. every set of one-place relations constitutes an object of type α. This type represents abstract objects.
- Type ν is defined as **datatype** $\nu = \omega\nu\ \omega \mid \alpha\nu\ \alpha$. This type represents individuals and can be either an ordinary urelement of type ω or an abstract object of type α (with the respective type constructors $\omega\nu$ and $\alpha\nu$).

- Type κ is defined as the set of all objects of type ν *option* and represents individual terms. The type *'a option* is part of Isabelle/HOL and consists of a type constructor *Some x* for an object x of type *'a* (in this case type ν) and an additional special element called *None*. *None* is used to represent individual terms that are definite descriptions that are not logically proper (i.e. they do not denote an individual).

Remark. *The Isabelle syntax* **typedef** o = *UNIV*::(*j*⇒*i*⇒*bool*) *set* **morphisms** *eval*o *make*o **..** *found in the theory source in the appendix introduces a new abstract type* o *that is represented by the full set (UNIV) of objects of type* $j \Rightarrow i \Rightarrow bool$. *The morphism eval*o *maps an object of abstract type* o *to its representative of type* $j \Rightarrow i \Rightarrow bool$, *whereas the morphism make*o *maps an object of type* $j \Rightarrow i \Rightarrow bool$ *to the object of type* o *that is represented by it. Defining these abstract types makes it possible to consider the defined types as primitives in later stages of the embedding, once their meta-logical properties are derived from the underlying representation. For a theoretical analysis of the representation layer the type* o *can be considered a synonym of* $j \Rightarrow i \Rightarrow bool$.
The Isabelle syntax **setup-lifting** *type-definition-*o *allows definitions for the abstract type* o *to be stated directly for its representation type* $j \Rightarrow i \Rightarrow bool$ *using the syntax* **lift-definition**.
For the sake of readability in the documentation of the embedding the morphisms are omitted and definitions are stated directly for the representation types[3].

3.4.2. Individual Terms and Definite Descriptions

There are two basic types of individual terms in PLM: definite descriptions and individual variables (and constants). Every logically proper definite description denotes an individual. A definite description is logically proper if its matrix is (actually) true for a unique individual.
In the embedding the type κ encompasses all individual terms, i.e. individual variables, constants *and* definite descriptions. An individual (i.e. a variable or constant of type ν) can be used in place of an individual term of type κ via the decoration $\_^P$ (see A.1.3):

$$x^P = Some\ x$$

The expression x^P (of type κ) is marked to be logically proper (it can only be substituted by objects that are internally of the form *Some x*) and to denote the individual x.
Definite descriptions are defined as follows:

$$\iota x\ .\ \varphi\ x = (if\ \exists!x.\ (\varphi\ x)\ dj\ dw\ then\ Some\ (THE\ x.\ (\varphi\ x)\ dj\ dw)\ else\ None)$$

If the propriety condition of a definite description $\exists!x.\ \varphi\ x\ dj\ dw$ holds, i.e. *there exists a unique x, such that* $\varphi\ x$ *holds for the actual state and the actual world*, the term

[3] The omission of the morphisms is achieved using custom *pretty printing* rules for the document generation facility of Isabelle. The full technical details without these minor omissions can be found in the raw Isabelle theory in the appendix.

$\iota x.\ \varphi\ x$ evaluates to *Some* (*THE* $x.\ \varphi\ x$ *dj dw*). Isabelle's *THE* operator evaluates to the unique object, for which the given condition holds, if there is such a unique object, and is undefined otherwise. If the propriety condition does not hold, the term evaluates to *None*.

The following meta-logical functions are defined to aid in handling individual terms:

- *proper* $x = (None \neq x)$
- *rep* $x = the\ x$

the maps an object of type *'a option* that is of the form *Some* x to x and is undefined for *None*. For an object of type κ the expression *proper* x is true, if the term is logically proper, and if this is the case, the expression *rep* x evaluates to the individual of type ν that the term denotes.

3.4.3. Mapping from Individuals to Urelements

To map abstract objects to urelements (for which relations can be evaluated), a constant $\alpha\sigma$ of type $\alpha \Rightarrow \sigma$ is introduced, which maps abstract objects (of type α) to special urelements (of type σ), see A.1.4.

To assure that every object in the full domain of urelements actually is an urelement for (one or more) individual objects, the constant $\alpha\sigma$ is axiomatized to be surjective.

Now the mapping $\nu\upsilon$ of type $\nu \Rightarrow \upsilon$ can be defined as follows:

$$\nu\upsilon \equiv \textit{case-}\nu\ \omega\upsilon\ (\sigma\upsilon \circ \alpha\sigma)$$

To clarify the syntax note that this is equivalent to the following:

$$(\forall x.\ \nu\upsilon\ (\omega\nu\ x) = \omega\upsilon\ x) \wedge (\forall x.\ \nu\upsilon\ (\alpha\nu\ x) = \sigma\upsilon\ (\alpha\sigma\ x))$$

So ordinary objects are simply converted to an urelements by the type constructor $\omega\upsilon$, whereas for abstract objects the corresponding special urelement under $\alpha\sigma$ is converted to an urelement using the type constructor $\sigma\upsilon$.

Remark. *Future versions of the embedding may introduce a dependency of the mapping from individuals to urelements on states (see 3.12).*

3.4.4. Exemplification of n-place relations

Exemplification of n-place relations can now be defined. Exemplification of zero-place relations is simply defined as the identity, whereas exemplification of n-place relations for $n \geq 1$ is defined to be true, if all individual terms are logically proper and the function application of the relation to the urelements corresponding to the individuals yields true for a given possible world and state (see A.1.5):

- $(\!|p|\!) = p$
- $(\!|F,x|\!) = (\lambda s\ w.\ proper\ x \wedge F\ (\nu\upsilon\ (rep\ x))\ s\ w)$
- $(\!|F,x,y|\!) = (\lambda s\ w.\ proper\ x \wedge proper\ y \wedge F\ (\nu\upsilon\ (rep\ x))\ (\nu\upsilon\ (rep\ y))\ s\ w)$
- $(\!|F,x,y,z|\!) =$
  $(\lambda s\ w.\ proper\ x \wedge$
  $proper\ y \wedge proper\ z \wedge F\ (\nu\upsilon\ (rep\ x))\ (\nu\upsilon\ (rep\ y))\ (\nu\upsilon\ (rep\ z))\ s\ w)$

3.4.5. Encoding

Encoding is defined as follows (see A.1.6):

$$\{\!|x,F|\!\} = (\lambda s\ w.\ proper\ x \wedge (\mathsf{case}\ rep\ x\ \mathsf{of}\ \omega\nu\ \omega \Rightarrow False \mid \alpha\nu\ \alpha \Rightarrow F \in \alpha))$$

For a given state s and a given possible world w it holds that an individual term x encodes F, if x is logically proper, the denoted individual $rep\ x$ is of the form $\alpha\nu\ \alpha$ for some object α (i.e. it is an abstract object) and F is contained in α (recall that abstract objects are defined to be sets of one-place relations).

Encoding is represented as a function of states and possible worlds to ensure type-correctness, but its evaluation does not depend on either. On the other hand whether F is contained in α does depend on the behavior of F in *all* states.

3.4.6. Connectives and Quantifiers

Following the model described in section 3.3.1 the connectives and quantifiers are defined in such a way that they behave classically if evaluated for the designated actual state dj, whereas their behavior is governed by uninterpreted constants in any other state[4].

For this purpose the following uninterpreted constants are introduced (see A.1.7):

- *I-NOT* of type $j \Rightarrow (i \Rightarrow bool) \Rightarrow i \Rightarrow bool$
- *I-IMPL* of type $j \Rightarrow (i \Rightarrow bool) \Rightarrow (i \Rightarrow bool) \Rightarrow i \Rightarrow bool$

Modality is represented using the dependency on primitive possible worlds using a standard Kripke semantics for a S5 modal logic.

The basic connectives and quantifiers are defined as follows (see A.1.7):

- $\neg p = (\lambda s\ w.\ s = dj \wedge \neg\ p\ dj\ w \vee s \neq dj \wedge I\text{-}NOT\ s\ (p\ s)\ w)$
- $p \rightarrow q =$
  $(\lambda s\ w.\ s = dj \wedge (p\ dj\ w \longrightarrow q\ dj\ w) \vee s \neq dj \wedge I\text{-}IMPL\ s\ (p\ s)\ (q\ s)\ w)$
- $\forall_{\nu}\ x\ .\ \varphi\ x = (\lambda s\ w.\ \forall x.\ (\varphi\ x)\ s\ w)$
- $\forall_{0}\ p\ .\ \varphi\ p = (\lambda s\ w.\ \forall p.\ (\varphi\ p)\ s\ w)$
- $\forall_{1}\ F\ .\ \varphi\ F = (\lambda s\ w.\ \forall F.\ (\varphi\ F)\ s\ w)$
- $\forall_{2}\ F\ .\ \varphi\ F = (\lambda s\ w.\ \forall F.\ (\varphi\ F)\ s\ w)$

[4] Early attempts in using an intuitionistic version of connectives and quantifiers based on [6] were found to be insufficient to capture the full hyperintensionality of PLM, but served as inspiration for the current construction.

- $\forall_3\ F\ .\ \varphi\ F = (\lambda s\ w.\ \forall F.\ (\varphi\ F)\ s\ w)$
- $\Box p = (\lambda s\ w.\ \forall v.\ p\ s\ v)$
- $\boldsymbol{\mathcal{A}} p = (\lambda s\ w.\ p\ s\ dw)$

Note in particular that negation and implication behave classically if evaluated for the actual state $s = dj$, but are governed by the uninterpreted constants *I-NOT* and *I-IMPL* for $s \neq dj$:

- $s = dj \Longrightarrow \neg p\ s\ w = (\neg\ p\ s\ w)$
- $s \neq dj \Longrightarrow \neg p\ s\ w = \textit{I-NOT}\ s\ (p\ s)\ w$
- $s = dj \Longrightarrow p \rightarrow q\ s\ w = (p\ s\ w \longrightarrow q\ s\ w)$
- $s \neq dj \Longrightarrow p \rightarrow q\ s\ w = \textit{I-IMPL}\ s\ (p\ s)\ (q\ s)\ w$

Remark. *Future research may conclude that non-classical behavior in states $s \neq dj$ for negation and implication is not sufficient for achieving the desired level of hyperintensionality for λ-expressions. It would be trivial to introduce additional uninterpreted constants to govern the behavior of the remaining connectives and quantifiers in such states as well, though. The remainder of the embedding would not be affected, i.e. no assumption about the behavior of connectives and quantifiers in states other than dj is made in the subsequent reasoning. At the time of writing non-classical behavior for negation and implication is considered sufficient.*

3.4.7. λ-Expressions

The bound variables of the λ-expressions of the embedded logic are individual variables, whereas relations are represented as functions acting on urelements. Therefore the definition of the λ-expressions of the embedded logic is non-trivial. The embedding defines them as follows (see A.1.8):

- $\boldsymbol{\lambda}^0\ p = p$
- $\boldsymbol{\lambda} x.\ \varphi\ x = (\lambda u\ s\ w.\ \exists x.\ \nu\upsilon\ x = u \wedge (\varphi\ x)\ s\ w)$
- $\boldsymbol{\lambda}^2\ (\lambda x\ y.\ \varphi\ x\ y) = (\lambda u\ v\ s\ w.\ \exists x\ y.\ \nu\upsilon\ x = u \wedge \nu\upsilon\ y = v \wedge (\varphi\ x\ y)\ s\ w)$
- $\boldsymbol{\lambda}^3\ (\lambda x\ y\ z.\ \varphi\ x\ y\ z) =$
  $(\lambda u\ v\ r\ s\ w.\ \exists x\ y\ z.\ \nu\upsilon\ x = u \wedge \nu\upsilon\ y = v \wedge \nu\upsilon\ z = r \wedge (\varphi\ x\ y\ z)\ s\ w)$

Remark. *For technical reasons Isabelle only allows λ-expressions for one-place relations to use a nice binder notation. Although better workarounds may be possible, for now the issue is avoided by the use of the primitive λ-expressions of the background logic in combination with the constants $\boldsymbol{\lambda}^2$ and $\boldsymbol{\lambda}^3$ as shown above.*

The representation of zero-place λ-expressions as the identity is straight-forward; the representation of n-place λ-expressions for $n \geq 1$ is illustrated for the case $n = 1$:
The matrix of the λ-expression φ is a function from individuals (of type ν) to truth values (of type o, resp. $j \Rightarrow i \Rightarrow bool$). One-place relations are represented as functions of type $\upsilon \Rightarrow j \Rightarrow i \Rightarrow bool$ though, where υ is the type of urelements.

The λ-expression $\boldsymbol{\lambda} x.\ \varphi\ x$ evaluates to *True* for an urelement u, a state s and a world w, if there is an individual x in the preimage of u under $\nu\upsilon$ and it holds that $\varphi\ x\ s\ w$.

$$\boldsymbol{\lambda} x.\ \varphi\ x\ u\ s\ w = (\exists\, x.\ \nu\upsilon\ x = u \wedge \varphi\ x\ s\ w)$$

If restricted to ordinary objects, the definition can be simplified, since $\nu\upsilon$ is bijective on the set of ordinary objects:

$$\boldsymbol{\lambda} x.\ \varphi\ x\ (\omega\upsilon\ u)\ s\ w = (\varphi\ (\omega\nu\ u))\ s\ w$$

However in general $\nu\upsilon$ can map several abstract objects to the same special urelement, so an analog statement for abstract objects does not hold for arbitrary φ. As described in section 3.2 such a statement would in fact not be desirable, since it would lead to inconsistencies.

Instead the embedding introduces the concept of *proper maps*. A map from individuals to propositions is defined to be proper if its truth evaluation for the actual state only depends on the urelements corresponding to the individuals (see A.1.9):

- $\mathit{IsProperInX}\ \varphi = (\forall\, x\ v.\ (\exists\, a.\ \nu\upsilon\ a = \nu\upsilon\ x \wedge (\varphi\ (a^P))\ dj\ v) = (\varphi\ (x^P))\ dj\ v)$
- $\mathit{IsProperInXY}\ \varphi =$
  $(\forall\, x\ y\ v.$
  $\quad(\exists\, a\ b.\ \nu\upsilon\ a = \nu\upsilon\ x \wedge \nu\upsilon\ b = \nu\upsilon\ y \wedge (\varphi\ (a^P)\ (b^P))\ dj\ v) =$
  $\quad(\varphi\ (x^P)\ (y^P))\ dj\ v)$
- $\mathit{IsProperInXYZ}\ \varphi =$
  $(\forall\, x\ y\ z\ v.$
  $\quad(\exists\, a\ b\ c.$
  $\quad\quad\nu\upsilon\ a = \nu\upsilon\ x \wedge \nu\upsilon\ b = \nu\upsilon\ y \wedge \nu\upsilon\ c = \nu\upsilon\ z \wedge (\varphi\ (a^P)\ (b^P)\ (c^P))\ dj\ v) =$
  $(\varphi\ (x^P)\ (y^P)\ (z^P))\ dj\ v)$

Now by the definition of proper maps the evaluation of λ-expressions behaves as expected for proper φ:

$$\mathit{IsProperInX}\ \varphi = (\forall\, w\ x.\ \boldsymbol{\lambda} x.\ \varphi\ (x^P)\ (\nu\upsilon\ x)\ dj\ w = \varphi\ (x^P)\ dj\ w)$$

Remark. *The right-hand side of the equation above does not quantify over all states, but is restricted to the actual state dj. This is sufficient given that truth evaluation only depends on the actual state and goes along with the desired semantics of λ-expressions (see 3.5.5).*

Maps that contain encoding formulas in their arguments are in general not proper and thereby the paradox mentioned in section 3.2 is prevented.

In fact proper maps are the most general kind of functions that may appear in a lambda-expression, such that β-conversion holds. In what way proper maps correspond to the formulas that PLM allows as the matrix of a λ-expression is a complex question and discussed separately in section 5.1.1.

3.4.8. Validity

Semantic validity is defined as follows (see A.1.10):

$$[\varphi \ in \ v] = \varphi \ dj \ v$$

A formula is considered semantically valid for a possible world v if it evaluates to *True* for the actual state dj and the given possible world v.

Remark. *The Isabelle Theory in the appendix defines the syntax $v \models p$ in the representation layer, following the syntax used in the formal semantics of PLM. The syntax $[p \ in \ v]$ that is easier to use in Isabelle due to bracketing the expression is only introduced after the semantics is derived in A.2.3. For simplicity only the latter syntax is used in this documentation.*

3.4.9. Concreteness

PLM defines concreteness as a one-place relation constant. For the embedding care has to be taken that concreteness actually matches the primitive distinction between ordinary and abstract objects. The following requirements have to be satisfied by the introduced notion of concreteness:

- Ordinary objects are possibly concrete. In the meta-logic this means that for every ordinary object there exists at least one possible world, in which the object is concrete.
- Abstract objects are not possibly concrete.

An additional requirement is enforced by axiom (32.4)[12], see 3.10.7. To satisfy this axiom the following has to be assured:

- Possibly contingent objects exist. In the meta-logic this means that there exists an ordinary object and two possible worlds, such that the ordinary object is concrete in one of the worlds, but not concrete in the other.
- Possibly no contingent objects exist. In the meta-logic this means that there exists a possible world, such that all objects that are concrete in this world, are concrete in all possible worlds.

In order to satisfy these requirements a constant *ConcreteInWorld* is introduced, that maps ordinary objects (of type ω) and possible worlds (of type i) to meta-logical truth values (of type *bool*). This constant is axiomatized in the following way (see A.1.11):

- $\forall x.\ \exists v.\ ConcreteInWorld\ x\ v$
- $\exists x\ v.\ ConcreteInWorld\ x\ v \wedge (\exists w.\ \neg\ ConcreteInWorld\ x\ w)$
- $\exists w.\ \forall x.\ ConcreteInWorld\ x\ w \longrightarrow (\forall v.\ ConcreteInWorld\ x\ v)$

Concreteness can now be defined as a one-place relation:

$$E! = (\lambda u\ s\ w.\ \mathit{case}\ u\ \mathit{of}\ \omega v\ x \Rightarrow ConcreteInWorld\ x\ w \mid \sigma v\ \sigma \Rightarrow False)$$

Whether an ordinary object is concrete is governed by the introduced constant, whereas abstract objects are never concrete.

3.4.10. The Syntax of the Embedded Logic

The embedding aims to provide a readable syntax for the embedded logic that is as close as possible to the syntax of PLM and clearly distinguishes between the embedded logic and the meta-logic. Some concessions have to be made due to the limitations of definable syntax in Isabelle, though. Moreover exemplification and encoding have to use a dedicated syntax in order to be distinguishable from function application.

The syntax for the basic formulas of PLM used in the embedding is summarized in the following table:

PLM	syntax in words	embedded logic	type
φ	it holds that φ	φ	o
$\neg\varphi$	not φ	$\neg\varphi$	o
$\varphi \rightarrow \psi$	φ implies ψ	$\varphi \rightarrow \psi$	o
$\Box\varphi$	necessarily φ	$\Box\varphi$	o
$\mathcal{A}\varphi$	actually φ	$\boldsymbol{\mathcal{A}}\varphi$	o
$\Pi\upsilon$	υ (an individual term) exemplifies Π	$(\!\vert\Pi,\upsilon\vert\!)$	o
Πx	x (an individual variable) exemplifies Π	$(\!\vert\Pi,x^P\vert\!)$	o
$\Pi\upsilon_1\upsilon_2$	υ_1 and υ_2 exemplify Π	$(\!\vert\Pi,\upsilon_1,\upsilon_2\vert\!)$	o
Πxy	x and y exemplify Π	$(\!\vert\Pi,x^P,y^P\vert\!)$	o
$\Pi\upsilon_1\upsilon_2\upsilon_3$	υ_1, υ_2 and υ_3 exemplify Π	$(\!\vert\Pi,\upsilon_1,\upsilon_2,\upsilon_3\vert\!)$	o
Πxyz	x, y and z exemplify Π	$(\!\vert\Pi,x^P,y^P,z^P\vert\!)$	o
$\upsilon\Pi$	υ encodes Π	$\{\!\vert\upsilon,\Pi\vert\!\}$	o
$\iota x\varphi$	*the* x, such that φ	$\boldsymbol{\iota} x.\ \varphi\ x$	κ
$\forall\, x(\varphi)$	for all individuals x it holds that φ	$\forall_\nu x.\ \varphi\ x$	o
$\forall\, p(\varphi)$	for all propositions p it holds that φ	$\forall_0 p.\ \varphi\ p$	o
$\forall\, F(\varphi)$	for all relations F it holds that φ	$\forall_1 F.\ \varphi\ F$	o
		$\forall_2 F.\ \varphi\ F$	
		$\forall_3 F.\ \varphi\ F$	
$[\lambda\ p]$	being such that p	$\boldsymbol{\lambda}^0\ p$	Π_0
$[\lambda x\ \varphi]$	being x such that φ	$\boldsymbol{\lambda} x.\ \varphi\ x$	Π_1
$[\lambda xy\ \varphi]$	being x and y such that φ	$\boldsymbol{\lambda}^2\ (\lambda x\ y.\ \varphi\ x\ y)$	Π_2
$[\lambda xyz\ \varphi]$	being x, y and z such that φ	$\boldsymbol{\lambda}^3\ (\lambda x\ y\ z.\ \varphi\ x\ y\ z)$	Π_3

Several subtleties have to be considered:

- n-place relations are only represented for $n \leq 3$. As the resulting language is already expressive enough to represent the most interesting parts of the theory and it would be trivial to add analog implementations for $n > 3$, this is considered to be sufficient. Future work may attempt to construct a general representation for n-place relations for arbitrary n.
- Individual terms (that can be descriptions) and individual variables, resp. constants have different types. Exemplification and encoding is defined for individual terms of type κ. Individual variables (i.e. variables of type ν) or individual constants (i.e. constants of type ν) can be converted to type κ using the decoration $\_^P$.
- In PLM a general term φ, as it occurs in definite descriptions, quantification formulas and λ-expressions above, can contain *free* variables. If such a term occurs within the scope of a variable binding operator, free occurrences of the variable are considered to be *bound* by the operator. In the embedding this concept is replaced by representing φ as a *function* acting on the bound variables and using the native concept of binding operators in Isabelle.
- The representation layer of the embedding defines a separate quantifier for every type of variable in PLM. This is done to assure that only quantification ranging over these types is part of the embedded language. The definition of a general quantifier in the representation layer could for example be used to quantify over individual *terms* (of type κ), whereas only quantification ranging over individuals (of type ν) is part of the language of PLM. After the semantics is introduced in section 3.5, a *type class* is constructed that is characterized by the semantics of quantification and instantiated for all variable types. This way a general binder that can be used for all variable types can be defined. The details of this approach are explained in section 3.6.

The syntax used for stating that a proposition is semantically valid is the following:

$$[\varphi \ in \ v]$$

Here φ and v are free variables (in the meta-logic). Therefore, stating the expression above as a lemma will implicitly be a quantified statement over all propositions φ and all possible worlds v (unless φ or v are explicitly restricted in the current scope or globally declared as constants).

3.5. Semantic Abstraction

The second layer of the embedding (see A.2) abstracts away from the technicalities of the representation layer and states the truth conditions for formulas of the embedded logic in a similar way as the (at the time of writing unpublished) semantics of object theory.

3.5.1. Domains and Denotation Functions

In order to do so the abstract types introduced in the representation layer κ, o resp. Π_0, Π_1, Π_2 and Π_3 are considered as primitive types and assigned semantic domains: R_κ, R_0, R_1, R_2 and R_3 (see A.2.1.1).

For the embedding the definition of these semantic domains is trivial, since the abstract types of the representation layer are already modeled using representation sets. Therefore the semantic domain for each type can simply be defined as the type of its representatives.

As a next step denotation functions are defined that assign semantic denotations to the objects of each abstract type (see A.2.1.2). The formal semantics of PLM does not a priori assume that every term has a denotation. Therefore, the denotation functions are represented as functions that map to the *option* type of the respective domain. This way they can either map a term to *Some* x, if the term denotes x, or to *None*, if the term does not denote.

In the embedding all relation terms always denote, therefore the denotation functions $d_0, \ldots, d_3$ for relations can simply be defined as the type constructor *Some*. Individual terms on the other hand are already represented by an *option* type, so the denotation function d_κ can be defined as the identity.

Moreover the primitive type of possible worlds i is used as the semantic domain of possible worlds W and the primitive actual world dw as the semantic actual world w_0 (see A.2.1.3).

Remark. *Although the definitions for semantic domains and denotations may seem redundant, conceptually the abstract types of the representation layer now have the role of primitive types. Although for simplicity the last section regarded the type* o *as synonym of* $j \Rightarrow i \Rightarrow$ *bool, it was introduced as a distinct type for which the set of all functions of type* $j \Rightarrow i \Rightarrow$ *bool merely serves as the underlying set of representatives. An object of type* o *cannot directly be substituted for a variable of type* $j \Rightarrow i \Rightarrow$ *bool. To do so it first has to be mapped to its representative of type* $j \Rightarrow i \Rightarrow$ *bool by the use of the morphism* $eval_o$ *that was introduced in the type definition and omitted in the last section for the sake of readability. Therefore although the definitions of the semantic domains and denotation functions may seem superfluous, the domains are different types than the corresponding abstract type and the denotation functions are functions between distinct types (note the use of* **lift-definition** *rather than* **definition** *for the denotation functions in A.2.1.2 that allows to define functions on abstract types in the terms of the underlying representation types).*

3.5.2. Exemplification and Encoding Extensions

Semantic truth conditions for exemplification formulas are defined using *exemplification extensions*. Exemplification extensions are functions relative to semantic possible worlds that map objects in the domain of n-place relations to meta-logical truth values in the case $n = 0$ and sets of n-tuples of objects in the domain of individuals in the case $n \geq 1$. Formally they are defined as follows (see A.2.1.4):

- $ex0\ p\ w = p\ dj\ w$
- $ex1\ F\ w = \{x \mid F\ (\nu\upsilon\ x)\ dj\ w\}$
- $ex2\ R\ w = \{(x, y) \mid R\ (\nu\upsilon\ x)\ (\nu\upsilon\ y)\ dj\ w\}$
- $ex3\ R\ w = \{(x, y, z) \mid R\ (\nu\upsilon\ x)\ (\nu\upsilon\ y)\ (\nu\upsilon\ z)\ dj\ w\}$

The exemplification extension of a 0-place relation is its evaluation for the actual state and the given possible world. The exemplification extension of n-place relations ($n \geq 1$) in a possible world is the set of all (tuples of) *individuals* that are mapped to *urelements* for which the relation evaluates to true for the given possible world and the actual state. This is in accordance with the constructed Aczel-model (see 3.3.1).

Conceptually, exemplification extensions as maps to sets of *individuals* are independent of the underlying model and in particular do not require the concept of *urelements* as they are present in an Aczel-model. Their use in the definition of truth conditions for exemplification formulas below is therefore an abstraction away from the technicalities of the representation layer.

Similarly to the exemplification extension for one-place relations an *encoding extension* is defined as follows (see A.2.1.5):

$$en\ F = \{x \mid \mathsf{case}\ x\ \mathsf{of}\ \omega\nu\ \omega \Rightarrow False \mid \alpha\nu\ y \Rightarrow F \in y\}$$

The encoding extension of a relation is defined as the set of all abstract objects that contain the relation. Since encoding is modally rigid the encoding extension does not need to be relativized for possible worlds.

3.5.3. Truth Conditions of Formulas

Based on the definitions above it is now possible to define truth conditions for the atomic formulas of the language.

For exemplification formulas of n-place relations it suffices to consider the case of one-place relations, for which the truth condition is defined as follows (see A.2.1.7):

$$[(|\Pi,\kappa|)\ in\ w] = (\exists\, r\ o_1.\ Some\ r = d_1\ \Pi \wedge Some\ o_1 = d_\kappa\ \kappa \wedge o_1 \in ex1\ r\ w)$$

The relation term Π is exemplified by an individual term κ in a possible world w if both terms have a denotation and the denoted individual is contained in the exemplification extension of the denoted relation in w. The definitions for n-place relations ($n > 1$) and 0-place relations are analog.

The truth condition for encoding formulas is defined in a similar manner (see A.2.1.8):

$$[\{\!|\kappa,\Pi|\!\}\ in\ w] = (\exists\, r\ o_1.\ Some\ r = d_1\ \Pi \wedge Some\ o_1 = d_\kappa\ \kappa \wedge o_1 \in en\ r)$$

The only difference to exemplification formulas is that the encoding extension does not depend on the possible world w.

The truth conditions for complex formulas are straightforward (see A.2.1.9):

- $[\neg\psi\ in\ w] = (\neg\ [\psi\ in\ w])$
- $[\psi \rightarrow \chi\ in\ w] = (\neg\ [\psi\ in\ w] \vee [\chi\ in\ w])$
- $[\Box\psi\ in\ w] = (\forall\, v.\ [\psi\ in\ v])$
- $[\mathcal{A}\psi\ in\ w] = [\psi\ in\ dw]$
- $[\forall_{\ \nu} x.\ \psi\ x\ in\ w] = (\forall\, x.\ [\psi\ x\ in\ w])$
- $[\forall_{\ 0} x.\ \psi\ x\ in\ w] = (\forall\, x.\ [\psi\ x\ in\ w])$
- $[\forall_{\ 1} x.\ \psi\ x\ in\ w] = (\forall\, x.\ [\psi\ x\ in\ w])$
- $[\forall_{\ 2} x.\ \psi\ x\ in\ w] = (\forall\, x.\ [\psi\ x\ in\ w])$
- $[\forall_{\ 3} x.\ \psi\ x\ in\ w] = (\forall\, x.\ [\psi\ x\ in\ w])$

A negation formula $\neg\psi$ is semantically true in a possible world, if and only if ψ is not semantically true in the given possible world. Similarly truth conditions for implication formulas and quantification formulas are defined canonically.

The truth condition of the modal box operator $\Box\psi$ as ψ being true in all possible worlds, shows that modality follows a S5 logic. A formula involving the actuality operator $\mathcal{A}\psi$ is defined to be semantically true, if and only if ψ is true in the designated actual world.

3.5.4. Denotation of Definite Descriptions

The definition of the denotation of description terms (see A.2.1.10) can be presented in a more readable form by splitting it into its two cases and by using the meta-logical quantifier for unique existence:

- $\exists!x.\ [\psi\ x\ in\ w_0] \Longrightarrow d_\kappa\ \iota x.\ \psi\ x = Some\ (THE\ x.\ [\psi\ x\ in\ w_0])$
- $\nexists!x.\ [\psi\ x\ in\ w_0] \Longrightarrow d_\kappa\ \iota x.\ \psi\ x = None$

If there exists a unique x, such that $\psi\ x$ is true in the actual world, the definite description denotes and its denotation is this unique x. Otherwise the definite description fails to denote.

It is important to consider what happens if a non-denoting definite description occurs in a formula: The only positions in which such a term could occur in a complex formula is in an exemplification expression or in an encoding expression. Given the above truth conditions it becomes clear, that the presence of non-denoting terms does *not* mean that there are formulas without truth conditions: Since exemplification and encoding formulas are defined to be true *only if* the contained individual terms have denotations, such formulas are *False* for non-denoting individual terms.

3.5.5. Denotation of λ-Expressions

The most complex part of the semantic abstraction is the definition of denotations for λ-expressions. The formal semantics of PLM is split into several cases and uses a special class of *Hilbert-Ackermann ε-terms* that are challenging to represent. Therefore a simplified formulation of the denotation criteria is used. Moreover the denotations of λ-expressions are coupled to syntactical conditions. This fact is represented using the notion of *proper maps* as a restriction for the matrix of a λ-expression that was introduced in section 3.4.7. The definitions are implemented as follows (see A.2.1.11):

- $d_1\ \boldsymbol{\lambda} x.\ (\!|\Pi, x^P|\!) = d_1\ \Pi$
- $IsProperInX\ \varphi \Longrightarrow$
  $Some\ r = d_1\ \boldsymbol{\lambda} x.\ \varphi\ (x^P) \wedge Some\ o_1 = d_\kappa\ x \longrightarrow (o_1 \in ex1\ r\ w) = [\varphi\ x\ in\ w]$
- $Some\ r = d_0\ \boldsymbol{\lambda}^0\ \varphi \longrightarrow ex0\ r\ w = [\varphi\ in\ w]$

The first condition for *elementary* λ-expressions is straightforward. The general case in the second condition is more complex: Given that the matrix φ is a proper map, the relation denoted by the λ-expression has the property, that for a denoting individual term x, the denoted individual is contained in its exemplification extension for a possible world w, if and only if $\varphi\ x$ holds in w. At a closer look this is the statement of β-conversion restricted to denoting individuals: the truth condition of the λ-expression being exemplified by some denoting individual term, is the same as the truth condition of the matrix of the term for the denoted individual. Therefore it is clear that the precondition that φ is a proper map is necessary and sufficient. Given this consideration the case for *0*-place relations is straightforward and the cases for $n \geq 2$ are analog to the case $n = 1$.

3.5.6. Properties of the Semantics

The formal semantics of PLM imposes several further restrictions some of which are derived as auxiliary lemmas. Furthermore some auxiliary statements that are specific to the underlying representation layer are proven.

The following auxiliary statements are derived (see A.2.1.12):

1. All relations denote, e.g.
   $\exists\, r.\ Some\ r = d_1\ F$
2. An individual term of the form x^P denotes x:
   $d_\kappa\ x^P = Some\ x$
3. Every ordinary object is contained in the extension of the concreteness property for some possible world:
   $Some\ r = d_1\ E! \Longrightarrow \forall\, x.\ \exists\, w.\ \omega\nu\ x \in ex1\ r\ w$
4. An object that is contained in the extension of the concreteness property in any world is an ordinary object:
   $Some\ r = d_1\ E! \Longrightarrow \forall\, x.\ x \in ex1\ r\ w \longrightarrow (\exists\, y.\ x = \omega\nu\ y)$
5. The denotation functions for relation terms are injective, e.g.

$d_1\ F = d_1\ G \Longrightarrow F = G$

6. The denotation function for individual terms is injective for denoting terms:
   $Some\ o_1 = d_\kappa\ x \wedge Some\ o_1 = d_\kappa\ y \Longrightarrow x = y$

Especially statements 5 and 6 are only derivable due to the specific construction of the representation layer: since the semantic domains were defined as the representation sets of the respective abstract types and denotations were defined canonically, objects that have the same denotation are identical as objects of the abstract type. 3 and 4 are necessary to connect concreteness with the underlying distinction between ordinary and abstract objects in the model.

3.5.7. Proper Maps

The definition of *proper maps* as described in section 3.4.7 is formulated in terms of the meta-logic. Since denotation conditions in the semantics and later some of the axioms have to be restricted to proper maps, a method has to be devised by which the propriety of a map can easily be shown without using meta-logical concepts.

Therefore introduction rules for *IsProperInX*, *IsProperInXY* and *IsProperInXYZ* are derived and a proving method *show-proper* is defined that can be used to proof the propriety of a map using these introduction rules (see A.2.2).

The rules themselves rely on the power of the *unifier* of Isabelle/HOL: Any map acting on individuals that can be expressed by another map that solely acts on exemplification expressions involving the individuals, is shown to be proper. This effectively means that all maps whose arguments only appear in exemplification expressions are proper. Using the provided introduction rules Isabelle's unifier can derive the propriety of such maps automatically.

For a discussion about the relation between this concept and admissible λ-expressions in PLM see section 5.1.1.

3.6. General All-Quantifier

Since the last section established the semantic truth conditions of the specific versions of the all-quantifier for all variable types of PLM, it is now possible to define a binding symbol for general all-quantification.

This is done using the concept of *type classes* in Isabelle/HOL. Type classes define constants that depend on a *type variable* and state assumptions about this constant. In subsequent reasoning the type of an object can be restricted to a type of the introduced type class. Thereby the reasoning can make use of all assumptions that have been stated about the constants of the type class. A priori it is not assumed that any type actually satisfies the requirements of the type class, so initially statements involving types restricted to a type class can not be applied to any specific type.

To allow that the type class has to be *instantiated* for the desired type. This is done by first providing definitions for the constants of the type class specific to the respective type. Then each assumption made by the type class has to be proven given the particular type and the provided definitions. After that any statement that was proven for the type class can be applied to the instantiated type.

In the case of general all-quantification for the embedding this concept can be utilized by introducing the type class *quantifiable* that is equipped with a constant that is used as the general all-quantification binder (see A.3.1). For this constant it can now be assumed that it satisfies the semantic property of all quantification: $[\forall\, x.\ \psi\ x\ in\ w] = (\forall\, x.\ [\psi\ x\ in\ w])$.

Since it was already shown in the last section that the specific all-quantifier for each variable type satisfies this property, the type class can immediately be instantiated for the types ν, Π_0, Π_1, Π_2 and Π_3 (see A.3.2). The instantiation proofs only need to refer to the statements derived in the semantics section for the respective version of the quantifier and are thereby independent of the representation layer.

From this point onward the general all-quantifier can completely replace the type specific quantifiers. This is true even if a quantification is meant to only range over objects of a particular type: In this case the desired type (if it can not implicitly be deduced from the context) can be stated explicitly while still using the general quantifier.

Remark. *Technically it would be possible to instantiate the type class quantifiable for any other type that satisfies the semantic criterion, thereby compromising the restriction of the all-quantifier to the primitive types of PLM. However, this is not done in the embedding and therefore the introduction of a general quantifier using a type class is considered a reasonable compromise.*

3.7. Derived Language Elements

The language of the embedded logic constructed so far is limited to a minimal set of primitive elements. This section introduces further derived language elements that are defined directly in the embedded logic.

Notably identity is not part of the primitive language, but introduced as a *defined* concept.

3.7.1. Connectives

The remaining classical connectives and the modal diamond operator are defined in the traditional manner (see A.4.1):

- $\varphi\ \&\ \psi = \neg(\varphi \rightarrow \neg\psi)$
- $\varphi \vee \psi = \neg\varphi \rightarrow \psi$
- $\varphi \equiv \psi = (\varphi \rightarrow \psi)\ \&\ (\psi \rightarrow \varphi)$
- $\Diamond\varphi = \neg\Box\neg\varphi$

Furthermore, the general all-quantifier is supplemented by an existential quantifier as follows:

- $\exists\ \alpha\ .\ \varphi\ \alpha = \neg(\forall\, \alpha.\ \neg\varphi\ \alpha)$

3.7.2. Identity

The definitions for identity are stated separately for each type of term (see A.4.3):

- $x =_E\ y = (\!|\boldsymbol{\lambda}^2\ (\lambda x\ y.\ (\!|O!,x^P|\!)\ \&\ (\!|O!,y^P|\!)\ \&\ \Box(\forall\, F.\ (\!|F,x^P|\!) \equiv (\!|F,y^P|\!))),x,y|\!)$
- $F =_1\ G = \Box(\forall\, x.\ \{\!|x^P,F|\!\} \equiv \{\!|x^P,G|\!\})$
- $F =_2\ G =$
  $\forall\, x.\ (\boldsymbol{\lambda} y.\ (\!|F,x^P,y^P|\!)) =_1 (\boldsymbol{\lambda} y.\ (\!|G,x^P,y^P|\!))\ \&\ (\boldsymbol{\lambda} y.\ (\!|F,y^P,x^P|\!)) =_1 (\boldsymbol{\lambda} y.\ (\!|G,y^P,x^P|\!))$
- $F =_3\ G = \forall\, x\ y.\ (\boldsymbol{\lambda} z.\ (\!|F,z^P,x^P,y^P|\!)) =_1 (\boldsymbol{\lambda} z.\ (\!|G,z^P,x^P,y^P|\!))\ \&\ (\boldsymbol{\lambda} z.\ (\!|F,x^P,z^P,y^P|\!))$
  $=_1 (\boldsymbol{\lambda} z.\ (\!|G,x^P,z^P,y^P|\!))\ \&\ (\boldsymbol{\lambda} z.\ (\!|F,x^P,y^P,z^P|\!)) =_1 (\boldsymbol{\lambda} z.\ (\!|G,x^P,y^P,z^P|\!))$
- $p =_0 q = (\boldsymbol{\lambda} x.\ p) =_1 (\boldsymbol{\lambda} x.\ q)$

Similarly to the general all-quantifier it makes sense to introduce a general identity relation for all types of terms (κ, o resp. Π_0, Π_1, Π_2, Π_3). However, whereas all-quantification is characterized by a semantic criterion that can be generalized in a type class, identity is defined independently for each type. Therefore a general identity symbol will only be introduced in section 3.9, since it will then be possible to formulate and prove a reasonable property shared by the identity of all types of terms.

3.8. The Proving Method meta_solver

3.8.1. General Concept

Since the semantics in section 3.5 constructed a first abstraction on top of the representation layer, it makes sense to revisit the general concept of the layered structure of the embedding.

The idea behind this structure is that reasoning in subsequent layers should - as far as possible - only rely on the previous layer. However, the restriction of proofs to a specific subset of the facts that are valid in the global context can be cumbersome for automated reasoning. While it is possible to restrict automated reasoning tools to only consider specific sets of facts, it is still an interesting question whether the process of automated reasoning in the layered approach can be made easier.

To that end the embedding utilizes the Isabelle package *Eisbach*. This package allows to conveniently define new proving methods that are based on the systematic application of existing methods.

Remark. *The Eisbach package even allows the construction of more complex proving methods that involve pattern matching. This functionality is utilized in the construction of a substitution method as described in section 3.11.5.*

The idea is to construct a simple resolution prover that can deconstruct complex formulas of the embedded logic to simpler formulas that are connected by a relation in the meta-logic as required by the semantics.

For example an implication formula can be deconstructed as follows:

$$[\varphi \rightarrow \psi \ in \ v] = ([\varphi \ in \ v] \longrightarrow [\psi \ in \ v])$$

Whereas the basic proving methods available in Isabelle cannot immediately prove $[\varphi \rightarrow \varphi \ in \ v]$ without any facts about the definitions of validity and implication, they *can* prove $[\varphi \ in \ v] \longrightarrow [\varphi \ in \ v]$ directly as an instance of $p \longrightarrow p$.

3.8.2. Implementation

Following this idea the method *meta-solver* is introduced (see A.5) that repeatedly applies rules like the above in order to translate complex formulas of the embedded logic to meta-logical statements involving simpler formulas.

The formulation of appropriate introduction, elimination and substitution rules for the logical connectives and quantifiers is straightforward. Beyond that the concept can be used to resolve exemplification and encoding formulas to their semantic truth conditions as well, e.g. (see A.5.10):

$$[(\!|F,x|\!) \ in \ v] = (\exists \, r \ o_1. \ Some \ r = d_1 \ F \wedge Some \ o_1 = d_\kappa \ x \wedge o_1 \in ex1 \ r \ v)$$

This way a large set of formulas can be decomposed to semantic expressions that can be automatically proven without having to rely on the meta-logical definitions directly.

Additionally the *meta-solver* is equipped with rules for being abstract and ordinary and for the defined identity.

Notably the representation layer has the property that the defined identities are equivalent to the identity in the meta-logic. Formally the following statements are true and derived as rules for the *meta-solver*:

- $[x =_E y \ in \ v] =$
  $(\exists \, o_1 \ o_2. \ Some \ (\omega\nu \ o_1) = d_\kappa \ x \wedge Some \ (\omega\nu \ o_2) = d_\kappa \ y \wedge o_1 = o_2)$
- $[x =_\kappa y \ in \ v] = (\exists \, o_1 \ o_2. \ Some \ o_1 = d_\kappa \ x \wedge Some \ o_2 = d_\kappa \ y \wedge o_1 = o_2)$
- $[F =_1 G \ in \ v] = (F = G)$
- $[F =_2 G \ in \ v] = (F = G)$
- $[F =_3 G \ in \ v] = (F = G)$
- $[F =_0 G \ in \ v] = (F = G)$

The proofs for these facts (see A.5.15) are complex and do not solely rely on the properties of the formal semantics of PLM.

The fact that they are derivable has a distinct advantage: since identical terms in the sense of PLM are identical in the meta-logic, proving the axiom of substitution (see 3.10.4) is trivial. A derivation that is solely based on the semantics on the other

hand, would require a complex induction proof. For this reason it is considered a reasonable compromise to include these statements as admissible rules for the *meta-solver*. However, future work may attempt to enforce the separation of layers more strictly and consequently abstain from these rules.

Remark. *Instead of introducing a custom proving method using the Eisbach package, a similar effect could be achieved by instead supplying the derived introduction, elimination and substitution rules directly to one of the existing proving methods like auto or clarsimp. In practice, however, we found that the custom meta-solver produces more reliable results, especially in the case that a proving objective cannot be solved completely by the supplied rules. Moreover the constructed custom proving method serves as a proof of concept and may inspire the development of further more complex proving methods that go beyond a simple resolution prover in the future.*

3.8.3. Applicability

Given the discussion above and keeping the layered structure of the embedding in mind, it is important to precisely determine for which purposes it is valid to use the constructed *meta-solver*.

The main application of the method in the embedding is to support the derivation of the axiom system as described in section 3.10. Furthermore the *meta-solver* can aid in examining the meta-logical properties of the embedding. The *meta-solver* is only supplied with rules that are *reversible*. Thereby it is justified to use it to simplify a statement before employing a tool like **nitpick** in order to look for models or counter-models for a statement.

However it is *not* justified to assume that a theorem that can be proven with the aid of the *meta-solver* method is derivable in the formal system of PLM, since the result still depends on the specific structure of the representation layer. However, based on the concept of the *meta-solver* another proving method is introduced in section 3.11.3, namely the *PLM-solver*. This proving method only employs rules that are derivable from the formal system of PLM itself. Thereby this method *can* be used in proofs without sacrificing the universality of the result.

3.9. General Identity Relation

As already mentioned in section 3.6 similarly to the general quantification binder it is desirable to introduce a general identity relation.

Since the identity of PLM is not directly characterized by semantic truth conditions, but instead *defined* using specific complex formulas in the embedded logic for each type of term, some other property has to be found that is shared by the respective definitions and can reasonably be used as the condition of a type class.

A natural choice for such a condition is the axiom of the substitution of identicals (see 3.10.4). The axiom states that if two objects are identical (in the sense of the defined identity of PLM), then a formula involving the first object implies the formula resulting from substituting the second object for the first object. This inspires the following condition for the type class *identifiable* (see A.6.1):

$$[\alpha = \beta \ in \ v] \wedge [\varphi \ \alpha \ in \ v] \Longrightarrow [\varphi \ \beta \ in \ v]$$

Using the fact that in the last section it was already derived, that the defined identity in the embedded-logic for each term implies the primitive identity of the meta-logical objects, this type class can be instantiated for all types of terms: κ, Π_0 resp. o, Π_1, Π_2, Π_3 (see A.6.2).

Since now general quantification and general identity are available, an additional quantifier for unique existence can be introduced (such a quantifier involves both quantification and identity). To that end a derived type class is introduced that is the combination of the *quantifiable* and the *identifiable* classes. Although this is straightforward for the relation types, this reveals a subtlety involving the distinction between individuals of type ν and individual terms of type κ: The type ν belongs to the class *quantifiable*, the type κ on the other hand does not: no quantification over individual *terms* (that may not denote) was defined. On the other hand the class *identifiable* was only instantiated for the type κ, but not for the type ν. This issue can be solved by noticing that it is straightforward and justified to define an identity for ν as follows:

$$x = y = x^P = y^P$$

This way type ν is equipped with both the general all-quantifier and the general identity relation and unique existence can be defined for all variable types as expected:

$$\exists ! \ \alpha \ . \ \varphi \ \alpha = \exists \, \alpha. \ \varphi \ \alpha \ \& \ (\forall \, \beta. \ \varphi \ \beta \rightarrow \beta = \alpha)$$

Another subtlety has to be considered: at times it is necessary to expand the definitions of identity for a specific type to derive statements in PLM. Since the defined identities were introduced prior to the general identity symbol, such an expansion is therefore so far not possible for a statement that uses the general identity, even if the types are fixed in the context.

To allow such an expansion the definitions of identity are equivalently restated for the general identity symbol and each specific type (see A.6.3). This way the general identity can from this point onward completely replace the type-specific identity symbols.

3.10. The Axiom System of PLM

The last step in abstracting away from the representation layer is the derivation of the axiom system of PLM. Conceptionally the derivation of the axioms is the last moment in which it is deemed admissible to rely on the meta-logical properties of the underlying model structure. Future work may even restrict this further to only allow the use of the properties of the semantics in the proofs (if this is found to be possible).

To be able to distinguish between the axioms and other statements and theorems in the embedded logic they are stated using a dedicated syntax (see A.7):

$$[\![\varphi]\!] = (\forall v.\ [\varphi\ in\ v])$$

Axioms are unconditionally true in all possible worlds. The only exceptions are *necessitation-averse*, resp. *modally-fragile* axioms[5]. Such axioms are stated using the following syntax:

$$[\varphi] = [\varphi\ in\ dw]$$

3.10.1. Axioms as Schemata

Most of the axioms in PLM are stated as *axiom schemata*. They use variables that range over and can therefore be instantiated for any formula and term. Furthermore PLM introduces the notion of *closures* (see [12, (20)]). Effectively this means that the statement of an axiom schema implies that the universal generalization of the schema, the actualization of the schema and (except for modally-fragile axioms) the necessitation of the schema is also an axiom.

Since in Isabelle/HOL free variables in a theorem already range over all terms of the same type no special measures have to be taken to allow instantiations for arbitrary terms. The concept of closures is introduced using the following rules (see A.7.1):

- $[\![\varphi]\!] \Longrightarrow [\varphi\ in\ v]$
- $(\bigwedge x.\ [\![\varphi\ x]\!]) \Longrightarrow [\![\forall x.\ \varphi\ x]\!]$
- $[\![\varphi]\!] \Longrightarrow [\![\mathcal{A}\varphi]\!]$
- $[\![\varphi]\!] \Longrightarrow [\![\Box\varphi]\!]$

For modally-fragile axioms only the following rules are introduced:

- $[\varphi] \Longrightarrow [\varphi\ in\ dw]$
- $(\bigwedge x.\ [\varphi\ x]) \Longrightarrow [\forall x.\ \varphi\ x]$

Remark. *To simplify the instantiation of the axioms in subsequent proofs, a set of* attributes *is defined that can be used to transform the statement of the axioms using the rules defined above.*

This way for example the axiom $[\![\Box\varphi \rightarrow \varphi]\!]$ *can be directly transformed to* $[\forall x.\ \Box\varphi\ x \rightarrow \varphi\ x\ in\ v]$ *by not referencing it directly as qml-2, but by applying the defined attributes to it: qml-2*[*axiom-universal*, *axiom-instance*]

[5] Currently PLM uses only one such axiom, see 3.10.6.

3.10.2. Derivation of the Axioms

To simplify the derivation of the axioms a proving method *axiom-meta-solver* is introduced, that unfolds the dedicated syntax, then applies the meta-solver and if possible resolves the proof objective automatically.

Most of the axioms can be derived by the *axiom-meta-solver* directly. Some axioms, however, require more verbose proofs or their representation in the functional setting of Isabelle/HOL requires special attention. Therefore in the following the complete axiom system is listed and discussed in detail where necessary. Additionally each axiom is associated with the numbering in the current draft of PLM[12].

3.10.3. Axioms for Negations and Conditionals

The axioms for negations and conditionals can be derived automatically and present no further issues (see A.7.2):

- $[\![\varphi \to (\psi \to \varphi)]\!]$ (21.1)
- $[\![\varphi \to (\psi \to \chi) \to (\varphi \to \psi \to (\varphi \to \chi))]\!]$ (21.2)
- $[\![\neg\varphi \to \neg\psi \to (\neg\varphi \to \psi \to \varphi)]\!]$ (21.3)

3.10.4. Axioms of Identity

The axiom of the substitution of identicals can be proven automatically, if additionally supplied with the defining assumption of the type class *identifiable*. The statement is the following (see A.7.3):

- $[\![\alpha = \beta \to (\varphi\ \alpha \to \varphi\ \beta)]\!]$ (25)

3.10.5. Axioms of Quantification

The axioms of quantification are formulated in a way that differs from the statements in PLM, as follows (see A.7.4):

- $[\![(\forall \alpha.\ \varphi\ \alpha) \to \varphi\ \tau]\!]$ (29.1a)
- $[\![(\forall \alpha.\ \varphi\ (\alpha^P)) \to ((\exists \beta.\ \beta^P = \tau) \to \varphi\ \tau)]\!]$ (29.1b)
- $[\![(\forall \alpha.\ \varphi\ \alpha \to \psi\ \alpha) \to ((\forall \alpha.\ \varphi\ \alpha) \to (\forall \alpha.\ \psi\ \alpha))]\!]$ (29.3)
- $[\![\varphi \to (\forall \alpha.\ \varphi)]\!]$ (29.4)
- *SimpleExOrEnc* $\psi \Longrightarrow [\![\psi\ (\iota x.\ \varphi\ x) \to (\exists \nu.\ \nu^P = (\iota x.\ \varphi\ x))]\!]$ (29.5a)
- *SimpleExOrEnc* $\psi \Longrightarrow [\![\psi\ \tau \to (\exists \nu.\ \nu^P = \tau)]\!]$ (29.5b)

The original axioms in PLM[6] are the following:

[6] Note that the axioms will in all likelihood be adjusted in future versions of PLM in order to prevent the paradox described in section 5.2.

- $\forall \alpha \varphi \rightarrow (\exists \beta(\beta = \tau) \rightarrow \varphi^{\tau}{}_{\alpha})$ (29.1)
- $\exists \beta(\beta = \tau)$, provided τ is not a description and β doesn't occur free in τ. (29.2)
- $\forall \alpha(\varphi \rightarrow \psi) \rightarrow (\forall \alpha\ \varphi \rightarrow \forall \alpha\ \psi)$ (29.3)
- $\varphi \rightarrow (\forall \alpha\ \varphi)$, provided α doesn't occur free in φ (29.4)
- $\psi^{\iota x \varphi}{}_{\mu} \rightarrow \exists \nu\ (\nu = \iota x \varphi)$, provided (a) ψ is either an exemplification formula $\Pi^n \kappa_1 \ldots \kappa_n$ $(n \geq 1)$ or an encoding formula $\kappa_1 \Pi^1$, (b) μ is an individual variable that occurs in ψ and only as one or more of the κ_i $(1 \leq i \leq n)$, and (c) ν is any individual variable that doesn't occur free in φ. (29.5)

In the embedding definite descriptions have the type κ that is different from the type for individuals ν. Quantification is only defined for ν, not for κ.

Therefore, the restriction of (29.2) does not apply, since the type restriction of quantification ensures that τ cannot be a definite description. Consequently the inner precondition of (29.1) can be dropped in (29.1a) - since a quantifier is used in the formulation, the problematic case of definite descriptions is excluded and the dropped precondition would always hold.

The second formulation (29.1b) for definite descriptions involves the type conversion $\_^P$ and keeps the inner precondition (since descriptions may not denote).

(29.5b) can be stated as a generalization of (29.5a) to general individual terms, since (29.2) already implies its right hand side for every term except descriptions.

Consequently (29.1b) and (29.5b) can replace the original axioms (29.1) and (29.5) for individual terms. For individual variables and constants as well as relations the simplified formulation (29.1a) can be used instead.

Future work may want to reconsider the reformulation of the axioms, especially considering the most recent developments described in section 5.2. At the time of writing the reformulation is considered a reasonable compromise, since due to the type restrictions of the embedding the reformulated version of the axioms is an equivalent representation of the original axioms.

The predicate *SimpleExOrEnc* used as the precondition for (29.5) is defined as an inductive predicate with the following introduction rules:

- *SimpleExOrEnc* $(\lambda x.\ (\!|F,x|\!))$
- *SimpleExOrEnc* $(\lambda x.\ (\!|F,x,\_|\!))$
- *SimpleExOrEnc* $(\lambda x.\ (\!|F,\_,x|\!))$
- *SimpleExOrEnc* $(\lambda x.\ (\!|F,x,\_,\_|\!))$
- *SimpleExOrEnc* $(\lambda x.\ (\!|F,\_,x,\_|\!))$
- *SimpleExOrEnc* $(\lambda x.\ (\!|F,\_,\_,x|\!))$
- *SimpleExOrEnc* $(\lambda x.\ \{\!|x,F|\!\})$

This corresponds exactly to the restriction of ψ to an exemplification or encoding formula in PLM.

3.10.6. Axioms of Actuality

As mentioned in the beginning of the section the modally-fragile axiom of actuality is stated using a different syntax (see A.7.5):

- $[\mathcal{A}\varphi \equiv \varphi]$ (30)

Note that the model finding tool **nitpick** can find a counter-model for the formulation as a regular axiom, as expected.

The remaining axioms of actuality are not modally-fragile and therefore stated as regular axioms:

- $[\![\ \mathcal{A}\neg\varphi \equiv \neg\mathcal{A}\varphi\ ]\!]$ (31.1)
- $[\![\ \mathcal{A}(\varphi \to \psi) \equiv (\mathcal{A}\varphi \to \mathcal{A}\psi)\ ]\!]$ (31.2)
- $[\![\ \mathcal{A}(\forall\,\alpha.\ \varphi\ \alpha) \equiv (\forall\,\alpha.\ \mathcal{A}\varphi\ \alpha)\ ]\!]$ (31.3)
- $[\![\ \mathcal{A}\varphi \equiv \mathcal{A}\mathcal{A}\varphi\ ]\!]$ (31.4)

All of the above can be proven automatically by the *axiom-meta-solver* method.

3.10.7. Axioms of Necessity

The axioms of necessity are the following (see A.7.6):

- $[\![\ \Box(\varphi \to \psi) \to (\Box\varphi \to \Box\psi)\ ]\!]$ (32.1)
- $[\![\ \Box\varphi \to \varphi\ ]\!]$ (32.2)
- $[\![\ \Diamond\varphi \to \Box\Diamond\varphi\ ]\!]$ (32.3)
- $[\![\ \Diamond(\exists\,x.\ (\!|E!,x^P|\!)\ \&\ \Diamond\neg(\!|E!,x^P|\!))\ \&\ \Diamond\neg(\exists\,x.\ (\!|E!,x^P|\!)\ \&\ \Diamond\neg(\!|E!,x^P|\!))\ ]\!]$ (32.4)

While the first three axioms can be derived automatically, the last axiom requires special attention. On a closer look the formulation may be familiar. The axiom was already mentioned in section 3.4.9 while constructing the representation of the constant $E!$. To be able to derive this axiom here the constant was specifically axiomatized. Consequently the derivation requires the use of these meta-logical axioms stated in the representation layer.

3.10.8. Axioms of Necessity and Actuality

The axioms of necessity and actuality can be derived automatically and require no further attention (see A.7.7):

- $[\![\ \mathcal{A}\varphi \to \Box\mathcal{A}\varphi\ ]\!]$ (33.1)
- $[\![\ \Box\varphi \equiv \mathcal{A}(\Box\varphi)\ ]\!]$ (33.2)

3.10.9. Axioms of Descriptions

There is only one axiom dedicated to descriptions (note, however, that descriptions play a role in the axioms of quantification). The statement is the following (see A.7.8):

- $[\![\ x^P = (\iota x.\ \varphi\ x) \equiv (\forall\, z.\ \mathcal{A}\varphi\ z \equiv z = x)\ ]\!]$ (34)

Given the technicalities of descriptions already discussed in section 3.10.5 it comes at no surprise that this statement requires a verbose proof.

3.10.10. Axioms of Complex Relation Terms

The axioms of complex relation terms deal with the properties of λ-expressions.

Since the *meta-solver* was not equipped with explicit rules for λ-expressions, the statements rely on their semantic properties as described in section 3.5 directly.

The statements are the following (see A.7.9):

- $\boldsymbol{\lambda} x.\ \varphi\ x = \boldsymbol{\lambda} y.\ \varphi\ y$ (36.1)
- *IsProperInX* $\varphi \Longrightarrow [\![\ (\!|\boldsymbol{\lambda} x.\ \varphi\ (x^P), x^P|\!) \equiv \varphi\ (x^P)\ ]\!]$ (36.2)
- *IsProperInXY* $\varphi \Longrightarrow [\![\ (\!|\boldsymbol{\lambda}^2\ (\lambda x\ y.\ \varphi\ (x^P)\ (y^P)), x^P, y^P|\!) \equiv \varphi\ (x^P)\ (y^P)\ ]\!]$ (36.2)
- *IsProperInXYZ* $\varphi \Longrightarrow$
  $[\![\ (\!|\boldsymbol{\lambda}^3\ (\lambda x\ y\ z.\ \varphi\ (x^P)\ (y^P)\ (z^P)), x^P, y^P, z^P|\!) \equiv \varphi\ (x^P)\ (y^P)\ (z^P)\ ]\!]$ (36.2)
- $[\![\ \boldsymbol{\lambda}^0\ \varphi = \varphi\ ]\!]$ (36.3)
- $[\![\ (\boldsymbol{\lambda} x.\ (\!|F, x^P|\!)) = F\ ]\!]$ (36.3)
- $[\![\ \boldsymbol{\lambda}^2\ (\lambda x\ y.\ (\!|F, x^P, y^P|\!)) = F\ ]\!]$ (36.3)
- $[\![\ \boldsymbol{\lambda}^3\ (\lambda x\ y\ z.\ (\!|F, x^P, y^P, z^P|\!)) = F\ ]\!]$ (36.3)
- $(\bigwedge x.\ [\boldsymbol{\mathcal{A}}(\varphi\ x \equiv \psi\ x)\ in\ v]) \Longrightarrow [\![\ \boldsymbol{\lambda}^0\ (\chi\ (\boldsymbol{\iota} x.\ \varphi\ x)) = \boldsymbol{\lambda}^0\ (\chi\ (\boldsymbol{\iota} x.\ \psi\ x))\ ]\!]$ (36.4)
- $(\bigwedge x.\ [\boldsymbol{\mathcal{A}}(\varphi\ x \equiv \psi\ x)\ in\ v]) \Longrightarrow [\![\ (\boldsymbol{\lambda} x.\ \chi\ (\boldsymbol{\iota} x.\ \varphi\ x)\ x) = (\boldsymbol{\lambda} x.\ \chi\ (\boldsymbol{\iota} x.\ \psi\ x)\ x)\ ]\!]$ (36.4)
- $(\bigwedge x.\ [\boldsymbol{\mathcal{A}}(\varphi\ x \equiv \psi\ x)\ in\ v]) \Longrightarrow [\![\ \boldsymbol{\lambda}^2\ (\chi\ (\boldsymbol{\iota} x.\ \varphi\ x)) = \boldsymbol{\lambda}^2\ (\chi\ (\boldsymbol{\iota} x.\ \psi\ x))\ ]\!]$ (36.4)
- $(\bigwedge x.\ [\boldsymbol{\mathcal{A}}(\varphi\ x \equiv \psi\ x)\ in\ v]) \Longrightarrow [\![\ \boldsymbol{\lambda}^3\ (\chi\ (\boldsymbol{\iota} x.\ \varphi\ x)) = \boldsymbol{\lambda}^3\ (\chi\ (\boldsymbol{\iota} x.\ \psi\ x))\ ]\!]$ (36.4)

The first axiom, α-conversion, could be omitted entirely. Since lambda-expressions are modeled using functions with bound variables and α-conversion is part of the logic of Isabelle/HOL, it already holds implicitly.

As explained in section 3.4.7 β-conversion has to be restricted to *proper maps*. In PLM this restriction is implicit due to the fact that λ-expressions are only well-formed if their matrix is a propositional formula.

The formulation of the last class of axioms ((36.4), ι-conversion) has to be adjusted to be representable in the functional setting. The original axiom is stated as follows in PLM:

$$\mathcal{A}(\varphi \equiv \psi) \rightarrow ([\lambda x_1 \cdots x_n\ \chi^*] = [\lambda x_1 \cdots x_n\ \chi^{*\prime}]$$

$\chi^{*\prime}$ is required to be the result of substituting $\iota x\psi$ for zero or more occurrences of $\iota x\varphi$ in χ^*. In the functional setting χ can be represented as function from individual terms of type κ to propositions of type o. Thereby substituting $\iota x\psi$ for occurrences of $\iota x\varphi$ can be expressed by comparing the function application of χ to $\boldsymbol{\iota} x.\ \varphi\ x$ with the function application of χ to $\boldsymbol{\iota} x.\ \psi\ x$.

Since in this representation φ and ψ are functions as well (from type ν to type o) the precondition has to be reformulated to hold for the application of φ and ψ to an arbitrary individual x to capture the concept of $\mathcal{A}(\varphi \equiv \psi)$ in PLM, where φ and ψ may contain x as a free variable.

3.10.11. Axioms of Encoding

The last class of axioms deals with encoding (see A.7.10):

- $[\![\{\!\{x,F\}\!\} \rightarrow \Box\{\!\{x,F\}\!\}]\!]$ (37)
- $[\![(\!|O!,x|\!) \rightarrow \neg(\exists\, F.\ \{\!\{x,F\}\!\})]\!]$ (38)
- $[\![\exists\, x.\ (\!|A!,x^P|\!)\ \&\ (\forall\, F.\ \{\!\{x^P,F\}\!\} \equiv \varphi\ F)]\!]$ (39)

Whereas the first statement, *encoding is modally rigid*, is a direct consequence of the semantics (recall that the encoding extension of a property was not relativized to possible worlds; see section 3.5), the second axiom, *ordinary objects do not encode*, is only derivable by expanding the definition of the encoding extension and the meta-logical distinction between ordinary and abstract objects.

Similarly the comprehension axiom for abstract objects depends on the model structure and follows from the representation of abstract objects as sets of one-place relations and the definition of encoding as set membership.

Furthermore in the functional setting φ has to be represented as a function and the condition it imposes on F is expressed as its application to F. The formulation in PLM on the other hand has to explicitly exclude a free occurrence of x in φ. In the functional setting this is not necessary. Since x is bound by the existential quantifier and not explicitly given to φ as an argument, the condition φ imposes on F cannot depend on x by construction.

3.10.12. Summary

Although some of the axioms have to be adjusted to be representable in the functional environment, the resulting formulation faithfully represents the original axiom system of PLM.

Furthermore a large part of the axioms can be derived independently of the technicalities of the representation layer with proofs that only depend on the representation of the semantics described in section 3.5. Future work may explore available options to further minimize the dependency on the underlying model structure.

To verify that the axiom system faithfully represents the reference system, the deductive system PLM as described in [12, Chap. 9] is derived solely based on the formulation of the axioms without falling back to the model structure or the semantics (see A.9).

3.11. The Deductive System PLM

The derivation of the deductive system PLM ([12, Chap. 9]) from the axiom system constitutes a major part of the Isabelle theory in the appendix (see A.9). Its extent of over one hundred pages makes it infeasible to discuss every aspect in full detail.

Nevertheless it is worthwhile to have a look at the mechanics of the derivation and to highlight some interesting concepts.

3.11.1. Modally Strict Proofs

PLM distinguishes between two sets of theorems: the theorems, that are derivable from the complete axiom system including the modally-fragile axiom, and the set of theorems, that have *modally-strict* proofs (see [12, (42)]).

A proof is modally-strict, if it does not depend on any modally-fragile axioms.

In the embedding modally-strict theorems are stated to be true for an arbitrary semantic possible world: $[\varphi\ in\ v]$

Here the variable v implicitly ranges over all semantic possible worlds of type i, including the designated actual world dw. Since modally-fragile axioms only hold in dw, they therefore cannot be used to prove a statement formulated this way, as desired.

Modally-fragile theorems on the other hand are stated to be true only for the designated actual world: $[\varphi\ in\ dw]$

This way necessary axioms, as well as modally-fragile axioms can be used in their proofs. However it is not possible to infer from a modally-fragile theorem that the same statement holds as a modally-strict theorem.

This representation of modally-strict and modally-fragile theorems is discussed in more detail in section 5.1.3.

3.11.2. Fundamental Metarules of PLM

The primitive rule of PLM is the modus ponens rule (see A.9.2):

- $[\varphi\ in\ v] \wedge [\varphi \rightarrow \psi\ in\ v] \Longrightarrow [\psi\ in\ v]$ (41)

This rule is a direct consequence of the semantics of the implication.

Additionally two fundamental Metarules are derived in PLM, *GEN* and *RN* (see A.9.5):

- $(\bigwedge\alpha.\ [\varphi\ \alpha\ in\ v]) \Longrightarrow [\forall\,\alpha.\ \varphi\ \alpha\ in\ v]$ (49)
- $[\![\bigwedge w.\ [\varphi\ in\ w] \Longrightarrow [\psi\ in\ w];\ [\Box\varphi\ in\ v]]\!] \Longrightarrow [\Box\psi\ in\ v]$ (51)

Although in PLM these rules can be derived by structural induction on the length of a derivation, this proving mechanism cannot be reproduced in Isabelle. However, the rules

are direct consequences of the semantics described in section 3.5. The same is true for the deduction rule (see A.9.6):

- $([\varphi \ in \ v] \Longrightarrow [\psi \ in \ v]) \Longrightarrow [\varphi \rightarrow \psi \ in \ v]$ (54)

Consequently this rule is derived from the semantics as well.

These rules are the *only* exceptions to the concept that the deductive system of PLM is derived solely from the axiom system without relying on the previous layers of the embedding.

3.11.3. PLM Solver

Similarly to the *meta-solver* described in section 3.8 another proving method is introduced, namely the *PLM-solver* (see A.9.1).

This proving method is initially not equipped with any rules. Throughout the derivation of the deductive system, whenever an appropriate rule is derived as part of PLM directly or becomes trivially derivable from the proven theorems, it is added to the *PLM-solver*.

Additionally the *PLM-solver* can instantiate any theorem of the deductive system PLM as well as any axiom, if doing so resolves the current proving goal.

By its construction the *PLM-solver* has the property, that it can *only* prove statements that are derivable from the deductive system PLM. Thereby it is safe to use to aid in any proof throughout the section. In practice it can automatically prove a variety of simple statements and aid in more complex proofs throughout the derivation of the deductive system.

3.11.4. Additional Type Classes

In PLM it is possible to derive statements involving the general identity symbol by case distinction: if such a statement is derivable for all types of terms in the language separately, it can be concluded that it is derivable for the identity symbol in general. Such a case distinction cannot be directly reproduced in the embedding, since it cannot be assumed that every instantiation of the type class *identifiable* is in fact one of the types of terms of PLM.

However, there is a simple way to still formulate such general statements. This is done by the introduction of additional type classes. A simple example is the type class *id-eq* (see A.9.7). This new type class assumes the following statements to be true:

- $[\alpha = \alpha \ in \ v]$ (71.1)
- $[\alpha = \beta \rightarrow \beta = \alpha \ in \ v]$ (71.2)
- $[\alpha = \beta \ \& \ \beta = \gamma \rightarrow \alpha = \gamma \ in \ v]$ (71.3)

Since these statements can be derived *separately* for the types ν, Π_0, Π_1, Π_2 and Π_3, the type class *id-eq* can be instantiated for each of these types.

3.11.5. The Rule of Substitution

A challenge in the derivation of the deductive system that is worth to examine in detail is the *rule of substitution*. The rule is stated in PLM as follows (see (113)[12]):

> If $\vdash_\Box \psi \equiv \chi$ and φ' is the result of substituting the formula χ for zero or more occurrences of ψ where the latter is a subformula of φ, then if $\Gamma \vdash \varphi$, then $\Gamma \vdash \varphi'$. [Variant: If $\vdash_\Box \psi \equiv \chi$, then $\varphi \vdash \varphi'$]

A naive representation of the rule would be the following:

$$(\bigwedge v.\ [\psi \equiv \chi\ \mathit{in}\ v]) \Longrightarrow [\varphi\ \psi\ \mathit{in}\ v] = [\varphi\ \chi\ \mathit{in}\ v]$$

However this statement is *not* derivable. The issue is connected to the restriction of ψ to be a *subformula* of φ in PLM. The formulation above would allow the rule to be instantiated for *any function* φ from formulas to formulas.

Formulas in the embedding have type o which is internally represented by functions of the type $j \Rightarrow i \Rightarrow \mathit{bool}$. Therefore the formulation above could be instantiated with a function φ that has the following internal representation: $\lambda\psi\ s\ w.\ \forall s.\ \psi\ s\ w$

So nothing prevents φ from evaluating its argument for a state different from the designated actual state *dj*. The condition $\bigwedge v.\ [\psi \equiv \chi\ \mathit{in}\ v]$ on the other hand only requires ψ and χ to be (necessarily) equivalent in the *actual state* - no statement about other states is implied.

Another issue arises if one considers one of the example cases of legitimate uses of the rule of substitution in PLM (see [12, (113)]):

> If $\vdash \exists x\ A!x$ and $\vdash_\Box A!x \equiv \neg\Diamond E!x$, then $\vdash \exists x\ \neg\Diamond E!x$.

This would not follow from the naive formulation above, even if it were derivable. Since x is *bound* by the existential quantifier, in the functional representation φ has to have a different type. In the example φ has to be $\lambda\psi.\ \exists x.\ \psi\ x$ which is of type $(\nu \Rightarrow \mathrm{o}) \Rightarrow \mathrm{o}$. ψ and χ have to be functions as well: $\psi = (\lambda x.\ (\!|A!,x|\!))$ and $\chi = (\lambda x.\ \neg\Diamond(\!|E!,x|\!))$. Consequently the equivalence condition for this case has to be reformulated to $\bigwedge x\ v.\ [\psi\ x \equiv \chi\ x\ \mathit{in}\ v]$[7].

Solution

The embedding employs a solution that is complex, but can successfully address the described issues.

The following definition is introduced (see A.9.10):

$$\mathit{Substable}\ \mathit{cond}\ \varphi = (\forall \psi\ \chi\ v.\ \mathit{cond}\ \psi\ \chi \longrightarrow [\varphi\ \psi \equiv \varphi\ \chi\ \mathit{in}\ v])$$

Given a condition *cond* a function φ is considered *Substable*, if and only if for all ψ and χ that satisfy *cond* it follows in each possible world v that $[\varphi\ \psi \equiv \varphi\ \chi\ \mathit{in}\ v]$[8].

[7] This is analog to the fact that x is a free variable in the condition $\vdash_\Box A!x \equiv \neg\Diamond E!x$ in PLM.

[8] ψ and χ can have an arbitrary type. φ is a function from this type to formulas.

Now several introduction rules for this property are derived. The idea is to capture the notion of *subformula* in PLM. A few examples are:

- *Substable cond* $(\lambda\varphi.\ \Theta)$
- *Substable cond* $\psi \Longrightarrow$ *Substable cond* $(\lambda\varphi.\ \neg\psi\ \varphi)$
- *Substable cond* $\psi \wedge$ *Substable cond* $\chi \Longrightarrow$ *Substable cond* $(\lambda\varphi.\ \psi\ \varphi \rightarrow \chi\ \varphi)$

These rules can be derived using theorems of PLM.

As illustrated above in the functional setting substitution has to be allowed not only for formulas, but also for *functions* to formulas. To that end the type class *Substable* is introduced that fixes a condition *Substable-Cond* to be used as *cond* in the definition above and assumes the following:

$$\textit{Substable Substable-Cond}\ \varphi \wedge \textit{Substable-Cond}\ \psi\ \chi \wedge \Theta\ [\varphi\ \psi\ \textit{in}\ v] \Longrightarrow \Theta\ [\varphi\ \chi\ \textit{in}\ v]$$

If φ is *Substable* (as per the definition above) under the condition *Substable-Cond* that was fixed in the type class, and ψ and χ satisfy the fixed condition *Substable-Cond*, then everything that is true for $[\varphi\ \psi\ \textit{in}\ v]$ is also true for $[\varphi\ \chi\ \textit{in}\ v]$.

As a base case this type class is *instantiated* for the type of formulas o with the following definition of *Substable-Cond*:

$$\textit{Substable-Cond}\ \psi\ \chi = (\forall v.\ [\psi \equiv \chi\ \textit{in}\ v])$$

Furthermore the type class is instantiated for *functions* from an arbitrary type to a type of the class *Substable* with the following definition of *Substable-Cond*:

$$\textit{Substable-Cond}\ \psi\ \chi = (\forall x.\ \textit{Substable-Cond}\ (\psi\ x)\ (\chi\ x))$$

Proving Methods

Although the construction above covers exactly the cases in which PLM allows substitutions, it does not yet have a form that allows to conveniently *apply* the rule of substitution. In order to apply the rule, it first has to be established that a formula can be decomposed into a function with the substituents as arguments and it further has to be shown that this function satisfies the appropriate *Substable* condition. This complexity prevents any reasonable use cases. This problem is mitigated by the introduction of proving methods. The main method is called *PLM-subst-method*.

This method uses a combination of pattern matching and automatic rule application to provide a convenient way to apply the rule of substitution in practice.

For example assume the current proof objective is $[\neg\neg\Diamond(\!|E!,x|\!)\ \textit{in}\ v]$. Now it is possible to apply *PLM-subst-method* as follows:

$$\textbf{apply}\ (\textit{PLM-subst-method}\ (\!|A!,x|\!)\ (\neg(\Diamond(\!|E!,x|\!))))$$

The method automatically analyzes the current proving goal, uses pattern matching to find an appropriate choice for a function φ, applies the substitution rule and resolves the substitutability claim about φ.

Consequently it can resolve the current proof objective by producing two new proving goals: $\forall\, v.\ [\!|A!,x|\!) \equiv \neg\Diamond(\!|E!,x|\!)\ in\ v]$ and $[\neg(\!|A!,x|\!)\ in\ v]$, as expected. The complexity of the construction above is hidden away entirely.

Similarly assume the proof objective is $[\exists\, x.\ \neg\Diamond(\!|E!,x^P|\!)\ in\ v]$. Now the method *PLM-subst-method* can be invoked as follows:

$$\textbf{apply}\ (\textit{PLM-subst-method}\ \lambda x\ .\ (\!|A!,x^P|\!)\ \lambda x\ .\ (\neg(\Diamond(\!|E!,x^P|\!))))$$

This will result in the new proving goals: $\forall\, x\ v.\ [(\!|A!,x^P|\!) \equiv \neg\Diamond(\!|E!,x^P|\!)\ in\ v]$ and $[\exists\, x.\ (\!|A!,x^P|\!)\ in\ v]$, as desired.

Conclusion

Although an adequate representation of the rule of substitution in the functional setting is challenging, the above construction allows a convenient use of the rule. Moreover it is important to note that despite the complexity of the representation no assumptions about the underlying model structure were made. The construction is completely derivable from the rules of PLM itself, so the devised rule is safe to use without compromising the provability claim of the layered structure of the embedding.

All statements that are proven using the constructed substitution methods, remain derivable from the deductive system of PLM.

3.11.6. An Example Proof

To illustrate how the derivation of theorems in the embedding works in practice, consider the following example[9]:

lemma $[\Box(\varphi \to \Box\varphi) \to ((\neg\Box\varphi) \equiv (\Box(\neg\varphi)))\ in\ v]$
proof (*rule CP*)
  assume $[\Box(\varphi \to \Box\varphi)\ in\ v]$
  hence $[(\neg\Box(\neg\varphi)) \equiv \Box\varphi\ in\ v]$
    by (*metis sc-eq-box-box-1 diamond-def vdash-properties-10*)
  thus $[((\neg\Box\varphi) \equiv (\Box(\neg\varphi)))\ in\ v]$
    by (*meson CP* $\equiv I$ $\equiv E$ $\neg\neg I$ $\neg\neg E$)
qed

Since the statement is an implication it is derived using a *conditional proof*. To that end the proof statement already applies the initial rule *CP*.

The proof objective inside the proof body is now $[\Box(\varphi \to \Box\varphi)\ in\ v] \Longrightarrow [\neg\Box\varphi \equiv \Box\neg\varphi\ in\ v]$, so $[\neg\Box\varphi \equiv \Box\neg\varphi\ in\ v]$ has to be shown under the assumption $[\Box(\varphi \to \Box\varphi)\ in\ v]$. Therefore the first step is to assume $[\Box(\varphi \to \Box\varphi)\ in\ v]$.

[9] Since the whole proof is stated as raw Isabelle code, unfortunately no color-coding can be applied.

The second statement can now be automatically derived using the previously proven theorem *sc-eq-box-box-1*, the definition of the diamond operator and a deduction rule. The final proof objective follows from a combination of introduction and elimination rules.

The automated reasoning tool **sledgehammer** can find proofs for the second and final statement automatically. It can even automatically find a proof for the entire theorem resulting in the following one-line proof:

lemma $[\Box(\varphi \rightarrow \Box\varphi) \rightarrow ((\neg\Box\varphi) \equiv (\Box(\neg\varphi)))\ in\ v]$
by (*metis* $\equiv I$ *CP* $\equiv E(1)$ $\equiv E(2)$ *raa-cor-1 sc-eq-box-box-1 diamond-def*)

So it can be seen that the embedding can be used to interactively prove statements with the support of automated reasoning tools and often even complete proofs for complex statements can be found automatically.

3.11.7. Summary

A full representation of the deductive system PLM, as described in [12, Chap. 9], could be derived without violating the layered structure of the embedding.

Although compromises affecting the degree of automation had to be made, the resulting representation can conveniently be used for the interactive construction of complex proofs while retaining the support of the automation facilities of Isabelle/HOL.

3.12. Artificial Theorems

The layered approach of the embedding provides the means to derive theorems independently of the representation layer and model structure. It is still interesting to consider some examples of theorems that are *not* part of PLM, but can be derived in the embedding using its meta-logical properties.

3.12.1. Non-Standard λ-Expressions

The following statement involves a λ-expressions that contains encoding subformulas and is consequently not part of PLM (see A.11):

$$[(\!|\boldsymbol{\lambda}x.\ \{\!\{F^P,y\}\!\},x^P|\!) \equiv \{\!\{F^P,y\}\!\}\ in\ v]$$

In this case traditional β-conversion still holds, since the λ-expression does not contain encoding expressions involving its bound variable[10]. On the other hand the following is *not* a theorem in the embedding (the tool **nitpick** can find a counter-model):

$$[(\!|\boldsymbol{\lambda}x.\ \{\!\{x^P,F\}\!\},x^P|\!) \rightarrow \{\!\{x^P,F\}\!\}\ in\ v]$$

[10] Consequently the matrix is a *proper map*.

Instead the following generalized versions of β-conversion are theorems:

- $[(\!|\boldsymbol{\lambda} x.\ \{\!|x^P,F|\!\},z^P|\!)\ in\ v] = (\exists\, y.\ \nu\upsilon\ y = \nu\upsilon\ z \wedge [\{\!|y^P,F|\!\}\ in\ v])$
- $[(\!|\boldsymbol{\lambda} x.\ \varphi\ (x^P),z^P|\!)\ in\ v] = (\exists\, y.\ \nu\upsilon\ y = \nu\upsilon\ z \wedge [\varphi\ (y^P)\ in\ v])$

These theorems can be equivalently stated purely in the embedded logic:

- $[(\!|\boldsymbol{\lambda} x.\ \{\!|x^P,F|\!\},z^P|\!) \equiv (\exists\, y.\ (\forall\, F.\ (\!|F,z^P|\!) \equiv (\!|F,y^P|\!))\ \&\ \{\!|y^P,F|\!\})\ in\ v]$
- $[(\!|\boldsymbol{\lambda} x.\ \varphi\ (x^P),z^P|\!) \equiv (\exists\, y.\ (\forall\, F.\ (\!|F,z^P|\!) \equiv (\!|F,y^P|\!))\ \&\ \varphi\ (y^P))\ in\ v]$

The second statement shows that in general λ-expressions in the embedding have a *non-standard* semantics. As a special case, however, the behavior of λ-expressions is classical if restricted to proper maps, which is due to the following theorem[11]:

$$IsProperInX\ \varphi \Longrightarrow [(\exists\, y.\ (\forall\, F.\ (\!|F,x^P|\!) \equiv (\!|F,y^P|\!))\ \&\ \varphi\ (y^P)) \equiv \varphi\ (x^P)\ in\ v]$$

As a consequence of the generalized β-conversion there are theorems in the embedding involving λ-expressions that *do* contain encoding subformulas in the bound variable, e.g.:

$$[(\!|\boldsymbol{\lambda} x.\ \{\!|x^P,F|\!\} \equiv \{\!|x^P,F|\!\},y^P|\!)\ in\ v]$$

This topic is discussed in more detail in section 5.1.1.

3.12.2. Consequences of the Aczel-model

Independently the following theorem is a consequence of the constructed Aczel-model:

$$[\forall\, F.\ (\!|F,a^P|\!) \equiv (\!|F,b^P|\!)\ in\ v] \Longrightarrow \boldsymbol{\lambda} x.\ (\!|R,x^P,a^P|\!) = \boldsymbol{\lambda} x.\ (\!|R,x^P,b^P|\!)$$

The reason for this theorem to hold is that the condition on a and b forces the embedding to map both objects to the same urelement. By the definition of exemplification the presented λ-expressions only depend on this urelement, therefore they are forced to be equal. Neither the deductive system of PLM nor its formal semantics require this equality.

Initial research suggests that this artificial theorem can be avoided by extending the embedding in the following way: the mapping from abstract objects to special urelements constructed in section 3.4.3 can be modified to depend on states. This way the condition used in the theorem only implies that a and b are mapped to the same urelement in the *actual state*. Since they can still be mapped to different urelements in different states, the derived equality no longer follows.

This extension of the embedding increases the complexity of the representation layer slightly, but its preliminary analysis suggests that it presents no further issues, so future versions of the embedding will in all likelihood include such a modification.

[11] Note that for propositional formulas an equivalent statement is derivable in PLM as well.

3.13. Sanity Tests

The consistency of the constructed embedding can be verified by the model-finding tool **nitpick** (see A.12.1). Since the main construction of the embedding is definitional and only a minimal set of meta-logical axioms is used, this is expected.

The hyperintensionality of the constructed model can be verified for some simple example cases. The following statements have counter-models (see A.12.2):

- $[(\boldsymbol{\lambda} y.\ q \vee \neg q) = (\boldsymbol{\lambda} y.\ p \vee \neg p)\ in\ v]$
- $[(\boldsymbol{\lambda} y.\ p \vee q) = (\boldsymbol{\lambda} y.\ q \vee p)\ in\ v]$

Furthermore the meta-logical axioms stated in section 3.4.9 can be justified (see A.12.4):

- $(\forall x.\ \exists v.\ ConcreteInWorld\ x\ v) =$
  $(\forall y.\ [(\!|\boldsymbol{\lambda} u.\ \neg\Box\neg(\!|E!,u^P|\!),y^P|\!)\ in\ v] = (\mathsf{case}\ y\ \mathsf{of}\ \omega\nu\ z \Rightarrow True \mid \alpha\nu\ z \Rightarrow False))$
- $(\forall x.\ \exists v.\ ConcreteInWorld\ x\ v) =$
  $(\forall y.\ [(\!|\boldsymbol{\lambda} u.\ \Box\neg(\!|E!,u^P|\!),y^P|\!)\ in\ v] = (\mathsf{case}\ y\ \mathsf{of}\ \omega\nu\ z \Rightarrow False \mid \alpha\nu\ z \Rightarrow True))$
- $(\exists x\ v.\ ConcreteInWorld\ x\ v \wedge (\exists w.\ \neg\ ConcreteInWorld\ x\ w)) =$
  $[\neg\Box(\forall x.\ (\!|E!,x^P|\!) \rightarrow \Box(\!|E!,x^P|\!))\ in\ v]$
- $(\exists w.\ \forall x.\ ConcreteInWorld\ x\ w \longrightarrow (\forall v.\ ConcreteInWorld\ x\ v)) =$
  $[\neg\Box\neg(\forall x.\ (\!|E!,x^P|\!) \rightarrow \Box(\!|E!,x^P|\!))\ in\ v]$

The first axiom is equivalent to the fact that concreteness matches the domains of ordinary, resp. abstract objects, whereas the second and third axiom correspond to the conjuncts of axiom (32.4)[12].

Remark. *Additionally some further desirable meta-logical properties of the embedding are verified in A.12.5 and A.12.6.*

4. Technical Limitations of Isabelle/HOL

Although the presented embedding shows that the generic proof assistant Isabelle/HOL offers a lot of flexibility in expressing even a very complex and challenging theory as the Theory of Abstract Objects, it has some limitations that required compromises in the formulation of the theory.

In this chapter some of these limitations and their consequences for the embedding are discussed. Future versions of Isabelle may allow a clearer implementation especially of the layered approach of the embedding.

4.1. Limitations of Type Classes and Locales

Isabelle provides a powerful tool for abstract reasoning called **locale**. Locales are used for *parametric* reasoning. Type classes, as already described briefly in section 3.6 and further mentioned in sections 3.9 and 3.11.4, are in fact special cases of locales that are additionally connected to Isabelle's internal type system.

The definition of a locale defines a set of constants that can use arbitrary type variables[1]. Assumptions about these constants can be postulated that can be used in the reasoning within the context of the locale. Similarly to the instantiation of a type class, a locale can be *interpreted* for specific definitions of the introduced constants, if it can be proven that the postulated assumptions are satisfied for the interpretation.

Thereby it is possible to reason about abstract structures that are solely characterized by a specific set of assumptions. Given that it can be shown that these assumptions are satisfied for a concrete case, an interpretation of the locale allows the use of all theorems shown for the abstract case in the concrete application.

Therefore in principle locales would be a perfect fit for the layered structure of the embedding: If the representation of the formal semantics and the axiom system could both be formulated as locales, it could first be shown that the axiom system is a *sublocale* of the formal semantics, i.e. every set of constants that satisfies the requirements of the formal semantics also satisfies the requirements of the axiom system, and further the formal semantics could be interpreted for a concrete model structure.

Since the reasoning within a locale cannot use further assumptions that are only satisfied by a specific interpretation, this way the universality of the reasoning based on the axiom system could be formally guaranteed - no proof that is solely based on the axiom locale

[1] Type classes on the other hand are restricted to only one type variable.

could use any meta-logical statement tied to the underlying representation layer and model structure[2].

However, a major issue arises when trying to formulate the axiom system as a locale. Constants in a locale have to be introduced with a fixed type. Although this type can use type variables, e.g. $'a \Rightarrow 'a \Rightarrow 'o$, the type variable $'a$ is fixed throughout the locale. This makes it impossible to introduce a general binder for all-quantification or a general identity symbol in a single axiom locale that could be used for the statement of the axioms of quantification and the substitution of identicals.

Several solutions to this problem could be considered: the identity relation could be introduced as a polymorphic constant *outside the locale* and the locale could assume some properties for this constant for specific type variables. Before interpreting the locale the polymorphic constant could then be *overloaded* for concrete types in order to be able to satisfy the assumptions. However, it would still be impossible to prove a general statement about identity: every statement would have to be restricted to a specific type, because in general no assumptions about the properties of identity could be made.

Another solution would be to refrain from using general quantifiers and identity relations altogether, but to introduce separate binders and identity symbols for the type of individuals and each relation type. However, this would add a significant amount of notational complexity and would require to duplicate all statements that hold for quantification and identity in general for every specific type. Statements ranging over multiple types would even have to be stated for every possible combination of types separately.

It could also be considered to introduce the axioms of quantification and identity separately from the axiom locale in a type class. An interpretation of the complete axiom system would then have to interpret the axiom locale, as well as instantiate the respective type classes. Since type classes can only use one type variable, this would make it impossible to use a type variable for truth values in the definition of the respective type classes, though. Consequently it is unclear how appropriate assumptions for such type classes could be formulated. Using separate locales instead of type classes would be connected with different issues.

Several other concepts were considered during the construction of the embedding, but no solution was found that would both accurately represent the axiom system and still be notationally convenient.

The most natural extension of Isabelle's locale system that would solve the described issues, would be the ability to introduce polymorphic constants in a locale that can be restricted to a type class (resp. a *sort*). The type class could potentially even be introduced simultaneously with the locale. However, such a construction is currently not possible in Isabelle and as of yet it is unknown whether the internal type system of Isabelle would allow such an extension in general.

[2] Although the construction of chapter 3 provides the means for universal reasoning that is independent of a model as well, it depends on *fair use* of the provided layer structure.

4.2. Case Distinctions by Type

Although a general identity relation can be represented using type classes as described in sections 3.6 and 3.9, this construction differs from the concept used in PLM. The identity relation of PLM is not determined by some set of properties, but by its definition for the specific concrete types.

Isabelle does not allow the restriction of a type variable in a statement to a specific set of types. Type variables can only be restricted to specific *sorts*, so effectively to type classes. As mentioned in section 3.11.4, this means that statements about the general identity relation, that depend on the specific definitions for the concrete types, cannot be proven as in PLM by case distinction on types. Instead additional type classes have to be introduced that *assume* the statements and then have to be instantiated for the concrete types.

Although this construction involves some technical overhead, the solution is elegant and provides a flexible representation for such general statements.

4.3. Structural Induction and Proof-Theoretic Reasoning

As mentioned in section 3.11.2, some of the meta-rules that PLM can derive by induction on the length of a derivation, have to be proven using the semantics instead in the embedding, e.g. the deduction theorem $([\varphi \ in \ v] \Longrightarrow [\psi \ in \ v]) \Longrightarrow [\varphi \rightarrow \psi \ in \ v]$.

While the derivation of these fundamental rules using the semantics is justified, it would be interesting to investigate whether the proof-theoretic reasoning PLM uses in these cases can be reproduced in Isabelle/HOL. A related topic is the representation of the concept of *modally-strict proofs* as described in sections 3.11.1 and 5.1.3.

5. Discussion and Results

5.1. Differences between the Embedding and PLM

Although the embedding attempts to represent the language and logic of PLM as precisely as possible, there remain some differences between PLM and its representation in Isabelle/HOL. Some of the known differences are discussed in the following sections. A complete analysis of the precise relation between PLM and the embedding unfortunately goes beyond the scope of this thesis and will only be possible after PLM has recovered from the discovered paradox (see 5.2). Such an analysis will be a highly interesting and relevant topic for future research.

5.1.1. Propositional Formulas and λ-Expressions

The main difference between the embedding and PLM is the fact that the embedding does not distinguish between propositional and non-propositional formulas.

This purely syntactic distinction is challenging to reproduce in a shallow embedding that does not introduce the complete term structure of the embedded language directly. Instead the embedding attempts to analyze the semantic reason for the syntactic distinction and to devise a semantic criterion that can be used as a replacement for the syntactic restriction.

The identified issue, that is addressed by the distinction in PLM, is described in section 3.2: Allowing non-propositional formulas in β-convertible λ-expressions without restriction leads to paradoxes.

Since the embedding is known to be consistent, the issue presents itself in a slightly different fashion: the paradox is constructed under the assumption that β-conversion holds unconditionally for all λ-expressions. In the embedding on the other hand in general λ-expressions have a *non-standard* semantics and β-conversion only follows as a special case (see 3.12.1). Thereby the consistency of the system is preserved.

With the definition of *proper maps* (see 3.4.7), the embedding constructs a necessary and sufficient condition on functions that may serve as matrix of a λ-expression while allowing β-conversion.

The idea is that every λ-expression that is syntactically well-formed in PLM should have a proper map as its matrix. Two subtleties have to be considered, though:

It was discovered that there are λ-expressions which are part of PLM, whose matrix does not correspond to a proper map in the embedding. The analysis of this issue led to the discovery of a paradox in the formulation of PLM and is discussed in more detail

in section 5.2. As a consequence these cases will not constitute proper λ-expressions in future versions of PLM.

The remaining subtlety is the fact that there are proper maps, that do not correspond to propositional formulas. Some examples have already been mentioned in section 3.12.1. Therefore the embedding suggests that the theory of PLM can be consistently extended to include a larger set of proper, β-convertible λ-expressions. Since the set of relations of PLM already has to be adjusted to prevent the discovered paradox, such an extension presents a viable option.

Once PLM has recovered from the paradox, future research can consider available options to align the set of relations present in the embedding with the resulting set of relations of the new version of PLM.

5.1.2. Terms and Variables

In PLM an individual term can be an individual variable, an individual constant or a definite description. A large number of statements is formulated using specific object-language variables instead of metavariables ranging over arbitrary terms. From such a statement its universal generalization can be derived using the rule GEN, which then can be instantiated for any individual term, given that it denotes ($\exists\beta\ \beta = \tau$).

As already mentioned in sections 3.4.2 and 3.10.5 the embedding uses a slightly different approach: In the embedding individuals and individual terms have different *types*.

The technicalities of this approach and a discussion about the accuracy of this representation were already given in the referenced sections, so at this point it suffices to summarize the resulting differences between the embedding and PLM:

- The individual variables of PLM are represented as variables of type ν in the embedding.
- Individual constants can be represented by declaring constants of type ν.
- Meta-level variables (like τ) ranging over all individual terms in PLM can be represented as variables of type κ.
- Objects of type ν have to be explicitly converted to objects of type κ using the decoration $\_^P$, if they are to be used in a context that allows general individual terms.
- The axioms of quantification are adjusted to go along with this representation (see 3.10.5).

In PLM the situation for relation variables, constants and terms is analog. However, the embedding uses the following simplification in order to avoid the additional complexity introduced for individuals:

Since at the time of writing PLM unconditionally asserts $\exists\beta\ \beta = \tau$ for any relation term by an axiom, the embedding uses only one type Π_n for each arity of relations. Therefore no special type conversion between variables and terms is necessary and every relation term can immediately be instantiated for a variable of type Π_n. This hides the additional

steps PLM employs for such instantiations (the generalization by GEN followed by an instantiation using quantification theory). Since $\exists \beta\ \beta = \tau$ holds unconditionally for relation terms, this simplification is justified.

However, the recent developments described in section 5.2 suggest that $\exists \beta\ \beta = \tau$ will in all likelihood no longer hold unconditionally for every relation term in future versions of PLM. Therefore, future versions of the embedding will have to include a distinction between relation terms and relation variables in a similar way as is already done for individuals. An alternative approach that could result in a more elegant representation would be to implement concepts of free logic based on the research in [4] for both individuals and relations.

5.1.3. Modally-strict Proofs and the Converse of RN

As described in section 3.11.1 modally-strict theorems in the embedding are stated in the form $[\varphi\ in\ v]$, so they are stated to be semantically true for an arbitrary possible world v.

Modally-strict theorems in PLM are defined using a proof-theoretic concept: modally-strict proofs are not allowed to use modally-fragile axioms. They are solely derived from axioms whose necessitations are axioms as well (see 3.10.1).

The metarule RN states in essence that if there is a modally-strict proof for φ, then $\Box\varphi$ is derivable as a theorem. PLM proves this fact by induction on the length of the derivation. Remark (185)[12] gives an example of a case in which the converse is false: if $\Box\varphi$ is derivable as a theorem, this does not imply that there is a modally-strict proof for φ.

However, in the embedding the following is derivable from the semantics of the box operator:

$$[\Box\varphi\ in\ dw] \Longrightarrow \forall v.\ [\varphi\ in\ v]$$

So although the converse of RN is not true in PLM, an equivalent statement for theorems of the form $[\varphi\ in\ v]$ in the embedding can be derived from the semantics.

The modally-strict theorems of PLM are a subset of a larger class of theorems, namely the theorems that are *necessarily true*. Semantically a statement of the form $[\varphi\ in\ v]$ in the embedding is derivable, whenever φ is a *necessary theorem*.

Unfortunately there is no semantic criterion that allows to decide whether a statement is a necessary theorem or a modally-strict theorem. Therefore, the embedding has to express modally-strict theorems as necessary theorems, for which the converse of RN is in fact true.

This still does not compromise the claim that any statement that is derived in A.9 is also derivable in PLM: the basis for this claim is that no proofs in this layer may rely on the meta-logical properties of the embedding, but only the fundamental meta-rules of PLM are allowed to derive theorems from the axioms. Since the converse of RN is

neither a fundamental meta-rule of PLM, nor derivable without using the semantics, it is not stated as an admissible rule for these proofs. Thereby it is guaranteed that no statement of the form $[\varphi\ in\ v]$ is derived that is not a modally-strict theorem of PLM.

Unfortunately this has the consequence that the proving method *PLM-solver* cannot be equipped with a reversible elimination rule for the box operator, which reduces its power as a proving method. However, preserving the claim that theorems derived in the embedding are also theorems of PLM even when restricting to modally-strict theorems was given preference over an increased level of automation.

5.2. A Paradox in PLM

During the analysis of the constructed embedding it was discovered that the formulation of the theory in PLM at the time of writing allowed paradoxical constructions.

This section first describes the process that led to the discovery of the paradox and the role the embedding played in it, after which the construction of the paradox is outlined in the language of PLM.

The paradox has since been confirmed by Edward Zalta and a vivid discussion about its repercussions and possible solutions has developed. At the time of writing it has become clear that there are several options to recover from the paradox while in essence retaining the full set of theorems of PLM. So far no final decision has been reached about which option will be implemented in future versions of PLM.

5.2.1. Discovery of the Paradox

The discovery of the paradox originates in the analysis of the concept of *proper maps* in the embedding and its relation to propositional formulas in PLM, which are the only formulas PLM allows as the matrix of λ-expressions (see 5.1.1).

While trying to verify the conjecture, that the matrix of every λ-expression allowed in PLM corresponds to a proper map in the embedding, it was discovered, that λ-expressions of the form $[\lambda y\ F\iota x(y[\lambda z\ Rxz])]$ in which the bound variable y occurs in an encoding formula inside the matrix of a definite description, were part of PLM, but their matrix was *not* a proper map in the embedding and therefore β-conversion was not derivable for these terms.

Further analysis showed that a modification of the embedding which would allow β-conversion for such expressions, would have to involve a restriction of the Aczel-model (in particular of the map from abstract objects to urelements).

In order to understand how the Aczel-model could be adequately restricted, the consequences of allowing β-conversion in the mentioned cases *by assumption* were studied in the embedding. This led to the first proof of inconsistency (see A.13.4):

$$(\bigwedge G\ \varphi.\ IsProperInX\ (\lambda x.\ (\!|G,\iota y.\ \varphi\ y\ x|\!))) \Longrightarrow False$$

Under the assumption that $\lambda x.\ (\!|G,\iota y.\ \varphi\ y\ x|\!)$ is a proper map for arbitrary G and φ, *False* is derivable in the embedding. However λ-expressions with the equivalent of such maps as matrix were in fact part of PLM.

Since the inconsistency can be derived without relying on the meta-logical properties of the embedding, it was immediately possible to translate the proof back to the language of PLM. The resulting formulation then served as the basis for further discussions with Edward Zalta.

Since then the issue leading to the paradox was identified as the *description backdoor* (see A.13.2) that can be used to construct a variety of paradoxical cases, e.g. the paradox described in section 3.2 can be reconstructed. This refined version of the paradox is used in the inconsistency proof in A.13.3 and is outlined in the language of PLM in the next section. The general situation leading to the paradox is repeated without referring to the particularities of the embedding.

5.2.2. Construction using the Language of PLM

Object theory distinguishes between propositional and non-propositional formulas. Propositional formulas are not allowed to contain encoding subformulas, so for example $\exists F\ xF$ is not propositional. Only propositional formulas can be the matrix of a λ-expression, so $[\lambda x\ \exists F\ xF]$ is not a valid term of the theory - it is excluded syntactically.

The reason for this is that considering $[\lambda x\ \exists F\ xF\ \&\ \neg Fx]$ a valid, denoting λ-expression for which β-conversion holds would result in a paradox as described in section 3.2.

Excluding non-propositional formulas in λ-expressions was believed to be sufficient to prevent such inconsistencies. This was shown to be incorrect, though.

The problem is the *description backdoor*. The term $[\lambda y\ F\iota x\psi]$ is well-formed, even if ψ is *not* propositional. This is due to the definition of *subformula*: ψ is *not* a subformula of $F\iota x\psi$, so ψ *may* contain encoding subformulas itself and $F\iota x\psi$ is still a propositional formula.

This was deemed to be no problem and for cases like $[\lambda y\ F\iota x(xG)]$ as they are mentioned and used in PLM this is indeed true.

It had not been considered that y may appear within the matrix of such a description and more so, it may appear in an encoding expression, for example $[\lambda y\ F\iota x(xG\ \&\ yG)]$ is still a propositional formula.

Therefore, the following construction is possible:

$$[\lambda y\ [\lambda z\ \forall p(p\rightarrow p)]\iota x(x = y\ \&\ \psi)] \tag{1}$$

Here ψ can be an arbitrary non-propositional formula in which x and y may be free and (1) is still a valid, denoting λ-expression for which β-conversion holds.

By β-conversion and description theory the following is derivable:

$$[\lambda y\ [\lambda z\ \forall p(p\rightarrow p)]\iota x(x = y\ \&\ \psi)]x \equiv \psi^{x}{}_{y} \tag{2}$$

Remark. *Using a modally-strict proof only the following is derivable:*
$[\lambda y\ [\lambda z\ \forall p(p \rightarrow p)]\iota x(x = y\ \&\ \psi)]x \equiv \mathcal{A}\psi^x{}_y$
For the construction of the paradox, the modally-fragile statement is sufficient. However, it is possible to construct similar paradoxical cases without appealing to any modally-fragile axioms or theorems as well.

This effectively undermines the intention of restricting λ-expressions to only propositional formulas:

Although $[\lambda x\ \exists F\ xF\ \&\ \neg Fx]$ is not part of the language, it is possible to formulate the following instead:

$$[\lambda y\ [\lambda z\ \forall p(p \rightarrow p)]\iota x(x = y\ \&\ (\exists F\ yF\ \&\ \neg Fy))] \tag{3}$$

If one considers (2) now, one can see that this λ-expressions behaves exactly the way that $[\lambda x\ \exists F\ xF\ \&\ \neg Fx]$ would, if it were part of the language, i.e. the result of β-reduction for $[\lambda x\ \exists F\ xF\ \&\ \neg Fx]$ would be the same as the right hand side of (2) when applied to (3). Therefore, the λ-expression in (3) can be used to reproduce the paradox described in section 3.2.

5.2.3. Possible Solutions

Fortunately no theorems were derived in PLM, that actually use problematic λ-expressions as described above. Therefore, it is possible to recover from the paradox without losing any theorems. At the time of writing, it seems likely that a concept of *proper* λ-expressions will be introduced to the theory and only *proper* λ-expressions will be forced to have denotations and allow β-conversion. Problematic λ-expressions that would lead to paradoxes, will not be considered *proper*. Several options are available to define the propriety of λ-*expressions* and to adjust PLM in detail.

As a consequence the purely syntactical distinction between propositional and non-propositional formulas is no longer sufficient to guarantee that every relation term has a denotation. The embedding of the theory shows that an adequate definition of *proper* λ-*expressions* can consistently replace this distinction entirely yielding a broader set of relations. The philosophical implications of such a radical modification of the theory have not yet been analyzed entirely though, and at the time of writing it is an open question whether such a modification may be implemented in future versions of PLM.

5.3. A Meta-Conjecture about Possible Worlds

A conversation between Bruno Woltzenlogel Paleo and Edward Zalta about the Theory of Abstract Objects led to the following meta-conjecture:

“ For every syntactic possible world w, there exists a semantic point p which is the denotation of w. ”[1]

Since the embedding constructs a representation of the semantics of PLM, it was possible to formally analyze the relationship between syntactic and semantic possible worlds and arrive at the following theorems (see A.10):

- $\forall x.\ [\mathit{PossibleWorld}\ (x^P)\ \mathit{in}\ w] \longrightarrow (\exists v.\ \forall p.\ [x^P \models p\ \mathit{in}\ w] = [p\ \mathit{in}\ v])$
- $\forall v.\ \exists x.\ [\mathit{PossibleWorld}\ (x^P)\ \mathit{in}\ w] \wedge (\forall p.\ [p\ \mathit{in}\ v] = [x^P \models p\ \mathit{in}\ w])$

The first statement shows that for every *syntactic* possible world x there is a *semantic* possible world v, such that a proposition is syntactically true in x, if and only if it is semantically true in v.

The second statement shows that for every *semantic* possible world v there is a *syntactic* possible world x, such that a proposition is semantically true in v, if and only if it is *syntactically* true in x.

This result extends the following theorems already derived syntactically in PLM (w is restricted to only range over syntactic possible worlds):

- $\Diamond p \equiv \exists w(w \models p)$ (433.1)
- $\Box p \equiv \forall w(w \models p)$ (433.2)

Whereas the syntactic statements of PLM already show the relation between the modal operators and syntactic possible worlds, the semantic statements derived in the embedding show that there is in fact a natural bijection between syntactic and semantic possible worlds.

This example shows that a semantical embedding allows a detailed analysis of the semantical properties of a theory and to arrive at interesting meta-logical results.

5.4. Functional Object Theory

The first and foremost goal of the presented work was to show that the second-order fragment of the Theory of Abstract Objects as described in PLM can be represented in functional higher-order logic using a shallow semantical embedding.

As a result a theory was constructed in Isabelle/HOL that - although its faithfulness is yet to be formally verified - is most likely able to represent and verify all reasoning in the target theory. A formal analysis of the faithfulness of the embedding is unfortunately not possible at this time, since the theory of PLM first has to be adjusted to prevent the discovered paradox. Depending on the precise modifications of PLM the embedding will have to be adjusted accordingly, after which the question can be revisited.

The embedding goes to great lengths to construct a restricted environment, in which it is possible to derive new theorems that can easily be translated back to the reference system of PLM. The fact that the construction of the paradox described in section 5.2

[1] This formulation originates in the resulting e-mail correspondence between Bruno Woltzenlogel Paleo and Christoph Benzmüller.

could be reproduced in the target logic, strongly indicates the merits and success of this approach.

Independently of the relation between the embedding and the target system, a byproduct of the embedding is a working functional variant of object theory that deserves to be studied in its own right. To that end future research may want to drop the layered structure of the embedding and dismiss all constructions that solely serve to restrict reasoning in the embedding in order to more closely reproduce the language of PLM. Automated reasoning in the resulting theory will be significantly more powerful and the interesting properties of the original theory, that result from the introduction of abstract objects and encoding, can still be preserved.

5.5. Relations vs. Functions

As mentioned in the introduction, Oppenheimer and Zalta argue that relational type theory is more fundamental than functional type theory (see [8]). One of their main arguments is that the Theory of Abstract Objects is not representable in functional type theory. The success of the presented embedding, however, suggests that the topic has to be examined more closely.

Their result is supported by the presented work in the following sense: it is impossible to represent the Theory of Abstract Objects by representing its λ-expressions directly as primitive λ-expressions in functional logic. Furthermore, exemplification cannot be represented classically as function application, while at the same time introducing encoding as a second mode of predication.

This already establishes that the traditional approach of translating relational type theory to functional type theory in fact fails for the Theory of Abstract Objects. A simple version of functional type theory, that only involves two primitive types (for individuals and propositions), is insufficient for a representation of the theory.

The embedding does not share several of the properties of the representative functional type theory constructed in [8, pp. 9-12]:

- Relations are *not* represented as functions from individuals to propositions.
- Exemplification is *not* represented as simple function application.
- The λ-expressions of object theory are *not* represented as primitive λ-expressions.

To illustrate the general schema that the embedding uses instead assume that there is a primitive type for each arity of relations R_n. Let further ι be the type of individuals and o be the type of propositions. The general construct is now the following:

- Exemplification (of an n-place relation) is a function of type $R_n \Rightarrow \iota \Rightarrow \ldots \Rightarrow \iota \Rightarrow \mathrm{o}$.
- Encoding is a function of type $\iota \Rightarrow R_1 \Rightarrow \mathrm{o}$.
- To represent λ-expressions functions Λ_n of type $(\iota \Rightarrow \ldots \Rightarrow \iota \Rightarrow \mathrm{o}) \Rightarrow R_n$ are introduced. The λ-expression $[\lambda x_1 \ldots x_n\ \varphi]$ of object theory is represented as $\Lambda_n[\lambda x_1 \ldots x_n\ \varphi]$.

The Theory of Abstract Objects restricts the matrix of λ-expressions to propositional formulas, so not all functions of type $\iota \Rightarrow \ldots \Rightarrow \iota \Rightarrow o$ are supposed to denote relations. However, since in classical functional type theory functions are total, Λ_n has to map all these functions to some object of type R_n. To solve this problem concepts used in the embedding of free logic can help[2]. The function Λ_n can map functions of type $\iota \Rightarrow \ldots \Rightarrow \iota \Rightarrow o$ that do not correspond to propositional formulas to objects of type R_n that represent invalid (resp. non-existing) relations. For invalid relations the functions used to represent encoding and exemplification can be defined to map to an object of type o that represents invalid propositions.

Oppenheimer and Zalta argue that using a free logic and letting non-propositional formulas fail to denote is not an option, since it prevents classical reasoning for non-propositional formulas[3]. Although this is true for the case of a simple functional type theory, it does not apply to the constructed theory: since only objects of type R_n may fail to denote, non-propositional reasoning is unaffected.

Remark. *Although the constructed functional type theory is based on the general structure of the presented embedding, instead of introducing concepts of free logic, λ-expressions involving non-propositional formulas are assigned* non-standard *denotations, i.e. they do denote, but β-conversion only holds under certain conditions (see 5.1.1). Although this concept has merits as well, future versions of the embedding may instead utilize the concepts described in [4] to replace this construction by a free logic implementation that will more closely reflect the concepts of propositional formulas and λ-expressions in object theory.*

The constructed theory can represent the relations and λ-expressions of object theory, as well as exemplification and encoding. Furthermore, the embedding shows that it has a model and that an adequate intensional interpretation of propositions can be used to preserve the desired hyperintensionality of relations in λ-expressions.

In summary it can be concluded that a representation of object theory in functional type theory is feasible, although it is connected with a fair amount of complexity (i.e. the introduction of additional primitive types and the usage of concepts of intensional and free logic). On the other hand, whether this result contradicts the philosophical claim that relations are more fundamental than functions, is still debatable considering the fact that the proposed construction has to introduce new primitive types for relations[4] and the construction is complex in general. Further it has to be noted that so far only the second-order fragment of object theory has been considered and the full type-theoretic version of the theory may present further challenges.

[2] See the embedding of free logic constructed in [4].

[3] See [8, pp. 30-31].

[4] Note, however, that the embedding can represent relations as functions acting on urelements following the Aczel-model.

5.6. Conclusion

The presented work shows that shallow semantical embeddings in HOL have the potential to represent even highly complex theories that originate in a fundamentally different tradition of logical reasoning (e.g. relational instead of functional type theory). The presented embedding represents the most ambitious project in this area so far and its success clearly shows the merits of the approach.

Not only could the embedding uncover a previously unknown paradox in the formulation of its target theory, but it could contribute to the understanding of the relation between functional and relational type theory and provide further insights into the general structure of the target theory, its semantics and possible models. It can even show that a consistent extension of the theory is possible that can increase its expressibility.

For the field of mathematics an analysis of chapters 14 and 15 of PLM, that construct natural numbers and theoretical mathematical objects and relations in object theory, is of particular interest. The embedding can be a significant aid in the study of these chapters, since the properties of the derived objects and relations can immediately be analyzed and verified using the extensive library for abstract mathematical reasoning already present in Isabelle/HOL as a reference.

The presented work introduces novel concepts that can benefit future endeavors of semantical embeddings in general: a layered structure allows the representation of a target theory without extensive prior results about its model structure and provides the means to comprehensively study potential models. Custom proving methods can benefit automated reasoning in an embedded logic and provide the means to reproduce even complex deductive rules of a target system in a user-friendly manner.

The fact that the embedding can construct a verified environment which allows to conveniently prove and verify theorems in the complex target system while retaining the support of automated reasoning tools, shows the great potential of semantical embeddings in providing the means for a productive interaction between humans and computer systems.

A. Isabelle Theory

A.1. Representation Layer

A.1.1. Primitives

typedecl i — possible worlds
typedecl j — states

consts $dw :: i$ — actual world
consts $dj :: j$ — actual state

typedecl ω — ordinary objects
typedecl σ — special urelements
datatype $\upsilon = \omega\upsilon\ \omega \mid \sigma\upsilon\ \sigma$ — urelements

A.1.2. Derived Types

typedef o = $UNIV$::$(j \Rightarrow i \Rightarrow bool)$ set
  morphisms $eval$o $make$o **..** — truth values

type-synonym Π_0 = o — zero place relations
typedef Π_1 = $UNIV$::$(\upsilon \Rightarrow j \Rightarrow i \Rightarrow bool)$ set
  morphisms $eval\Pi_1$ $make\Pi_1$ **..** — one place relations
typedef Π_2 = $UNIV$::$(\upsilon \Rightarrow \upsilon \Rightarrow j \Rightarrow i \Rightarrow bool)$ set
  morphisms $eval\Pi_2$ $make\Pi_2$ **..** — two place relations
typedef Π_3 = $UNIV$::$(\upsilon \Rightarrow \upsilon \Rightarrow \upsilon \Rightarrow j \Rightarrow i \Rightarrow bool)$ set
  morphisms $eval\Pi_3$ $make\Pi_3$ **..** — three place relations

type-synonym $\alpha = \Pi_1$ set — abstract objects

datatype $\nu = \omega\nu\ \omega \mid \alpha\nu\ \alpha$ — individuals

typedef κ = $UNIV$::$(\nu\ option)$ set
  morphisms $eval\kappa$ $make\kappa$ **..** — individual terms

setup-lifting *type-definition*-o
setup-lifting *type-definition*-κ
setup-lifting *type-definition*-Π_1
setup-lifting *type-definition*-Π_2
setup-lifting *type-definition*-Π_3

A.1.3. Individual Terms and Definite Descriptions

lift-definition $\nu\kappa :: \nu \Rightarrow \kappa$ $(-^P\ [90]\ 90)$ **is** $Some$ **.**
lift-definition $proper :: \kappa \Rightarrow bool$ **is** $op \neq\ None$ **.**
lift-definition $rep :: \kappa \Rightarrow \nu$ **is** the **.**

lift-definition $that$::$(\nu \Rightarrow$o$) \Rightarrow \kappa$ (**binder** $\boldsymbol{\iota}$ $[8]\ 9$) **is**
  $\lambda\ \varphi\ .\ if\ (\exists!\ x\ .\ (\varphi\ x)\ dj\ dw)$
    $then\ Some\ (THE\ x\ .\ (\varphi\ x)\ dj\ dw)$

else None .

A.1.4. Mapping from Individuals to Urelements

consts $\alpha\sigma :: \alpha \Rightarrow \sigma$
axiomatization where $\alpha\sigma$*-surj*: *surj* $\alpha\sigma$
definition $\nu\upsilon :: \nu \Rightarrow \upsilon$ **where** $\nu\upsilon \equiv$ *case-*ν $\omega\upsilon$ $(\sigma\upsilon \circ \alpha\sigma)$

A.1.5. Exemplification of n-place-Relations.

lift-definition *exe0*$::\Pi_0 \Rightarrow o$ $(\langle\!|\text{-}|\!\rangle)$ **is** *id* .
lift-definition *exe1*$::\Pi_1 \Rightarrow \kappa \Rightarrow o$ $(\langle\!|\text{-},\text{-}|\!\rangle)$ **is**
$\lambda\ F\ x\ s\ w$. $(proper\ x) \wedge F\ (\nu\upsilon\ (rep\ x))\ s\ w$.
lift-definition *exe2*$::\Pi_2 \Rightarrow \kappa \Rightarrow \kappa \Rightarrow o$ $(\langle\!|\text{-},\text{-},\text{-}|\!\rangle)$ **is**
$\lambda\ F\ x\ y\ s\ w$. $(proper\ x) \wedge (proper\ y) \wedge$
$F\ (\nu\upsilon\ (rep\ x))\ (\nu\upsilon\ (rep\ y))\ s\ w$.
lift-definition *exe3*$::\Pi_3 \Rightarrow \kappa \Rightarrow \kappa \Rightarrow \kappa \Rightarrow o$ $(\langle\!|\text{-},\text{-},\text{-},\text{-}|\!\rangle)$ **is**
$\lambda\ F\ x\ y\ z\ s\ w$. $(proper\ x) \wedge (proper\ y) \wedge (proper\ z) \wedge$
$F\ (\nu\upsilon\ (rep\ x))\ (\nu\upsilon\ (rep\ y))\ (\nu\upsilon\ (rep\ z))\ s\ w$.

A.1.6. Encoding

lift-definition *enc* $:: \kappa \Rightarrow \Pi_1 \Rightarrow o$ $(\{\!|\text{-},\text{-}|\!\})$ **is**
$\lambda\ x\ F\ s\ w$. $(proper\ x) \wedge$ *case-*ν $(\lambda\ \omega$. *False*$)$ $(\lambda\ \alpha$. $F \in \alpha)$ $(rep\ x)$.

A.1.7. Connectives and Quantifiers

consts *I-NOT* $:: j \Rightarrow (i \Rightarrow bool) \Rightarrow i \Rightarrow bool$
consts *I-IMPL* $:: j \Rightarrow (i \Rightarrow bool) \Rightarrow (i \Rightarrow bool) \Rightarrow (i \Rightarrow bool)$

lift-definition *not* $::$ o$\Rightarrow$o $(\neg\text{-}\ [54]\ 70)$ **is**
$\lambda\ p\ s\ w$. $s = dj \wedge \neg p\ dj\ w \vee s \neq dj \wedge (I\text{-}NOT\ s\ (p\ s)\ w)$.
lift-definition *impl* $::$ o$\Rightarrow$o$\Rightarrow$o (**infixl** $\rightarrow$ *51*) **is**
$\lambda\ p\ q\ s\ w$. $s = dj \wedge (p\ dj\ w \longrightarrow q\ dj\ w) \vee s \neq dj \wedge (I\text{-}IMPL\ s\ (p\ s)\ (q\ s)\ w)$.
lift-definition $forall_\nu$ $:: (\nu \Rightarrow o) \Rightarrow o$ (**binder** $\forall_\nu$ $[8]$ *9*) **is**
$\lambda\ \varphi\ s\ w$. $\forall\ x :: \nu$. $(\varphi\ x)\ s\ w$.
lift-definition $forall_0$ $:: (\Pi_0 \Rightarrow o) \Rightarrow o$ (**binder** $\forall_0$ $[8]$ *9*) **is**
$\lambda\ \varphi\ s\ w$. $\forall\ x :: \Pi_0$. $(\varphi\ x)\ s\ w$.
lift-definition $forall_1$ $:: (\Pi_1 \Rightarrow o) \Rightarrow o$ (**binder** $\forall_1$ $[8]$ *9*) **is**
$\lambda\ \varphi\ s\ w$. $\forall\ x :: \Pi_1$. $(\varphi\ x)\ s\ w$.
lift-definition $forall_2$ $:: (\Pi_2 \Rightarrow o) \Rightarrow o$ (**binder** $\forall_2$ $[8]$ *9*) **is**
$\lambda\ \varphi\ s\ w$. $\forall\ x :: \Pi_2$. $(\varphi\ x)\ s\ w$.
lift-definition $forall_3$ $:: (\Pi_3 \Rightarrow o) \Rightarrow o$ (**binder** $\forall_3$ $[8]$ *9*) **is**
$\lambda\ \varphi\ s\ w$. $\forall\ x :: \Pi_3$. $(\varphi\ x)\ s\ w$.
lift-definition $forall_o$ $:: (o \Rightarrow o) \Rightarrow o$ (**binder** $\forall_o$ $[8]$ *9*) **is**
$\lambda\ \varphi\ s\ w$. $\forall\ x :: o$. $(\varphi\ x)\ s\ w$.
lift-definition *box* $::$ o$\Rightarrow$o $(\Box\text{-}\ [62]\ 63)$ **is**
$\lambda\ p\ s\ w$. $\forall\ v$. $p\ s\ v$.
lift-definition *actual* $::$ o$\Rightarrow$o $(\mathcal{A}\text{-}\ [64]\ 65)$ **is**
$\lambda\ p\ s\ w$. $p\ s\ dw$.

Remark. *The connectives behave classically if evaluated for the actual state dj, whereas their behavior is governed by uninterpreted constants for any other state.*

A.1.8. Lambda Expressions

Remark. *Lambda expressions have to convert maps from individuals to propositions to relations that are represented by maps from urelements to truth values.*

lift-definition *lambdabinder0* :: $o \Rightarrow \Pi_0$ $(\boldsymbol{\lambda}^0)$ **is** *id* **.**
lift-definition *lambdabinder1* :: $(\nu \Rightarrow o) \Rightarrow \Pi_1$ (**binder** $\boldsymbol{\lambda}$ $[8]$ 9) **is**
$\lambda\ \varphi\ u\ s\ w\ .\ \exists\ x\ .\ \nu\upsilon\ x = u \wedge \varphi\ x\ s\ w$ **.**
lift-definition *lambdabinder2* :: $(\nu \Rightarrow \nu \Rightarrow o) \Rightarrow \Pi_2$ $(\boldsymbol{\lambda}^2)$ **is**
$\lambda\ \varphi\ u\ v\ s\ w\ .\ \exists\ x\ y\ .\ \nu\upsilon\ x = u \wedge \nu\upsilon\ y = v \wedge \varphi\ x\ y\ s\ w$ **.**
lift-definition *lambdabinder3* :: $(\nu \Rightarrow \nu \Rightarrow \nu \Rightarrow o) \Rightarrow \Pi_3$ $(\boldsymbol{\lambda}^3)$ **is**
$\lambda\ \varphi\ u\ v\ r\ s\ w\ .\ \exists\ x\ y\ z\ .\ \nu\upsilon\ x = u \wedge \nu\upsilon\ y = v \wedge \nu\upsilon\ z = r \wedge \varphi\ x\ y\ z\ s\ w$ **.**

A.1.9. Proper Maps

Remark. *The embedding introduces the notion of* proper *maps from individual terms to propositions.*
Such a map is proper if and only if for all proper individual terms its truth evaluation in the actual state only depends on the urelements corresponding to the individuals the terms denote.
Proper maps are exactly those maps that - when used as matrix of a lambda-expression - unconditionally allow beta-reduction.

lift-definition *IsProperInX* :: $(\kappa \Rightarrow o) \Rightarrow$ *bool* **is**
$\lambda\ \varphi\ .\ \forall\ x\ v\ .\ (\exists\ a\ .\ \nu\upsilon\ a = \nu\upsilon\ x \wedge (\varphi\ (a^P)\ dj\ v)) = (\varphi\ (x^P)\ dj\ v)$ **.**
lift-definition *IsProperInXY* :: $(\kappa \Rightarrow \kappa \Rightarrow o) \Rightarrow$ *bool* **is**
$\lambda\ \varphi\ .\ \forall\ x\ y\ v\ .\ (\exists\ a\ b\ .\ \nu\upsilon\ a = \nu\upsilon\ x \wedge \nu\upsilon\ b = \nu\upsilon\ y$
$\wedge\ (\varphi\ (a^P)\ (b^P)\ dj\ v)) = (\varphi\ (x^P)\ (y^P)\ dj\ v)$ **.**
lift-definition *IsProperInXYZ* :: $(\kappa \Rightarrow \kappa \Rightarrow \kappa \Rightarrow o) \Rightarrow$ *bool* **is**
$\lambda\ \varphi\ .\ \forall\ x\ y\ z\ v\ .\ (\exists\ a\ b\ c\ .\ \nu\upsilon\ a = \nu\upsilon\ x \wedge \nu\upsilon\ b = \nu\upsilon\ y \wedge \nu\upsilon\ c = \nu\upsilon\ z$
$\wedge\ (\varphi\ (a^P)\ (b^P)\ (c^P)\ dj\ v)) = (\varphi\ (x^P)\ (y^P)\ (z^P)\ dj\ v)$ **.**

A.1.10. Validity

lift-definition *valid-in* :: $i \Rightarrow o \Rightarrow$ *bool* (**infixl** $\models$ 5) **is**
$\lambda\ v\ \varphi\ .\ \varphi\ dj\ v$ **.**

Remark. *A formula is considered semantically valid for a possible world, if it evaluates to True for the actual state dj and the given possible world.*

A.1.11. Concreteness

consts *ConcreteInWorld* :: $\omega \Rightarrow i \Rightarrow$ *bool*

abbreviation (*input*) *OrdinaryObjectsPossiblyConcrete* **where**
OrdinaryObjectsPossiblyConcrete $\equiv \forall\ x\ .\ \exists\ v\ .$ *ConcreteInWorld* $x\ v$
abbreviation (*input*) *PossiblyContingentObjectExists* **where**
PossiblyContingentObjectExists $\equiv \exists\ x\ v\ .$ *ConcreteInWorld* $x\ v$
$\wedge\ (\exists\ w\ .\ \neg$ *ConcreteInWorld* $x\ w)$
abbreviation (*input*) *PossiblyNoContingentObjectExists* **where**
PossiblyNoContingentObjectExists $\equiv \exists\ w\ .\ \forall\ x\ .$ *ConcreteInWorld* $x\ w$
$\longrightarrow (\forall\ v\ .$ *ConcreteInWorld* $x\ v)$
axiomatization where
OrdinaryObjectsPossiblyConcreteAxiom:

OrdinaryObjectsPossiblyConcrete
and *PossiblyContingentObjectExistsAxiom*:
PossiblyContingentObjectExists
and *PossiblyNoContingentObjectExistsAxiom*:
PossiblyNoContingentObjectExists

Remark. *Care has to be taken that the defined notion of concreteness coincides with the meta-logical distinction between abstract objects and ordinary objects. Furthermore the axioms about concreteness have to be satisfied. This is achieved by introducing an uninterpreted constant ConcreteInWorld that determines whether an ordinary object is concrete in a given possible world. This constant is axiomatized, such that all ordinary objects are possibly concrete, contingent objects possibly exist and possibly no contingent objects exist.*

lift-definition *Concrete*::Π_1 ($E!$) **is**
$\lambda\ u\ s\ w$. *case* u *of* $\omega\upsilon\ x \Rightarrow$ *ConcreteInWorld* $x\ w$ | - $\Rightarrow$ *False* **.**

Remark. *Concreteness of ordinary objects is now defined using this axiomatized uninterpreted constant. Abstract objects on the other hand are never concrete.*

A.1.12. Collection of Meta-Definitions

named-theorems *meta-defs*

declare *not-def*[*meta-defs*] *impl-def*[*meta-defs*] *forall*$_\nu$*-def*[*meta-defs*]
forall$_0$*-def*[*meta-defs*] *forall*$_1$*-def*[*meta-defs*]
forall$_2$*-def*[*meta-defs*] *forall*$_3$*-def*[*meta-defs*] *forall*$_\mathrm{o}$*-def*[*meta-defs*]
box-def[*meta-defs*] *actual-def*[*meta-defs*] *that-def*[*meta-defs*]
lambdabinder0-def[*meta-defs*] *lambdabinder1-def*[*meta-defs*]
lambdabinder2-def[*meta-defs*] *lambdabinder3-def*[*meta-defs*]
exe0-def[*meta-defs*] *exe1-def*[*meta-defs*] *exe2-def*[*meta-defs*]
exe3-def[*meta-defs*] *enc-def*[*meta-defs*] *inv-def*[*meta-defs*]
that-def[*meta-defs*] *valid-in-def*[*meta-defs*] *Concrete-def*[*meta-defs*]

declare [[*smt-solver* = *cvc4*]]
declare [[*simp-depth-limit* = *10*]]
declare [[*unify-search-bound* = *40*]]

A.1.13. Auxiliary Lemmata

named-theorems *meta-aux*

declare *make*κ*-inverse*[*meta-aux*] *eval*κ*-inverse*[*meta-aux*]
*make*o*-inverse*[*meta-aux*] *eval*o*-inverse*[*meta-aux*]
makeΠ_1*-inverse*[*meta-aux*] *eval*Π_1*-inverse*[*meta-aux*]
makeΠ_2*-inverse*[*meta-aux*] *eval*Π_2*-inverse*[*meta-aux*]
makeΠ_3*-inverse*[*meta-aux*] *eval*Π_3*-inverse*[*meta-aux*]
lemma $\nu\upsilon$*-*$\omega\nu$*-is-*$\omega\upsilon$[*meta-aux*]: $\nu\upsilon\ (\omega\nu\ x) = \omega\upsilon\ x$ **by** (*simp add*: $\nu\upsilon$*-def*)
lemma *rep-proper-id*[*meta-aux*]: *rep* $(x^P) = x$
by (*simp add*: *meta-aux* $\nu\kappa$*-def rep-def*)
lemma $\nu\kappa$*-proper*[*meta-aux*]: *proper* (x^P)
by (*simp add*: *meta-aux* $\nu\kappa$*-def proper-def*)
lemma *no-*$\alpha\omega$[*meta-aux*]: $\neg(\nu\upsilon\ (\alpha\nu\ x) = \omega\upsilon\ y)$ **by** (*simp add*: $\nu\upsilon$*-def*)
lemma *no-*$\sigma\omega$[*meta-aux*]: $\neg(\sigma\upsilon\ x = \omega\upsilon\ y)$ **by** *blast*
lemma $\nu\upsilon$*-surj*[*meta-aux*]: *surj* $\nu\upsilon$
using $\alpha\sigma$*-surj* **unfolding** $\nu\upsilon$*-def surj-def*

by (*metis* ν*.simps*(5) ν*.simps*(6) υ*.exhaust comp-apply*)
lemma *lambda*Π_1*-aux*[*meta-aux*]:
makeΠ_1 ($\lambda u\ s\ w.\ \exists x.\ \nu\upsilon\ x = u \wedge$ *eval*$\Pi_1\ F\ (\nu\upsilon\ x)\ s\ w) = F$
proof −
have $\bigwedge u\ s\ w\ \varphi\ .\ (\exists\ x\ .\ \nu\upsilon\ x = u \wedge \varphi\ (\nu\upsilon\ x)\ (s::j)\ (w::i)) \longleftrightarrow \varphi\ u\ s\ w$
using $\nu\upsilon$*-surj* **unfolding** *surj-def* **by** *metis*
thus *?thesis* **apply** *transfer* **by** *simp*
qed
lemma *lambda*Π_2*-aux*[*meta-aux*]:
makeΠ_2 ($\lambda u\ v\ s\ w.\ \exists x\ .\ \nu\upsilon\ x = u \wedge (\exists\ y\ .\ \nu\upsilon\ y = v \wedge$ *eval*$\Pi_2\ F\ (\nu\upsilon\ x)\ (\nu\upsilon\ y)\ s\ w)) = F$
proof −
have $\bigwedge u\ v\ (s::j)\ (w::i)\ \varphi\ .$
$(\exists\ x\ .\ \nu\upsilon\ x = u \wedge (\exists\ y\ .\ \nu\upsilon\ y = v \wedge \varphi\ (\nu\upsilon\ x)\ (\nu\upsilon\ y)\ s\ w))$
$\longleftrightarrow \varphi\ u\ v\ s\ w$
using $\nu\upsilon$*-surj* **unfolding** *surj-def* **by** *metis*
thus *?thesis* **apply** *transfer* **by** *simp*
qed
lemma *lambda*Π_3*-aux*[*meta-aux*]:
makeΠ_3 ($\lambda u\ v\ r\ s\ w.\ \exists x.\ \nu\upsilon\ x = u \wedge (\exists y.\ \nu\upsilon\ y = v \wedge$
$(\exists z.\ \nu\upsilon\ z = r \wedge$ *eval*$\Pi_3\ F\ (\nu\upsilon\ x)\ (\nu\upsilon\ y)\ (\nu\upsilon\ z)\ s\ w))) = F$
proof −
have $\bigwedge u\ v\ r\ (s::j)\ (w::i)\ \varphi\ .\ \exists x.\ \nu\upsilon\ x = u \wedge (\exists y.\ \nu\upsilon\ y = v$
$\wedge (\exists z.\ \nu\upsilon\ z = r \wedge \varphi\ (\nu\upsilon\ x)\ (\nu\upsilon\ y)\ (\nu\upsilon\ z)\ s\ w)) = \varphi\ u\ v\ r\ s\ w$
using $\nu\upsilon$*-surj* **unfolding** *surj-def* **by** *metis*
thus *?thesis* **apply** *transfer* **apply** (*rule ext*)+ **by** *metis*
qed

A.2. Semantic Abstraction

A.2.1. Semantics

locale *Semantics*
begin
named-theorems *semantics*

A.2.1.1. Semantic Domains

type-synonym $R_\kappa = \nu$
type-synonym $R_0 = j \Rightarrow i \Rightarrow bool$
type-synonym $R_1 = \upsilon \Rightarrow R_0$
type-synonym $R_2 = \upsilon \Rightarrow \upsilon \Rightarrow R_0$
type-synonym $R_3 = \upsilon \Rightarrow \upsilon \Rightarrow \upsilon \Rightarrow R_0$
type-synonym $W = i$

A.2.1.2. Denotation Functions

lift-definition $d_\kappa :: \kappa \Rightarrow R_\kappa$ *option* **is** *id* .
lift-definition $d_0 :: \Pi_0 \Rightarrow R_0$ *option* **is** *Some* .
lift-definition $d_1 :: \Pi_1 \Rightarrow R_1$ *option* **is** *Some* .
lift-definition $d_2 :: \Pi_2 \Rightarrow R_2$ *option* **is** *Some* .
lift-definition $d_3 :: \Pi_3 \Rightarrow R_3$ *option* **is** *Some* .

A.2.1.3. Actual World

definition w_0 **where** $w_0 \equiv dw$

A.2.1.4. Exemplification Extensions

definition *ex0* :: $R_0 \Rightarrow W \Rightarrow bool$
  where *ex0* $\equiv \lambda\ F\ .\ F\ dj$
definition *ex1* :: $R_1 \Rightarrow W \Rightarrow (R_\kappa\ set)$
  where *ex1* $\equiv \lambda\ F\ w\ .\ \{\ x\ .\ F\ (\nu\upsilon\ x)\ dj\ w\ \}$
definition *ex2* :: $R_2 \Rightarrow W \Rightarrow ((R_\kappa \times R_\kappa)\ set)$
  where *ex2* $\equiv \lambda\ F\ w\ .\ \{\ (x,y)\ .\ F\ (\nu\upsilon\ x)\ (\nu\upsilon\ y)\ dj\ w\ \}$
definition *ex3* :: $R_3 \Rightarrow W \Rightarrow ((R_\kappa \times R_\kappa \times R_\kappa)\ set)$
  where *ex3* $\equiv \lambda\ F\ w\ .\ \{\ (x,y,z)\ .\ F\ (\nu\upsilon\ x)\ (\nu\upsilon\ y)\ (\nu\upsilon\ z)\ dj\ w\ \}$

A.2.1.5. Encoding Extensions

definition *en* :: $R_1 \Rightarrow (R_\kappa\ set)$
  where *en* $\equiv \lambda\ F\ .\ \{\ x\ .\ case\ x\ of\ \alpha\nu\ y \Rightarrow make\Pi_1\ (\lambda\ x\ .\ F\ x) \in y$
    $\mid$ - $\Rightarrow$ *False* $\}$

A.2.1.6. Collection of Semantic Definitions

named-theorems *semantics-defs*
declare d_0*-def*[*semantics-defs*] d_1*-def*[*semantics-defs*]
    d_2*-def*[*semantics-defs*] d_3*-def*[*semantics-defs*]
    ex0-def[*semantics-defs*] *ex1-def*[*semantics-defs*]
    ex2-def[*semantics-defs*] *ex3-def*[*semantics-defs*]
    en-def[*semantics-defs*] d_κ*-def*[*semantics-defs*]
    w_0*-def*[*semantics-defs*]

A.2.1.7. Truth Conditions of Exemplification Formulas

lemma *T1-1*[*semantics*]:
  $(w \models (\!|F,x|\!)) = (\exists\ r\ o_1\ .\ Some\ r = d_1\ F \wedge Some\ o_1 = d_\kappa\ x \wedge o_1 \in ex1\ r\ w)$
  unfolding *semantics-defs*
  apply (*simp add*: *meta-defs meta-aux rep-def proper-def*)
  by (*metis option.discI option.exhaust option.sel*)

lemma *T1-2*[*semantics*]:
  $(w \models (\!|F,x,y|\!)) = (\exists\ r\ o_1\ o_2\ .\ Some\ r = d_2\ F \wedge Some\ o_1 = d_\kappa\ x$
    $\wedge\ Some\ o_2 = d_\kappa\ y \wedge (o_1,\ o_2) \in ex2\ r\ w)$
  unfolding *semantics-defs*
  apply (*simp add*: *meta-defs meta-aux rep-def proper-def*)
  by (*metis option.discI option.exhaust option.sel*)

lemma *T1-3*[*semantics*]:
  $(w \models (\!|F,x,y,z|\!)) = (\exists\ r\ o_1\ o_2\ o_3\ .\ Some\ r = d_3\ F \wedge Some\ o_1 = d_\kappa\ x$
    $\wedge\ Some\ o_2 = d_\kappa\ y \wedge Some\ o_3 = d_\kappa\ z$
    $\wedge\ (o_1,\ o_2,\ o_3) \in ex3\ r\ w)$
  unfolding *semantics-defs*
  apply (*simp add*: *meta-defs meta-aux rep-def proper-def*)
  by (*metis option.discI option.exhaust option.sel*)

lemma *T3*[*semantics*]:
  $(w \models (\!|F|\!)) = (\exists\ r\ .\ Some\ r = d_0\ F \wedge ex0\ r\ w)$
  unfolding *semantics-defs*
  by (*simp add*: *meta-defs meta-aux*)

A.2.1.8. Truth Conditions of Encoding Formulas

lemma *T2*[*semantics*]:

$(w \models \{\!|x,F|\!\}) = (\exists\ r\ o_1\ .\ Some\ r = d_1\ F \wedge Some\ o_1 = d_\kappa\ x \wedge o_1 \in en\ r)$
unfolding *semantics-defs*
apply (*simp add*: *meta-defs meta-aux rep-def proper-def split*: ν.*split*)
by (*metis* ν.*exhaust* ν.*inject*(2) ν.*simps*(4) $\nu\kappa$.*rep-eq option.collapse*
option.discI rep.rep-eq rep-proper-id)

A.2.1.9. Truth Conditions of Complex Formulas

lemma *T4*[*semantics*]: $(w \models \neg\psi) = (\neg(w \models \psi))$
by (*simp add*: *meta-defs meta-aux*)

lemma *T5*[*semantics*]: $(w \models \psi \rightarrow \chi) = (\neg(w \models \psi) \vee (w \models \chi))$
by (*simp add*: *meta-defs meta-aux*)

lemma *T6*[*semantics*]: $(w \models \Box\psi) = (\forall\ v\ .\ (v \models \psi))$
by (*simp add*: *meta-defs meta-aux*)

lemma *T7*[*semantics*]: $(w \models \boldsymbol{\mathcal{A}}\psi) = (dw \models \psi)$
by (*simp add*: *meta-defs meta-aux*)

lemma *T8-*ν[*semantics*]: $(w \models \forall_\nu\ x.\ \psi\ x) = (\forall\ x\ .\ (w \models \psi\ x))$
by (*simp add*: *meta-defs meta-aux*)

lemma *T8-0*[*semantics*]: $(w \models \forall_0\ x.\ \psi\ x) = (\forall\ x\ .\ (w \models \psi\ x))$
by (*simp add*: *meta-defs meta-aux*)

lemma *T8-1*[*semantics*]: $(w \models \forall_1\ x.\ \psi\ x) = (\forall\ x\ .\ (w \models \psi\ x))$
by (*simp add*: *meta-defs meta-aux*)

lemma *T8-2*[*semantics*]: $(w \models \forall_2\ x.\ \psi\ x) = (\forall\ x\ .\ (w \models \psi\ x))$
by (*simp add*: *meta-defs meta-aux*)

lemma *T8-3*[*semantics*]: $(w \models \forall_3\ x.\ \psi\ x) = (\forall\ x\ .\ (w \models \psi\ x))$
by (*simp add*: *meta-defs meta-aux*)

lemma *T8-*o[*semantics*]: $(w \models \forall_o\ x.\ \psi\ x) = (\forall\ x\ .\ (w \models \psi\ x))$
by (*simp add*: *meta-defs meta-aux*)

A.2.1.10. Denotations of Descriptions

lemma *D3*[*semantics*]:
$d_\kappa\ (\boldsymbol{\iota}x\ .\ \psi\ x) = (if\ (\exists x\ .\ (w_0 \models \psi\ x) \wedge (\forall\ y\ .\ (w_0 \models \psi\ y) \longrightarrow y = x))$
$then\ (Some\ (THE\ x\ .\ (w_0 \models \psi\ x)))\ else\ None)$
unfolding *semantics-defs*
by (*auto simp*: *meta-defs meta-aux*)

A.2.1.11. Denotations of Lambda Expressions

lemma *D4-1*[*semantics*]: $d_1\ (\boldsymbol{\lambda}\ x\ .\ (\!|F,\ x^P|\!)) = d_1\ F$
by (*simp add*: *meta-defs meta-aux*)

lemma *D4-2*[*semantics*]: $d_2\ (\boldsymbol{\lambda}^2\ (\lambda\ x\ y\ .\ (\!|F,\ x^P,\ y^P|\!))) = d_2\ F$
by (*simp add*: *meta-defs meta-aux*)

lemma *D4-3*[*semantics*]: $d_3\ (\boldsymbol{\lambda}^3\ (\lambda\ x\ y\ z\ .\ (\!|F,\ x^P,\ y^P,\ z^P|\!))) = d_3\ F$
by (*simp add*: *meta-defs meta-aux*)

lemma *D5-1*[*semantics*]:
assumes *IsProperInX* φ
shows $\bigwedge$ *w* o_1 *r* . *Some r* = d_1 ($\boldsymbol{\lambda}$ *x* . (φ (x^P))) $\wedge$ *Some* o_1 = d_κ *x*
$\longrightarrow$ (o_1 $\in$ *ex1 r w*) = (*w* $\models$ φ *x*)
using *assms* **unfolding** *IsProperInX-def semantics-defs*
by (*auto simp*: *meta-defs meta-aux rep-def proper-def* $\nu\kappa$*.abs-eq*)

lemma *D5-2*[*semantics*]:
assumes *IsProperInXY* φ
shows $\bigwedge$ *w* o_1 o_2 *r* . *Some r* = d_2 ($\boldsymbol{\lambda}^2$ (λ *x y* . φ (x^P) (y^P)))
$\wedge$ *Some* o_1 = d_κ *x* $\wedge$ *Some* o_2 = d_κ *y*
$\longrightarrow$ ((o_1,o_2) $\in$ *ex2 r w*) = (*w* $\models$ φ *x y*)
using *assms* **unfolding** *IsProperInXY-def semantics-defs*
by (*auto simp*: *meta-defs meta-aux rep-def proper-def* $\nu\kappa$*.abs-eq*)

lemma *D5-3*[*semantics*]:
assumes *IsProperInXYZ* φ
shows $\bigwedge$ *w* o_1 o_2 o_3 *r* . *Some r* = d_3 ($\boldsymbol{\lambda}^3$ (λ *x y z* . φ (x^P) (y^P) (z^P)))
$\wedge$ *Some* o_1 = d_κ *x* $\wedge$ *Some* o_2 = d_κ *y* $\wedge$ *Some* o_3 = d_κ *z*
$\longrightarrow$ ((o_1,o_2,o_3) $\in$ *ex3 r w*) = (*w* $\models$ φ *x y z*)
using *assms* **unfolding** *IsProperInXYZ-def semantics-defs*
by (*auto simp*: *meta-defs meta-aux rep-def proper-def* $\nu\kappa$*.abs-eq*)

lemma *D6*[*semantics*]: ($\bigwedge$ *w r* . *Some r* = d_0 ($\boldsymbol{\lambda}^0$ φ) $\longrightarrow$ *ex0 r w* = (*w* $\models$ φ))
by (*auto simp*: *meta-defs meta-aux semantics-defs*)

A.2.1.12. Auxiliary Lemmas

lemma *propex*$_0$: $\exists$ *r* . *Some r* = d_0 *F*
unfolding d_0*-def* **by** *simp*
lemma *propex*$_1$: $\exists$ *r* . *Some r* = d_1 *F*
unfolding d_1*-def* **by** *simp*
lemma *propex*$_2$: $\exists$ *r* . *Some r* = d_2 *F*
unfolding d_2*-def* **by** *simp*
lemma *propex*$_3$: $\exists$ *r* . *Some r* = d_3 *F*
unfolding d_3*-def* **by** *simp*
lemma d_κ*-proper*: d_κ (u^P) = *Some u*
unfolding d_κ*-def* **by** (*simp add*: $\nu\kappa$*-def meta-aux*)
lemma *ConcretenessSemantics1*:
Some r = d_1 *E*! $\Longrightarrow$ ($\exists$ *w* . $\omega\nu$ *x* $\in$ *ex1 r w*)
unfolding *semantics-defs* **apply** *transfer*
by (*simp add*: *OrdinaryObjectsPossiblyConcreteAxiom* $\nu\upsilon$*-*$\omega\nu$*-is-*$\omega\upsilon$)
lemma *ConcretenessSemantics2*:
Some r = d_1 *E*! $\Longrightarrow$ (*x* $\in$ *ex1 r w* $\longrightarrow$ ($\exists$ *y*. *x* = $\omega\nu$ *y*))
unfolding *semantics-defs* **apply** *transfer* **apply** *simp*
by (*metis* ν*.exhaust* υ*.exhaust* υ*.simps*(6) *no-*$\alpha\omega$)
lemma d_0*-inject*: $\bigwedge$*x y*. d_0 *x* = d_0 *y* $\Longrightarrow$ *x* = *y*
unfolding d_0*-def* **by** (*simp add*: *evalo-inject*)
lemma d_1*-inject*: $\bigwedge$*x y*. d_1 *x* = d_1 *y* $\Longrightarrow$ *x* = *y*
unfolding d_1*-def* **by** (*simp add*: *eval*Π_1*-inject*)
lemma d_2*-inject*: $\bigwedge$*x y*. d_2 *x* = d_2 *y* $\Longrightarrow$ *x* = *y*
unfolding d_2*-def* **by** (*simp add*: *eval*Π_2*-inject*)
lemma d_3*-inject*: $\bigwedge$*x y*. d_3 *x* = d_3 *y* $\Longrightarrow$ *x* = *y*
unfolding d_3*-def* **by** (*simp add*: *eval*Π_3*-inject*)
lemma d_κ*-inject*: $\bigwedge$*x y* o_1. *Some* o_1 = d_κ *x* $\wedge$ *Some* o_1 = d_κ *y* $\Longrightarrow$ *x* = *y*
proof −
fix *x* :: κ **and** *y* :: κ **and** o_1 :: ν
assume *Some* o_1 = d_κ *x* $\wedge$ *Some* o_1 = d_κ *y*

thus $x = y$ **apply** *transfer* **by** *auto*
qed
end

A.2.2. Introduction Rules for Proper Maps

Remark. *Every map whose arguments only occur in exemplification expressions is proper.*

named-theorems *IsProper-intros*

lemma *IsProperInX-intro*[*IsProper-intros*]:
IsProperInX (λ x . χ
($*$ *one place* $*$) (λ F . ($\!|F,x|\!$))
($*$ *two place* $*$) (λ F . ($\!|F,x,x|\!$)) (λ F a . ($\!|F,x,a|\!$)) (λ F a . ($\!|F,a,x|\!$))
($*$ *three place three x* $*$) (λ F . ($\!|F,x,x,x|\!$))
($*$ *three place two x* $*$) (λ F a . ($\!|F,x,x,a|\!$)) (λ F a . ($\!|F,x,a,x|\!$))
(λ F a . ($\!|F,a,x,x|\!$))
($*$ *three place one x* $*$) (λ F a b. ($\!|F,x,a,b|\!$)) (λ F a b. ($\!|F,a,x,b|\!$))
(λ F a b . ($\!|F,a,b,x|\!$)))
unfolding *IsProperInX-def*
by (*auto simp*: *meta-defs meta-aux*)

lemma *IsProperInXY-intro*[*IsProper-intros*]:
IsProperInXY (λ x y . χ
($*$ *only x* $*$)
($*$ *one place* $*$) (λ F . ($\!|F,x|\!$))
($*$ *two place* $*$) (λ F . ($\!|F,x,x|\!$)) (λ F a . ($\!|F,x,a|\!$)) (λ F a . ($\!|F,a,x|\!$))
($*$ *three place three x* $*$) (λ F . ($\!|F,x,x,x|\!$))
($*$ *three place two x* $*$) (λ F a . ($\!|F,x,x,a|\!$)) (λ F a . ($\!|F,x,a,x|\!$))
(λ F a . ($\!|F,a,x,x|\!$))
($*$ *three place one x* $*$) (λ F a b. ($\!|F,x,a,b|\!$)) (λ F a b. ($\!|F,a,x,b|\!$))
(λ F a b . ($\!|F,a,b,x|\!$))
($*$ *only y* $*$)
($*$ *one place* $*$) (λ F . ($\!|F,y|\!$))
($*$ *two place* $*$) (λ F . ($\!|F,y,y|\!$)) (λ F a . ($\!|F,y,a|\!$)) (λ F a . ($\!|F,a,y|\!$))
($*$ *three place three y* $*$) (λ F . ($\!|F,y,y,y|\!$))
($*$ *three place two y* $*$) (λ F a . ($\!|F,y,y,a|\!$)) (λ F a . ($\!|F,y,a,y|\!$))
(λ F a . ($\!|F,a,y,y|\!$))
($*$ *three place one y* $*$) (λ F a b. ($\!|F,y,a,b|\!$)) (λ F a b. ($\!|F,a,y,b|\!$))
(λ F a b . ($\!|F,a,b,y|\!$))
($*$ *x and y* $*$)
($*$ *two place* $*$) (λ F . ($\!|F,x,y|\!$)) (λ F . ($\!|F,y,x|\!$))
($*$ *three place* (x,y) $*$) (λ F a . ($\!|F,x,y,a|\!$)) (λ F a . ($\!|F,x,a,y|\!$))
(λ F a . ($\!|F,a,x,y|\!$))
($*$ *three place* (y,x) $*$) (λ F a . ($\!|F,y,x,a|\!$)) (λ F a . ($\!|F,y,a,x|\!$))
(λ F a . ($\!|F,a,y,x|\!$))
($*$ *three place* (x,x,y) $*$) (λ F . ($\!|F,x,x,y|\!$)) (λ F . ($\!|F,x,y,x|\!$))
(λ F . ($\!|F,y,x,x|\!$))
($*$ *three place* (x,y,y) $*$) (λ F . ($\!|F,x,y,y|\!$)) (λ F . ($\!|F,y,x,y|\!$))
(λ F . ($\!|F,y,y,x|\!$))
($*$ *three place* (x,x,x) $*$) (λ F . ($\!|F,x,x,x|\!$))
($*$ *three place* (y,y,y) $*$) (λ F . ($\!|F,y,y,y|\!$)))
unfolding *IsProperInXY-def* **by** (*auto simp*: *meta-defs meta-aux*)

lemma *IsProperInXYZ-intro*[*IsProper-intros*]:
IsProperInXYZ (λ x y z . χ
($*$ *only x* $*$)

$(*\ one\ place\ *)\ (\lambda\ F\ .\ (\!|F,x|\!))$
$(*\ two\ place\ *)\ (\lambda\ F\ .\ (\!|F,x,x|\!))\ (\lambda\ F\ a\ .\ (\!|F,x,a|\!))\ (\lambda\ F\ a\ .\ (\!|F,a,x|\!))$
$(*\ three\ place\ three\ x\ *)\ (\lambda\ F\ .\ (\!|F,x,x,x|\!))$
$(*\ three\ place\ two\ x\ *)\ (\lambda\ F\ a\ .\ (\!|F,x,x,a|\!))\ (\lambda\ F\ a\ .\ (\!|F,x,a,x|\!))$
$(\lambda\ F\ a\ .\ (\!|F,a,x,x|\!))$
$(*\ three\ place\ one\ x\ *)\ (\lambda\ F\ a\ b.\ (\!|F,x,a,b|\!))\ (\lambda\ F\ a\ b.\ (\!|F,a,x,b|\!))$
$(\lambda\ F\ a\ b\ .\ (\!|F,a,b,x|\!))$
$(*\ only\ y\ *)$
$(*\ one\ place\ *)\ (\lambda\ F\ .\ (\!|F,y|\!))$
$(*\ two\ place\ *)\ (\lambda\ F\ .\ (\!|F,y,y|\!))\ (\lambda\ F\ a\ .\ (\!|F,y,a|\!))\ (\lambda\ F\ a\ .\ (\!|F,a,y|\!))$
$(*\ three\ place\ three\ y\ *)\ (\lambda\ F\ .\ (\!|F,y,y,y|\!))$
$(*\ three\ place\ two\ y\ *)\ (\lambda\ F\ a\ .\ (\!|F,y,y,a|\!))\ (\lambda\ F\ a\ .\ (\!|F,y,a,y|\!))$
$(\lambda\ F\ a\ .\ (\!|F,a,y,y|\!))$
$(*\ three\ place\ one\ y\ *)\ (\lambda\ F\ a\ b.\ (\!|F,y,a,b|\!))\ (\lambda\ F\ a\ b.\ (\!|F,a,y,b|\!))$
$(\lambda\ F\ a\ b\ .\ (\!|F,a,b,y|\!))$
$(*\ only\ z\ *)$
$(*\ one\ place\ *)\ (\lambda\ F\ .\ (\!|F,z|\!))$
$(*\ two\ place\ *)\ (\lambda\ F\ .\ (\!|F,z,z|\!))\ (\lambda\ F\ a\ .\ (\!|F,z,a|\!))\ (\lambda\ F\ a\ .\ (\!|F,a,z|\!))$
$(*\ three\ place\ three\ z\ *)\ (\lambda\ F\ .\ (\!|F,z,z,z|\!))$
$(*\ three\ place\ two\ z\ *)\ (\lambda\ F\ a\ .\ (\!|F,z,z,a|\!))\ (\lambda\ F\ a\ .\ (\!|F,z,a,z|\!))$
$(\lambda\ F\ a\ .\ (\!|F,a,z,z|\!))$
$(*\ three\ place\ one\ z\ *)\ (\lambda\ F\ a\ b.\ (\!|F,z,a,b|\!))\ (\lambda\ F\ a\ b.\ (\!|F,a,z,b|\!))$
$(\lambda\ F\ a\ b\ .\ (\!|F,a,b,z|\!))$
$(*\ x\ and\ y\ *)$
$(*\ two\ place\ *)\ (\lambda\ F\ .\ (\!|F,x,y|\!))\ (\lambda\ F\ .\ (\!|F,y,x|\!))$
$(*\ three\ place\ (x,y)\ *)\ (\lambda\ F\ a\ .\ (\!|F,x,y,a|\!))\ (\lambda\ F\ a\ .\ (\!|F,x,a,y|\!))$
$(\lambda\ F\ a\ .\ (\!|F,a,x,y|\!))$
$(*\ three\ place\ (y,x)\ *)\ (\lambda\ F\ a\ .\ (\!|F,y,x,a|\!))\ (\lambda\ F\ a\ .\ (\!|F,y,a,x|\!))$
$(\lambda\ F\ a\ .\ (\!|F,a,y,x|\!))$
$(*\ three\ place\ (x,x,y)\ *)\ (\lambda\ F\ .\ (\!|F,x,x,y|\!))\ (\lambda\ F\ .\ (\!|F,x,y,x|\!))$
$(\lambda\ F\ .\ (\!|F,y,x,x|\!))$
$(*\ three\ place\ (x,y,y)\ *)\ (\lambda\ F\ .\ (\!|F,x,y,y|\!))\ (\lambda\ F\ .\ (\!|F,y,x,y|\!))$
$(\lambda\ F\ .\ (\!|F,y,y,x|\!))$
$(*\ three\ place\ (x,x,x)\ *)\ (\lambda\ F\ .\ (\!|F,x,x,x|\!))$
$(*\ three\ place\ (y,y,y)\ *)\ (\lambda\ F\ .\ (\!|F,y,y,y|\!))$
$(*\ x\ and\ z\ *)$
$(*\ two\ place\ *)\ (\lambda\ F\ .\ (\!|F,x,z|\!))\ (\lambda\ F\ .\ (\!|F,z,x|\!))$
$(*\ three\ place\ (x,z)\ *)\ (\lambda\ F\ a\ .\ (\!|F,x,z,a|\!))\ (\lambda\ F\ a\ .\ (\!|F,x,a,z|\!))$
$(\lambda\ F\ a\ .\ (\!|F,a,x,z|\!))$
$(*\ three\ place\ (z,x)\ *)\ (\lambda\ F\ a\ .\ (\!|F,z,x,a|\!))\ (\lambda\ F\ a\ .\ (\!|F,z,a,x|\!))$
$(\lambda\ F\ a\ .\ (\!|F,a,z,x|\!))$
$(*\ three\ place\ (x,x,z)\ *)\ (\lambda\ F\ .\ (\!|F,x,x,z|\!))\ (\lambda\ F\ .\ (\!|F,x,z,x|\!))$
$(\lambda\ F\ .\ (\!|F,z,x,x|\!))$
$(*\ three\ place\ (x,z,z)\ *)\ (\lambda\ F\ .\ (\!|F,x,z,z|\!))\ (\lambda\ F\ .\ (\!|F,z,x,z|\!))$
$(\lambda\ F\ .\ (\!|F,z,z,x|\!))$
$(*\ three\ place\ (x,x,x)\ *)\ (\lambda\ F\ .\ (\!|F,x,x,x|\!))$
$(*\ three\ place\ (z,z,z)\ *)\ (\lambda\ F\ .\ (\!|F,z,z,z|\!))$
$(*\ y\ and\ z\ *)$
$(*\ two\ place\ *)\ (\lambda\ F\ .\ (\!|F,y,z|\!))\ (\lambda\ F\ .\ (\!|F,z,y|\!))$
$(*\ three\ place\ (y,z)\ *)\ (\lambda\ F\ a\ .\ (\!|F,y,z,a|\!))\ (\lambda\ F\ a\ .\ (\!|F,y,a,z|\!))$
$(\lambda\ F\ a\ .\ (\!|F,a,y,z|\!))$
$(*\ three\ place\ (z,y)\ *)\ (\lambda\ F\ a\ .\ (\!|F,z,y,a|\!))\ (\lambda\ F\ a\ .\ (\!|F,z,a,y|\!))$
$(\lambda\ F\ a\ .\ (\!|F,a,z,y|\!))$
$(*\ three\ place\ (y,y,z)\ *)\ (\lambda\ F\ .\ (\!|F,y,y,z|\!))\ (\lambda\ F\ .\ (\!|F,y,z,y|\!))$
$(\lambda\ F\ .\ (\!|F,z,y,y|\!))$
$(*\ three\ place\ (y,z,z)\ *)\ (\lambda\ F\ .\ (\!|F,y,z,z|\!))\ (\lambda\ F\ .\ (\!|F,z,y,z|\!))$
$(\lambda\ F\ .\ (\!|F,z,z,y|\!))$
$(*\ three\ place\ (y,y,y)\ *)\ (\lambda\ F\ .\ (\!|F,y,y,y|\!))$

```
    (* three place (z,z,z) *) (λ F . (|F,z,z,z|))
  (* x y z *)
    (* three place (x,...) *) (λ F . (|F,x,y,z|)) (λ F . (|F,x,z,y|))
    (* three place (y,...) *) (λ F . (|F,y,x,z|)) (λ F . (|F,y,z,x|))
    (* three place (z,...) *) (λ F . (|F,z,x,y|)) (λ F . (|F,z,y,x|)))
 unfolding IsProperInXYZ-def
 by (auto simp: meta-defs meta-aux)

method show-proper = (fast intro: IsProper-intros)
```

A.2.3. Validity Syntax

```
abbreviation validity-in :: o⇒i⇒bool ([- in -] [1]) where
  validity-in ≡ λ φ v . v ⊨ φ
definition actual-validity :: o⇒bool ([-] [1]) where
  actual-validity ≡ λ φ . dw ⊨ φ
definition necessary-validity :: o⇒bool (□[-] [1]) where
  necessary-validity ≡ λ φ . ∀ v . (v ⊨ φ)
```

A.3. General Quantification

Remark. *In order to define general quantifiers that can act on individuals as well as relations a type class is introduced which assumes the semantics of the all quantifier. This type class is then instantiated for individuals and relations.*

A.3.1. Type Class

```
class quantifiable = fixes forall :: ('a⇒o)⇒o (binder ∀ [8] 9)
  assumes quantifiable-T8: (w ⊨ (∀ x . ψ x)) = (∀ x . (w ⊨ (ψ x)))
begin
end

lemma (in Semantics) T8: shows (w ⊨ ∀ x . ψ x) = (∀ x . (w ⊨ ψ x))
  using quantifiable-T8 .
```

A.3.2. Instantiations

```
instantiation ν :: quantifiable
begin
  definition forall-ν :: (ν⇒o)⇒o where forall-ν ≡ forall_ν
  instance proof
    fix w :: i and ψ :: ν⇒o
    show (w ⊨ ∀x. ψ x) = (∀x. (w ⊨ ψ x))
      unfolding forall-ν-def using Semantics.T8-ν .
  qed
end

instantiation o :: quantifiable
begin
  definition forall-o :: (o⇒o)⇒o where forall-o ≡ forall_o
  instance proof
    fix w :: i and ψ :: o⇒o
```

```
    show (w ⊨ ∀ x. ψ x) = (∀ x. (w ⊨ ψ x))
      unfolding forall-o-def using Semantics.T8-o .
  qed
end

instantiation Π1 :: quantifiable
begin
  definition forall-Π1 :: (Π1⇒o)⇒o where forall-Π1 ≡ forall1
  instance proof
    fix w :: i and ψ :: Π1⇒o
    show (w ⊨ ∀ x. ψ x) = (∀ x. (w ⊨ ψ x))
      unfolding forall-Π1-def using Semantics.T8-1 .
  qed
end

instantiation Π2 :: quantifiable
begin
  definition forall-Π2 :: (Π2⇒o)⇒o where forall-Π2 ≡ forall2
  instance proof
    fix w :: i and ψ :: Π2⇒o
    show (w ⊨ ∀ x. ψ x) = (∀ x. (w ⊨ ψ x))
      unfolding forall-Π2-def using Semantics.T8-2 .
  qed
end

instantiation Π3 :: quantifiable
begin
  definition forall-Π3 :: (Π3⇒o)⇒o where forall-Π3 ≡ forall3
  instance proof
    fix w :: i and ψ :: Π3⇒o
    show (w ⊨ ∀ x. ψ x) = (∀ x. (w ⊨ ψ x))
      unfolding forall-Π3-def using Semantics.T8-3 .
  qed
end
```

A.4. Basic Definitions

A.4.1. Derived Connectives

```
definition conj::o⇒o⇒o (infixl & 53) where
  conj ≡ λ x y . ¬(x → ¬y)
definition disj::o⇒o⇒o (infixl ∨ 52) where
  disj ≡ λ x y . ¬x → y
definition equiv::o⇒o⇒o (infixl ≡ 51) where
  equiv ≡ λ x y . (x → y) & (y → x)
definition diamond::o⇒o (◇- [62] 63) where
  diamond ≡ λ φ . ¬□¬φ
definition (in quantifiable) exists :: ('a⇒o)⇒o (binder ∃ [8] 9) where
    exists ≡ λ φ . ¬(∀ x . ¬φ x)

named-theorems conn-defs
declare diamond-def[conn-defs] conj-def[conn-defs]
        disj-def[conn-defs] equiv-def[conn-defs]
        exists-def[conn-defs]
```

A.4.2. Abstract and Ordinary Objects

definition *Ordinary* :: Π_1 $(O!)$ **where** $Ordinary \equiv \boldsymbol{\lambda} x.\ \Diamond(\!|E!,x^P|\!)$
definition *Abstract* :: Π_1 $(A!)$ **where** $Abstract \equiv \boldsymbol{\lambda} x.\ \neg\Diamond(\!|E!,x^P|\!)$

A.4.3. Identity Definitions

definition $basic\text{-}identity_E$::Π_2 **where**
$basic\text{-}identity_E \equiv \boldsymbol{\lambda}^2\ (\lambda\ x\ y\ .\ (\!|O!,x^P|\!)\ \&\ (\!|O!,y^P|\!)$
$\&\ \Box(\forall\ F.\ (\!|F,x^P|\!) \equiv (\!|F,y^P|\!)))$

definition $basic\text{-}identity_E\text{-}infix$::$\kappa\Rightarrow\kappa\Rightarrow$o (**infixl** $=_E$ *63*) **where**
$x =_E y \equiv (\!|basic\text{-}identity_E,\ x,\ y|\!)$

definition $basic\text{-}identity_\kappa$ (**infixl** $=_\kappa$ *63*) **where**
$basic\text{-}identity_\kappa \equiv \lambda\ x\ y\ .\ (x =_E y) \vee (\!|A!,x|\!)\ \&\ (\!|A!,y|\!)$
$\&\ \Box(\forall\ F.\ \{\!|x,F|\!\} \equiv \{\!|y,F|\!\})$

definition $basic\text{-}identity_1$ (**infixl** $=_1$ *63*) **where**
$basic\text{-}identity_1 \equiv \lambda\ F\ G\ .\ \Box(\forall\ x.\ \{\!|x^P,F|\!\} \equiv \{\!|x^P,G|\!\})$

definition $basic\text{-}identity_2$ :: $\Pi_2\Rightarrow\Pi_2\Rightarrow$o (**infixl** $=_2$ *63*) **where**
$basic\text{-}identity_2 \equiv \lambda\ F\ G\ .\ \forall\ x.\ ((\boldsymbol{\lambda} y.\ (\!|F,x^P,y^P|\!)) =_1 (\boldsymbol{\lambda} y.\ (\!|G,x^P,y^P|\!)))$
$\&\ ((\boldsymbol{\lambda} y.\ (\!|F,y^P,x^P|\!)) =_1 (\boldsymbol{\lambda} y.\ (\!|G,y^P,x^P|\!)))$

definition $basic\text{-}identity_3$::$\Pi_3\Rightarrow\Pi_3\Rightarrow$o (**infixl** $=_3$ *63*) **where**
$basic\text{-}identity_3 \equiv \lambda\ F\ G\ .\ \forall\ x\ y.\ (\boldsymbol{\lambda} z.\ (\!|F,z^P,x^P,y^P|\!)) =_1 (\boldsymbol{\lambda} z.\ (\!|G,z^P,x^P,y^P|\!))$
$\&\ (\boldsymbol{\lambda} z.\ (\!|F,x^P,z^P,y^P|\!)) =_1 (\boldsymbol{\lambda} z.\ (\!|G,x^P,z^P,y^P|\!))$
$\&\ (\boldsymbol{\lambda} z.\ (\!|F,x^P,y^P,z^P|\!)) =_1 (\boldsymbol{\lambda} z.\ (\!|G,x^P,y^P,z^P|\!))$

definition $basic\text{-}identity_0$::o$\Rightarrow$o$\Rightarrow$o (**infixl** $=_0$ *63*) **where**
$basic\text{-}identity_0 \equiv \lambda\ F\ G\ .\ (\boldsymbol{\lambda} y.\ F) =_1 (\boldsymbol{\lambda} y.\ G)$

A.5. MetaSolver

Remark. *meta-solver is a resolution prover that translates expressions in the embedded logic to expressions in the meta-logic, resp. semantic expressions. The rules for connectives, quantifiers, exemplification and encoding are straightforward. Furthermore, rules for the defined identities are derived. The defined identities in the embedded logic coincide with the meta-logical equality.*

locale *MetaSolver*
begin
interpretation *Semantics* **.**

named-theorems *meta-intro*
named-theorems *meta-elim*
named-theorems *meta-subst*
named-theorems *meta-cong*

method *meta-solver* = (*assumption* | *rule meta-intro*
| *erule meta-elim* | *drule meta-elim* | *subst meta-subst*
| *subst* (*asm*) *meta-subst* | (*erule notE*; (*meta-solver*; *fail*))
)+

A.5.1. Rules for Implication

lemma *ImplI*[*meta-intro*]: $([\varphi\ in\ v] \Longrightarrow [\psi\ in\ v]) \Longrightarrow ([\varphi \rightarrow \psi\ in\ v])$
by (*simp add*: *Semantics.T5*)
lemma *ImplE*[*meta-elim*]: $([\varphi \rightarrow \psi\ in\ v]) \Longrightarrow ([\varphi\ in\ v] \longrightarrow [\psi\ in\ v])$
by (*simp add*: *Semantics.T5*)
lemma *ImplS*[*meta-subst*]: $([\varphi \rightarrow \psi\ in\ v]) = ([\varphi\ in\ v] \longrightarrow [\psi\ in\ v])$
by (*simp add*: *Semantics.T5*)

A.5.2. Rules for Negation

lemma *NotI*[*meta-intro*]: $\neg[\varphi\ in\ v] \Longrightarrow [\neg\varphi\ in\ v]$
by (*simp add*: *Semantics.T4*)
lemma *NotE*[*meta-elim*]: $[\neg\varphi\ in\ v] \Longrightarrow \neg[\varphi\ in\ v]$
by (*simp add*: *Semantics.T4*)
lemma *NotS*[*meta-subst*]: $[\neg\varphi\ in\ v] = (\neg[\varphi\ in\ v])$
by (*simp add*: *Semantics.T4*)

A.5.3. Rules for Conjunction

lemma *ConjI*[*meta-intro*]: $([\varphi\ in\ v] \wedge [\psi\ in\ v]) \Longrightarrow [\varphi\ \&\ \psi\ in\ v]$
by (*simp add*: *conj-def NotS ImplS*)
lemma *ConjE*[*meta-elim*]: $[\varphi\ \&\ \psi\ in\ v] \Longrightarrow ([\varphi\ in\ v] \wedge [\psi\ in\ v])$
by (*simp add*: *conj-def NotS ImplS*)
lemma *ConjS*[*meta-subst*]: $[\varphi\ \&\ \psi\ in\ v] = ([\varphi\ in\ v] \wedge [\psi\ in\ v])$
by (*simp add*: *conj-def NotS ImplS*)

A.5.4. Rules for Equivalence

lemma *EquivI*[*meta-intro*]: $([\varphi\ in\ v] \longleftrightarrow [\psi\ in\ v]) \Longrightarrow [\varphi \equiv \psi\ in\ v]$
by (*simp add*: *equiv-def NotS ImplS ConjS*)
lemma *EquivE*[*meta-elim*]: $[\varphi \equiv \psi\ in\ v] \Longrightarrow ([\varphi\ in\ v] \longleftrightarrow [\psi\ in\ v])$
by (*auto simp*: *equiv-def NotS ImplS ConjS*)
lemma *EquivS*[*meta-subst*]: $[\varphi \equiv \psi\ in\ v] = ([\varphi\ in\ v] \longleftrightarrow [\psi\ in\ v])$
by (*auto simp*: *equiv-def NotS ImplS ConjS*)

A.5.5. Rules for Disjunction

lemma *DisjI*[*meta-intro*]: $([\varphi\ in\ v] \vee [\psi\ in\ v]) \Longrightarrow [\varphi \vee \psi\ in\ v]$
by (*auto simp*: *disj-def NotS ImplS*)
lemma *DisjE*[*meta-elim*]: $[\varphi \vee \psi\ in\ v] \Longrightarrow ([\varphi\ in\ v] \vee [\psi\ in\ v])$
by (*auto simp*: *disj-def NotS ImplS*)
lemma *DisjS*[*meta-subst*]: $[\varphi \vee \psi\ in\ v] = ([\varphi\ in\ v] \vee [\psi\ in\ v])$
by (*auto simp*: *disj-def NotS ImplS*)

A.5.6. Rules for Necessity

lemma *BoxI*[*meta-intro*]: $(\bigwedge v.[\varphi\ in\ v]) \Longrightarrow [\Box\varphi\ in\ v]$
by (*simp add*: *Semantics.T6*)
lemma *BoxE*[*meta-elim*]: $[\Box\varphi\ in\ v] \Longrightarrow (\bigwedge v.[\varphi\ in\ v])$
by (*simp add*: *Semantics.T6*)
lemma *BoxS*[*meta-subst*]: $[\Box\varphi\ in\ v] = (\forall\ v.[\varphi\ in\ v])$
by (*simp add*: *Semantics.T6*)

A.5.7. Rules for Possibility

lemma *DiaI*[*meta-intro*]: $(\exists\, v.[\varphi\ in\ v]) \Longrightarrow [\Diamond\varphi\ in\ v]$
by (*metis BoxS NotS diamond-def*)
lemma *DiaE*[*meta-elim*]: $[\Diamond\varphi\ in\ v] \Longrightarrow (\exists\, v.[\varphi\ in\ v])$
by (*metis BoxS NotS diamond-def*)
lemma *DiaS*[*meta-subst*]: $[\Diamond\varphi\ in\ v] = (\exists\ v.[\varphi\ in\ v])$
by (*metis BoxS NotS diamond-def*)

A.5.8. Rules for Quantification

lemma *AllI*[*meta-intro*]: $(\bigwedge x.\ [\varphi\ x\ in\ v]) \Longrightarrow [\forall\ x.\ \varphi\ x\ in\ v]$
by (*auto simp*: *T8*)
lemma *AllE*[*meta-elim*]: $[\forall\, x.\ \varphi\ x\ in\ v] \Longrightarrow (\bigwedge x.[\varphi\ x\ in\ v])$
by (*auto simp*: *T8*)
lemma *AllS*[*meta-subst*]: $[\forall\, x.\ \varphi\ x\ in\ v] = (\forall\, x.[\varphi\ x\ in\ v])$
by (*auto simp*: *T8*)

A.5.8.1. Rules for Existence

lemma *ExIRule*: $([\varphi\ y\ in\ v]) \Longrightarrow [\exists\, x.\ \varphi\ x\ in\ v]$
by (*auto simp*: *exists-def Semantics.T8 Semantics.T4*)
lemma *ExI*[*meta-intro*]: $(\exists\ y\ .\ [\varphi\ y\ in\ v]) \Longrightarrow [\exists\, x.\ \varphi\ x\ in\ v]$
by (*auto simp*: *exists-def Semantics.T8 Semantics.T4*)
lemma *ExE*[*meta-elim*]: $[\exists\, x.\ \varphi\ x\ in\ v] \Longrightarrow (\exists\ y\ .\ [\varphi\ y\ in\ v])$
by (*auto simp*: *exists-def Semantics.T8 Semantics.T4*)
lemma *ExS*[*meta-subst*]: $[\exists\, x.\ \varphi\ x\ in\ v] = (\exists\ y\ .\ [\varphi\ y\ in\ v])$
by (*auto simp*: *exists-def Semantics.T8 Semantics.T4*)
lemma *ExERule*: **assumes** $[\exists\, x.\ \varphi\ x\ in\ v]$ **obtains** x **where** $[\varphi\ x\ in\ v]$
using *ExE assms* **by** *auto*

A.5.9. Rules for Actuality

lemma *ActualI*[*meta-intro*]: $[\varphi\ in\ dw] \Longrightarrow [\mathcal{A}\varphi\ in\ v]$
by (*auto simp*: *Semantics.T7*)
lemma *ActualE*[*meta-elim*]: $[\mathcal{A}\varphi\ in\ v] \Longrightarrow [\varphi\ in\ dw]$
by (*auto simp*: *Semantics.T7*)
lemma *ActualS*[*meta-subst*]: $[\mathcal{A}\varphi\ in\ v] = [\varphi\ in\ dw]$
by (*auto simp*: *Semantics.T7*)

A.5.10. Rules for Encoding

lemma *EncI*[*meta-intro*]:
assumes $\exists\ r\ o_1\ .\ Some\ r = d_1\ F \wedge Some\ o_1 = d_\kappa\ x \wedge o_1 \in en\ r$
shows $[\!\{x,F\}\!]\ in\ v$
using *assms* **by** (*auto simp*: *Semantics.T2*)
lemma *EncE*[*meta-elim*]:
assumes $[\!\{x,F\}\!]\ in\ v$
shows $\exists\ r\ o_1\ .\ Some\ r = d_1\ F \wedge Some\ o_1 = d_\kappa\ x \wedge o_1 \in en\ r$
using *assms* **by** (*auto simp*: *Semantics.T2*)
lemma *EncS*[*meta-subst*]:
$[\!\{x,F\}\!]\ in\ v] = (\exists\ r\ o_1\ .\ Some\ r = d_1\ F \wedge Some\ o_1 = d_\kappa\ x \wedge o_1 \in en\ r)$
by (*auto simp*: *Semantics.T2*)

A.5.11. Rules for Exemplification

A.5.11.1. Zero-place Relations

lemma *Exe0I*[*meta-intro*]:
assumes $\exists\ r\ .\ Some\ r = d_0\ p \wedge ex0\ r\ v$
shows [(|p|) *in* v]
using *assms* **by** (*auto simp*: *Semantics.T3*)
lemma *Exe0E*[*meta-elim*]:
assumes [(|p|) *in* v]
shows $\exists\ r\ .\ Some\ r = d_0\ p \wedge ex0\ r\ v$
using *assms* **by** (*auto simp*: *Semantics.T3*)
lemma *Exe0S*[*meta-subst*]:
[(|p|) *in* v] = ($\exists\ r\ .\ Some\ r = d_0\ p \wedge ex0\ r\ v$)
by (*auto simp*: *Semantics.T3*)

A.5.11.2. One-Place Relations

lemma *Exe1I*[*meta-intro*]:
assumes $\exists\ r\ o_1\ .\ Some\ r = d_1\ F \wedge Some\ o_1 = d_\kappa\ x \wedge o_1 \in ex1\ r\ v$
shows [(|F,x|) *in* v]
using *assms* **by** (*auto simp*: *Semantics.T1-1*)
lemma *Exe1E*[*meta-elim*]:
assumes [(|F,x|) *in* v]
shows $\exists\ r\ o_1\ .\ Some\ r = d_1\ F \wedge Some\ o_1 = d_\kappa\ x \wedge o_1 \in ex1\ r\ v$
using *assms* **by** (*auto simp*: *Semantics.T1-1*)
lemma *Exe1S*[*meta-subst*]:
[(|F,x|) *in* v] = ($\exists\ r\ o_1\ .\ Some\ r = d_1\ F \wedge Some\ o_1 = d_\kappa\ x \wedge o_1 \in ex1\ r\ v$)
by (*auto simp*: *Semantics.T1-1*)

A.5.11.3. Two-Place Relations

lemma *Exe2I*[*meta-intro*]:
assumes $\exists\ r\ o_1\ o_2\ .\ Some\ r = d_2\ F \wedge Some\ o_1 = d_\kappa\ x$
$\wedge\ Some\ o_2 = d_\kappa\ y \wedge (o_1,\ o_2) \in ex2\ r\ v$
shows [(|F,x,y|) *in* v]
using *assms* **by** (*auto simp*: *Semantics.T1-2*)
lemma *Exe2E*[*meta-elim*]:
assumes [(|F,x,y|) *in* v]
shows $\exists\ r\ o_1\ o_2\ .\ Some\ r = d_2\ F \wedge Some\ o_1 = d_\kappa\ x$
$\wedge\ Some\ o_2 = d_\kappa\ y \wedge (o_1,\ o_2) \in ex2\ r\ v$
using *assms* **by** (*auto simp*: *Semantics.T1-2*)
lemma *Exe2S*[*meta-subst*]:
[(|F,x,y|) *in* v] = ($\exists\ r\ o_1\ o_2\ .\ Some\ r = d_2\ F \wedge Some\ o_1 = d_\kappa\ x$
$\wedge\ Some\ o_2 = d_\kappa\ y \wedge (o_1,\ o_2) \in ex2\ r\ v$)
by (*auto simp*: *Semantics.T1-2*)

A.5.11.4. Three-Place Relations

lemma *Exe3I*[*meta-intro*]:
assumes $\exists\ r\ o_1\ o_2\ o_3\ .\ Some\ r = d_3\ F \wedge Some\ o_1 = d_\kappa\ x$
$\wedge\ Some\ o_2 = d_\kappa\ y \wedge Some\ o_3 = d_\kappa\ z$
$\wedge\ (o_1,\ o_2,\ o_3) \in ex3\ r\ v$
shows [(|F,x,y,z|) *in* v]
using *assms* **by** (*auto simp*: *Semantics.T1-3*)
lemma *Exe3E*[*meta-elim*]:
assumes [(|F,x,y,z|) *in* v]
shows $\exists\ r\ o_1\ o_2\ o_3\ .\ Some\ r = d_3\ F \wedge Some\ o_1 = d_\kappa\ x$
$\wedge\ Some\ o_2 = d_\kappa\ y \wedge Some\ o_3 = d_\kappa\ z$

$\wedge\ (o_1,\ o_2,\ o_3) \in ex3\ r\ v$
using *assms* **by** (*auto simp*: *Semantics.T1-3*)
lemma *Exe3S*[*meta-subst*]:
$[(\!|F,x,y,z|\!)\ in\ v] = (\exists\ r\ o_1\ o_2\ o_3\ .\ Some\ r = d_3\ F \wedge Some\ o_1 = d_\kappa\ x \wedge Some\ o_2 = d_\kappa\ y \wedge Some\ o_3 = d_\kappa\ z \wedge (o_1,\ o_2,\ o_3) \in ex3\ r\ v)$
by (*auto simp*: *Semantics.T1-3*)

A.5.12. Rules for Being Ordinary

lemma *OrdI*[*meta-intro*]:
assumes $\exists\ o_1\ y.\ Some\ o_1 = d_\kappa\ x \wedge o_1 = \omega\nu\ y$
shows $[(\!|O!,x|\!)\ in\ v]$
proof −
have *IsProperInX* $(\lambda x.\ \Diamond(\!|E!,x|\!))$
by *show-proper*
moreover have $[\Diamond(\!|E!,x|\!)\ in\ v]$
apply *meta-solver*
using *ConcretenessSemantics1 propex*$_1$ *assms* **by** *fast*
ultimately show $[(\!|O!,x|\!)\ in\ v]$
unfolding *Ordinary-def*
using *D5-1 propex*$_1$ *assms ConcretenessSemantics1 Exe1S*
by *blast*
qed
lemma *OrdE*[*meta-elim*]:
assumes $[(\!|O!,x|\!)\ in\ v]$
shows $\exists\ o_1\ y.\ Some\ o_1 = d_\kappa\ x \wedge o_1 = \omega\nu\ y$
proof −
have $\exists r\ o_1.\ Some\ r = d_1\ O! \wedge Some\ o_1 = d_\kappa\ x \wedge o_1 \in ex1\ r\ v$
using *assms Exe1E* **by** *simp*
moreover have *IsProperInX* $(\lambda x.\ \Diamond(\!|E!,x|\!))$
by *show-proper*
ultimately have $[\Diamond(\!|E!,x|\!)\ in\ v]$
using *D5-1* **unfolding** *Ordinary-def* **by** *fast*
thus *?thesis*
apply − **apply** *meta-solver*
using *ConcretenessSemantics2* **by** *blast*
qed
lemma *OrdS*[*meta-cong*]:
$[(\!|O!,x|\!)\ in\ v] = (\exists\ o_1\ y.\ Some\ o_1 = d_\kappa\ x \wedge o_1 = \omega\nu\ y)$
using *OrdI OrdE* **by** *blast*

A.5.13. Rules for Being Abstract

lemma *AbsI*[*meta-intro*]:
assumes $\exists\ o_1\ y.\ Some\ o_1 = d_\kappa\ x \wedge o_1 = \alpha\nu\ y$
shows $[(\!|A!,x|\!)\ in\ v]$
proof −
have *IsProperInX* $(\lambda x.\ \neg\Diamond(\!|E!,x|\!))$
by *show-proper*
moreover have $[\neg\Diamond(\!|E!,x|\!)\ in\ v]$
apply *meta-solver*
using *ConcretenessSemantics2 propex*$_1$ *assms*
by (*metis* ν.*distinct*(*1*) *option.sel*)
ultimately show $[(\!|A!,x|\!)\ in\ v]$
unfolding *Abstract-def*
using *D5-1 propex*$_1$ *assms ConcretenessSemantics1 Exe1S*
by *blast*

qed
lemma *AbsE*[*meta-elim*]:
assumes [(|*A*!,*x*|) *in v*]
shows ∃ *o*$_1$ *y*. *Some o*$_1$ = *d*$_\kappa$ *x* ∧ *o*$_1$ = *α*ν *y*
proof −
have *1*: *IsProperInX* (*λx*. ¬◇(|*E*!,*x*|))
by *show-proper*
have ∃ *r o*$_1$. *Some r* = *d*$_1$ *A*! ∧ *Some o*$_1$ = *d*$_\kappa$ *x* ∧ *o*$_1$ ∈ *ex1 r v*
using *assms Exe1E* **by** *simp*
moreover hence [¬◇(|*E*!,*x*|) *in v*]
using *D5-1*[*OF 1*]
unfolding *Abstract-def* **by** *fast*
ultimately show *?thesis*
apply − **apply** *meta-solver*
using *ConcretenessSemantics1 propex*$_1$
by (*metis ν.exhaust*)
qed
lemma *AbsS*[*meta-cong*]:
[(|*A*!,*x*|) *in v*] = (∃ *o*$_1$ *y*. *Some o*$_1$ = *d*$_\kappa$ *x* ∧ *o*$_1$ = *α*ν *y*)
using *AbsI AbsE* **by** *blast*

A.5.14. Rules for Definite Descriptions

lemma *TheEqI*:
assumes ⋀*x*. [*φ x in dw*] = [*ψ x in dw*]
shows (*ι*x. *φ x*) = (*ι*x. *ψ x*)
proof −
have *1*: *d*$_\kappa$ (*ι*x. *φ x*) = *d*$_\kappa$ (*ι*x. *ψ x*)
using *assms D3* **unfolding** *w*$_0$*-def* **by** *simp*
{
assume ∃ *o*$_1$. *Some o*$_1$ = *d*$_\kappa$ (*ι*x. *φ x*)
hence *?thesis* **using** *1 d*$_\kappa$*-inject* **by** *force*
}
moreover {
assume ¬(∃ *o*$_1$. *Some o*$_1$ = *d*$_\kappa$ (*ι*x. *φ x*))
hence *?thesis* **using** *1 D3*
by (*metis d*$_\kappa$*.rep-eq evalκ-inverse*)
}
ultimately show *?thesis* **by** *blast*
qed

A.5.15. Rules for Identity

A.5.15.1. Ordinary Objects

lemma *Eq*$_E$*I*[*meta-intro*]:
assumes ∃ *o*$_1$ *o*$_2$. *Some* (*ων o*$_1$) = *d*$_\kappa$ *x* ∧ *Some* (*ων o*$_2$) = *d*$_\kappa$ *y* ∧ *o*$_1$ = *o*$_2$
shows [*x* =$_E$ *y in v*]
proof −
obtain *o*$_1$ *o*$_2$ **where** *1*:
Some (*ων o*$_1$) = *d*$_\kappa$ *x* ∧ *Some* (*ων o*$_2$) = *d*$_\kappa$ *y* ∧ *o*$_1$ = *o*$_2$
using *assms* **by** *auto*
obtain *r* **where** *2*:
Some r = *d*$_2$ *basic-identity*$_E$
using *propex*$_2$ **by** *auto*
have [(|*O*!,*x*|) & (|*O*!,*y*|) & □(∀ *F*. (|*F*,*x*|) ≡ (|*F*,*y*|)) *in v*]
proof −

```
    have [(|O!,x|) in v] ∧ [(|O!,y|) in v]
      using OrdI 1 by blast
    moreover have [□(∀ F. (|F,x|) ≡ (|F,y|)) in v]
      apply meta-solver using 1 by force
    ultimately show ?thesis using ConjI by simp
  qed
  moreover have IsProperInXY (λ x y . (|O!,x|) & (|O!,y|) & □(∀ F. (|F,x|) ≡ (|F,y|)))
    by show-proper
  ultimately have (ων o1, ων o2) ∈ ex2 r v
    using D5-2 1 2
    unfolding basic-identity_E-def by fast
  thus [x =_E y in v]
    using Exe2I 1 2
    unfolding basic-identity_E-infix-def basic-identity_E-def
    by blast
qed
lemma Eq_EE[meta-elim]:
  assumes [x =_E y in v]
  shows ∃ o1 o2. Some (ων o1) = d_κ x ∧ Some (ων o2) = d_κ y ∧ o1 = o2
proof −
  have IsProperInXY (λ x y . (|O!,x|) & (|O!,y|) & □(∀ F. (|F,x|) ≡ (|F,y|)))
    by show-proper
  hence 1: [(|O!,x|) & (|O!,y|) & □(∀ F. (|F,x|) ≡ (|F,y|)) in v]
    using assms unfolding basic-identity_E-def basic-identity_E-infix-def
    using D4-2 T1-2 D5-2 by meson
  hence 2: ∃ o1 o2 . Some (ων o1) = d_κ x
                   ∧ Some (ων o2) = d_κ y
    apply (subst (asm) ConjS)
    apply (subst (asm) ConjS)
    using OrdE by auto
  then obtain o1 o2 where 3:
    Some (ων o1) = d_κ x ∧ Some (ων o2) = d_κ y
    by auto
  have ∃ r . Some r = d1 (λ z . makeo (λ w s . d_κ (z^P) = Some (ων o1)))
    using prope x1 by auto
  then obtain r where 4:
    Some r = d1 (λ z . makeo (λ w s . d_κ (z^P) = Some (ων o1)))
    by auto
  hence 5: r = (λu s w. ∃ x . νυ x = u ∧ Some x = Some (ων o1))
    unfolding lambdabinder1-def d1-def d_κ-proper
    apply transfer
    by simp
  have [□(∀ F. (|F,x|) ≡ (|F,y|)) in v]
    using 1 using ConjE by blast
  hence 6: ∀ v F . [(|F,x|) in v] ⟷ [(|F,y|) in v]
    using BoxE EquivE AllE by fast
  hence ∀ v . ((ων o1) ∈ ex1 r v) = ((ων o2) ∈ ex1 r v)
    using 2 4 unfolding valid-in-def
    by (metis 3 6 d1.rep-eq d_κ-inject d_κ-proper ex1-def evalo-inverse exe1.rep-eq
        mem-Collect-eq option.sel rep-proper-id νκ-proper valid-in.abs-eq)
  moreover have (ων o1) ∈ ex1 r v
    unfolding 5 ex1-def by simp
  ultimately have (ων o2) ∈ ex1 r v
    by auto
  hence o1 = o2 unfolding 5 ex1-def by (auto simp: meta-aux)
  thus ?thesis
    using 3 by auto
qed
```

lemma *Eq$_E$S[meta-subst]*:
  $[x =_E y \ in \ v] = (\exists\ o_1\ o_2.\ Some\ (\omega\nu\ o_1) = d_\kappa\ x \wedge Some\ (\omega\nu\ o_2) = d_\kappa\ y \wedge o_1 = o_2)$
  using *Eq$_E$I Eq$_E$E* **by** *blast*

A.5.15.2. Individuals

lemma *EqκI[meta-intro]*:
  assumes $\exists\ o_1\ o_2.\ Some\ o_1 = d_\kappa\ x \wedge Some\ o_2 = d_\kappa\ y \wedge o_1 = o_2$
  shows $[x =_\kappa y \ in \ v]$
proof −
  have $x = y$ **using** *assms d$_\kappa$-inject* **by** *meson*
  moreover have $[x =_\kappa x \ in \ v]$
    unfolding *basic-identity$_\kappa$-def*
    apply *meta-solver*
    by (*metis* (*no-types*, *lifting*) *assms AbsI Exe1E ν.exhaust*)
  ultimately show *?thesis* **by** *auto*
qed
lemma *Eqκ-prop*:
  assumes $[x =_\kappa y \ in \ v]$
  shows $[\varphi\ x \ in \ v] = [\varphi\ y \ in \ v]$
proof −
  have $[x =_E y \vee (\!|A!,x|\!) \ \&\ (\!|A!,y|\!) \ \&\ \Box(\forall\ F.\ \{\!|x,F|\!\} \equiv \{\!|y,F|\!})\ in\ v]$
    using *assms* **unfolding** *basic-identity$_\kappa$-def* **by** *simp*
  moreover {
    assume $[x =_E y \ in \ v]$
    hence $(\exists\ o_1\ o_2.\ Some\ o_1 = d_\kappa\ x \wedge Some\ o_2 = d_\kappa\ y \wedge o_1 = o_2)$
      using *Eq$_E$E* **by** *fast*
  }
  moreover {
    assume *1*: $[(\!|A!,x|\!) \ \&\ (\!|A!,y|\!) \ \&\ \Box(\forall\ F.\ \{\!|x,F|\!\} \equiv \{\!|y,F|\!})\ in\ v]$
    hence *2*: $(\exists\ o_1\ o_2\ X\ Y.\ Some\ o_1 = d_\kappa\ x \wedge Some\ o_2 = d_\kappa\ y \wedge o_1 = \alpha\nu\ X \wedge o_2 = \alpha\nu\ Y)$
      using *AbsE ConjE* **by** *meson*
    moreover then obtain $o_1\ o_2\ X\ Y$ **where** *3*:
      $Some\ o_1 = d_\kappa\ x \wedge Some\ o_2 = d_\kappa\ y \wedge o_1 = \alpha\nu\ X \wedge o_2 = \alpha\nu\ Y$
      by *auto*
    moreover have *4*: $[\Box(\forall\ F.\ \{\!|x,F|\!\} \equiv \{\!|y,F|\!})\ in\ v]$
      using *1 ConjE* **by** *blast*
    hence *6*: $\forall\ v\ F\ .\ [\{\!|x,F|\!\}\ in\ v] \longleftrightarrow [\{\!|y,F|\!\}\ in\ v]$
      using *BoxE AllE EquivE* **by** *fast*
    hence *7*: $\forall\, v\ r.\ (\exists\ o_1.\ Some\ o_1 = d_\kappa\ x \wedge o_1 \in en\ r) = (\exists\ o_1.\ Some\ o_1 = d_\kappa\ y \wedge o_1 \in en\ r)$
      apply − **apply** *meta-solver*
      using *propex$_1$ d$_1$-inject* **apply** *simp*
      apply *transfer* **by** *simp*
    hence *8*: $\forall\ r.\ (o_1 \in en\ r) = (o_2 \in en\ r)$
      using *3 d$_\kappa$-inject d$_\kappa$-proper* **apply** *simp*
      by (*metis option.inject*)
    hence $\forall\, r.\ (o_1 \in r) = (o_2 \in r)$
      unfolding *en-def* **using** *3*
      by (*metis Collect-cong Collect-mem-eq ν.simps*(*6*) *mem-Collect-eq make*Π_1*-cases*)
    hence $(o_1 \in \{\ x\ .\ o_1 = x\ \}) = (o_2 \in \{\ x\ .\ o_1 = x\ \})$
      by *metis*
    hence $o_1 = o_2$ **by** *simp*
    hence $(\exists\ o_1\ o_2.\ Some\ o_1 = d_\kappa\ x \wedge Some\ o_2 = d_\kappa\ y \wedge o_1 = o_2)$
      using *3* **by** *auto*

}
ultimately have $x = y$
using *DisjS* **using** *Semantics.*d_κ*-inject* **by** *auto*
thus $(v \models (\varphi\ x)) = (v \models (\varphi\ y))$ **by** *simp*
qed
lemma *Eq*κ*E*[*meta-elim*]:
assumes $[x =_\kappa y\ in\ v]$
shows $\exists\ o_1\ o_2.\ Some\ o_1 = d_\kappa\ x \wedge Some\ o_2 = d_\kappa\ y \wedge o_1 = o_2$
proof −
have $\forall\ \varphi\ .\ (v \models \varphi\ x) = (v \models \varphi\ y)$
using *assms Eq*κ*-prop* **by** *blast*
moreover obtain φ **where** φ*-prop*:
$\varphi = (\lambda\ \alpha\ .\ makeo\ (\lambda\ w\ s\ .\ (\exists\ o_1\ o_2.\ Some\ o_1 = d_\kappa\ x$
$\wedge\ Some\ o_2 = d_\kappa\ \alpha \wedge o_1 = o_2)))$
by *auto*
ultimately have $(v \models \varphi\ x) = (v \models \varphi\ y)$ **by** *metis*
moreover have $(v \models \varphi\ x)$
using *assms* **unfolding** φ*-prop basic-identity*$_\kappa$*-def*
by (*metis* (*mono-tags*, *lifting*) *AbsS ConjE DisjS*
$Eq_E S$ *valid-in.abs-eq*)
ultimately have $(v \models \varphi\ y)$ **by** *auto*
thus *?thesis*
unfolding φ*-prop*
by (*simp add*: *valid-in-def meta-aux*)
qed
lemma *Eq*κ*S*[*meta-subst*]:
$[x =_\kappa y\ in\ v] = (\exists\ o_1\ o_2.\ Some\ o_1 = d_\kappa\ x \wedge Some\ o_2 = d_\kappa\ y \wedge o_1 = o_2)$
using *Eq*κ*I Eq*κ*E* **by** *blast*

A.5.15.3. One-Place Relations

lemma $Eq_1 I$[*meta-intro*]: $F = G \Longrightarrow [F =_1 G\ in\ v]$
unfolding *basic-identity*$_1$*-def*
apply (*rule BoxI*, *rule AllI*, *rule EquivI*)
by *simp*
lemma $Eq_1 E$[*meta-elim*]: $[F =_1 G\ in\ v] \Longrightarrow F = G$
unfolding *basic-identity*$_1$*-def*
apply (*drule BoxE*, *drule-tac* x=($\alpha\nu$ { F }) **in** *AllE*, *drule EquivE*)
apply (*simp add*: *Semantics.T2*)
unfolding *en-def* d_κ*-def* d_1*-def*
using $\nu\kappa$*-proper rep-proper-id*
by (*simp add*: *rep-def proper-def meta-aux* $\nu\kappa$*.rep-eq*)
lemma $Eq_1 S$[*meta-subst*]: $[F =_1 G\ in\ v] = (F = G)$
using $Eq_1 I$ $Eq_1 E$ **by** *auto*
lemma Eq_1*-prop*: $[F =_1 G\ in\ v] \Longrightarrow [\varphi\ F\ in\ v] = [\varphi\ G\ in\ v]$
using $Eq_1 E$ **by** *blast*

A.5.15.4. Two-Place Relations

lemma $Eq_2 I$[*meta-intro*]: $F = G \Longrightarrow [F =_2 G\ in\ v]$
unfolding *basic-identity*$_2$*-def*
apply (*rule AllI*, *rule ConjI*, (*subst* $Eq_1 S$)+)
by *simp*
lemma $Eq_2 E$[*meta-elim*]: $[F =_2 G\ in\ v] \Longrightarrow F = G$
proof −
assume $[F =_2 G\ in\ v]$
hence *1*: $[\forall\ x.\ (\boldsymbol{\lambda} y.\ (\!|F,x^P,y^P|\!)) =_1 (\boldsymbol{\lambda} y.\ (\!|G,x^P,y^P|\!))\ in\ v]$
unfolding *basic-identity*$_2$*-def*

```
    apply − apply meta-solver by auto
  {
    fix u v s w
    obtain x where x-def: νυ x = v by (metis νυ-surj surj-def)
    obtain a where a-def:
      a = (λu s w. ∃xa. νυ xa = u ∧ evalΠ2 F (νυ x) (νυ xa) s w)
      by auto
    obtain b where b-def:
      b = (λu s w. ∃xa. νυ xa = u ∧ evalΠ2 G (νυ x) (νυ xa) s w)
      by auto
    have a = b unfolding a-def b-def
      using 1 apply − apply meta-solver
      by (auto simp: meta-defs meta-aux makeΠ1-inject)
    hence a u s w = b u s w by auto
    hence (evalΠ2 F (νυ x) u s w) = (evalΠ2 G (νυ x) u s w)
      unfolding a-def b-def
      by (metis (no-types, hide-lams) νυ-surj surj-def)
    hence (evalΠ2 F v u s w) = (evalΠ2 G v u s w)
      unfolding x-def by auto
  }
  hence (evalΠ2 F) = (evalΠ2 G) by blast
  thus F = G by (simp add: evalΠ2-inject)
qed
lemma Eq2S[meta-subst]: [F =2 G in v] = (F = G)
  using Eq2I Eq2E by auto
lemma Eq2-prop: [F =2 G in v] ⟹ [φ F in v] = [φ G in v]
  using Eq2E by blast
```

A.5.15.5. Three-Place Relations

```
lemma Eq3I[meta-intro]: F = G ⟹ [F =3 G in v]
  apply (simp add: meta-defs meta-aux conn-defs forall-ν-def basic-identity3-def)
  using MetaSolver.Eq1I valid-in.rep-eq by auto
lemma Eq3E[meta-elim]: [F =3 G in v] ⟹ F = G
proof −

  assume [F =3 G in v]
  hence 1: [∀ x y. (λz. (|F,x^P,y^P,z^P|)) =1 (λz. (|G,x^P,y^P,z^P|)) in v]
    unfolding basic-identity3-def
    apply − apply meta-solver by auto
  {
    fix u v r s w
    obtain x where x-def: νυ x = v by (metis νυ-surj surj-def)
    obtain y where y-def: νυ y = r by (metis νυ-surj surj-def)
    obtain a where a-def:
      a = (λu s w. ∃xa. νυ xa = u ∧ evalΠ3 F (νυ x) (νυ y) (νυ xa) s w)
      by auto
    obtain b where b-def:
      b = (λu s w. ∃xa. νυ xa = u ∧ evalΠ3 G (νυ x) (νυ y) (νυ xa) s w)
      by auto
    have a = b unfolding a-def b-def
      using 1 apply − apply meta-solver
      by (auto simp: meta-defs meta-aux makeΠ1-inject)
    hence a u s w = b u s w by auto
    hence (evalΠ3 F (νυ x) (νυ y) u s w) = (evalΠ3 G (νυ x) (νυ y) u s w)
      unfolding a-def b-def
      by (metis (no-types, hide-lams) νυ-surj surj-def)
    hence (evalΠ3 F v r u s w) = (evalΠ3 G v r u s w)
```

```
    unfolding x-def y-def by auto
  }
  hence (evalΠ3 F) = (evalΠ3 G) by blast
  thus F = G by (simp add: evalΠ3-inject)
qed
lemma Eq3S[meta-subst]: [F =3 G in v] = (F = G)
  using Eq3I Eq3E by auto
lemma Eq3-prop: [F =3 G in v] ⟹ [φ F in v] = [φ G in v]
  using Eq3E by blast
```

A.5.15.6. Propositions

```
lemma Eq0I[meta-intro]: x = y ⟹ [x =0 y in v]
  unfolding basic-identity0-def by (simp add: Eq1S)
lemma Eq0E[meta-elim]: [F =0 G in v] ⟹ F = G
  proof −
    assume [F =0 G in v]
    hence [(λy. F) =1 (λy. G) in v]
      unfolding basic-identity0-def by simp
    hence (λy. F) = (λy. G)
      using Eq1S by simp
    hence (λu s w. (∃ x. νυ x = u) ∧ evalo F s w)
        = (λu s w. (∃ x. νυ x = u) ∧ evalo G s w)
      apply (simp add: meta-defs meta-aux)
      by (metis (no-types, lifting) UNIV-I makeΠ1-inverse)
    hence ⋀s w.(evalo F s w) = (evalo G s w)
      by metis
    hence (evalo F) = (evalo G) by blast
    thus F = G
    by (metis evalo-inverse)
  qed
lemma Eq0S[meta-subst]: [F =0 G in v] = (F = G)
  using Eq0I Eq0E by auto
lemma Eq0-prop: [F =0 G in v] ⟹ [φ F in v] = [φ G in v]
  using Eq0E by blast

end
```

A.6. General Identity

Remark. *In order to define a general identity symbol that can act on all types of terms a type class is introduced which assumes the substitution property which is needed to derive the corresponding axiom. This type class is instantiated for all relation types, individual terms and individuals.*

A.6.1. Type Classes

```
class identifiable =
fixes identity :: 'a⇒'a⇒o (infixl = 63)
assumes l-identity:
  w ⊨ x = y ⟹ w ⊨ φ x ⟹ w ⊨ φ y
begin
  abbreviation notequal (infixl ≠ 63) where
    notequal ≡ λ x y . ¬(x = y)
```

end

class *quantifiable-and-identifiable* = *quantifiable* + *identifiable*
begin
  definition *exists-unique*::('*a*⇒o)⇒o (**binder** ∃! [*8*] *9*) **where**
    exists-unique ≡ λ φ . ∃ α . φ α **&** (∀ β. φ β → β = α)

  declare *exists-unique-def*[*conn-defs*]
end

A.6.2. Instantiations

instantiation κ :: *identifiable*
begin
  definition *identity-κ* **where** *identity-κ* ≡ *basic-identity*$_\kappa$
  instance proof
    fix *x y* :: κ **and** *w* φ
    show [*x* = *y in w*] ⟹ [φ *x in w*] ⟹ [φ *y in w*]
      unfolding *identity-κ-def*
      using *MetaSolver.Eqκ-prop* **..**
  qed
end

instantiation ν :: *identifiable*
begin
  definition *identity-ν* **where** *identity-ν* ≡ λ *x y* . $x^P = y^P$
  instance proof
    fix α :: ν **and** β :: ν **and** *v* φ
    assume *v* ⊨ α = β
    hence *v* ⊨ $\alpha^P = \beta^P$
      unfolding *identity-ν-def* **by** *auto*
    hence ⋀φ.(*v* ⊨ φ (α^P)) ⟹ (*v* ⊨ φ (β^P))
      using *l-identity* **by** *auto*
    hence (*v* ⊨ φ (*rep* (α^P))) ⟹ (*v* ⊨ φ (*rep* (β^P)))
      by *meson*
    thus (*v* ⊨ φ α) ⟹ (*v* ⊨ φ β)
      by (*simp only*: *rep-proper-id*)
  qed
end

instantiation Π_1 :: *identifiable*
begin
  definition *identity*-Π_1 **where** *identity*-Π_1 ≡ *basic-identity*$_1$
  instance proof
    fix *F G* :: Π_1 **and** *w* φ
    show (*w* ⊨ *F* = *G*) ⟹ (*w* ⊨ φ *F*) ⟹ (*w* ⊨ φ *G*)
      unfolding *identity*-Π_1-*def* **using** *MetaSolver.Eq*$_1$-*prop* **..**
  qed
end

instantiation Π_2 :: *identifiable*
begin
  definition *identity*-Π_2 **where** *identity*-Π_2 ≡ *basic-identity*$_2$
  instance proof
    fix *F G* :: Π_2 **and** *w* φ
    show (*w* ⊨ *F* = *G*) ⟹ (*w* ⊨ φ *F*) ⟹ (*w* ⊨ φ *G*)
      unfolding *identity*-Π_2-*def* **using** *MetaSolver.Eq*$_2$-*prop* **..**
  qed

end

instantiation Π_3 :: *identifiable*
begin
definition *identity-*Π_3 **where** *identity-*$\Pi_3 \equiv$ *basic-identity*$_3$
instance proof
fix $F\ G :: \Pi_3$ **and** $w\ \varphi$
show $(w \models F = G) \Longrightarrow (w \models \varphi\ F) \Longrightarrow (w \models \varphi\ G)$
unfolding *identity-*Π_3*-def* **using** *MetaSolver.Eq*$_3$*-prop* **..**
qed
end

instantiation o :: *identifiable*
begin
definition *identity-*o **where** *identity-*o $\equiv$ *basic-identity*$_0$
instance proof
fix $F\ G$ :: o **and** $w\ \varphi$
show $(w \models F = G) \Longrightarrow (w \models \varphi\ F) \Longrightarrow (w \models \varphi\ G)$
unfolding *identity-*o*-def* **using** *MetaSolver.Eq*$_0$*-prop* **..**
qed
end

instance ν :: *quantifiable-and-identifiable* **..**
instance Π_1 :: *quantifiable-and-identifiable* **..**
instance Π_2 :: *quantifiable-and-identifiable* **..**
instance Π_3 :: *quantifiable-and-identifiable* **..**
instance o :: *quantifiable-and-identifiable* **..**

A.6.3. New Identity Definitions

Remark. *The basic definitions of identity use type specific quantifiers and identity symbols. Equivalent definitions that use the general identity symbol and general quantifiers are provided.*

named-theorems *identity-defs*
lemma *identity*$_E$*-def*[*identity-defs*]:
$\mathit{basic\text{-}identity}_E \equiv \boldsymbol{\lambda}^2\ (\lambda x\ y.\ (\!|O!,x^P|\!)\ \&\ (\!|O!,y^P|\!)\ \&\ \Box(\forall\ F.\ (\!|F,x^P|\!) \equiv (\!|F,y^P|\!)))$
unfolding *basic-identity*$_E$*-def forall-*Π_1*-def* **by** *simp*
lemma *identity*$_E$*-infix-def*[*identity-defs*]:
$x =_E y \equiv (\!|\mathit{basic\text{-}identity}_E,x,y|\!)$ **using** *basic-identity*$_E$*-infix-def* **.**
lemma *identity*$_\kappa$*-def*[*identity-defs*]:
$op = \equiv \lambda x\ y.\ x =_E y \vee (\!|A!,x|\!)\ \&\ (\!|A!,y|\!)\ \&\ \Box(\forall\ F.\ \{\!|x,F|\!\} \equiv \{\!|y,F|\!\})$
unfolding *identity-*κ*-def basic-identity*$_\kappa$*-def forall-*Π_1*-def* **by** *simp*
lemma *identity*$_\nu$*-def*[*identity-defs*]:
$op = \equiv \lambda x\ y.\ (x^P) =_E (y^P) \vee (\!|A!,x^P|\!)\ \&\ (\!|A!,y^P|\!)\ \&\ \Box(\forall\ F.\ \{\!|x^P,F|\!\} \equiv \{\!|y^P,F|\!\})$
unfolding *identity-*ν*-def identity*$_\kappa$*-def* **by** *simp*
lemma *identity*$_1$*-def*[*identity-defs*]:
$op = \equiv \lambda F\ G.\ \Box(\forall\ x\ .\ \{\!|x^P,F|\!\} \equiv \{\!|x^P,G|\!\})$
unfolding *identity-*Π_1*-def basic-identity*$_1$*-def forall-*ν*-def* **by** *simp*
lemma *identity*$_2$*-def*[*identity-defs*]:
$op = \equiv \lambda F\ G.\ \forall\ x.\ (\boldsymbol{\lambda}y.\ (\!|F,x^P,y^P|\!)) = (\boldsymbol{\lambda}y.\ (\!|G,x^P,y^P|\!))$
$\&\ (\boldsymbol{\lambda}y.\ (\!|F,y^P,x^P|\!)) = (\boldsymbol{\lambda}y.\ (\!|G,y^P,x^P|\!))$
unfolding *identity-*Π_2*-def identity-*Π_1*-def basic-identity*$_2$*-def forall-*ν*-def* **by** *simp*
lemma *identity*$_3$*-def*[*identity-defs*]:
$op = \equiv \lambda F\ G.\ \forall\ x\ y.\ (\boldsymbol{\lambda}z.\ (\!|F,z^P,x^P,y^P|\!)) = (\boldsymbol{\lambda}z.\ (\!|G,z^P,x^P,y^P|\!))$
$\&\ (\boldsymbol{\lambda}z.\ (\!|F,x^P,z^P,y^P|\!)) = (\boldsymbol{\lambda}z.\ (\!|G,x^P,z^P,y^P|\!))$
$\&\ (\boldsymbol{\lambda}z.\ (\!|F,x^P,y^P,z^P|\!)) = (\boldsymbol{\lambda}z.\ (\!|G,x^P,y^P,z^P|\!))$
unfolding *identity-*Π_3*-def identity-*Π_1*-def basic-identity*$_3$*-def forall-*ν*-def* **by** *simp*

lemma *identity$_o$-def*[*identity-defs*]: $op = \equiv \lambda F\ G.\ (\boldsymbol{\lambda} y.\ F) = (\boldsymbol{\lambda} y.\ G)$
unfolding *identity-o-def identity-*Π_1*-def basic-identity$_0$-def* **by** *simp*

A.7. The Axioms of PLM

Remark. *The axioms of PLM can now be derived from the Semantics and the model structure.*

locale *Axioms*
begin
interpretation *MetaSolver* .
interpretation *Semantics* .
named-theorems *axiom*

Remark. *The special syntax* [[-]] *is introduced for stating the axioms. Modally-fragile axioms are stated with the syntax for actual validity* [-].

definition *axiom* :: o⇒*bool* ([[-]]) **where** *axiom* ≡ λ φ . ∀ *v* . [φ *in v*]

method *axiom-meta-solver* = ((((*unfold axiom-def*)*?*, *rule allI*) | (*unfold actual-validity-def*)*?*), *meta-solver*,
(*simp* | (*auto*; *fail*))*?*)

A.7.1. Closures

Remark. *Rules resembling the concepts of closures in PLM are derived. Theorem attributes are introduced to aid in the instantiation of the axioms.*

lemma *axiom-instance*[*axiom*]: [[φ]] ⟹ [φ *in v*]
unfolding *axiom-def* **by** *simp*
lemma *closures-universal*[*axiom*]: (⋀*x*.[[φ *x*]]) ⟹ [[∀ *x*. φ *x*]]
by *axiom-meta-solver*
lemma *closures-actualization*[*axiom*]: [[φ]] ⟹ [[$\mathcal{A}$ φ]]
by *axiom-meta-solver*
lemma *closures-necessitation*[*axiom*]: [[φ]] ⟹ [[□ φ]]
by *axiom-meta-solver*
lemma *necessitation-averse-axiom-instance*[*axiom*]: [φ] ⟹ [φ *in dw*]
by *axiom-meta-solver*
lemma *necessitation-averse-closures-universal*[*axiom*]: (⋀*x*.[φ *x*]) ⟹ [∀ *x*. φ *x*]
by *axiom-meta-solver*

attribute-setup *axiom-instance* = ⟪
Scan.succeed (*Thm.rule-attribute* []
(*fn - => fn thm => thm RS* @{*thm axiom-instance*}))
⟫

attribute-setup *necessitation-averse-axiom-instance* = ⟪
Scan.succeed (*Thm.rule-attribute* []
(*fn - => fn thm => thm RS* @{*thm necessitation-averse-axiom-instance*}))
⟫

attribute-setup *axiom-necessitation* = ⟪
Scan.succeed (*Thm.rule-attribute* []
(*fn - => fn thm => thm RS* @{*thm closures-necessitation*}))

⟫

attribute-setup *axiom-actualization* = ⟪
Scan.succeed (*Thm.rule-attribute* []
(*fn - => fn thm => thm RS* @{*thm closures-actualization*}))
⟫

attribute-setup *axiom-universal* = ⟪
Scan.succeed (*Thm.rule-attribute* []
(*fn - => fn thm => thm RS* @{*thm closures-universal*}))
⟫

A.7.2. Axioms for Negations and Conditionals

lemma *pl-1*[*axiom*]:
[[$\varphi \to (\psi \to \varphi)$]]
by *axiom-meta-solver*
lemma *pl-2*[*axiom*]:
[[$(\varphi \to (\psi \to \chi)) \to ((\varphi \to \psi) \to (\varphi \to \chi))$]]
by *axiom-meta-solver*
lemma *pl-3*[*axiom*]:
[[$(\neg\varphi \to \neg\psi) \to ((\neg\varphi \to \psi) \to \varphi)$]]
by *axiom-meta-solver*

A.7.3. Axioms of Identity

lemma *l-identity*[*axiom*]:
[[$\alpha = \beta \to (\varphi\ \alpha \to \varphi\ \beta)$]]
using *l-identity* **apply** − **by** *axiom-meta-solver*

A.7.4. Axioms of Quantification

lemma *cqt-1*[*axiom*]:
[[$(\forall\ \alpha.\ \varphi\ \alpha) \to \varphi\ \alpha$]]
by *axiom-meta-solver*
lemma *cqt-1-κ*[*axiom*]:
[[$(\forall\ \alpha.\ \varphi\ (\alpha^P)) \to ((\exists\ \beta\ .\ (\beta^P) = \alpha) \to \varphi\ \alpha)$]]
proof −
{
fix *v*
assume *1*: [$(\forall\ \alpha.\ \varphi\ (\alpha^P))$ *in v*]
assume [$(\exists\ \beta\ .\ (\beta^P) = \alpha)$ *in v*]
then obtain β **where** *2*:
[$(\beta^P) = \alpha$ *in v*] **by** (*rule ExERule*)
hence [$\varphi\ (\beta^P)$ *in v*] **using** *1 AllE* **by** *fast*
hence [$\varphi\ \alpha$ *in v*]
using *l-identity*[**where** φ=φ, *axiom-instance*]
ImplS 2 **by** *simp*
}
thus [[$(\forall\ \alpha.\ \varphi\ (\alpha^P)) \to ((\exists\ \beta\ .\ (\beta^P) = \alpha) \to \varphi\ \alpha)$]]
unfolding *axiom-def* **using** *ImplI* **by** *blast*
qed
lemma *cqt-3*[*axiom*]:
[[$(\forall\,\alpha.\ \varphi\ \alpha \to \psi\ \alpha) \to ((\forall\,\alpha.\ \varphi\ \alpha) \to (\forall\,\alpha.\ \psi\ \alpha))$]]
by *axiom-meta-solver*
lemma *cqt-4*[*axiom*]:
[[$\varphi \to (\forall\,\alpha.\ \varphi)$]]

by *axiom-meta-solver*

inductive *SimpleExOrEnc*
where *SimpleExOrEnc* ($\lambda\ x\ .\ (\!|F,x|\!)$)
| *SimpleExOrEnc* ($\lambda\ x\ .\ (\!|F,x,y|\!)$)
| *SimpleExOrEnc* ($\lambda\ x\ .\ (\!|F,y,x|\!)$)
| *SimpleExOrEnc* ($\lambda\ x\ .\ (\!|F,x,y,z|\!)$)
| *SimpleExOrEnc* ($\lambda\ x\ .\ (\!|F,y,x,z|\!)$)
| *SimpleExOrEnc* ($\lambda\ x\ .\ (\!|F,y,z,x|\!)$)
| *SimpleExOrEnc* ($\lambda\ x\ .\ \{\!|x,F|\!\}$)

lemma *cqt-5*[*axiom*]:
assumes *SimpleExOrEnc* ψ
shows $[[(\psi\ (\boldsymbol{\iota}x\ .\ \varphi\ x)) \rightarrow (\exists\ \alpha.\ (\alpha^P) = (\boldsymbol{\iota}x\ .\ \varphi\ x))]]$
proof −
have $\forall\ w\ .\ ([(\psi\ (\boldsymbol{\iota}x\ .\ \varphi\ x))\ in\ w] \longrightarrow (\exists\ o_1\ .\ Some\ o_1 = d_\kappa\ (\boldsymbol{\iota}x\ .\ \varphi\ x)))$
using *assms* **apply** *induct* **by** (*meta-solver*;*metis*)+
thus *?thesis*
apply − **unfolding** *identity-κ-def*
apply *axiom-meta-solver*
using d_κ*-proper* **by** *auto*
qed

lemma *cqt-5-mod*[*axiom*]:
assumes *SimpleExOrEnc* ψ
shows $[[\psi\ \tau \rightarrow (\exists\ \alpha\ .\ (\alpha^P) = \tau)]]$
proof −
have $\forall\ w\ .\ ([(\psi\ \tau)\ in\ w] \longrightarrow (\exists\ o_1\ .\ Some\ o_1 = d_\kappa\ \tau))$
using *assms* **apply** *induct* **by** (*meta-solver*;*metis*)+
thus *?thesis*
apply − **unfolding** *identity-κ-def*
apply *axiom-meta-solver*
using d_κ*-proper* **by** *auto*
qed

A.7.5. Axioms of Actuality

lemma *logic-actual*[*axiom*]: $[(\boldsymbol{\mathcal{A}}\varphi) \equiv \varphi]$
by *axiom-meta-solver*
lemma $[[(\boldsymbol{\mathcal{A}}\varphi) \equiv \varphi]]$
nitpick[*user-axioms*, *expect* = *genuine*, *card* = *1*, *card i* = *2*]
oops — Counter-model by nitpick

lemma *logic-actual-nec-1*[*axiom*]:
$[[\boldsymbol{\mathcal{A}}\neg\varphi \equiv \neg\boldsymbol{\mathcal{A}}\varphi]]$
by *axiom-meta-solver*
lemma *logic-actual-nec-2*[*axiom*]:
$[[(\boldsymbol{\mathcal{A}}(\varphi \rightarrow \psi)) \equiv (\boldsymbol{\mathcal{A}}\varphi \rightarrow \boldsymbol{\mathcal{A}}\psi)]]$
by *axiom-meta-solver*
lemma *logic-actual-nec-3*[*axiom*]:
$[[\boldsymbol{\mathcal{A}}(\forall\,\alpha.\ \varphi\ \alpha) \equiv (\forall\,\alpha.\ \boldsymbol{\mathcal{A}}(\varphi\ \alpha))]]$
by *axiom-meta-solver*
lemma *logic-actual-nec-4*[*axiom*]:
$[[\boldsymbol{\mathcal{A}}\varphi \equiv \boldsymbol{\mathcal{A}}\boldsymbol{\mathcal{A}}\varphi]]$
by *axiom-meta-solver*

A.7.6. Axioms of Necessity

lemma *qml-1*[*axiom*]:
[[$\Box(\varphi \to \psi) \to (\Box\varphi \to \Box\psi)$]]
by *axiom-meta-solver*
lemma *qml-2*[*axiom*]:
[[$\Box\varphi \to \varphi$]]
by *axiom-meta-solver*
lemma *qml-3*[*axiom*]:
[[$\Diamond\varphi \to \Box\Diamond\varphi$]]
by *axiom-meta-solver*
lemma *qml-4*[*axiom*]:
[[$\Diamond(\exists\, x.\ (\!|E!,x^P|\!)$ **&** $\Diamond\neg(\!|E!,x^P|\!))$ **&** $\Diamond\neg(\exists\, x.\ (\!|E!,x^P|\!)$ **&** $\Diamond\neg(\!|E!,x^P|\!))$]]
unfolding *axiom-def*
using *PossiblyContingentObjectExistsAxiom*
PossiblyNoContingentObjectExistsAxiom
apply (*simp add*: *meta-defs meta-aux conn-defs forall-ν-def*
split: *ν.split υ.split*)
by (*metis νυ-ων-is-ωυ υ.distinct*(*1*) *υ.inject*(*1*))

A.7.7. Axioms of Necessity and Actuality

lemma *qml-act-1*[*axiom*]:
[[$\mathcal{A}\varphi \to \Box\mathcal{A}\varphi$]]
by *axiom-meta-solver*
lemma *qml-act-2*[*axiom*]:
[[$\Box\varphi \equiv \mathcal{A}(\Box\varphi)$]]
by *axiom-meta-solver*

A.7.8. Axioms of Descriptions

lemma *descriptions*[*axiom*]:
[[$x^P = (\iota x.\ \varphi\ x) \equiv (\forall\, z.(\mathcal{A}(\varphi\ z) \equiv z = x))$]]
unfolding *axiom-def*
proof (*rule allI*, *rule EquivI*; *rule*)
fix *v*
assume [$x^P = (\iota x.\ \varphi\ x)$ *in v*]
moreover hence *1*:
$\exists\, o_1\ o_2.$ *Some* $o_1 = d_\kappa\ (x^P) \wedge$ *Some* $o_2 = d_\kappa\ (\iota x.\ \varphi\ x) \wedge o_1 = o_2$
apply − **unfolding** *identity-κ-def* **by** *meta-solver*
then obtain $o_1\ o_2$ **where** *2*:
Some $o_1 = d_\kappa\ (x^P) \wedge$ *Some* $o_2 = d_\kappa\ (\iota x.\ \varphi\ x) \wedge o_1 = o_2$
by *auto*
hence *3*:
$(\exists\ x\ .((w_0 \models \varphi\ x) \wedge (\forall\, y.\ (w_0 \models \varphi\ y) \longrightarrow y = x)))$
$\wedge\ d_\kappa\ (\iota x.\ \varphi\ x) =$ *Some* (*THE x*. $(w_0 \models \varphi\ x)$)
using *D3* **by** (*metis option.distinct*(*1*))
then obtain *X* **where** *4*:
$((w_0 \models \varphi\ X) \wedge (\forall\, y.\ (w_0 \models \varphi\ y) \longrightarrow y = X))$
by *auto*
moreover have $o_1 =$ (*THE x*. $(w_0 \models \varphi\ x)$)
using *2 3* **by** *auto*
ultimately have *5*: $X = o_1$
by (*metis* (*mono-tags*) *theI*)
have $\forall\ z\ .$ [$\mathcal{A}\varphi\ z$ *in v*] = [$(z^P) = (x^P)$ *in v*]
proof
fix *z*

```
    have [𝓐φ z in v] ⟹ [(z^P) = (x^P) in v]
      unfolding identity-κ-def apply meta-solver
      using 4 5 2 d_κ-proper w_0-def by auto
    moreover have [(z^P) = (x^P) in v] ⟹ [𝓐φ z in v]
      unfolding identity-κ-def apply meta-solver
      using 2 4 5
      by (simp add: d_κ-proper w_0-def)
    ultimately show [𝓐φ z in v] = [(z^P) = (x^P) in v]
      by auto
  qed
  thus [∀ z. 𝓐φ z ≡ (z) = (x) in v]
    unfolding identity-ν-def
    by (simp add: AllI EquivS)
next
  fix v
  assume [∀ z. 𝓐φ z ≡ (z) = (x) in v]
  hence ⋀z. (dw ⊨ φ z) = (∃ o_1 o_2. Some o_1 = d_κ (z^P)
          ∧ Some o_2 = d_κ (x^P) ∧ o_1 = o_2)
    apply − unfolding identity-ν-def identity-κ-def by meta-solver
  hence ∀ z . (dw ⊨ φ z) = (z = x)
    by (simp add: d_κ-proper)
  moreover hence x = (THE z . (dw ⊨ φ z)) by simp
  ultimately have x^P = (ιx. φ x)
    using D3 d_κ-inject d_κ-proper w_0-def by presburger
  thus [x^P = (ιx. φ x) in v]
    using Eqκ S unfolding identity-κ-def by (metis d_κ-proper)
qed
```

A.7.9. Axioms for Complex Relation Terms

```
lemma lambda-predicates-1[axiom]:
  (λ x . φ x) = (λ y . φ y) ..

lemma lambda-predicates-2-1[axiom]:
  assumes IsProperInX φ
  shows [[(|λ x . φ (x^P), x^P|) ≡ φ (x^P)]]
  apply axiom-meta-solver
  using D5-1[OF assms] d_κ-proper propex_1
  by metis

lemma lambda-predicates-2-2[axiom]:
  assumes IsProperInXY φ
  shows [[(|(λ^2 (λ x y . φ (x^P) (y^P))), x^P, y^P|) ≡ φ (x^P) (y^P)]]
  apply axiom-meta-solver
  using D5-2[OF assms] d_κ-proper propex_2
  by metis

lemma lambda-predicates-2-3[axiom]:
  assumes IsProperInXYZ φ
  shows [[(|(λ^3 (λ x y z . φ (x^P) (y^P) (z^P))),x^P,y^P,z^P|) ≡ φ (x^P) (y^P) (z^P)]]
  proof −
    have [[(|(λ^3 (λ x y z . φ (x^P) (y^P) (z^P))),x^P,y^P,z^P|) → φ (x^P) (y^P) (z^P)]]
      apply axiom-meta-solver using D5-3[OF assms] by auto
    moreover have
      [[φ (x^P) (y^P) (z^P) → (|(λ^3 (λ x y z . φ (x^P) (y^P) (z^P))),x^P,y^P,z^P|)]]
      apply axiom-meta-solver
      using D5-3[OF assms] d_κ-proper propex_3
      by (metis (no-types, lifting))
```

ultimately show *?thesis* **unfolding** *axiom-def equiv-def ConjS* **by** *blast*
qed

lemma *lambda-predicates-3-0*[*axiom*]:
[[($\boldsymbol{\lambda}^0$ φ) = φ]]
unfolding *identity-defs*
apply *axiom-meta-solver*
by (*simp add*: *meta-defs meta-aux*)

lemma *lambda-predicates-3-1*[*axiom*]:
[[($\boldsymbol{\lambda}$ x . (|F, x^P|)) = F]]
unfolding *axiom-def*
apply (*rule allI*)
unfolding *identity*-Π_1-*def* **apply** (*rule* Eq_1I)
using *D4-1* d_1-*inject* **by** *simp*

lemma *lambda-predicates-3-2*[*axiom*]:
[[($\boldsymbol{\lambda}^2$ (λ x y . (|F, x^P, y^P|))) = F]]
unfolding *axiom-def*
apply (*rule allI*)
unfolding *identity*-Π_2-*def* **apply** (*rule* Eq_2I)
using *D4-2* d_2-*inject* **by** *simp*

lemma *lambda-predicates-3-3*[*axiom*]:
[[($\boldsymbol{\lambda}^3$ (λ x y z . (|F, x^P, y^P, z^P|))) = F]]
unfolding *axiom-def*
apply (*rule allI*)
unfolding *identity*-Π_3-*def* **apply** (*rule* Eq_3I)
using *D4-3* d_3-*inject* **by** *simp*

lemma *lambda-predicates-4-0*[*axiom*]:
assumes $\bigwedge x.$[($\boldsymbol{\mathcal{A}}$(φ $x \equiv \psi$ x)) *in* v]
shows [[($\boldsymbol{\lambda}^0$ (χ ($\boldsymbol{\iota}x.$ φ x)) = $\boldsymbol{\lambda}^0$ (χ ($\boldsymbol{\iota}x.$ ψ x)))]]
unfolding *axiom-def identity*-o-*def* **apply** − **apply** (*rule allI*; *rule* Eq_0I)
using *TheEqI*[*OF assms*[*THEN ActualE*, *THEN EquivE*]] **by** *auto*

lemma *lambda-predicates-4-1*[*axiom*]:
assumes $\bigwedge x.$[($\boldsymbol{\mathcal{A}}$(φ $x \equiv \psi$ x)) *in* v]
shows [[(($\boldsymbol{\lambda}$ x . χ ($\boldsymbol{\iota}x.$ φ x) x) = ($\boldsymbol{\lambda}$ x . χ ($\boldsymbol{\iota}x.$ ψ x) x))]]
unfolding *axiom-def identity*-Π_1-*def* **apply** − **apply** (*rule allI*; *rule* Eq_1I)
using *TheEqI*[*OF assms*[*THEN ActualE*, *THEN EquivE*]] **by** *auto*

lemma *lambda-predicates-4-2*[*axiom*]:
assumes $\bigwedge x.$[($\boldsymbol{\mathcal{A}}$(φ $x \equiv \psi$ x)) *in* v]
shows [[($\boldsymbol{\lambda}^2$ (λ x y . χ ($\boldsymbol{\iota}x.$ φ x) x y)) = ($\boldsymbol{\lambda}^2$ (λ x y . χ ($\boldsymbol{\iota}x.$ ψ x) x y)))]]
unfolding *axiom-def identity*-Π_2-*def* **apply** − **apply** (*rule allI*; *rule* Eq_2I)
using *TheEqI*[*OF assms*[*THEN ActualE*, *THEN EquivE*]] **by** *auto*

lemma *lambda-predicates-4-3*[*axiom*]:
assumes $\bigwedge x.$[($\boldsymbol{\mathcal{A}}$(φ $x \equiv \psi$ x)) *in* v]
shows [[($\boldsymbol{\lambda}^3$ (λ x y z . χ ($\boldsymbol{\iota}x.$ φ x) x y z)) = ($\boldsymbol{\lambda}^3$ (λ x y z . χ ($\boldsymbol{\iota}x.$ ψ x) x y z))]]
unfolding *axiom-def identity*-Π_3-*def* **apply** − **apply** (*rule allI*; *rule* Eq_3I)
using *TheEqI*[*OF assms*[*THEN ActualE*, *THEN EquivE*]] **by** *auto*

A.7.10. Axioms of Encoding

lemma *encoding*[*axiom*]:
[[{|x,F|} → □{|x,F|}]]

by *axiom-meta-solver*
lemma *nocoder*[*axiom*]:
[[⦇*O*!,*x*⦈ → ¬(∃ *F* . ⦃*x*,*F*⦄)]]
unfolding *axiom-def*
apply (*rule allI*, *rule ImplI*, *subst* (*asm*) *OrdS*)
apply *meta-solver* **unfolding** *en-def*
by (*metis* ν.*simps*(5) *mem-Collect-eq option.sel*)
lemma *A-objects*[*axiom*]:
[[∃ *x*. (⦇*A*!,x^P⦈ & (∀ *F* . (⦃x^P,*F*⦄ ≡ φ *F*))]]
unfolding *axiom-def*
proof (*rule allI*, *rule ExIRule*)
fix *v*
let *?x* = αν { *F* . [φ *F in v*]}
have [⦇*A*!,*?x*^{P}⦈ *in v*] **by** (*simp add*: *AbsS* d_κ-*proper*)
moreover have [(∀ *F*. ⦃*?x*^{P},*F*⦄ ≡ φ *F*) *in v*]
apply *meta-solver* **unfolding** *en-def*
using d_1.*rep-eq* d_κ-*def* d_κ-*proper eval*Π_1-*inverse* **by** *auto*
ultimately show [⦇*A*!,*?x*^{P}⦈ & (∀ *F*. ⦃*?x*^{P},*F*⦄ ≡ φ *F*) *in v*]
by (*simp only*: *ConjS*)
qed

end

A.8. Definitions

A.8.1. Property Negations

consts *propnot* :: *'a*⇒*'a* (-$^-$ [90] 90)
overloading $propnot_0 \equiv propnot :: \Pi_0 \Rightarrow \Pi_0$
$propnot_1 \equiv propnot :: \Pi_1 \Rightarrow \Pi_1$
$propnot_2 \equiv propnot :: \Pi_2 \Rightarrow \Pi_2$
$propnot_3 \equiv propnot :: \Pi_3 \Rightarrow \Pi_3$
begin
definition $propnot_0 :: \Pi_0 \Rightarrow \Pi_0$ **where**
$propnot_0 \equiv \lambda\ p\ .\ \boldsymbol{\lambda}^0\ (\neg p)$
definition $propnot_1$ **where**
$propnot_1 \equiv \lambda\ F\ .\ \boldsymbol{\lambda}\ x\ .\ \neg(\!|F,\ x^P|\!)$
definition $propnot_2$ **where**
$propnot_2 \equiv \lambda\ F\ .\ \boldsymbol{\lambda}^2\ (\lambda\ x\ y\ .\ \neg(\!|F,\ x^P,\ y^P|\!))$
definition $propnot_3$ **where**
$propnot_3 \equiv \lambda\ F\ .\ \boldsymbol{\lambda}^3\ (\lambda\ x\ y\ z\ .\ \neg(\!|F,\ x^P,\ y^P,\ z^P|\!))$
end

named-theorems *propnot-defs*
declare $propnot_0$-*def*[*propnot-defs*] $propnot_1$-*def*[*propnot-defs*]
$propnot_2$-*def*[*propnot-defs*] $propnot_3$-*def*[*propnot-defs*]

A.8.2. Noncontingent and Contingent Relations

consts *Necessary* :: *'a*⇒o
overloading $Necessary_0 \equiv Necessary :: \Pi_0 \Rightarrow o$
$Necessary_1 \equiv Necessary :: \Pi_1 \Rightarrow o$
$Necessary_2 \equiv Necessary :: \Pi_2 \Rightarrow o$
$Necessary_3 \equiv Necessary :: \Pi_3 \Rightarrow o$
begin
definition $Necessary_0$ **where**
$Necessary_0 \equiv \lambda\ p\ .\ \Box p$

definition *Necessary*$_1$ :: $\Pi_1 \Rightarrow o$ **where**
Necessary$_1 \equiv \lambda\ F\ .\ \Box(\forall\ x\ .\ (\!|F,x^P|\!))$
definition *Necessary*$_2$ **where**
Necessary$_2 \equiv \lambda\ F\ .\ \Box(\forall\ x\ y\ .\ (\!|F,x^P,y^P|\!))$
definition *Necessary*$_3$ **where**
Necessary$_3 \equiv \lambda\ F\ .\ \Box(\forall\ x\ y\ z\ .\ (\!|F,x^P,y^P,z^P|\!))$
end

named-theorems *Necessary-defs*
declare *Necessary*$_0$*-def*[*Necessary-defs*] *Necessary*$_1$*-def*[*Necessary-defs*]
Necessary$_2$*-def*[*Necessary-defs*] *Necessary*$_3$*-def*[*Necessary-defs*]

consts *Impossible* :: $'a \Rightarrow o$
overloading *Impossible*$_0 \equiv$ *Impossible* :: $\Pi_0 \Rightarrow o$
Impossible$_1 \equiv$ *Impossible* :: $\Pi_1 \Rightarrow o$
Impossible$_2 \equiv$ *Impossible* :: $\Pi_2 \Rightarrow o$
Impossible$_3 \equiv$ *Impossible* :: $\Pi_3 \Rightarrow o$
begin
definition *Impossible*$_0$ **where**
Impossible$_0 \equiv \lambda\ p\ .\ \Box\neg p$
definition *Impossible*$_1$ **where**
Impossible$_1 \equiv \lambda\ F\ .\ \Box(\forall\ x.\ \neg(\!|F,x^P|\!))$
definition *Impossible*$_2$ **where**
Impossible$_2 \equiv \lambda\ F\ .\ \Box(\forall\ x\ y.\ \neg(\!|F,x^P,y^P|\!))$
definition *Impossible*$_3$ **where**
Impossible$_3 \equiv \lambda\ F\ .\ \Box(\forall\ x\ y\ z.\ \neg(\!|F,x^P,y^P,z^P|\!))$
end

named-theorems *Impossible-defs*
declare *Impossible*$_0$*-def*[*Impossible-defs*] *Impossible*$_1$*-def*[*Impossible-defs*]
Impossible$_2$*-def*[*Impossible-defs*] *Impossible*$_3$*-def*[*Impossible-defs*]

definition *NonContingent* **where**
NonContingent $\equiv \lambda\ F\ .\ (Necessary\ F) \vee (Impossible\ F)$
definition *Contingent* **where**
Contingent $\equiv \lambda\ F\ .\ \neg(Necessary\ F \vee Impossible\ F)$

definition *ContingentlyTrue* :: $o \Rightarrow o$ **where**
ContingentlyTrue $\equiv \lambda\ p\ .\ p\ \&\ \Diamond\neg p$
definition *ContingentlyFalse* :: $o \Rightarrow o$ **where**
ContingentlyFalse $\equiv \lambda\ p\ .\ \neg p\ \&\ \Diamond p$

definition *WeaklyContingent* **where**
WeaklyContingent $\equiv \lambda\ F\ .\ Contingent\ F\ \&\ (\forall\ x.\ \Diamond(\!|F,x^P|\!) \rightarrow \Box(\!|F,x^P|\!))$

A.8.3. Null and Universal Objects

definition *Null* :: $\kappa \Rightarrow o$ **where**
Null $\equiv \lambda\ x\ .\ (\!|A!,x|\!)\ \&\ \neg(\exists\ F\ .\ \{\!|x,\ F|\!\})$
definition *Universal* :: $\kappa \Rightarrow o$ **where**
Universal $\equiv \lambda\ x\ .\ (\!|A!,x|\!)\ \&\ (\forall\ F\ .\ \{\!|x,\ F|\!\})$

definition *NullObject* :: $\kappa\ (\mathbf{a}_\emptyset)$ **where**
NullObject $\equiv (\boldsymbol{\iota}x\ .\ Null\ (x^P))$
definition *UniversalObject* :: $\kappa\ (\mathbf{a}_V)$ **where**
UniversalObject $\equiv (\boldsymbol{\iota}x\ .\ Universal\ (x^P))$

A.8.4. Propositional Properties

definition *Propositional* **where**
Propositional $F \equiv \exists\ \ p\ .\ F = (\boldsymbol{\lambda}\ x\ .\ p)$

A.8.5. Indiscriminate Properties

definition *Indiscriminate* :: $\Pi_1 \Rightarrow o$ **where**
Indiscriminate $\equiv \lambda\ F\ .\ \Box((\exists\ \ x\ .\ (\!|F,x^P|\!)) \rightarrow (\forall\ \ x\ .\ (\!|F,x^P|\!)))$

A.8.6. Miscellaneous

definition *not-identical*$_E$ :: $\kappa \Rightarrow \kappa \Rightarrow o$ (**infixl** $\neq_E$ *63*)
where *not-identical*$_E \equiv \lambda\ x\ y\ .\ (\!|(\boldsymbol{\lambda}^2\ (\lambda\ x\ y\ .\ x^P =_E y^P))^-,\ x,\ y|\!)$

A.9. The Deductive System PLM

declare *meta-defs*[*no-atp*] *meta-aux*[*no-atp*]

locale *PLM* = *Axioms*
begin

A.9.1. Automatic Solver

named-theorems *PLM*
named-theorems *PLM-intro*
named-theorems *PLM-elim*
named-theorems *PLM-dest*
named-theorems *PLM-subst*

method *PLM-solver* **declares** *PLM-intro PLM-elim PLM-subst PLM-dest PLM*
= ((*assumption* | (*match axiom* **in** *A*: [[φ]] **for** $\varphi \Rightarrow$ ‹*fact A*[*axiom-instance*]›)
| *fact PLM* | *rule PLM-intro* | *subst PLM-subst* | *subst* (*asm*) *PLM-subst*
| *fastforce* | *safe* | *drule PLM-dest* | *erule PLM-elim*); (*PLM-solver*)*?*)

A.9.2. Modus Ponens

lemma *modus-ponens*[*PLM*]:
$[\![[\varphi\ in\ v];\ [\varphi \rightarrow \psi\ in\ v]]\!] \Longrightarrow [\psi\ in\ v]$
by (*simp add*: *Semantics.T5*)

A.9.3. Axioms

interpretation *Axioms* **.**
declare *axiom*[*PLM*]
declare *conn-defs*[*PLM*]

A.9.4. (Modally Strict) Proofs and Derivations

lemma *vdash-properties-6*[*no-atp*]:
$[\![[\varphi\ in\ v];\ [\varphi \rightarrow \psi\ in\ v]]\!] \Longrightarrow [\psi\ in\ v]$
using *modus-ponens* **.**
lemma *vdash-properties-9*[*PLM*]:
$[\varphi\ in\ v] \Longrightarrow [\psi \rightarrow \varphi\ in\ v]$
using *modus-ponens pl-1*[*axiom-instance*] **by** *blast*

lemma *vdash-properties-10*[*PLM*]:
  [$\varphi \rightarrow \psi$ *in* v] $\Longrightarrow$ ([φ *in* v] $\Longrightarrow$ [ψ *in* v])
  using *vdash-properties-6* **.**

attribute-setup *deduction* = ⟪
  Scan.succeed (*Thm.rule-attribute* []
    (*fn - => fn thm => thm RS* @{*thm vdash-properties-10*}))
⟫

A.9.5. GEN and RN

lemma *rule-gen*[*PLM*]:
  ⟦$\bigwedge\alpha$. [φ α *in* v]⟧ $\Longrightarrow$ [$\forall\,\alpha$. φ α *in* v]
  by (*simp add*: *Semantics.T8*)

lemma *RN-2*[*PLM*]:
  ($\bigwedge$ v . [ψ *in* v] $\Longrightarrow$ [φ *in* v]) $\Longrightarrow$ ([$\Box\psi$ *in* v] $\Longrightarrow$ [$\Box\varphi$ *in* v])
  by (*simp add*: *Semantics.T6*)

lemma *RN*[*PLM*]:
  ($\bigwedge$ v . [φ *in* v]) $\Longrightarrow$ [$\Box\varphi$ *in* v]
  using *qml-3*[*axiom-necessitation*, *axiom-instance*] *RN-2* **by** *blast*

A.9.6. Negations and Conditionals

lemma *if-p-then-p*[*PLM*]:
  [$\varphi \rightarrow \varphi$ *in* v]
  using *pl-1 pl-2 vdash-properties-10 axiom-instance* **by** *blast*

lemma *deduction-theorem*[*PLM*,*PLM-intro*]:
  ⟦[φ *in* v] $\Longrightarrow$ [ψ *in* v]⟧ $\Longrightarrow$ [$\varphi \rightarrow \psi$ *in* v]
  by (*simp add*: *Semantics.T5*)
lemmas *CP* = *deduction-theorem*

lemma *ded-thm-cor-3*[*PLM*]:
  ⟦[$\varphi \rightarrow \psi$ *in* v]; [$\psi \rightarrow \chi$ *in* v]⟧ $\Longrightarrow$ [$\varphi \rightarrow \chi$ *in* v]
  by (*meson pl-2 vdash-properties-10 vdash-properties-9 axiom-instance*)
lemma *ded-thm-cor-4*[*PLM*]:
  ⟦[$\varphi \rightarrow (\psi \rightarrow \chi)$ *in* v]; [ψ *in* v]⟧ $\Longrightarrow$ [$\varphi \rightarrow \chi$ *in* v]
  by (*meson pl-2 vdash-properties-10 vdash-properties-9 axiom-instance*)

lemma *useful-tautologies-1*[*PLM*]:
  [$\neg\neg\varphi \rightarrow \varphi$ *in* v]
  by (*meson pl-1 pl-3 ded-thm-cor-3 ded-thm-cor-4 axiom-instance*)
lemma *useful-tautologies-2*[*PLM*]:
  [$\varphi \rightarrow \neg\neg\varphi$ *in* v]
  by (*meson pl-1 pl-3 ded-thm-cor-3 useful-tautologies-1*
        vdash-properties-10 axiom-instance)
lemma *useful-tautologies-3*[*PLM*]:
  [$\neg\varphi \rightarrow (\varphi \rightarrow \psi)$ *in* v]
  by (*meson pl-1 pl-2 pl-3 ded-thm-cor-3 ded-thm-cor-4 axiom-instance*)
lemma *useful-tautologies-4*[*PLM*]:
  [$(\neg\psi \rightarrow \neg\varphi) \rightarrow (\varphi \rightarrow \psi)$ *in* v]
  by (*meson pl-1 pl-2 pl-3 ded-thm-cor-3 ded-thm-cor-4 axiom-instance*)
lemma *useful-tautologies-5*[*PLM*]:
  [$(\varphi \rightarrow \psi) \rightarrow (\neg\psi \rightarrow \neg\varphi)$ *in* v]
  by (*metis CP useful-tautologies-4 vdash-properties-10*)
lemma *useful-tautologies-6*[*PLM*]:

$[(\varphi \rightarrow \neg\psi) \rightarrow (\psi \rightarrow \neg\varphi)\ in\ v]$
by (*metis CP useful-tautologies-4 vdash-properties-10*)
lemma *useful-tautologies-7*[*PLM*]:
$[(\neg\varphi \rightarrow \psi) \rightarrow (\neg\psi \rightarrow \varphi)\ in\ v]$
using *ded-thm-cor-3 useful-tautologies-4 useful-tautologies-5*
useful-tautologies-6 **by** *blast*
lemma *useful-tautologies-8*[*PLM*]:
$[\varphi \rightarrow (\neg\psi \rightarrow \neg(\varphi \rightarrow \psi))\ in\ v]$
by (*meson ded-thm-cor-3 CP useful-tautologies-5*)
lemma *useful-tautologies-9*[*PLM*]:
$[(\varphi \rightarrow \psi) \rightarrow ((\neg\varphi \rightarrow \psi) \rightarrow \psi)\ in\ v]$
by (*metis CP useful-tautologies-4 vdash-properties-10*)
lemma *useful-tautologies-10*[*PLM*]:
$[(\varphi \rightarrow \neg\psi) \rightarrow ((\varphi \rightarrow \psi) \rightarrow \neg\varphi)\ in\ v]$
by (*metis ded-thm-cor-3 CP useful-tautologies-6*)

lemma *modus-tollens-1*[*PLM*]:
$[\![[\varphi \rightarrow \psi\ in\ v]; [\neg\psi\ in\ v]]\!] \Longrightarrow [\neg\varphi\ in\ v]$
by (*metis ded-thm-cor-3 ded-thm-cor-4 useful-tautologies-3*
useful-tautologies-7 vdash-properties-10)
lemma *modus-tollens-2*[*PLM*]:
$[\![[\varphi \rightarrow \neg\psi\ in\ v]; [\psi\ in\ v]]\!] \Longrightarrow [\neg\varphi\ in\ v]$
using *modus-tollens-1 useful-tautologies-2*
vdash-properties-10 **by** *blast*

lemma *contraposition-1*[*PLM*]:
$[\varphi \rightarrow \psi\ in\ v] = [\neg\psi \rightarrow \neg\varphi\ in\ v]$
using *useful-tautologies-4 useful-tautologies-5*
vdash-properties-10 **by** *blast*
lemma *contraposition-2*[*PLM*]:
$[\varphi \rightarrow \neg\psi\ in\ v] = [\psi \rightarrow \neg\varphi\ in\ v]$
using *contraposition-1 ded-thm-cor-3*
useful-tautologies-1 **by** *blast*

lemma *reductio-aa-1*[*PLM*]:
$[\![[\neg\varphi\ in\ v] \Longrightarrow [\neg\psi\ in\ v]; [\neg\varphi\ in\ v] \Longrightarrow [\psi\ in\ v]]\!] \Longrightarrow [\varphi\ in\ v]$
using *CP modus-tollens-2 useful-tautologies-1*
vdash-properties-10 **by** *blast*
lemma *reductio-aa-2*[*PLM*]:
$[\![[\varphi\ in\ v] \Longrightarrow [\neg\psi\ in\ v]; [\varphi\ in\ v] \Longrightarrow [\psi\ in\ v]]\!] \Longrightarrow [\neg\varphi\ in\ v]$
by (*meson contraposition-1 reductio-aa-1*)
lemma *reductio-aa-3*[*PLM*]:
$[\![[\neg\varphi \rightarrow \neg\psi\ in\ v]; [\neg\varphi \rightarrow \psi\ in\ v]]\!] \Longrightarrow [\varphi\ in\ v]$
using *reductio-aa-1 vdash-properties-10* **by** *blast*
lemma *reductio-aa-4*[*PLM*]:
$[\![[\varphi \rightarrow \neg\psi\ in\ v]; [\varphi \rightarrow \psi\ in\ v]]\!] \Longrightarrow [\neg\varphi\ in\ v]$
using *reductio-aa-2 vdash-properties-10* **by** *blast*

lemma *raa-cor-1*[*PLM*]:
$[\![[\varphi\ in\ v]; [\neg\psi\ in\ v] \Longrightarrow [\neg\varphi\ in\ v]]\!] \Longrightarrow ([\varphi\ in\ v] \Longrightarrow [\psi\ in\ v])$
using *reductio-aa-1 vdash-properties-9* **by** *blast*
lemma *raa-cor-2*[*PLM*]:
$[\![[\neg\varphi\ in\ v]; [\neg\psi\ in\ v] \Longrightarrow [\varphi\ in\ v]]\!] \Longrightarrow ([\neg\varphi\ in\ v] \Longrightarrow [\psi\ in\ v])$
using *reductio-aa-1 vdash-properties-9* **by** *blast*
lemma *raa-cor-3*[*PLM*]:
$[\![[\varphi\ in\ v]; [\neg\psi \rightarrow \neg\varphi\ in\ v]]\!] \Longrightarrow ([\varphi\ in\ v] \Longrightarrow [\psi\ in\ v])$
using *raa-cor-1 vdash-properties-10* **by** *blast*
lemma *raa-cor-4*[*PLM*]:

$[\![[\neg\varphi \ in \ v]; [\neg\psi \rightarrow \varphi \ in \ v]]\!] \Longrightarrow ([\neg\varphi \ in \ v] \Longrightarrow [\psi \ in \ v])$
using *raa-cor-2 vdash-properties-10* **by** *blast*

Remark. *In contrast to PLM the classical introduction and elimination rules are proven before the tautologies. The statements proven so far are sufficient for the proofs and using the derived rules the tautologies can be derived automatically.*

lemma *intro-elim-1*[*PLM*]:
$[\![[\varphi \ in \ v]; [\psi \ in \ v]]\!] \Longrightarrow [\varphi \ \& \ \psi \ in \ v]$
unfolding *conj-def* **using** *ded-thm-cor-4 if-p-then-p modus-tollens-2* **by** *blast*
lemmas &I = *intro-elim-1*
lemma *intro-elim-2-a*[*PLM*]:
$[\varphi \ \& \ \psi \ in \ v] \Longrightarrow [\varphi \ in \ v]$
unfolding *conj-def* **using** *CP reductio-aa-1* **by** *blast*
lemma *intro-elim-2-b*[*PLM*]:
$[\varphi \ \& \ \psi \ in \ v] \Longrightarrow [\psi \ in \ v]$
unfolding *conj-def* **using** *pl-1 CP reductio-aa-1 axiom-instance* **by** *blast*
lemmas &E = *intro-elim-2-a intro-elim-2-b*
lemma *intro-elim-3-a*[*PLM*]:
$[\varphi \ in \ v] \Longrightarrow [\varphi \vee \psi \ in \ v]$
unfolding *disj-def* **using** *ded-thm-cor-4 useful-tautologies-3* **by** *blast*
lemma *intro-elim-3-b*[*PLM*]:
$[\psi \ in \ v] \Longrightarrow [\varphi \vee \psi \ in \ v]$
by (*simp only*: *disj-def vdash-properties-9*)
lemmas $\vee I$ = *intro-elim-3-a intro-elim-3-b*
lemma *intro-elim-4-a*[*PLM*]:
$[\![[\varphi \vee \psi \ in \ v]; [\varphi \rightarrow \chi \ in \ v]; [\psi \rightarrow \chi \ in \ v]]\!] \Longrightarrow [\chi \ in \ v]$
unfolding *disj-def* **by** (*meson reductio-aa-2 vdash-properties-10*)
lemma *intro-elim-4-b*[*PLM*]:
$[\![[\varphi \vee \psi \ in \ v]; [\neg\varphi \ in \ v]]\!] \Longrightarrow [\psi \ in \ v]$
unfolding *disj-def* **using** *vdash-properties-10* **by** *blast*
lemma *intro-elim-4-c*[*PLM*]:
$[\![[\varphi \vee \psi \ in \ v]; [\neg\psi \ in \ v]]\!] \Longrightarrow [\varphi \ in \ v]$
unfolding *disj-def* **using** *raa-cor-2 vdash-properties-10* **by** *blast*
lemma *intro-elim-4-d*[*PLM*]:
$[\![[\varphi \vee \psi \ in \ v]; [\varphi \rightarrow \chi \ in \ v]; [\psi \rightarrow \Theta \ in \ v]]\!] \Longrightarrow [\chi \vee \Theta \ in \ v]$
unfolding *disj-def* **using** *contraposition-1 ded-thm-cor-3* **by** *blast*
lemma *intro-elim-4-e*[*PLM*]:
$[\![[\varphi \vee \psi \ in \ v]; [\varphi \equiv \chi \ in \ v]; [\psi \equiv \Theta \ in \ v]]\!] \Longrightarrow [\chi \vee \Theta \ in \ v]$
unfolding *equiv-def* **using** &$E(1)$ *intro-elim-4-d* **by** *blast*
lemmas $\vee E$ = *intro-elim-4-a intro-elim-4-b intro-elim-4-c intro-elim-4-d*
lemma *intro-elim-5*[*PLM*]:
$[\![[\varphi \rightarrow \psi \ in \ v]; [\psi \rightarrow \varphi \ in \ v]]\!] \Longrightarrow [\varphi \equiv \psi \ in \ v]$
by (*simp only*: *equiv-def* &I)
lemmas $\equiv I$ = *intro-elim-5*
lemma *intro-elim-6-a*[*PLM*]:
$[\![[\varphi \equiv \psi \ in \ v]; [\varphi \ in \ v]]\!] \Longrightarrow [\psi \ in \ v]$
unfolding *equiv-def* **using** &$E(1)$ *vdash-properties-10* **by** *blast*
lemma *intro-elim-6-b*[*PLM*]:
$[\![[\varphi \equiv \psi \ in \ v]; [\psi \ in \ v]]\!] \Longrightarrow [\varphi \ in \ v]$
unfolding *equiv-def* **using** &$E(2)$ *vdash-properties-10* **by** *blast*
lemma *intro-elim-6-c*[*PLM*]:
$[\![[\varphi \equiv \psi \ in \ v]; [\neg\varphi \ in \ v]]\!] \Longrightarrow [\neg\psi \ in \ v]$
unfolding *equiv-def* **using** &$E(2)$ *modus-tollens-1* **by** *blast*
lemma *intro-elim-6-d*[*PLM*]:
$[\![[\varphi \equiv \psi \ in \ v]; [\neg\psi \ in \ v]]\!] \Longrightarrow [\neg\varphi \ in \ v]$
unfolding *equiv-def* **using** &$E(1)$ *modus-tollens-1* **by** *blast*
lemma *intro-elim-6-e*[*PLM*]:

$[\![[\varphi \equiv \psi \ in \ v]; [\psi \equiv \chi \ in \ v]]\!] \Longrightarrow [\varphi \equiv \chi \ in \ v]$
by (*metis equiv-def ded-thm-cor-3* &E ≡I)
lemma *intro-elim-6-f*[*PLM*]:
$[\![[\varphi \equiv \psi \ in \ v]; [\varphi \equiv \chi \ in \ v]]\!] \Longrightarrow [\chi \equiv \psi \ in \ v]$
by (*metis equiv-def ded-thm-cor-3* &E ≡I)
lemmas ≡E = *intro-elim-6-a intro-elim-6-b intro-elim-6-c*
intro-elim-6-d intro-elim-6-e intro-elim-6-f
lemma *intro-elim-7*[*PLM*]:
$[\varphi \ in \ v] \Longrightarrow [\neg\neg\varphi \ in \ v]$
using *if-p-then-p modus-tollens-2* **by** *blast*
lemmas ¬¬I = *intro-elim-7*
lemma *intro-elim-8*[*PLM*]:
$[\neg\neg\varphi \ in \ v] \Longrightarrow [\varphi \ in \ v]$
using *if-p-then-p raa-cor-2* **by** *blast*
lemmas ¬¬E = *intro-elim-8*

context
begin
private lemma *NotNotI*[*PLM-intro*]:
$[\varphi \ in \ v] \Longrightarrow [\neg(\neg\varphi) \ in \ v]$
by (*simp add*: ¬¬I)
private lemma *NotNotD*[*PLM-dest*]:
$[\neg(\neg\varphi) \ in \ v] \Longrightarrow [\varphi \ in \ v]$
using ¬¬E **by** *blast*

private lemma *ImplI*[*PLM-intro*]:
$([\varphi \ in \ v] \Longrightarrow [\psi \ in \ v]) \Longrightarrow [\varphi \rightarrow \psi \ in \ v]$
using *CP* **.**
private lemma *ImplE*[*PLM-elim*, *PLM-dest*]:
$[\varphi \rightarrow \psi \ in \ v] \Longrightarrow ([\varphi \ in \ v] \Longrightarrow [\psi \ in \ v])$
using *modus-ponens* **.**
private lemma *ImplS*[*PLM-subst*]:
$[\varphi \rightarrow \psi \ in \ v] = ([\varphi \ in \ v] \longrightarrow [\psi \ in \ v])$
using *ImplI ImplE* **by** *blast*

private lemma *NotI*[*PLM-intro*]:
$([\varphi \ in \ v] \Longrightarrow (\bigwedge\psi \ .[\psi \ in \ v])) \Longrightarrow [\neg\varphi \ in \ v]$
using *CP modus-tollens-2* **by** *blast*
private lemma *NotE*[*PLM-elim*,*PLM-dest*]:
$[\neg\varphi \ in \ v] \Longrightarrow ([\varphi \ in \ v] \longrightarrow (\forall\,\psi \ .[\psi \ in \ v]))$
using ∨$I(2)$ ∨$E(3)$ **by** *blast*
private lemma *NotS*[*PLM-subst*]:
$[\neg\varphi \ in \ v] = ([\varphi \ in \ v] \longrightarrow (\forall\,\psi \ .[\psi \ in \ v]))$
using *NotI NotE* **by** *blast*

private lemma *ConjI*[*PLM-intro*]:
$[\![[\varphi \ in \ v]; [\psi \ in \ v]]\!] \Longrightarrow [\varphi \ \& \ \psi \ in \ v]$
using &I **by** *blast*
private lemma *ConjE*[*PLM-elim*,*PLM-dest*]:
$[\varphi \ \& \ \psi \ in \ v] \Longrightarrow (([\varphi \ in \ v] \wedge [\psi \ in \ v]))$
using *CP* &E **by** *blast*
private lemma *ConjS*[*PLM-subst*]:
$[\varphi \ \& \ \psi \ in \ v] = (([\varphi \ in \ v] \wedge [\psi \ in \ v]))$
using *ConjI ConjE* **by** *blast*

private lemma *DisjI*[*PLM-intro*]:
$[\varphi \ in \ v] \vee [\psi \ in \ v] \Longrightarrow [\varphi \vee \psi \ in \ v]$
using ∨I **by** *blast*

private lemma *DisjE*[*PLM-elim*,*PLM-dest*]:
$[\varphi \lor \psi\ in\ v] \Longrightarrow [\varphi\ in\ v] \lor [\psi\ in\ v]$
using *CP* $\lor E(1)$ **by** *blast*
private lemma *DisjS*[*PLM-subst*]:
$[\varphi \lor \psi\ in\ v] = ([\varphi\ in\ v] \lor [\psi\ in\ v])$
using *DisjI DisjE* **by** *blast*

private lemma *EquivI*[*PLM-intro*]:
$[\![[\varphi\ in\ v] \Longrightarrow [\psi\ in\ v]; [\psi\ in\ v] \Longrightarrow [\varphi\ in\ v]]\!] \Longrightarrow [\varphi \equiv \psi\ in\ v]$
using *CP* $\equiv I$ **by** *blast*
private lemma *EquivE*[*PLM-elim*,*PLM-dest*]:
$[\varphi \equiv \psi\ in\ v] \Longrightarrow (([\varphi\ in\ v] \longrightarrow [\psi\ in\ v]) \land ([\psi\ in\ v] \longrightarrow [\varphi\ in\ v]))$
using $\equiv E(1)$ $\equiv E(2)$ **by** *blast*
private lemma *EquivS*[*PLM-subst*]:
$[\varphi \equiv \psi\ in\ v] = ([\varphi\ in\ v] \longleftrightarrow [\psi\ in\ v])$
using *EquivI EquivE* **by** *blast*

private lemma *NotOrD*[*PLM-dest*]:
$\neg[\varphi \lor \psi\ in\ v] \Longrightarrow \neg[\varphi\ in\ v] \land \neg[\psi\ in\ v]$
using $\lor I$ **by** *blast*
private lemma *NotAndD*[*PLM-dest*]:
$\neg[\varphi\ \&\ \psi\ in\ v] \Longrightarrow \neg[\varphi\ in\ v] \lor \neg[\psi\ in\ v]$
using $\&I$ **by** *blast*
private lemma *NotEquivD*[*PLM-dest*]:
$\neg[\varphi \equiv \psi\ in\ v] \Longrightarrow [\varphi\ in\ v] \neq [\psi\ in\ v]$
by (*meson NotI contraposition-1* $\equiv I$ *vdash-properties-9*)

private lemma *BoxI*[*PLM-intro*]:
$(\bigwedge v\ .\ [\varphi\ in\ v]) \Longrightarrow [\Box\varphi\ in\ v]$
using *RN* **by** *blast*
private lemma *NotBoxD*[*PLM-dest*]:
$\neg[\Box\varphi\ in\ v] \Longrightarrow (\exists\ v\ .\ \neg[\varphi\ in\ v])$
using *BoxI* **by** *blast*

private lemma *AllI*[*PLM-intro*]:
$(\bigwedge x\ .\ [\varphi\ x\ in\ v]) \Longrightarrow [\forall\ x\ .\ \varphi\ x\ in\ v]$
using *rule-gen* **by** *blast*
lemma *NotAllD*[*PLM-dest*]:
$\neg[\forall\ x\ .\ \varphi\ x\ in\ v] \Longrightarrow (\exists\ x\ .\ \neg[\varphi\ x\ in\ v])$
using *AllI* **by** *fastforce*
end

lemma *oth-class-taut-1-a*[*PLM*]:
$[\neg(\varphi\ \&\ \neg\varphi)\ in\ v]$
by *PLM-solver*
lemma *oth-class-taut-1-b*[*PLM*]:
$[\neg(\varphi \equiv \neg\varphi)\ in\ v]$
by *PLM-solver*
lemma *oth-class-taut-2*[*PLM*]:
$[\varphi \lor \neg\varphi\ in\ v]$
by *PLM-solver*
lemma *oth-class-taut-3-a*[*PLM*]:
$[(\varphi\ \&\ \varphi) \equiv \varphi\ in\ v]$
by *PLM-solver*
lemma *oth-class-taut-3-b*[*PLM*]:
$[(\varphi\ \&\ \psi) \equiv (\psi\ \&\ \varphi)\ in\ v]$
by *PLM-solver*
lemma *oth-class-taut-3-c*[*PLM*]:

$[(\varphi \;\&\; (\psi \;\&\; \chi)) \equiv ((\varphi \;\&\; \psi) \;\&\; \chi)\ in\ v]$
by *PLM-solver*
lemma *oth-class-taut-3-d*[*PLM*]:
$[(\varphi \vee \varphi) \equiv \varphi\ in\ v]$
by *PLM-solver*
lemma *oth-class-taut-3-e*[*PLM*]:
$[(\varphi \vee \psi) \equiv (\psi \vee \varphi)\ in\ v]$
by *PLM-solver*
lemma *oth-class-taut-3-f*[*PLM*]:
$[(\varphi \vee (\psi \vee \chi)) \equiv ((\varphi \vee \psi) \vee \chi)\ in\ v]$
by *PLM-solver*
lemma *oth-class-taut-3-g*[*PLM*]:
$[(\varphi \equiv \psi) \equiv (\psi \equiv \varphi)\ in\ v]$
by *PLM-solver*
lemma *oth-class-taut-3-i*[*PLM*]:
$[(\varphi \equiv (\psi \equiv \chi)) \equiv ((\varphi \equiv \psi) \equiv \chi)\ in\ v]$
by *PLM-solver*
lemma *oth-class-taut-4-a*[*PLM*]:
$[\varphi \equiv \varphi\ in\ v]$
by *PLM-solver*
lemma *oth-class-taut-4-b*[*PLM*]:
$[\varphi \equiv \neg\neg\varphi\ in\ v]$
by *PLM-solver*
lemma *oth-class-taut-5-a*[*PLM*]:
$[(\varphi \rightarrow \psi) \equiv \neg(\varphi \;\&\; \neg\psi)\ in\ v]$
by *PLM-solver*
lemma *oth-class-taut-5-b*[*PLM*]:
$[\neg(\varphi \rightarrow \psi) \equiv (\varphi \;\&\; \neg\psi)\ in\ v]$
by *PLM-solver*
lemma *oth-class-taut-5-c*[*PLM*]:
$[(\varphi \rightarrow \psi) \rightarrow ((\psi \rightarrow \chi) \rightarrow (\varphi \rightarrow \chi))\ in\ v]$
by *PLM-solver*
lemma *oth-class-taut-5-d*[*PLM*]:
$[(\varphi \equiv \psi) \equiv (\neg\varphi \equiv \neg\psi)\ in\ v]$
by *PLM-solver*
lemma *oth-class-taut-5-e*[*PLM*]:
$[(\varphi \equiv \psi) \rightarrow ((\varphi \rightarrow \chi) \equiv (\psi \rightarrow \chi))\ in\ v]$
by *PLM-solver*
lemma *oth-class-taut-5-f*[*PLM*]:
$[(\varphi \equiv \psi) \rightarrow ((\chi \rightarrow \varphi) \equiv (\chi \rightarrow \psi))\ in\ v]$
by *PLM-solver*
lemma *oth-class-taut-5-g*[*PLM*]:
$[(\varphi \equiv \psi) \rightarrow ((\varphi \equiv \chi) \equiv (\psi \equiv \chi))\ in\ v]$
by *PLM-solver*
lemma *oth-class-taut-5-h*[*PLM*]:
$[(\varphi \equiv \psi) \rightarrow ((\chi \equiv \varphi) \equiv (\chi \equiv \psi))\ in\ v]$
by *PLM-solver*
lemma *oth-class-taut-5-i*[*PLM*]:
$[(\varphi \equiv \psi) \equiv ((\varphi \;\&\; \psi) \vee (\neg\varphi \;\&\; \neg\psi))\ in\ v]$
by *PLM-solver*
lemma *oth-class-taut-5-j*[*PLM*]:
$[(\neg(\varphi \equiv \psi)) \equiv ((\varphi \;\&\; \neg\psi) \vee (\neg\varphi \;\&\; \psi))\ in\ v]$
by *PLM-solver*
lemma *oth-class-taut-5-k*[*PLM*]:
$[(\varphi \rightarrow \psi) \equiv (\neg\varphi \vee \psi)\ in\ v]$
by *PLM-solver*

lemma *oth-class-taut-6-a*[*PLM*]:

$[(\varphi \;\&\; \psi) \equiv \neg(\neg\varphi \vee \neg\psi)\ in\ v]$
by *PLM-solver*
lemma *oth-class-taut-6-b*[*PLM*]:
$[(\varphi \vee \psi) \equiv \neg(\neg\varphi \;\&\; \neg\psi)\ in\ v]$
by *PLM-solver*
lemma *oth-class-taut-6-c*[*PLM*]:
$[\neg(\varphi \;\&\; \psi) \equiv (\neg\varphi \vee \neg\psi)\ in\ v]$
by *PLM-solver*
lemma *oth-class-taut-6-d*[*PLM*]:
$[\neg(\varphi \vee \psi) \equiv (\neg\varphi \;\&\; \neg\psi)\ in\ v]$
by *PLM-solver*

lemma *oth-class-taut-7-a*[*PLM*]:
$[(\varphi \;\&\; (\psi \vee \chi)) \equiv ((\varphi \;\&\; \psi) \vee (\varphi \;\&\; \chi))\ in\ v]$
by *PLM-solver*
lemma *oth-class-taut-7-b*[*PLM*]:
$[(\varphi \vee (\psi \;\&\; \chi)) \equiv ((\varphi \vee \psi) \;\&\; (\varphi \vee \chi))\ in\ v]$
by *PLM-solver*

lemma *oth-class-taut-8-a*[*PLM*]:
$[((\varphi \;\&\; \psi) \to \chi) \to (\varphi \to (\psi \to \chi))\ in\ v]$
by *PLM-solver*
lemma *oth-class-taut-8-b*[*PLM*]:
$[(\varphi \to (\psi \to \chi)) \to ((\varphi \;\&\; \psi) \to \chi)\ in\ v]$
by *PLM-solver*

lemma *oth-class-taut-9-a*[*PLM*]:
$[(\varphi \;\&\; \psi) \to \varphi\ in\ v]$
by *PLM-solver*
lemma *oth-class-taut-9-b*[*PLM*]:
$[(\varphi \;\&\; \psi) \to \psi\ in\ v]$
by *PLM-solver*

lemma *oth-class-taut-10-a*[*PLM*]:
$[\varphi \to (\psi \to (\varphi \;\&\; \psi))\ in\ v]$
by *PLM-solver*
lemma *oth-class-taut-10-b*[*PLM*]:
$[(\varphi \to (\psi \to \chi)) \equiv (\psi \to (\varphi \to \chi))\ in\ v]$
by *PLM-solver*
lemma *oth-class-taut-10-c*[*PLM*]:
$[(\varphi \to \psi) \to ((\varphi \to \chi) \to (\varphi \to (\psi \;\&\; \chi)))\ in\ v]$
by *PLM-solver*
lemma *oth-class-taut-10-d*[*PLM*]:
$[(\varphi \to \chi) \to ((\psi \to \chi) \to ((\varphi \vee \psi) \to \chi))\ in\ v]$
by *PLM-solver*
lemma *oth-class-taut-10-e*[*PLM*]:
$[(\varphi \to \psi) \to ((\chi \to \Theta) \to ((\varphi \;\&\; \chi) \to (\psi \;\&\; \Theta)))\ in\ v]$
by *PLM-solver*
lemma *oth-class-taut-10-f*[*PLM*]:
$[((\varphi \;\&\; \psi) \equiv (\varphi \;\&\; \chi)) \equiv (\varphi \to (\psi \equiv \chi))\ in\ v]$
by *PLM-solver*
lemma *oth-class-taut-10-g*[*PLM*]:
$[((\varphi \;\&\; \psi) \equiv (\chi \;\&\; \psi)) \equiv (\psi \to (\varphi \equiv \chi))\ in\ v]$
by *PLM-solver*

attribute-setup *equiv-lr* = ⟪
Scan.succeed (*Thm.rule-attribute* []
(*fn - => fn thm => thm RS* @{*thm* ≡*E*(*1*)}))

⟫

attribute-setup *equiv-rl* = ⟪
Scan.succeed (*Thm.rule-attribute* []
(*fn* - => *fn thm* => *thm RS* @{*thm* ≡*E*(*2*)}))
⟫

attribute-setup *equiv-sym* = ⟪
Scan.succeed (*Thm.rule-attribute* []
(*fn* - => *fn thm* => *thm RS* @{*thm oth-class-taut-3-g*[*equiv-lr*]}))
⟫

attribute-setup *conj1* = ⟪
Scan.succeed (*Thm.rule-attribute* []
(*fn* - => *fn thm* => *thm RS* @{*thm* &*E*(*1*)}))
⟫

attribute-setup *conj2* = ⟪
Scan.succeed (*Thm.rule-attribute* []
(*fn* - => *fn thm* => *thm RS* @{*thm* &*E*(*2*)}))
⟫

attribute-setup *conj-sym* = ⟪
Scan.succeed (*Thm.rule-attribute* []
(*fn* - => *fn thm* => *thm RS* @{*thm oth-class-taut-3-b*[*equiv-lr*]}))
⟫

A.9.7. Identity

lemma *id-eq-prop-prop-1*[*PLM*]:
[(F::Π_1) = F *in* v]
unfolding *identity-defs* **by** *PLM-solver*
lemma *id-eq-prop-prop-2*[*PLM*]:
[((F::Π_1) = G) → (G = F) *in* v]
by (*meson id-eq-prop-prop-1 CP ded-thm-cor-3 l-identity*[*axiom-instance*])
lemma *id-eq-prop-prop-3*[*PLM*]:
[(((F::Π_1) = G) & (G = H)) → (F = H) *in* v]
by (*metis l-identity*[*axiom-instance*] *ded-thm-cor-4 CP* &*E*)
lemma *id-eq-prop-prop-4-a*[*PLM*]:
[(F::Π_2) = F *in* v]
unfolding *identity-defs* **by** *PLM-solver*
lemma *id-eq-prop-prop-4-b*[*PLM*]:
[(F::Π_3) = F *in* v]
unfolding *identity-defs* **by** *PLM-solver*
lemma *id-eq-prop-prop-5-a*[*PLM*]:
[((F::Π_2) = G) → (G = F) *in* v]
by (*meson id-eq-prop-prop-4-a CP ded-thm-cor-3 l-identity*[*axiom-instance*])
lemma *id-eq-prop-prop-5-b*[*PLM*]:
[((F::Π_3) = G) → (G = F) *in* v]
by (*meson id-eq-prop-prop-4-b CP ded-thm-cor-3 l-identity*[*axiom-instance*])
lemma *id-eq-prop-prop-6-a*[*PLM*]:
[(((F::Π_2) = G) & (G = H)) → (F = H) *in* v]
by (*metis l-identity*[*axiom-instance*] *ded-thm-cor-4 CP* &*E*)
lemma *id-eq-prop-prop-6-b*[*PLM*]:
[(((F::Π_3) = G) & (G = H)) → (F = H) *in* v]
by (*metis l-identity*[*axiom-instance*] *ded-thm-cor-4 CP* &*E*)
lemma *id-eq-prop-prop-7*[*PLM*]:
[(p::Π_0) = p *in* v]

unfolding *identity-defs* **by** *PLM-solver*
lemma *id-eq-prop-prop-7-b*[*PLM*]:
[(*p*::o) = *p in v*]
unfolding *identity-defs* **by** *PLM-solver*
lemma *id-eq-prop-prop-8*[*PLM*]:
[((*p*::Π_0) = *q*) → (*q* = *p*) *in v*]
by (*meson id-eq-prop-prop-7 CP ded-thm-cor-3 l-identity*[*axiom-instance*])
lemma *id-eq-prop-prop-8-b*[*PLM*]:
[((*p*::o) = *q*) → (*q* = *p*) *in v*]
by (*meson id-eq-prop-prop-7-b CP ded-thm-cor-3 l-identity*[*axiom-instance*])
lemma *id-eq-prop-prop-9*[*PLM*]:
[(((*p*::Π_0) = *q*) & (*q* = *r*)) → (*p* = *r*) *in v*]
by (*metis l-identity*[*axiom-instance*] *ded-thm-cor-4 CP* &*E*)
lemma *id-eq-prop-prop-9-b*[*PLM*]:
[(((*p*::o) = *q*) & (*q* = *r*)) → (*p* = *r*) *in v*]
by (*metis l-identity*[*axiom-instance*] *ded-thm-cor-4 CP* &*E*)

lemma *eq-E-simple-1*[*PLM*]:
[(*x* =$_E$ *y*) ≡ ((|*O*!,*x*|) & (|*O*!,*y*|) & □(∀ *F* . (|*F*,*x*|) ≡ (|*F*,*y*|))) *in v*]
proof (*rule* ≡*I*; *rule CP*)
assume *1*: [*x* =$_E$ *y in v*]
have [∀ *x y* . ((*x*^{P}) =$_E$ (*y*^{P})) ≡ ((|*O*!,*x*^{P}|) & (|*O*!,*y*^{P}|)
& □(∀ *F* . (|*F*,*x*^{P}|) ≡ (|*F*,*y*^{P}|))) *in v*]
unfolding *identity*$_E$*-infix-def identity*$_E$*-def*
apply (*rule lambda-predicates-2-2*[*axiom-universal*, *axiom-universal*, *axiom-instance*])
by *show-proper*
moreover have [∃ α . (α^P) = *x in v*]
apply (*rule cqt-5-mod*[**where** ψ=λ *x* . *x* =$_E$ *y*, *axiom-instance*,*deduction*])
unfolding *identity*$_E$*-infix-def*
apply (*rule SimpleExOrEnc.intros*)
using *1* **unfolding** *identity*$_E$*-infix-def* **by** *auto*
moreover have [∃ β . (β^P) = *y in v*]
apply (*rule cqt-5-mod*[**where** ψ=λ *y* . *x* =$_E$ *y*,*axiom-instance*,*deduction*])
unfolding *identity*$_E$*-infix-def*
apply (*rule SimpleExOrEnc.intros*) **using** *1*
unfolding *identity*$_E$*-infix-def* **by** *auto*
ultimately have [(*x* =$_E$ *y*) ≡ ((|*O*!,*x*|) & (|*O*!,*y*|)
& □(∀ *F* . (|*F*,*x*|) ≡ (|*F*,*y*|))) *in v*]
using *cqt-1-*κ[*axiom-instance*,*deduction*, *deduction*] **by** *meson*
thus [((|*O*!,*x*|) & (|*O*!,*y*|) & □(∀ *F* . (|*F*,*x*|) ≡ (|*F*,*y*|))) *in v*]
using *1* ≡*E*(*1*) **by** *blast*
next
assume *1*: [(|*O*!,*x*|) & (|*O*!,*y*|) & □(∀ *F*. (|*F*,*x*|) ≡ (|*F*,*y*|)) *in v*]
have [∀ *x y* . ((*x*^{P}) =$_E$ (*y*^{P})) ≡ ((|*O*!,*x*^{P}|) & (|*O*!,*y*^{P}|)
& □(∀ *F* . (|*F*,*x*^{P}|) ≡ (|*F*,*y*^{P}|))) *in v*]
unfolding *identity*$_E$*-def identity*$_E$*-infix-def*
apply (*rule lambda-predicates-2-2*[*axiom-universal*, *axiom-universal*, *axiom-instance*])
by *show-proper*
moreover have [∃ α . (α^P) = *x in v*]
apply (*rule cqt-5-mod*[**where** ψ=λ *x* . (|*O*!,*x*|),*axiom-instance*,*deduction*])
apply (*rule SimpleExOrEnc.intros*)
using *1*[*conj1*,*conj1*] **by** *auto*
moreover have [∃ β . (β^P) = *y in v*]
apply (*rule cqt-5-mod*[**where** ψ=λ *y* . (|*O*!,*y*|),*axiom-instance*,*deduction*])
apply (*rule SimpleExOrEnc.intros*)
using *1*[*conj1*,*conj2*] **by** *auto*
ultimately have [(*x* =$_E$ *y*) ≡ ((|*O*!,*x*|) & (|*O*!,*y*|)
& □(∀ *F* . (|*F*,*x*|) ≡ (|*F*,*y*|))) *in v*]

```
    using cqt-1-κ[axiom-instance,deduction, deduction] by meson
    thus [(x =_E y) in v] using 1 ≡E(2) by blast
  qed
 lemma eq-E-simple-2[PLM]:
  [(x =_E y) → (x = y) in v]
  unfolding identity-defs by PLM-solver
 lemma eq-E-simple-3[PLM]:
  [(x = y) ≡ (((|O!,x|) & (|O!,y|) & □(∀ F . (|F,x|) ≡ (|F,y|)))
          ∨ ((|A!,x|) & (|A!,y|) & □(∀ F. {|x,F|} ≡ {|y,F|}))) in v]
  using eq-E-simple-1
  apply − unfolding identity-defs
  by PLM-solver

 lemma id-eq-obj-1[PLM]: [(x^P) = (x^P) in v]
  proof −
   have [(◇(|E!, x^P|)) ∨ (¬◇(|E!, x^P|)) in v]
    using PLM.oth-class-taut-2 by simp
   hence [(◇(|E!, x^P|)) in v] ∨ [(¬◇(|E!, x^P|)) in v]
    using CP ∨E(1) by blast
   moreover {
    assume [(◇(|E!, x^P|)) in v]
    hence [(|λx. ◇(|E!,x^P|),x^P|) in v]
     apply (rule lambda-predicates-2-1[axiom-instance, equiv-rl, rotated])
     by show-proper
    hence [(|λx. ◇(|E!,x^P|),x^P|) & (|λx. ◇(|E!,x^P|),x^P|)
          & □(∀ F. (|F,x^P|) ≡ (|F,x^P|)) in v]
     apply − by PLM-solver
    hence [(x^P) =_E (x^P) in v]
     using eq-E-simple-1[equiv-rl] unfolding Ordinary-def by fast
   }
   moreover {
    assume [(¬◇(|E!, x^P|)) in v]
    hence [(|λx. ¬◇(|E!,x^P|),x^P|) in v]
     apply (rule lambda-predicates-2-1[axiom-instance, equiv-rl, rotated])
     by show-proper
    hence [(|λx. ¬◇(|E!,x^P|),x^P|) & (|λx. ¬◇(|E!,x^P|),x^P|)
          & □(∀ F. {|x^P,F|} ≡ {|x^P,F|}) in v]
     apply − by PLM-solver
   }
   ultimately show ?thesis unfolding identity-defs Ordinary-def Abstract-def
    using ∨I by blast
  qed
 lemma id-eq-obj-2[PLM]:
  [((x^P) = (y^P)) → ((y^P) = (x^P)) in v]
  by (meson l-identity[axiom-instance] id-eq-obj-1 CP ded-thm-cor-3)
 lemma id-eq-obj-3[PLM]:
  [((x^P) = (y^P)) & ((y^P) = (z^P)) → ((x^P) = (z^P)) in v]
  by (metis l-identity[axiom-instance] ded-thm-cor-4 CP &E)
end
```

Remark. *To unify the statements of the properties of equality a type class is introduced.*

```
class id-eq = quantifiable-and-identifiable +
 assumes id-eq-1: [(x :: 'a) = x in v]
 assumes id-eq-2: [((x :: 'a) = y) → (y = x) in v]
 assumes id-eq-3: [((x :: 'a) = y) & (y = z) → (x = z) in v]

instantiation ν :: id-eq
```

```
begin
  instance proof
    fix x :: ν and v
    show [x = x in v]
      using PLM.id-eq-obj-1
      by (simp add: identity-ν-def)
  next
    fix x y::ν and v
    show [x = y → y = x in v]
      using PLM.id-eq-obj-2
      by (simp add: identity-ν-def)
  next
    fix x y z::ν and v
    show [((x = y) & (y = z)) → x = z in v]
      using PLM.id-eq-obj-3
      by (simp add: identity-ν-def)
  qed
end

instantiation o :: id-eq
begin
  instance proof
    fix x :: o and v
    show [x = x in v]
      using PLM.id-eq-prop-prop-7 .
  next
    fix x y :: o and v
    show [x = y → y = x in v]
      using PLM.id-eq-prop-prop-8 .
  next
    fix x y z :: o and v
    show [((x = y) & (y = z)) → x = z in v]
      using PLM.id-eq-prop-prop-9 .
  qed
end

instantiation Π₁ :: id-eq
begin
  instance proof
    fix x :: Π₁ and v
    show [x = x in v]
      using PLM.id-eq-prop-prop-1 .
  next
    fix x y :: Π₁ and v
    show [x = y → y = x in v]
      using PLM.id-eq-prop-prop-2 .
  next
    fix x y z :: Π₁ and v
    show [((x = y) & (y = z)) → x = z in v]
      using PLM.id-eq-prop-prop-3 .
  qed
end

instantiation Π₂ :: id-eq
begin
  instance proof
    fix x :: Π₂ and v
    show [x = x in v]
```

```
    using PLM.id-eq-prop-prop-4-a .
  next
    fix x y :: Π2 and v
    show [x = y → y = x in v]
      using PLM.id-eq-prop-prop-5-a .
  next
    fix x y z :: Π2 and v
    show [((x = y) & (y = z)) → x = z in v]
      using PLM.id-eq-prop-prop-6-a .
  qed
end

instantiation Π3 :: id-eq
begin
  instance proof
    fix x :: Π3 and v
    show [x = x in v]
      using PLM.id-eq-prop-prop-4-b .
  next
    fix x y :: Π3 and v
    show [x = y → y = x in v]
      using PLM.id-eq-prop-prop-5-b .
  next
    fix x y z :: Π3 and v
    show [((x = y) & (y = z)) → x = z in v]
      using PLM.id-eq-prop-prop-6-b .
  qed
end

context PLM
begin
  lemma id-eq-1[PLM]:
    [(x::'a::id-eq) = x in v]
    using id-eq-1 .
  lemma id-eq-2[PLM]:
    [((x::'a::id-eq) = y) → (y = x) in v]
    using id-eq-2 .
  lemma id-eq-3[PLM]:
    [((x::'a::id-eq) = y) & (y = z) → (x = z) in v]
    using id-eq-3 .

  attribute-setup eq-sym = ⟪
    Scan.succeed (Thm.rule-attribute []
      (fn - => fn thm => thm RS @{thm id-eq-2[deduction]}))
  ⟫


  lemma all-self-eq-1[PLM]:
    [□(∀ α :: 'a::id-eq . α = α) in v]
    by PLM-solver
  lemma all-self-eq-2[PLM]:
    [∀ α :: 'a::id-eq . □(α = α) in v]
    by PLM-solver

  lemma t-id-t-proper-1[PLM]:
    [τ = τ' → (∃ β . (β^P) = τ) in v]
    proof (rule CP)
      assume [τ = τ' in v]
```

```
    moreover {
      assume [τ =_E τ' in v]
      hence [∃ β . (β^P) = τ in v]
        apply −
        apply (rule cqt-5-mod[where ψ=λ τ . τ =_E τ', axiom-instance, deduction])
         subgoal unfolding identity-defs by (rule SimpleExOrEnc.intros)
        by simp
    }
    moreover {
      assume [(|A!,τ|) & (|A!,τ'|) & □(∀ F. {|τ,F|} ≡ {|τ',F|}) in v]
      hence [∃ β . (β^P) = τ in v]
        apply −
        apply (rule cqt-5-mod[where ψ=λ τ . (|A!,τ|), axiom-instance, deduction])
         subgoal unfolding identity-defs by (rule SimpleExOrEnc.intros)
        by PLM-solver
    }
    ultimately show [∃ β . (β^P) = τ in v] unfolding identity_κ-def
      using intro-elim-4-b reductio-aa-1 by blast
  qed

lemma t-id-t-proper-2[PLM]: [τ = τ' → (∃ β . (β^P) = τ') in v]
proof (rule CP)
  assume [τ = τ' in v]
  moreover {
    assume [τ =_E τ' in v]
    hence [∃ β . (β^P) = τ' in v]
      apply −
      apply (rule cqt-5-mod[where ψ=λ τ' . τ =_E τ', axiom-instance, deduction])
       subgoal unfolding identity-defs by (rule SimpleExOrEnc.intros)
      by simp
  }
  moreover {
    assume [(|A!,τ|) & (|A!,τ'|) & □(∀ F. {|τ,F|} ≡ {|τ',F|}) in v]
    hence [∃ β . (β^P) = τ' in v]
      apply −
      apply (rule cqt-5-mod[where ψ=λ τ . (|A!,τ|), axiom-instance, deduction])
       subgoal unfolding identity-defs by (rule SimpleExOrEnc.intros)
      by PLM-solver
  }
  ultimately show [∃ β . (β^P) = τ' in v] unfolding identity_κ-def
    using intro-elim-4-b reductio-aa-1 by blast
qed

lemma id-nec[PLM]: [((α::'a::id-eq) = (β)) ≡ □((α) = (β)) in v]
  apply (rule ≡I)
   using l-identity[where φ = (λ β . □((α) = (β))), axiom-instance]
         id-eq-1 RN ded-thm-cor-4 unfolding identity-ν-def
   apply blast
  using qml-2[axiom-instance] by blast

lemma id-nec-desc[PLM]:
  [((ιx. φ x) = (ιx. ψ x)) ≡ □((ιx. φ x) = (ιx. ψ x)) in v]
  proof (cases [(∃ α. (α^P) = (ιx . φ x)) in v] ∧ [(∃ β. (β^P) = (ιx . ψ x)) in v])
    assume [(∃ α. (α^P) = (ιx . φ x)) in v] ∧ [(∃ β. (β^P) = (ιx . ψ x)) in v]
    then obtain α and β where
      [(α^P) = (ιx . φ x) in v] ∧ [(β^P) = (ιx . ψ x) in v]
      apply − unfolding conn-defs by PLM-solver
    moreover {
```

moreover have $[(\alpha) = (\beta) \equiv \Box((\alpha) = (\beta))\ in\ v]$ **by** *PLM-solver*
ultimately have $[((\iota x.\ \varphi\ x) = (\beta^P) \equiv \Box((\iota x.\ \varphi\ x) = (\beta^P)))\ in\ v]$
using *l-identity*[**where** $\varphi{=}\lambda\ \alpha\ .\ (\alpha) = (\beta^P) \equiv \Box((\alpha) = (\beta^P))$, *axiom-instance*]
modus-ponens **unfolding** *identity-ν-def* **by** *metis*
}
ultimately show *?thesis*
using *l-identity*[**where** $\varphi{=}\lambda\ \alpha\ .\ (\iota x\ .\ \varphi\ x) = (\alpha)$
$\equiv \Box((\iota x\ .\ \varphi\ x) = (\alpha))$, *axiom-instance*]
modus-ponens **by** *metis*
next
assume $\neg([(\exists\ \alpha.\ (\alpha^P) = (\iota x\ .\ \varphi\ x))\ in\ v] \wedge [(\exists\ \beta.\ (\beta^P) = (\iota x\ .\ \psi\ x))\ in\ v])$
hence $\neg[(\!|A!,(\iota x\ .\ \varphi\ x)|\!)\ in\ v] \wedge \neg[(\iota x\ .\ \varphi\ x) =_E (\iota x\ .\ \psi\ x)\ in\ v]$
$\vee\ \neg[(\!|A!,(\iota x\ .\ \psi\ x)|\!)\ in\ v] \wedge \neg[(\iota x\ .\ \varphi\ x) =_E (\iota x\ .\ \psi\ x)\ in\ v]$
unfolding *identity*$_E$*-infix-def*
using *cqt-5*[*axiom-instance*] *PLM.contraposition-1 SimpleExOrEnc.intros*
vdash-properties-10 **by** *meson*
hence $\neg[(\iota x\ .\ \varphi\ x) = (\iota x\ .\ \psi\ x)\ in\ v]$
apply − **unfolding** *identity-defs* **by** *PLM-solver*
thus *?thesis* **apply** − **apply** *PLM-solver*
using *qml-2*[*axiom-instance*, *deduction*] **by** *auto*
qed

A.9.8. Quantification

lemma *rule-ui*[*PLM*,*PLM-elim*,*PLM-dest*]:
$[\forall\,\alpha\ .\ \varphi\ \alpha\ in\ v] \Longrightarrow [\varphi\ \beta\ in\ v]$
by (*meson cqt-1*[*axiom-instance*, *deduction*])
lemmas $\forall\,E$ = *rule-ui*

lemma *rule-ui-2*[*PLM*,*PLM-elim*,*PLM-dest*]:
$[\![[\forall\,\alpha\ .\ \varphi\ (\alpha^P)\ in\ v]; [\exists\ \alpha\ .\ (\alpha)^P = \beta\ in\ v]]\!] \Longrightarrow [\varphi\ \beta\ in\ v]$
using *cqt-1-κ*[*axiom-instance*, *deduction*, *deduction*] **by** *blast*

lemma *cqt-orig-1*[*PLM*]:
$[(\forall\,\alpha.\ \varphi\ \alpha) \rightarrow \varphi\ \beta\ in\ v]$
by *PLM-solver*
lemma *cqt-orig-2*[*PLM*]:
$[(\forall\,\alpha.\ \varphi \rightarrow \psi\ \alpha) \rightarrow (\varphi \rightarrow (\forall\,\alpha.\ \psi\ \alpha))\ in\ v]$
by *PLM-solver*

lemma *universal*[*PLM*]:
$(\bigwedge\alpha\ .\ [\varphi\ \alpha\ in\ v]) \Longrightarrow [\forall\ \alpha\ .\ \varphi\ \alpha\ in\ v]$
using *rule-gen* **.**
lemmas $\forall\,I$ = *universal*

lemma *cqt-basic-1*[*PLM*]:
$[(\forall\,\alpha.\ (\forall\,\beta\ .\ \varphi\ \alpha\ \beta)) \equiv (\forall\,\beta.\ (\forall\,\alpha.\ \varphi\ \alpha\ \beta))\ in\ v]$
by *PLM-solver*
lemma *cqt-basic-2*[*PLM*]:
$[(\forall\,\alpha.\ \varphi\ \alpha \equiv \psi\ \alpha) \equiv ((\forall\,\alpha.\ \varphi\ \alpha \rightarrow \psi\ \alpha)\ \&\ (\forall\,\alpha.\ \psi\ \alpha \rightarrow \varphi\ \alpha))\ in\ v]$
by *PLM-solver*
lemma *cqt-basic-3*[*PLM*]:
$[(\forall\,\alpha.\ \varphi\ \alpha \equiv \psi\ \alpha) \rightarrow ((\forall\,\alpha.\ \varphi\ \alpha) \equiv (\forall\,\alpha.\ \psi\ \alpha))\ in\ v]$
by *PLM-solver*
lemma *cqt-basic-4*[*PLM*]:
$[(\forall\,\alpha.\ \varphi\ \alpha\ \&\ \psi\ \alpha) \equiv ((\forall\,\alpha.\ \varphi\ \alpha)\ \&\ (\forall\,\alpha.\ \psi\ \alpha))\ in\ v]$
by *PLM-solver*
lemma *cqt-basic-6*[*PLM*]:

$[(\forall\,\alpha.\ (\forall\,\alpha.\ \varphi\ \alpha))\equiv(\forall\,\alpha.\ \varphi\ \alpha)\ in\ v]$
by *PLM-solver*
lemma *cqt-basic-7*[*PLM*]:
$[(\varphi\rightarrow(\forall\,\alpha\ .\ \psi\ \alpha))\equiv(\forall\,\alpha.(\varphi\rightarrow\psi\ \alpha))\ in\ v]$
by *PLM-solver*
lemma *cqt-basic-8*[*PLM*]:
$[((\forall\,\alpha.\ \varphi\ \alpha)\vee(\forall\,\alpha.\ \psi\ \alpha))\rightarrow(\forall\,\alpha.\ (\varphi\ \alpha\vee\psi\ \alpha))\ in\ v]$
by *PLM-solver*
lemma *cqt-basic-9*[*PLM*]:
$[((\forall\,\alpha.\ \varphi\ \alpha\rightarrow\psi\ \alpha)\ \&\ (\forall\,\alpha.\ \psi\ \alpha\rightarrow\chi\ \alpha))\rightarrow(\forall\,\alpha.\ \varphi\ \alpha\rightarrow\chi\ \alpha)\ in\ v]$
by *PLM-solver*
lemma *cqt-basic-10*[*PLM*]:
$[((\forall\,\alpha.\ \varphi\ \alpha\equiv\psi\ \alpha)\ \&\ (\forall\,\alpha.\ \psi\ \alpha\equiv\chi\ \alpha))\rightarrow(\forall\,\alpha.\ \varphi\ \alpha\equiv\chi\ \alpha)\ in\ v]$
by *PLM-solver*
lemma *cqt-basic-11*[*PLM*]:
$[(\forall\,\alpha.\ \varphi\ \alpha\equiv\psi\ \alpha)\equiv(\forall\,\alpha.\ \psi\ \alpha\equiv\varphi\ \alpha)\ in\ v]$
by *PLM-solver*
lemma *cqt-basic-12*[*PLM*]:
$[(\forall\,\alpha.\ \varphi\ \alpha)\equiv(\forall\,\beta.\ \varphi\ \beta)\ in\ v]$
by *PLM-solver*

lemma *existential*[*PLM*,*PLM-intro*]:
$[\varphi\ \alpha\ in\ v]\Longrightarrow[\exists\ \alpha.\ \varphi\ \alpha\ in\ v]$
unfolding *exists-def* **by** *PLM-solver*
lemmas $\exists\,I$ = *existential*
lemma *instantiation*-[*PLM*,*PLM-elim*,*PLM-dest*]:
$[\![[\exists\,\alpha\ .\ \varphi\ \alpha\ in\ v];\ (\bigwedge\alpha.[\varphi\ \alpha\ in\ v]\Longrightarrow[\psi\ in\ v])]\!]\Longrightarrow[\psi\ in\ v]$
unfolding *exists-def* **by** *PLM-solver*

lemma *Instantiate*:
assumes $[\exists\ \ x\ .\ \varphi\ x\ \ in\ v]$
obtains x **where** $[\varphi\ x\ in\ v]$
apply (*insert assms*) **unfolding** *exists-def* **by** *PLM-solver*
lemmas $\exists\,E$ = *Instantiate*

lemma *cqt-further-1*[*PLM*]:
$[(\forall\,\alpha.\ \varphi\ \alpha)\rightarrow(\exists\,\alpha.\ \varphi\ \alpha)\ in\ v]$
by *PLM-solver*
lemma *cqt-further-2*[*PLM*]:
$[(\neg(\forall\,\alpha.\ \varphi\ \alpha))\equiv(\exists\,\alpha.\ \neg\varphi\ \alpha)\ in\ v]$
unfolding *exists-def* **by** *PLM-solver*
lemma *cqt-further-3*[*PLM*]:
$[(\forall\,\alpha.\ \varphi\ \alpha)\equiv\neg(\exists\,\alpha.\ \neg\varphi\ \alpha)\ in\ v]$
unfolding *exists-def* **by** *PLM-solver*
lemma *cqt-further-4*[*PLM*]:
$[(\neg(\exists\,\alpha.\ \varphi\ \alpha))\equiv(\forall\,\alpha.\ \neg\varphi\ \alpha)\ in\ v]$
unfolding *exists-def* **by** *PLM-solver*
lemma *cqt-further-5*[*PLM*]:
$[(\exists\,\alpha.\ \varphi\ \alpha\ \&\ \psi\ \alpha)\rightarrow((\exists\,\alpha.\ \varphi\ \alpha)\ \&\ (\exists\,\alpha.\ \psi\ \alpha))\ in\ v]$
unfolding *exists-def* **by** *PLM-solver*
lemma *cqt-further-6*[*PLM*]:
$[(\exists\,\alpha.\ \varphi\ \alpha\vee\psi\ \alpha)\equiv((\exists\,\alpha.\ \varphi\ \alpha)\vee(\exists\,\alpha.\ \psi\ \alpha))\ in\ v]$
unfolding *exists-def* **by** *PLM-solver*
lemma *cqt-further-10*[*PLM*]:
$[(\varphi\ (\alpha::'a::id\text{-}eq)\ \&\ (\forall\ \ \beta\ .\ \varphi\ \beta\rightarrow\beta=\alpha))\equiv(\forall\ \ \beta\ .\ \varphi\ \beta\equiv\beta=\alpha)\ in\ v]$
apply *PLM-solver*
using *l-identity*[*axiom-instance*, *deduction*, *deduction*] *id-eq-2*[*deduction*]
apply *blast*

```
  using id-eq-1 by auto
lemma cqt-further-11[PLM]:
  [((∀ α. φ α) & (∀ α. ψ α)) → (∀ α. φ α ≡ ψ α) in v]
  by PLM-solver
lemma cqt-further-12[PLM]:
  [((¬(∃ α. φ α)) & (¬(∃ α. ψ α))) → (∀ α. φ α ≡ ψ α) in v]
  unfolding exists-def by PLM-solver
lemma cqt-further-13[PLM]:
  [((∃ α. φ α) & (¬(∃ α. ψ α))) → (¬(∀ α. φ α ≡ ψ α)) in v]
  unfolding exists-def by PLM-solver
lemma cqt-further-14[PLM]:
  [(∃ α. ∃ β. φ α β) ≡ (∃ β. ∃ α. φ α β) in v]
  unfolding exists-def by PLM-solver

lemma nec-exist-unique[PLM]:
  [(∀ x. φ x → □(φ x)) → ((∃ !x. φ x) → (∃ !x. □(φ x))) in v]
  proof (rule CP)
    assume a: [∀ x. φ x → □φ x in v]
    show [(∃ !x. φ x) → (∃ !x. □φ x) in v]
    proof (rule CP)
      assume [(∃ !x. φ x) in v]
      hence [∃ α. φ α & (∀ β. φ β → β = α) in v]
        by (simp only: exists-unique-def)
      then obtain α where 1:
        [φ α & (∀ β. φ β → β = α) in v]
        by (rule ∃ E)
      {
        fix β
        have [□φ β → β = α in v]
          using 1 &E(2) qml-2[axiom-instance]
            ded-thm-cor-3 ∀ E by fastforce
      }
      hence [∀ β. □φ β → β = α in v] by (rule ∀ I)
      moreover have [□(φ α) in v]
        using 1 &E(1) a vdash-properties-10 cqt-orig-1[deduction]
        by fast
      ultimately have [∃ α. □(φ α) & (∀ β. □φ β → β = α) in v]
        using &I ∃ I by fast
      thus [(∃ !x. □φ x) in v]
        unfolding exists-unique-def by assumption
    qed
  qed
```

A.9.9. Actuality and Descriptions

```
lemma nec-imp-act[PLM]: [□φ → 𝒜φ in v]
  apply (rule CP)
  using qml-act-2[axiom-instance, equiv-lr]
        qml-2[axiom-actualization, axiom-instance]
        logic-actual-nec-2[axiom-instance, equiv-lr, deduction]
  by blast
lemma act-conj-act-1[PLM]:
  [𝒜(𝒜φ → φ) in v]
  using equiv-def logic-actual-nec-2[axiom-instance]
        logic-actual-nec-4[axiom-instance] &E(2) ≡E(2)
  by metis
lemma act-conj-act-2[PLM]:
  [𝒜(φ → 𝒜φ) in v]
```

using *logic-actual-nec-2*[*axiom-instance*] *qml-act-1*[*axiom-instance*]
ded-thm-cor-3 ≡*E*(*2*) *nec-imp-act*
by *blast*
lemma *act-conj-act-3*[*PLM*]:
[(𝒜φ & 𝒜ψ) → 𝒜(φ & ψ) *in v*]
unfolding *conn-defs*
by (*metis logic-actual-nec-2*[*axiom-instance*]
logic-actual-nec-1[*axiom-instance*]
≡*E*(*2*) *CP* ≡*E*(*4*) *reductio-aa-2*
vdash-properties-10)
lemma *act-conj-act-4*[*PLM*]:
[𝒜(𝒜φ ≡ φ) *in v*]
unfolding *equiv-def*
by (*PLM-solver PLM-intro*: *act-conj-act-3*[**where** φ=𝒜φ → φ
and ψ=φ → 𝒜φ, *deduction*])
lemma *closure-act-1a*[*PLM*]:
[𝒜𝒜(𝒜φ ≡ φ) *in v*]
using *logic-actual-nec-4*[*axiom-instance*]
act-conj-act-4 ≡*E*(*1*)
by *blast*
lemma *closure-act-1b*[*PLM*]:
[𝒜𝒜𝒜(𝒜φ ≡ φ) *in v*]
using *logic-actual-nec-4*[*axiom-instance*]
act-conj-act-4 ≡*E*(*1*)
by *blast*
lemma *closure-act-1c*[*PLM*]:
[𝒜𝒜𝒜𝒜(𝒜φ ≡ φ) *in v*]
using *logic-actual-nec-4*[*axiom-instance*]
act-conj-act-4 ≡*E*(*1*)
by *blast*
lemma *closure-act-2*[*PLM*]:
[∀ α. 𝒜(𝒜(φ α) ≡ φ α) *in v*]
by *PLM-solver*

lemma *closure-act-3*[*PLM*]:
[𝒜(∀ α. 𝒜(φ α) ≡ φ α) *in v*]
by (*PLM-solver PLM-intro*: *logic-actual-nec-3*[*axiom-instance*, *equiv-rl*])
lemma *closure-act-4*[*PLM*]:
[𝒜(∀ α₁ α₂. 𝒜(φ α₁ α₂) ≡ φ α₁ α₂) *in v*]
by (*PLM-solver PLM-intro*: *logic-actual-nec-3*[*axiom-instance*, *equiv-rl*])
lemma *closure-act-4-b*[*PLM*]:
[𝒜(∀ α₁ α₂ α₃. 𝒜(φ α₁ α₂ α₃) ≡ φ α₁ α₂ α₃) *in v*]
by (*PLM-solver PLM-intro*: *logic-actual-nec-3*[*axiom-instance*, *equiv-rl*])
lemma *closure-act-4-c*[*PLM*]:
[𝒜(∀ α₁ α₂ α₃ α₄. 𝒜(φ α₁ α₂ α₃ α₄) ≡ φ α₁ α₂ α₃ α₄) *in v*]
by (*PLM-solver PLM-intro*: *logic-actual-nec-3*[*axiom-instance*, *equiv-rl*])

lemma *RA*[*PLM*,*PLM-intro*]:
([φ *in dw*]) ⟹ [𝒜φ *in dw*]
using *logic-actual*[*necessitation-averse-axiom-instance*, *equiv-rl*] **.**

lemma *RA-2*[*PLM*,*PLM-intro*]:
([ψ *in dw*] ⟹ [φ *in dw*]) ⟹ ([𝒜ψ *in dw*] ⟹ [𝒜φ *in dw*])
using *RA logic-actual*[*necessitation-averse-axiom-instance*] *intro-elim-6-a* **by** *blast*

context
begin
private lemma *ActualE*[*PLM*,*PLM-elim*,*PLM-dest*]:

$[\mathcal{A}\varphi \ in \ dw] \Longrightarrow [\varphi \ in \ dw]$
using *logic-actual*[*necessitation-averse-axiom-instance*, *equiv-lr*] **.**

private lemma *NotActualD*[*PLM-dest*]:
$\neg[\mathcal{A}\varphi \ in \ dw] \Longrightarrow \neg[\varphi \ in \ dw]$
using *RA* **by** *metis*

private lemma *ActualImplI*[*PLM-intro*]:
$[\mathcal{A}\varphi \rightarrow \mathcal{A}\psi \ in \ v] \Longrightarrow [\mathcal{A}(\varphi \rightarrow \psi) \ in \ v]$
using *logic-actual-nec-2*[*axiom-instance*, *equiv-rl*] **.**
private lemma *ActualImplE*[*PLM-dest*, *PLM-elim*]:
$[\mathcal{A}(\varphi \rightarrow \psi) \ in \ v] \Longrightarrow [\mathcal{A}\varphi \rightarrow \mathcal{A}\psi \ in \ v]$
using *logic-actual-nec-2*[*axiom-instance*, *equiv-lr*] **.**
private lemma *NotActualImplD*[*PLM-dest*]:
$\neg[\mathcal{A}(\varphi \rightarrow \psi) \ in \ v] \Longrightarrow \neg[\mathcal{A}\varphi \rightarrow \mathcal{A}\psi \ in \ v]$
using *ActualImplI* **by** *blast*

private lemma *ActualNotI*[*PLM-intro*]:
$[\neg\mathcal{A}\varphi \ in \ v] \Longrightarrow [\mathcal{A}\neg\varphi \ in \ v]$
using *logic-actual-nec-1*[*axiom-instance*, *equiv-rl*] **.**
lemma *ActualNotE*[*PLM-elim*,*PLM-dest*]:
$[\mathcal{A}\neg\varphi \ in \ v] \Longrightarrow [\neg\mathcal{A}\varphi \ in \ v]$
using *logic-actual-nec-1*[*axiom-instance*, *equiv-lr*] **.**
lemma *NotActualNotD*[*PLM-dest*]:
$\neg[\mathcal{A}\neg\varphi \ in \ v] \Longrightarrow \neg[\neg\mathcal{A}\varphi \ in \ v]$
using *ActualNotI* **by** *blast*

private lemma *ActualConjI*[*PLM-intro*]:
$[\mathcal{A}\varphi \ \& \ \mathcal{A}\psi \ in \ v] \Longrightarrow [\mathcal{A}(\varphi \ \& \ \psi) \ in \ v]$
unfolding *equiv-def*
by (*PLM-solver PLM-intro*: *act-conj-act-3*[*deduction*])
private lemma *ActualConjE*[*PLM-elim*,*PLM-dest*]:
$[\mathcal{A}(\varphi \ \& \ \psi) \ in \ v] \Longrightarrow [\mathcal{A}\varphi \ \& \ \mathcal{A}\psi \ in \ v]$
unfolding *conj-def* **by** *PLM-solver*

private lemma *ActualEquivI*[*PLM-intro*]:
$[\mathcal{A}\varphi \equiv \mathcal{A}\psi \ in \ v] \Longrightarrow [\mathcal{A}(\varphi \equiv \psi) \ in \ v]$
unfolding *equiv-def*
by (*PLM-solver PLM-intro*: *act-conj-act-3*[*deduction*])
private lemma *ActualEquivE*[*PLM-elim*, *PLM-dest*]:
$[\mathcal{A}(\varphi \equiv \psi) \ in \ v] \Longrightarrow [\mathcal{A}\varphi \equiv \mathcal{A}\psi \ in \ v]$
unfolding *equiv-def* **by** *PLM-solver*

private lemma *ActualBoxI*[*PLM-intro*]:
$[\Box\varphi \ in \ v] \Longrightarrow [\mathcal{A}(\Box\varphi) \ in \ v]$
using *qml-act-2*[*axiom-instance*, *equiv-lr*] **.**
private lemma *ActualBoxE*[*PLM-elim*, *PLM-dest*]:
$[\mathcal{A}(\Box\varphi) \ in \ v] \Longrightarrow [\Box\varphi \ in \ v]$
using *qml-act-2*[*axiom-instance*, *equiv-rl*] **.**
private lemma *NotActualBoxD*[*PLM-dest*]:
$\neg[\mathcal{A}(\Box\varphi) \ in \ v] \Longrightarrow \neg[\Box\varphi \ in \ v]$
using *ActualBoxI* **by** *blast*

private lemma *ActualDisjI*[*PLM-intro*]:
$[\mathcal{A}\varphi \vee \mathcal{A}\psi \ in \ v] \Longrightarrow [\mathcal{A}(\varphi \vee \psi) \ in \ v]$
unfolding *disj-def* **by** *PLM-solver*
private lemma *ActualDisjE*[*PLM-elim*,*PLM-dest*]:
$[\mathcal{A}(\varphi \vee \psi) \ in \ v] \Longrightarrow [\mathcal{A}\varphi \vee \mathcal{A}\psi \ in \ v]$

unfolding *disj-def* **by** *PLM-solver*
private lemma *NotActualDisjD*[*PLM-dest*]:
$\neg[\mathcal{A}(\varphi \vee \psi)\ in\ v] \Longrightarrow \neg[\mathcal{A}\varphi \vee \mathcal{A}\psi\ in\ v]$
using *ActualDisjI* **by** *blast*

private lemma *ActualForallI*[*PLM-intro*]:
$[\forall\ x\ .\ \mathcal{A}(\varphi\ x)\ in\ v] \Longrightarrow [\mathcal{A}(\forall\ x\ .\ \varphi\ x)\ in\ v]$
using *logic-actual-nec-3*[*axiom-instance*, *equiv-rl*] .
lemma *ActualForallE*[*PLM-elim*,*PLM-dest*]:
$[\mathcal{A}(\forall\ x\ .\ \varphi\ x)\ in\ v] \Longrightarrow [\forall\ x\ .\ \mathcal{A}(\varphi\ x)\ in\ v]$
using *logic-actual-nec-3*[*axiom-instance*, *equiv-lr*] .
lemma *NotActualForallD*[*PLM-dest*]:
$\neg[\mathcal{A}(\forall\ x\ .\ \varphi\ x)\ in\ v] \Longrightarrow \neg[\forall\ x\ .\ \mathcal{A}(\varphi\ x)\ in\ v]$
using *ActualForallI* **by** *blast*

lemma *ActualActualI*[*PLM-intro*]:
$[\mathcal{A}\varphi\ in\ v] \Longrightarrow [\mathcal{A}\mathcal{A}\varphi\ in\ v]$
using *logic-actual-nec-4*[*axiom-instance*, *equiv-lr*] .
lemma *ActualActualE*[*PLM-elim*,*PLM-dest*]:
$[\mathcal{A}\mathcal{A}\varphi\ in\ v] \Longrightarrow [\mathcal{A}\varphi\ in\ v]$
using *logic-actual-nec-4*[*axiom-instance*, *equiv-rl*] .
lemma *NotActualActualD*[*PLM-dest*]:
$\neg[\mathcal{A}\mathcal{A}\varphi\ in\ v] \Longrightarrow \neg[\mathcal{A}\varphi\ in\ v]$
using *ActualActualI* **by** *blast*
end

lemma *ANeg-1*[*PLM*]:
$[\neg\mathcal{A}\varphi \equiv \neg\varphi\ in\ dw]$
by *PLM-solver*
lemma *ANeg-2*[*PLM*]:
$[\neg\mathcal{A}\neg\varphi \equiv \varphi\ in\ dw]$
by *PLM-solver*
lemma *Act-Basic-1*[*PLM*]:
$[\mathcal{A}\varphi \vee \mathcal{A}\neg\varphi\ in\ v]$
by *PLM-solver*
lemma *Act-Basic-2*[*PLM*]:
$[\mathcal{A}(\varphi\ \&\ \psi) \equiv (\mathcal{A}\varphi\ \&\ \mathcal{A}\psi)\ in\ v]$
by *PLM-solver*
lemma *Act-Basic-3*[*PLM*]:
$[\mathcal{A}(\varphi \equiv \psi) \equiv ((\mathcal{A}(\varphi \rightarrow \psi))\ \&\ (\mathcal{A}(\psi \rightarrow \varphi)))\ in\ v]$
by *PLM-solver*
lemma *Act-Basic-4*[*PLM*]:
$[(\mathcal{A}(\varphi \rightarrow \psi)\ \&\ \mathcal{A}(\psi \rightarrow \varphi)) \equiv (\mathcal{A}\varphi \equiv \mathcal{A}\psi)\ in\ v]$
by *PLM-solver*
lemma *Act-Basic-5*[*PLM*]:
$[\mathcal{A}(\varphi \equiv \psi) \equiv (\mathcal{A}\varphi \equiv \mathcal{A}\psi)\ in\ v]$
by *PLM-solver*
lemma *Act-Basic-6*[*PLM*]:
$[\Diamond\varphi \equiv \mathcal{A}(\Diamond\varphi)\ in\ v]$
unfolding *diamond-def* **by** *PLM-solver*
lemma *Act-Basic-7*[*PLM*]:
$[\mathcal{A}\varphi \equiv \Box\mathcal{A}\varphi\ in\ v]$
by (*simp add*: *qml-2*[*axiom-instance*] *qml-act-1*[*axiom-instance*] $\equiv I$)
lemma *Act-Basic-8*[*PLM*]:
$[\mathcal{A}(\Box\varphi) \rightarrow \Box\mathcal{A}\varphi\ in\ v]$
by (*metis qml-act-2*[*axiom-instance*] *CP Act-Basic-7* $\equiv E(1)$
$\equiv E(2)$ *nec-imp-act vdash-properties-10*)
lemma *Act-Basic-9*[*PLM*]:

```
  [□φ → □𝒜φ in v]
  using qml-act-1[axiom-instance] ded-thm-cor-3 nec-imp-act by blast
lemma Act-Basic-10[PLM]:
  [𝒜(φ ∨ ψ) ≡ 𝒜φ ∨ 𝒜ψ in v]
  by PLM-solver

lemma Act-Basic-11[PLM]:
  [𝒜(∃ α. φ α) ≡ (∃ α.𝒜(φ α)) in v]
  proof −
    have [𝒜(∀ α . ¬φ α) ≡ (∀ α . 𝒜¬φ α) in v]
      using logic-actual-nec-3[axiom-instance] by blast
    hence [¬𝒜(∀ α . ¬φ α) ≡ ¬(∀ α . 𝒜¬φ α) in v]
      using oth-class-taut-5-d[equiv-lr] by blast
    moreover have [𝒜¬(∀ α . ¬φ α) ≡ ¬𝒜(∀ α . ¬φ α) in v]
      using logic-actual-nec-1[axiom-instance] by blast
    ultimately have [𝒜¬(∀ α . ¬φ α) ≡ ¬(∀ α . 𝒜¬φ α) in v]
      using ≡E(5) by auto
    moreover {
      have [∀ α . 𝒜¬φ α ≡ ¬𝒜φ α in v]
        using logic-actual-nec-1[axiom-universal, axiom-instance] by blast
      hence [(∀ α . 𝒜¬φ α) ≡ (∀ α . ¬𝒜φ α) in v]
        using cqt-basic-3[deduction] by fast
      hence [(¬(∀ α . 𝒜¬φ α)) ≡ ¬(∀ α . ¬𝒜φ α) in v]
        using oth-class-taut-5-d[equiv-lr] by blast
    }
    ultimately show ?thesis unfolding exists-def using ≡E(5) by auto
  qed

lemma act-quant-uniq[PLM]:
  [(∀ z . 𝒜φ z ≡ z = x) ≡ (∀ z . φ z ≡ z = x) in dw]
  by PLM-solver

lemma fund-cont-desc[PLM]:
  [(x^P = (ιx. φ x)) ≡ (∀ z . φ z ≡ (z = x)) in dw]
  using descriptions[axiom-instance] act-quant-uniq ≡E(5) by fast

lemma hintikka[PLM]:
  [(x^P = (ιx. φ x)) ≡ (φ x & (∀ z . φ z → z = x)) in dw]
  proof −
    have [(∀ z . φ z ≡ z = x) ≡ (φ x & (∀ z . φ z → z = x)) in dw]
      unfolding identity-ν-def apply PLM-solver using id-eq-obj-1 apply simp
      using l-identity[where φ=λ x . φ x, axiom-instance,
                       deduction, deduction]
      using id-eq-obj-2[deduction] unfolding identity-ν-def by fastforce
    thus ?thesis using ≡E(5) fund-cont-desc by blast
  qed

lemma russell-axiom-a[PLM]:
  [(⦇F, ιx. φ x⦈) ≡ (∃ x . φ x & (∀ z . φ z → z = x) & ⦇F, x^P⦈) in dw]
  (is [?lhs ≡ ?rhs in dw])
  proof −
    {
      assume 1: [?lhs in dw]
      hence [∃ α. α^P = (ιx. φ x) in dw]
      using cqt-5[axiom-instance, deduction]
            SimpleExOrEnc.intros
      by blast
      then obtain α where 2:
```

```
    [α^P = (ιx. φ x) in dw]
    using ∃E by auto
  hence 3: [φ α & (∀ z . φ z → z = α) in dw]
    using hintikka[equiv-lr] by simp
  from 2 have [(ιx. φ x) = (α^P) in dw]
    using l-identity[where α=α^P and β=ιx. φ x and φ=λ x . x = α^P,
          axiom-instance, deduction, deduction]
          id-eq-obj-1[where x=α] by auto
  hence [(|F, α^P|) in dw]
  using 1 l-identity[where β=α^P and α=ιx. φ x and φ=λ x . (|F,x|),
                     axiom-instance, deduction, deduction] by auto
  with 3 have [φ α & (∀ z . φ z → z = α) & (|F, α^P|) in dw] by (rule &I)
  hence [?rhs in dw] using ∃I[where α=α] by simp
 }
 moreover {
  assume [?rhs in dw]
  then obtain α where 4:
    [φ α & (∀ z . φ z → z = α) & (|F, α^P|) in dw]
    using ∃E by auto
  hence [α^P = (ιx . φ x) in dw] ∧ [(|F, α^P|) in dw]
    using hintikka[equiv-rl] &E by blast
  hence [?lhs in dw]
    using l-identity[axiom-instance, deduction, deduction]
    by blast
 }
 ultimately show ?thesis by PLM-solver
qed

lemma russell-axiom-g[PLM]:
 [{|ιx. φ x,F|} ≡ (∃ x . φ x & (∀ z . φ z → z = x) & {|x^P, F|}) in dw]
 (is [?lhs ≡ ?rhs in dw])
 proof −
  {
   assume 1: [?lhs in dw]
   hence [∃α. α^P = (ιx. φ x) in dw]
   using cqt-5[axiom-instance, deduction] SimpleExOrEnc.intros by blast
   then obtain α where 2: [α^P = (ιx. φ x) in dw] by (rule ∃E)
   hence 3: [(φ α & (∀ z . φ z → z = α)) in dw]
     using hintikka[equiv-lr] by simp
   from 2 have [(ιx. φ x) = α^P in dw]
     using l-identity[where α=α^P and β=ιx. φ x and φ=λ x . x = α^P,
           axiom-instance, deduction, deduction]
           id-eq-obj-1[where x=α] by auto
   hence [{|α^P, F|} in dw]
   using 1 l-identity[where β=α^P and α=ιx. φ x and φ=λ x . {|x,F|},
                      axiom-instance, deduction, deduction] by auto
   with 3 have [(φ α & (∀ z . φ z → z = α)) & {|α^P, F|} in dw]
     using &I by auto
   hence [?rhs in dw] using ∃I[where α=α] by (simp add: identity-defs)
  }
  moreover {
   assume [?rhs in dw]
   then obtain α where 4:
     [φ α & (∀ z . φ z → z = α) & {|α^P, F|} in dw]
     using ∃E by auto
   hence [α^P = (ιx . φ x) in dw] ∧ [{|α^P, F|} in dw]
     using hintikka[equiv-rl] &E by blast
   hence [?lhs in dw]
```

```
      using l-identity[axiom-instance, deduction, deduction]
      by fast
    }
    ultimately show ?thesis by PLM-solver
  qed

lemma russell-axiom[PLM]:
  assumes SimpleExOrEnc ψ
  shows [ψ (ιx. φ x) ≡ (∃ x . φ x & (∀ z . φ z → z = x) & ψ (x^P)) in dw]
  (is [?lhs ≡ ?rhs in dw])
  proof −
    {
      assume 1: [?lhs in dw]
      hence [∃ α. α^P = (ιx. φ x) in dw]
      using cqt-5[axiom-instance, deduction] assms by blast
      then obtain α where 2: [α^P = (ιx. φ x) in dw] by (rule ∃E)
      hence 3: [(φ α & (∀ z . φ z → z = α)) in dw]
        using hintikka[equiv-lr] by simp
      from 2 have [(ιx. φ x) = (α^P) in dw]
        using l-identity[where α=α^P and β=ιx. φ x and φ=λ x . x = α^P,
              axiom-instance, deduction, deduction]
              id-eq-obj-1[where x=α] by auto
      hence [ψ (α^P) in dw]
        using 1 l-identity[where β=α^P and α=ιx. φ x and φ=λ x . ψ x,
                           axiom-instance, deduction, deduction] by auto
      with 3 have [φ α & (∀ z . φ z → z = α) & ψ (α^P) in dw]
        using &I by auto
      hence [?rhs in dw] using ∃I[where α=α] by (simp add: identity-defs)
    }
    moreover {
      assume [?rhs in dw]
      then obtain α where 4:
        [φ α & (∀ z . φ z → z = α) & ψ (α^P) in dw]
        using ∃E by auto
      hence [α^P = (ιx . φ x) in dw] ∧ [ψ (α^P) in dw]
        using hintikka[equiv-rl] &E by blast
      hence [?lhs in dw]
        using l-identity[axiom-instance, deduction, deduction]
        by fast
    }
    ultimately show ?thesis by PLM-solver
  qed

lemma unique-exists[PLM]:
  [(∃ y . y^P = (ιx. φ x)) ≡ (∃!x . φ x) in dw]
  proof((rule ≡I, rule CP, rule-tac[2] CP))
    assume [∃ y. y^P = (ιx. φ x) in dw]
    then obtain α where
      [α^P = (ιx. φ x) in dw]
      by (rule ∃E)
    hence [φ α & (∀β. φ β → β = α) in dw]
      using hintikka[equiv-lr] by auto
    thus [∃!x . φ x in dw]
      unfolding exists-unique-def using ∃I by fast
  next
    assume [∃!x . φ x in dw]
    then obtain α where
      [φ α & (∀β. φ β → β = α) in dw]
```

unfolding *exists-unique-def* **by** (*rule* $\exists E$)
hence $[\alpha^P = (\boldsymbol{\iota}x.\ \varphi\ x)\ in\ dw]$
using *hintikka*[*equiv-rl*] **by** *auto*
thus $[\exists y.\ y^P = (\boldsymbol{\iota}x.\ \varphi\ x)\ in\ dw]$
using $\exists I$ **by** *fast*
qed

lemma *y-in-1*[*PLM*]:
$[x^P = (\boldsymbol{\iota}x\ .\ \varphi) \rightarrow \varphi\ in\ dw]$
using *hintikka*[*equiv-lr*, *conj1*] **by** (*rule CP*)

lemma *y-in-2*[*PLM*]:
$[z^P = (\boldsymbol{\iota}x\ .\ \varphi\ x) \rightarrow \varphi\ z\ in\ dw]$
using *hintikka*[*equiv-lr*, *conj1*] **by** (*rule CP*)

lemma *y-in-3*[*PLM*]:
$[(\exists\ y\ .\ y^P = (\boldsymbol{\iota}x\ .\ \varphi\ (x^P))) \rightarrow \varphi\ (\boldsymbol{\iota}x\ .\ \varphi\ (x^P))\ in\ dw]$
proof (*rule CP*)
assume $[(\exists\ y\ .\ y^P = (\boldsymbol{\iota}x\ .\ \varphi\ (x^P)))\ in\ dw]$
then obtain *y* **where** *1*:
$[y^P = (\boldsymbol{\iota}x.\ \varphi\ (x^P))\ in\ dw]$
by (*rule* $\exists E$)
hence $[\varphi\ (y^P)\ in\ dw]$
using *y-in-2*[*deduction*] **unfolding** *identity-ν-def* **by** *blast*
thus $[\varphi\ (\boldsymbol{\iota}x.\ \varphi\ (x^P))\ in\ dw]$
using *l-identity*[*axiom-instance*, *deduction*,
deduction] *1* **by** *fast*
qed

lemma *act-quant-nec*[*PLM*]:
$[(\forall z\ .\ (\boldsymbol{\mathcal{A}}\varphi\ z \equiv z = x)) \equiv (\forall z.\ \boldsymbol{\mathcal{A}}\boldsymbol{\mathcal{A}}\varphi\ z \equiv z = x)\ in\ v]$
by *PLM-solver*

lemma *equi-desc-descA-1*[*PLM*]:
$[(x^P = (\boldsymbol{\iota}x\ .\ \varphi\ x)) \equiv (x^P = (\boldsymbol{\iota}x\ .\ \boldsymbol{\mathcal{A}}\varphi\ x))\ in\ v]$
using *descriptions*[*axiom-instance*] **apply** (*rule* $\equiv E(5)$)
using *act-quant-nec* **apply** (*rule* $\equiv E(5)$)
using *descriptions*[*axiom-instance*]
by (*meson* $\equiv E(6)$ *oth-class-taut-4-a*)

lemma *equi-desc-descA-2*[*PLM*]:
$[(\exists\ y\ .\ y^P = (\boldsymbol{\iota}x.\ \varphi\ x)) \rightarrow ((\boldsymbol{\iota}x\ .\ \varphi\ x) = (\boldsymbol{\iota}x\ .\ \boldsymbol{\mathcal{A}}\varphi\ x))\ in\ v]$
proof (*rule CP*)
assume $[\exists y.\ y^P = (\boldsymbol{\iota}x.\ \varphi\ x)\ in\ v]$
then obtain *y* **where**
$[y^P = (\boldsymbol{\iota}x.\ \varphi\ x)\ in\ v]$
by (*rule* $\exists E$)
moreover hence $[y^P = (\boldsymbol{\iota}x.\ \boldsymbol{\mathcal{A}}\varphi\ x)\ in\ v]$
using *equi-desc-descA-1*[*equiv-lr*] **by** *auto*
ultimately show $[(\boldsymbol{\iota}x.\ \varphi\ x) = (\boldsymbol{\iota}x.\ \boldsymbol{\mathcal{A}}\varphi\ x)\ in\ v]$
using *l-identity*[*axiom-instance*, *deduction*, *deduction*]
by *fast*
qed

lemma *equi-desc-descA-3*[*PLM*]:
assumes *SimpleExOrEnc* ψ
shows $[\psi\ (\boldsymbol{\iota}x.\ \varphi\ x) \rightarrow (\exists\ y\ .\ y^P = (\boldsymbol{\iota}x.\ \boldsymbol{\mathcal{A}}\varphi\ x))\ in\ v]$
proof (*rule CP*)

```
    assume [ψ (ιx. φ x) in v]
    hence [∃α. α^P = (ιx. φ x) in v]
      using cqt-5[OF assms, axiom-instance, deduction] by auto
    then obtain α where [α^P = (ιx. φ x) in v] by (rule ∃E)
    hence [α^P = (ιx . 𝒜φ x) in v]
      using equi-desc-descA-1[equiv-lr] by auto
    thus [∃y. y^P = (ιx. 𝒜φ x) in v]
      using ∃I by fast
  qed

lemma equi-desc-descA-4[PLM]:
  assumes SimpleExOrEnc ψ
  shows [ψ (ιx. φ x) → ((ιx. φ x) = (ιx. 𝒜φ x)) in v]
  proof (rule CP)
    assume [ψ (ιx. φ x) in v]
    hence [∃α. α^P = (ιx. φ x) in v]
      using cqt-5[OF assms, axiom-instance, deduction] by auto
    then obtain α where [α^P = (ιx. φ x) in v] by (rule ∃E)
    moreover hence [α^P  = (ιx . 𝒜φ x) in v]
      using equi-desc-descA-1[equiv-lr] by auto
    ultimately show [(ιx. φ x)  = (ιx . 𝒜φ x) in v]
      using l-identity[axiom-instance, deduction, deduction] by fast
  qed

lemma nec-hintikka-scheme[PLM]:
  [(x^P = (ιx. φ x)) ≡ (𝒜φ x & (∀ z . 𝒜φ z → z = x)) in v]
  using descriptions[axiom-instance]
  apply (rule ≡E(5))
  apply PLM-solver
   using id-eq-obj-1 apply simp
   using id-eq-obj-2[deduction]
         l-identity[where α=x, axiom-instance, deduction, deduction]
   unfolding identity-ν-def
   apply blast
  using l-identity[where α=x, axiom-instance, deduction, deduction]
  id-eq-2[where 'a=ν, deduction] unfolding identity-ν-def by meson

lemma equiv-desc-eq[PLM]:
  assumes ⋀x.[𝒜(φ x ≡ ψ x) in v]
  shows [(∀ x . ((x^P = (ιx . φ x)) ≡ (x^P = (ιx . ψ x)))) in v]
  proof(rule ∀I)
    fix x
    {
      assume [x^P = (ιx . φ x) in v]
      hence 1: [𝒜φ x & (∀z. 𝒜φ z → z = x) in v]
        using nec-hintikka-scheme[equiv-lr] by auto
      hence 2: [𝒜φ x in v] ∧ [(∀z. 𝒜φ z → z = x) in v]
        using &E by blast
      {
        fix z
        {
          assume [𝒜ψ z in v]
          hence [𝒜φ z in v]
           using assms[where x=z] apply − by PLM-solver
          moreover have [𝒜φ z → z = x in v]
            using 2 cqt-1[axiom-instance,deduction] by auto
          ultimately have [z = x in v]
           using vdash-properties-10 by auto
```

```
      }
      hence [𝒜ψ z → z = x in v] by (rule CP)
    }
    hence [(∀ z . 𝒜ψ z → z = x) in v] by (rule ∀I)
    moreover have [𝒜ψ x in v]
      using 1[conj1] assms[where x=x]
      apply − by PLM-solver
    ultimately have [𝒜ψ x & (∀z. 𝒜ψ z → z = x) in v]
      by PLM-solver
    hence [x^P = (ιx. ψ x) in v]
     using nec-hintikka-scheme[where φ=ψ, equiv-rl] by auto
  }
  moreover {
    assume [x^P = (ιx . ψ x) in v]
    hence 1: [𝒜ψ x & (∀z. 𝒜ψ z → z = x) in v]
      using nec-hintikka-scheme[equiv-lr] by auto
    hence 2: [𝒜ψ x in v] ∧ [(∀z. 𝒜ψ z → z = x) in v]
      using &E by blast
    {
      fix z
      {
        assume [𝒜φ z in v]
        hence [𝒜ψ z in v]
          using assms[where x=z]
          apply − by PLM-solver
        moreover have [𝒜ψ z → z = x in v]
          using 2 cqt-1[axiom-instance,deduction] by auto
        ultimately have [z = x in v]
          using vdash-properties-10 by auto
      }
      hence [𝒜φ z → z = x in v] by (rule CP)
    }
    hence [(∀z. 𝒜φ z → z = x) in v] by (rule ∀I)
    moreover have [𝒜φ x in v]
      using 1[conj1] assms[where x=x]
      apply − by PLM-solver
    ultimately have [𝒜φ x & (∀z. 𝒜φ z → z = x) in v]
      by PLM-solver
    hence [x^P = (ιx. φ x) in v]
      using nec-hintikka-scheme[where φ=φ,equiv-rl]
      by auto
  }
  ultimately show [x^P = (ιx. φ x) ≡ (x^P) = (ιx. ψ x) in v]
    using ≡I CP by auto
qed

lemma UniqueAux:
  assumes [(𝒜φ (α::ν) & (∀ z . 𝒜(φ z) → z = α)) in v]
  shows [(∀ z . (𝒜(φ z) ≡ (z = α))) in v]
  proof −
    {
      fix z
      {
        assume [𝒜(φ z) in v]
        hence [z = α in v]
          using assms[conj2, THEN cqt-1[where α=z,
                      axiom-instance, deduction],
                    deduction] by auto
```

```
    }
    moreover {
      assume [z = α in v]
      hence [α = z in v]
        unfolding identity-ν-def
        using id-eq-obj-2[deduction] by fast
      hence [𝒜(φ z) in v] using assms[conj1]
        using l-identity[axiom-instance, deduction,
                         deduction] by fast
    }
    ultimately have [(𝒜(φ z) ≡ (z = α)) in v]
      using ≡I CP by auto
  }
  thus [(∀ z . (𝒜(φ z) ≡ (z = α))) in v]
  by (rule ∀I)
qed

lemma nec-russell-axiom[PLM]:
  assumes SimpleExOrEnc ψ
  shows [(ψ (ιx. φ x)) ≡ (∃ x . (𝒜φ x & (∀ z . 𝒜(φ z) → z = x))
                           & ψ (x^P)) in v]
  (is [?lhs ≡ ?rhs in v])
  proof −
    {
      assume 1: [?lhs in v]
      hence [∃α. (α^P) = (ιx. φ x) in v]
        using cqt-5[axiom-instance, deduction] assms by blast
      then obtain α where 2: [(α^P) = (ιx. φ x) in v] by (rule ∃E)
      hence [(∀ z . (𝒜(φ z) ≡ (z = α))) in v]
        using descriptions[axiom-instance, equiv-lr] by auto
      hence 3: [(𝒜φ α) & (∀ z . (𝒜(φ z) → (z = α))) in v]
        using cqt-1[where α=α and φ=λ z . (𝒜(φ z) ≡ (z = α)),
                    axiom-instance, deduction, equiv-rl]
        using id-eq-obj-1[where x=α] unfolding identity-ν-def
        using hintikka[equiv-lr] cqt-basic-2[equiv-lr,conj1]
        &I by fast
      from 2 have [(ιx. φ x) = (α^P) in v]
        using l-identity[where β=(ιx. φ x) and φ=λ x . x = (α^P),
              axiom-instance, deduction, deduction]
              id-eq-obj-1[where x=α] by auto
      hence [ψ (α^P) in v]
        using 1 l-identity[where α=(ιx. φ x) and φ=λ x . ψ x,
                           axiom-instance, deduction,
                           deduction] by auto
      with 3 have [(𝒜φ α & (∀ z . 𝒜(φ z) → (z = α))) & ψ (α^P) in v]
        using &I by simp
      hence [?rhs in v]
        using ∃I[where α=α]
        by (simp add: identity-defs)
    }
    moreover {
      assume [?rhs in v]
      then obtain α where 4:
        [(𝒜φ α & (∀ z . 𝒜(φ z) → z = α)) & ψ (α^P) in v]
        using ∃E by auto
      hence [(∀ z . (𝒜(φ z) ≡ (z = α))) in v]
        using UniqueAux &E(1) by auto
      hence [(α^P) = (ιx . φ x) in v] ∧ [ψ (α^P) in v]
```

```
      using descriptions[axiom-instance, equiv-rl]
          4[conj2] by blast
    hence [?lhs in v]
      using l-identity[axiom-instance, deduction,
                       deduction]
      by fast
  }
  ultimately show ?thesis by PLM-solver
qed

lemma actual-desc-1[PLM]:
[(∃ y . (y^P) = (ιx. φ x)) ≡ (∃! x . 𝒜(φ x)) in v] (is [?lhs ≡ ?rhs in v])
proof −
  {
    assume [?lhs in v]
    then obtain α where
      [((α^P) = (ιx. φ x)) in v]
      by (rule ∃E)
    hence [(|A!,(ιx. φ x)|) in v] ∨ [(α^P) =_E (ιx. φ x) in v]
      apply − unfolding identity-defs by PLM-solver
    then obtain x where
      [((𝒜φ x & (∀ z . 𝒜(φ z) → z = x))) in v]
      using nec-russell-axiom[where ψ=λx . (|A!,x|), equiv-lr, THEN ∃E]
      using nec-russell-axiom[where ψ=λx . (α^P) =_E x, equiv-lr, THEN ∃E]
      using SimpleExOrEnc.intros unfolding identity_E-infix-def
      by (meson &E)
    hence [?rhs in v] unfolding exists-unique-def by (rule ∃I)
  }
  moreover {
    assume [?rhs in v]
    then obtain x where
      [((𝒜φ x & (∀ z . 𝒜(φ z) → z = x))) in v]
      unfolding exists-unique-def by (rule ∃E)
    hence [∀z. 𝒜φ z ≡ z = x in v]
      using UniqueAux by auto
    hence [(x^P) = (ιx. φ x) in v]
      using descriptions[axiom-instance, equiv-rl] by auto
    hence [?lhs in v] by (rule ∃I)
  }
  ultimately show ?thesis
    using ≡I CP by auto
qed

lemma actual-desc-2[PLM]:
[(x^P) = (ιx. φ) → 𝒜φ in v]
using nec-hintikka-scheme[equiv-lr, conj1]
by (rule CP)

lemma actual-desc-3[PLM]:
[(z^P) = (ιx. φ x) → 𝒜(φ z) in v]
using nec-hintikka-scheme[equiv-lr, conj1]
by (rule CP)

lemma actual-desc-4[PLM]:
[(∃ y . ((y^P) = (ιx. φ (x^P)))) → 𝒜(φ (ιx. φ (x^P))) in v]
proof (rule CP)
  assume [(∃ y . (y^P) = (ιx . φ (x^P))) in v]
  then obtain y where 1:
```

```
    [y^P = (ιx. φ (x^P)) in v]
    by (rule ∃ E)
  hence [A(φ (y^P)) in v] using actual-desc-3[deduction] by fast
  thus [A(φ (ιx. φ (x^P))) in v]
    using l-identity[axiom-instance, deduction,
                     deduction] 1 by fast
qed

lemma unique-box-desc-1[PLM]:
  [(∃!x . □(φ x)) → (∀ y . (y^P) = (ιx. φ x) → φ y) in v]
  proof (rule CP)
    assume [(∃!x . □(φ x)) in v]
    then obtain α where 1:
      [□φ α & (∀β. □(φ β) → β = α) in v]
      unfolding exists-unique-def by (rule ∃ E)
    {
      fix y
      {
        assume [(y^P) = (ιx. φ x) in v]
        hence [Aφ α → α = y in v]
          using nec-hintikka-scheme[where x=y and φ=φ, equiv-lr, conj2,
                       THEN cqt-1[where α=α,axiom-instance, deduction]] by simp
        hence [α = y in v]
          using 1[conj1] nec-imp-act vdash-properties-10 by blast
        hence [φ y in v]
          using 1[conj1] qml-2[axiom-instance, deduction]
                l-identity[axiom-instance, deduction, deduction]
          by fast
      }
      hence [(y^P) = (ιx. φ x) → φ y in v]
        by (rule CP)
    }
    thus [∀ y . (y^P) = (ιx. φ x) → φ y in v]
      by (rule ∀ I)
qed

lemma unique-box-desc[PLM]:
  [(∀ x . (φ x → □(φ x))) → ((∃!x . φ x)
    → (∀ y . (y^P = (ιx . φ x)) → φ y)) in v]
  apply (rule CP, rule CP)
  using nec-exist-unique[deduction, deduction]
        unique-box-desc-1[deduction] by blast
```

A.9.10. Necessity

```
lemma RM-1[PLM]:
  (⋀v.[φ → ψ in v]) ⟹ [□φ → □ψ in v]
  using RN qml-1[axiom-instance] vdash-properties-10 by blast

lemma RM-1-b[PLM]:
  (⋀v.[χ in v] ⟹ [φ → ψ in v]) ⟹ ([□χ in v] ⟹ [□φ → □ψ in v])
  using RN-2 qml-1[axiom-instance] vdash-properties-10 by blast

lemma RM-2[PLM]:
  (⋀v.[φ → ψ in v]) ⟹ [◇φ → ◇ψ in v]
  unfolding diamond-def
  using RM-1 contraposition-1 by auto
```

lemma *RM-2-b*[*PLM*]:
(⋀*v*.[*χ in v*] ⟹ [*φ* → *ψ in v*]) ⟹ ([□*χ in v*] ⟹ [◇*φ* → ◇*ψ in v*])
unfolding *diamond-def*
using *RM-1-b contraposition-1* **by** *blast*

lemma *KBasic-1*[*PLM*]:
[□*φ* → □(*ψ* → *φ*) *in v*]
by (*simp only*: *pl-1*[*axiom-instance*] *RM-1*)
lemma *KBasic-2*[*PLM*]:
[□(¬*φ*) → □(*φ* → *ψ*) *in v*]
by (*simp only*: *RM-1 useful-tautologies-3*)
lemma *KBasic-3*[*PLM*]:
[□(*φ* & *ψ*) ≡ □*φ* & □*ψ in v*]
apply (*rule* ≡*I*)
apply (*rule CP*)
apply (*rule* &*I*)
using *RM-1 oth-class-taut-9-a vdash-properties-6* **apply** *blast*
using *RM-1 oth-class-taut-9-b vdash-properties-6* **apply** *blast*
using *qml-1*[*axiom-instance*] *RM-1 ded-thm-cor-3 oth-class-taut-10-a*
oth-class-taut-8-b vdash-properties-10
by *blast*
lemma *KBasic-4*[*PLM*]:
[□(*φ* ≡ *ψ*) ≡ (□(*φ* → *ψ*) & □(*ψ* → *φ*)) *in v*]
apply (*rule* ≡*I*)
unfolding *equiv-def* **using** *KBasic-3 PLM.CP* ≡*E*(*1*)
apply *blast*
using *KBasic-3 PLM.CP* ≡*E*(*2*)
by *blast*
lemma *KBasic-5*[*PLM*]:
[(□(*φ* → *ψ*) & □(*ψ* → *φ*)) → (□*φ* ≡ □*ψ*) *in v*]
by (*metis qml-1*[*axiom-instance*] *CP* &*E* ≡*I vdash-properties-10*)
lemma *KBasic-6*[*PLM*]:
[□(*φ* ≡ *ψ*) → (□*φ* ≡ □*ψ*) *in v*]
using *KBasic-4 KBasic-5* **by** (*metis equiv-def ded-thm-cor-3* &*E*(*1*))
lemma [(□*φ* ≡ □*ψ*) → □(*φ* ≡ *ψ*) *in v*]
nitpick[*expect*=*genuine*, *user-axioms*, *card* = *1*, *card i* = *2*]
oops — countermodel as desired
lemma *KBasic-7*[*PLM*]:
[(□*φ* & □*ψ*) → □(*φ* ≡ *ψ*) *in v*]
proof (*rule CP*)
assume [□*φ* & □*ψ in v*]
hence [□(*ψ* → *φ*) *in v*] ∧ [□(*φ* → *ψ*) *in v*]
using &*E KBasic-1 vdash-properties-10* **by** *blast*
thus [□(*φ* ≡ *ψ*) *in v*]
using *KBasic-4* ≡*E*(*2*) *intro-elim-1* **by** *blast*
qed

lemma *KBasic-8*[*PLM*]:
[□(*φ* & *ψ*) → □(*φ* ≡ *ψ*) *in v*]
using *KBasic-7 KBasic-3*
by (*metis equiv-def PLM.ded-thm-cor-3* &*E*(*1*))
lemma *KBasic-9*[*PLM*]:
[□((¬*φ*) & (¬*ψ*)) → □(*φ* ≡ *ψ*) *in v*]
proof (*rule CP*)
assume [□((¬*φ*) & (¬*ψ*)) *in v*]
hence [□((¬*φ*) ≡ (¬*ψ*)) *in v*]
using *KBasic-8 vdash-properties-10* **by** *blast*
moreover have ⋀*v*.[((¬*φ*) ≡ (¬*ψ*)) → (*φ* ≡ *ψ*) *in v*]

using *CP* ≡*E*(*2*) *oth-class-taut-5-d* **by** *blast*
ultimately show [□(*φ* ≡ *ψ*) *in v*]
using *RM-1 PLM.vdash-properties-10* **by** *blast*
qed

lemma *rule-sub-lem-1-a*[*PLM*]:
[□(*ψ* ≡ *χ*) *in v*] ⟹ [(¬*ψ*) ≡ (¬*χ*) *in v*]
using *qml-2*[*axiom-instance*] ≡*E*(*1*) *oth-class-taut-5-d*
vdash-properties-10
by *blast*
lemma *rule-sub-lem-1-b*[*PLM*]:
[□(*ψ* ≡ *χ*) *in v*] ⟹ [(*ψ* → Θ) ≡ (*χ* → Θ) *in v*]
by (*metis equiv-def contraposition-1 CP* &*E*(*2*) ≡*I*
≡*E*(*1*) *rule-sub-lem-1-a*)
lemma *rule-sub-lem-1-c*[*PLM*]:
[□(*ψ* ≡ *χ*) *in v*] ⟹ [(Θ → *ψ*) ≡ (Θ → *χ*) *in v*]
by (*metis CP* ≡*I* ≡*E*(*3*) ≡*E*(*4*) ¬¬*I*
¬¬*E rule-sub-lem-1-a*)
lemma *rule-sub-lem-1-d*[*PLM*]:
(⋀*x*.[□(*ψ x* ≡ *χ x*) *in v*]) ⟹ [(∀ *α*. *ψ α*) ≡ (∀ *α*. *χ α*) *in v*]
by (*metis equiv-def* ∀ *I CP* &*E* ≡*I raa-cor-1*
vdash-properties-10 rule-sub-lem-1-a ∀ *E*)
lemma *rule-sub-lem-1-e*[*PLM*]:
[□(*ψ* ≡ *χ*) *in v*] ⟹ [𝒜*ψ* ≡ 𝒜*χ in v*]
using *Act-Basic-5* ≡*E*(*1*) *nec-imp-act*
vdash-properties-10
by *blast*
lemma *rule-sub-lem-1-f*[*PLM*]:
[□(*ψ* ≡ *χ*) *in v*] ⟹ [□*ψ* ≡ □*χ in v*]
using *KBasic-6* ≡*I* ≡*E*(*1*) *vdash-properties-9*
by *blast*

named-theorems *Substable-intros*

definition *Substable* :: (*'a*⇒*'a*⇒*bool*)⇒(*'a*⇒o) ⇒ *bool*
where *Substable* ≡ (λ *cond φ* . ∀ *ψ χ v* . (*cond ψ χ*) ⟶ [*φ ψ* ≡ *φ χ in v*])

lemma *Substable-intro-const*[*Substable-intros*]:
Substable cond (λ *φ* . Θ)
unfolding *Substable-def* **using** *oth-class-taut-4-a* **by** *blast*

lemma *Substable-intro-not*[*Substable-intros*]:
assumes *Substable cond ψ*
shows *Substable cond* (λ *φ* . ¬(*ψ φ*))
using *assms* **unfolding** *Substable-def*
using *rule-sub-lem-1-a RN-2* ≡*E oth-class-taut-5-d* **by** *metis*
lemma *Substable-intro-impl*[*Substable-intros*]:
assumes *Substable cond ψ*
and *Substable cond χ*
shows *Substable cond* (λ *φ* . *ψ φ* → *χ φ*)
using *assms* **unfolding** *Substable-def*
by (*metis* ≡*I CP intro-elim-6-a intro-elim-6-b*)
lemma *Substable-intro-box*[*Substable-intros*]:
assumes *Substable cond ψ*
shows *Substable cond* (λ *φ* . □(*ψ φ*))
using *assms* **unfolding** *Substable-def*
using *rule-sub-lem-1-f RN* **by** *meson*

lemma *Substable-intro-actual*[*Substable-intros*]:
assumes *Substable cond* ψ
shows *Substable cond* $(\lambda\ \varphi\ .\ \mathcal{A}(\psi\ \varphi))$
using *assms* **unfolding** *Substable-def*
using *rule-sub-lem-1-e RN* **by** *meson*
lemma *Substable-intro-all*[*Substable-intros*]:
assumes $\forall\ x\ .$ *Substable cond* $(\psi\ x)$
shows *Substable cond* $(\lambda\ \varphi\ .\ \forall\ x\ .\ \psi\ x\ \varphi)$
using *assms* **unfolding** *Substable-def*
by (*simp add*: *RN rule-sub-lem-1-d*)

named-theorems *Substable-Cond-defs*
end

class *Substable* =
fixes *Substable-Cond* :: $'a \Rightarrow 'a \Rightarrow$ *bool*
assumes *rule-sub-nec*:
$\bigwedge\ \varphi\ \psi\ \chi\ \Theta\ v\ .\ [\![PLM.Substable\ Substable\text{-}Cond\ \varphi;\ Substable\text{-}Cond\ \psi\ \chi]\!]$
$\Longrightarrow \Theta\ [\varphi\ \psi\ in\ v] \Longrightarrow \Theta\ [\varphi\ \chi\ in\ v]$

instantiation o :: *Substable*
begin
definition *Substable-Cond*-o **where** [*PLM.Substable-Cond-defs*]:
Substable-Cond-o $\equiv \lambda\ \varphi\ \psi\ .\ \forall\ v\ .\ [\varphi \equiv \psi\ in\ v]$
instance proof
interpret *PLM* **.**
fix φ :: o $\Rightarrow$ o **and** $\psi\ \chi$ :: o **and** Θ :: *bool* $\Rightarrow$ *bool* **and** v::*i*
assume *Substable Substable-Cond* φ
moreover assume *Substable-Cond* $\psi\ \chi$
ultimately have $[\varphi\ \psi \equiv \varphi\ \chi\ in\ v]$
unfolding *Substable-def* **by** *blast*
hence $[\varphi\ \psi\ in\ v] = [\varphi\ \chi\ in\ v]$ **using** $\equiv E$ **by** *blast*
moreover assume $\Theta\ [\varphi\ \psi\ in\ v]$
ultimately show $\Theta\ [\varphi\ \chi\ in\ v]$ **by** *simp*
qed
end

instantiation *fun* :: (*type*, *Substable*) *Substable*
begin
definition *Substable-Cond-fun* **where** [*PLM.Substable-Cond-defs*]:
Substable-Cond-fun $\equiv \lambda\ \varphi\ \psi\ .\ \forall\ x\ .$ *Substable-Cond* $(\varphi\ x)\ (\psi\ x)$
instance proof
interpret *PLM* **.**
fix φ:: $('a \Rightarrow 'b) \Rightarrow$ o **and** $\psi\ \chi :: 'a \Rightarrow 'b$ **and** $\Theta\ v$
assume *Substable Substable-Cond* φ
moreover assume *Substable-Cond* $\psi\ \chi$
ultimately have $[\varphi\ \psi \equiv \varphi\ \chi\ in\ v]$
unfolding *Substable-def* **by** *blast*
hence $[\varphi\ \psi\ in\ v] = [\varphi\ \chi\ in\ v]$ **using** $\equiv E$ **by** *blast*
moreover assume $\Theta\ [\varphi\ \psi\ in\ v]$
ultimately show $\Theta\ [\varphi\ \chi\ in\ v]$ **by** *simp*
qed
end

context *PLM*
begin

lemma *Substable-intro-equiv*[*Substable-intros*]:

assumes *Substable cond* ψ
and *Substable cond* χ
shows *Substable cond* $(\lambda\ \varphi\ .\ \psi\ \varphi \equiv \chi\ \varphi)$
unfolding *conn-defs* **by** (*simp add*: *assms Substable-intros*)
lemma *Substable-intro-conj*[*Substable-intros*]:
assumes *Substable cond* ψ
and *Substable cond* χ
shows *Substable cond* $(\lambda\ \varphi\ .\ \psi\ \varphi\ \&\ \chi\ \varphi)$
unfolding *conn-defs* **by** (*simp add*: *assms Substable-intros*)
lemma *Substable-intro-disj*[*Substable-intros*]:
assumes *Substable cond* ψ
and *Substable cond* χ
shows *Substable cond* $(\lambda\ \varphi\ .\ \psi\ \varphi \vee \chi\ \varphi)$
unfolding *conn-defs* **by** (*simp add*: *assms Substable-intros*)
lemma *Substable-intro-diamond*[*Substable-intros*]:
assumes *Substable cond* ψ
shows *Substable cond* $(\lambda\ \varphi\ .\ \Diamond(\psi\ \varphi))$
unfolding *conn-defs* **by** (*simp add*: *assms Substable-intros*)
lemma *Substable-intro-exist*[*Substable-intros*]:
assumes $\forall\ x\ .$ *Substable cond* $(\psi\ x)$
shows *Substable cond* $(\lambda\ \varphi\ .\ \exists\ x\ .\ \psi\ x\ \varphi)$
unfolding *conn-defs* **by** (*simp add*: *assms Substable-intros*)

lemma *Substable-intro-id-*o[*Substable-intros*]:
Substable Substable-Cond $(\lambda\ \varphi\ .\ \varphi)$
unfolding *Substable-def Substable-Cond-*o*-def* **by** *blast*
lemma *Substable-intro-id-fun*[*Substable-intros*]:
assumes *Substable Substable-Cond* ψ
shows *Substable Substable-Cond* $(\lambda\ \varphi\ .\ \psi\ (\varphi\ x))$
using *assms* **unfolding** *Substable-def Substable-Cond-fun-def*
by *blast*

method *PLM-subst-method* **for** ψ::$'a$::*Substable* **and** χ::$'a$::*Substable* =
(*match* **conclusion in** Θ $[\varphi\ \chi\ in\ v]$ **for** Θ **and** φ **and** v $\Rightarrow$
‹(*rule rule-sub-nec*[*where* Θ=Θ *and* χ=χ *and* ψ=ψ *and* φ=φ *and* v=v],
((*fast intro*: *Substable-intros*, ((*assumption*)+)*?*)+; *fail*),
unfold Substable-Cond-defs)›)

method *PLM-autosubst* =
(*match* **premises in** $\bigwedge v\ .\ [\psi \equiv \chi\ in\ v]$ **for** ψ **and** χ $\Rightarrow$
‹ *match conclusion in* Θ $[\varphi\ \chi\ in\ v]$ *for* Θ φ *and* v $\Rightarrow$
‹(*rule rule-sub-nec*[*where* Θ=Θ *and* χ=χ *and* ψ=ψ *and* φ=φ *and* v=v],
((*fast intro*: *Substable-intros*, ((*assumption*)+)*?*)+; *fail*),
unfold Substable-Cond-defs)› ›)

method *PLM-autosubst1* =
(*match* **premises in** $\bigwedge v\ x\ .\ [\psi\ x \equiv \chi\ x\ in\ v]$
for ψ::$'a$::*type*$\Rightarrow$o **and** χ::$'a\Rightarrow$o $\Rightarrow$
‹ *match conclusion in* Θ $[\varphi\ \chi\ in\ v]$ *for* Θ φ *and* v $\Rightarrow$
‹(*rule rule-sub-nec*[*where* Θ=Θ *and* χ=χ *and* ψ=ψ *and* φ=φ *and* v=v],
((*fast intro*: *Substable-intros*, ((*assumption*)+)*?*)+; *fail*),
unfold Substable-Cond-defs)› ›)

method *PLM-autosubst2* =
(*match* **premises in** $\bigwedge v\ x\ y\ .\ [\psi\ x\ y \equiv \chi\ x\ y\ in\ v]$
for ψ::$'a$::*type*$\Rightarrow{'a}\Rightarrow$o **and** χ::$'a$::*type*$\Rightarrow{'a}\Rightarrow$o $\Rightarrow$
‹ *match conclusion in* Θ $[\varphi\ \chi\ in\ v]$ *for* Θ φ *and* v $\Rightarrow$
‹(*rule rule-sub-nec*[*where* Θ=Θ *and* χ=χ *and* ψ=ψ *and* φ=φ *and* v=v],

((*fast intro*: *Substable-intros*, ((*assumption*)+)*?*)+; *fail*),
unfold Substable-Cond-defs)›)

method *PLM-subst-goal-method* **for** φ::'*a*::*Substable*⇒o **and** ψ::'*a* =
(*match* **conclusion in** Θ [φ χ *in v*] **for** Θ **and** χ **and** *v* ⇒
‹(*rule rule-sub-nec*[*where* Θ=Θ *and* χ=χ *and* ψ=ψ *and* φ=φ *and v*=*v*],
((*fast intro*: *Substable-intros*, ((*assumption*)+)*?*)+; *fail*),
unfold Substable-Cond-defs)›)

lemma *rule-sub-nec*[*PLM*]:
assumes *Substable Substable-Cond* φ
shows (⋀*v*.[(ψ ≡ χ) *in v*]) ⟹ Θ [φ ψ *in v*] ⟹ Θ [φ χ *in v*]
proof −
assume (⋀*v*.[(ψ ≡ χ) *in v*])
hence [φ ψ *in v*] = [φ χ *in v*]
using *assms RN* **unfolding** *Substable-def Substable-Cond-defs*
using ≡*I CP* ≡*E*(*1*) ≡*E*(*2*) **by** *meson*
thus Θ [φ ψ *in v*] ⟹ Θ [φ χ *in v*] **by** *auto*
qed

lemma *rule-sub-nec1*[*PLM*]:
assumes *Substable Substable-Cond* φ
shows (⋀*v x* .[(ψ *x* ≡ χ *x*) *in v*]) ⟹ Θ [φ ψ *in v*] ⟹ Θ [φ χ *in v*]
proof −
assume (⋀*v x*.[(ψ *x* ≡ χ *x*) *in v*])
hence [φ ψ *in v*] = [φ χ *in v*]
using *assms RN* **unfolding** *Substable-def Substable-Cond-defs*
using ≡*I CP* ≡*E*(*1*) ≡*E*(*2*) **by** *metis*
thus Θ [φ ψ *in v*] ⟹ Θ [φ χ *in v*] **by** *auto*
qed

lemma *rule-sub-nec2*[*PLM*]:
assumes *Substable Substable-Cond* φ
shows (⋀*v x y* .[ψ *x y* ≡ χ *x y in v*]) ⟹ Θ [φ ψ *in v*] ⟹ Θ [φ χ *in v*]
proof −
assume (⋀*v x y* .[ψ *x y* ≡ χ *x y in v*])
hence [φ ψ *in v*] = [φ χ *in v*]
using *assms RN* **unfolding** *Substable-def Substable-Cond-defs*
using ≡*I CP* ≡*E*(*1*) ≡*E*(*2*) **by** *metis*
thus Θ [φ ψ *in v*] ⟹ Θ [φ χ *in v*] **by** *auto*
qed

lemma *rule-sub-remark-1-autosubst*:
assumes (⋀*v*.[(|*A*!,*x*|) ≡ (¬(◇(|*E*!,*x*|))) *in v*])
and [¬(|*A*!,*x*|) *in v*]
shows[¬¬◇(|*E*!,*x*|) *in v*]
apply (*insert assms*) **apply** *PLM-autosubst* **by** *auto*

lemma *rule-sub-remark-1*:
assumes (⋀*v*.[(|*A*!,*x*|) ≡ (¬(◇(|*E*!,*x*|))) *in v*])
and [¬(|*A*!,*x*|) *in v*]
shows[¬¬◇(|*E*!,*x*|) *in v*]
apply (*PLM-subst-method* (|*A*!,*x*|) (¬(◇(|*E*!,*x*|))))
apply (*simp add*: *assms*(*1*))
by (*simp add*: *assms*(*2*))

lemma *rule-sub-remark-2*:
assumes $(\bigwedge v.[(\!|R,x,y|\!) \equiv ((\!|R,x,y|\!) \;\&\; ((\!|Q,a|\!) \vee (\neg(\!|Q,a|\!))))\ in\ v])$
and $[p \rightarrow (\!|R,x,y|\!)\ in\ v]$
shows$[p \rightarrow ((\!|R,x,y|\!) \;\&\; ((\!|Q,a|\!) \vee (\neg(\!|Q,a|\!))))\ \ in\ v]$
apply (*insert assms*) **apply** *PLM-autosubst* **by** *auto*

lemma *rule-sub-remark-3-autosubst*:
assumes $(\bigwedge v\ x.[(\!|A!,x^P|\!) \equiv (\neg(\Diamond(\!|E!,x^P|\!)))\ in\ v])$
and $[\exists\ x\ .\ (\!|A!,x^P|\!)\ in\ v]$
shows$[\exists\ x\ .\ (\neg(\Diamond(\!|E!,x^P|\!)))\ \ in\ v]$
apply (*insert assms*) **apply** *PLM-autosubst1* **by** *auto*

lemma *rule-sub-remark-3*:
assumes $(\bigwedge v\ x.[(\!|A!,x^P|\!) \equiv (\neg(\Diamond(\!|E!,x^P|\!)))\ in\ v])$
and $[\exists\ x\ .\ (\!|A!,x^P|\!)\ in\ v]$
shows $[\exists\ x\ .\ (\neg(\Diamond(\!|E!,x^P|\!)))\ \ in\ v]$
apply (*PLM-subst-method* $\lambda x\ .\ (\!|A!,x^P|\!)\ \lambda x\ .\ (\neg(\Diamond(\!|E!,x^P|\!)))$)
apply (*simp add*: *assms(1)*)
by (*simp add*: *assms(2)*)

lemma *rule-sub-remark-4*:
assumes $\bigwedge v\ x.[(\neg(\neg(\!|P,x^P|\!))) \equiv (\!|P,x^P|\!)\ in\ v]$
and $[\boldsymbol{\mathcal{A}}(\neg(\neg(\!|P,x^P|\!)))\ in\ v]$
shows $[\boldsymbol{\mathcal{A}}(\!|P,x^P|\!)\ in\ v]$
apply (*insert assms*) **apply** *PLM-autosubst1* **by** *auto*

lemma *rule-sub-remark-5*:
assumes $\bigwedge v.[(\varphi \rightarrow \psi) \equiv ((\neg\psi) \rightarrow (\neg\varphi))\ in\ v]$
and $[\Box(\varphi \rightarrow \psi)\ in\ v]$
shows $[\Box((\neg\psi) \rightarrow (\neg\varphi))\ in\ v]$
apply (*insert assms*) **apply** *PLM-autosubst* **by** *auto*

lemma *rule-sub-remark-6*:
assumes $\bigwedge v.[\psi \equiv \chi\ in\ v]$
and $[\Box(\varphi \rightarrow \psi)\ in\ v]$
shows $[\Box(\varphi \rightarrow \chi)\ in\ v]$
apply (*insert assms*) **apply** *PLM-autosubst* **by** *auto*

lemma *rule-sub-remark-7*:
assumes $\bigwedge v.[\varphi \equiv (\neg(\neg\varphi))\ in\ v]$
and $[\Box(\varphi \rightarrow \varphi)\ in\ v]$
shows $[\Box((\neg(\neg\varphi)) \rightarrow \varphi)\ in\ v]$
apply (*insert assms*) **apply** *PLM-autosubst* **by** *auto*

lemma *rule-sub-remark-8*:
assumes $\bigwedge v.[\boldsymbol{\mathcal{A}}\varphi \equiv \varphi\ in\ v]$
and $[\Box(\boldsymbol{\mathcal{A}}\varphi)\ in\ v]$
shows $[\Box(\varphi)\ in\ v]$
apply (*insert assms*) **apply** *PLM-autosubst* **by** *auto*

lemma *rule-sub-remark-9*:
assumes $\bigwedge v.[(\!|P,a|\!) \equiv ((\!|P,a|\!) \;\&\; ((\!|Q,b|\!) \vee (\neg(\!|Q,b|\!))))\ in\ v]$
and $[(\!|P,a|\!) = (\!|P,a|\!)\ in\ v]$
shows $[(\!|P,a|\!) = ((\!|P,a|\!) \;\&\; ((\!|Q,b|\!) \vee (\neg(\!|Q,b|\!))))\ in\ v]$
unfolding *identity-defs* **apply** (*insert assms*)
apply *PLM-autosubst* **oops** — no match as desired

— *dr-alphabetic-rules* implicitly holds

— *dr-alphabetic-thm* implicitly holds

lemma *KBasic2-1*[*PLM*]:
[$\Box\varphi \equiv \Box(\neg(\neg\varphi))$ *in v*]
apply (*PLM-subst-method* φ $(\neg(\neg\varphi))$)
by *PLM-solver+*

lemma *KBasic2-2*[*PLM*]:
[$(\neg(\Box\varphi)) \equiv \Diamond(\neg\varphi)$ *in v*]
unfolding *diamond-def*
apply (*PLM-subst-method* φ $\neg(\neg\varphi)$)
by *PLM-solver+*

lemma *KBasic2-3*[*PLM*]:
[$\Box\varphi \equiv (\neg(\Diamond(\neg\varphi)))$ *in v*]
unfolding *diamond-def*
apply (*PLM-subst-method* φ $\neg(\neg\varphi)$)
apply *PLM-solver*
by (*simp add*: *oth-class-taut-4-b*)
lemmas *Df*$\Box$ = *KBasic2-3*

lemma *KBasic2-4*[*PLM*]:
[$\Box(\neg(\varphi)) \equiv (\neg(\Diamond\varphi))$ *in v*]
unfolding *diamond-def*
by (*simp add*: *oth-class-taut-4-b*)

lemma *KBasic2-5*[*PLM*]:
[$\Box(\varphi \rightarrow \psi) \rightarrow (\Diamond\varphi \rightarrow \Diamond\psi)$ *in v*]
by (*simp only*: *CP RM-2-b*)
lemmas *K*$\Diamond$ = *KBasic2-5*

lemma *KBasic2-6*[*PLM*]:
[$\Diamond(\varphi \vee \psi) \equiv (\Diamond\varphi \vee \Diamond\psi)$ *in v*]
proof −
have [$\Box((\neg\varphi)$ & $(\neg\psi)) \equiv (\Box(\neg\varphi)$ & $\Box(\neg\psi))$ *in v*]
using *KBasic-3* **by** *blast*
hence [$(\neg(\Diamond(\neg((\neg\varphi)$ & $(\neg\psi))))) \equiv (\Box(\neg\varphi)$ & $\Box(\neg\psi))$ *in v*]
using *Df*$\Box$ **by** (*rule* $\equiv E(6)$)
hence [$(\neg(\Diamond(\neg((\neg\varphi)$ & $(\neg\psi))))) \equiv ((\neg(\Diamond\varphi))$ & $(\neg(\Diamond\psi)))$ *in v*]
apply − **apply** (*PLM-subst-method* $\Box(\neg\varphi)$ $\neg(\Diamond\varphi)$)
apply (*simp add*: *KBasic2-4*)
apply (*PLM-subst-method* $\Box(\neg\psi)$ $\neg(\Diamond\psi)$)
apply (*simp add*: *KBasic2-4*)
unfolding *diamond-def* **by** *assumption*
hence [$(\neg(\Diamond(\varphi \vee \psi))) \equiv ((\neg(\Diamond\varphi))$ & $(\neg(\Diamond\psi)))$ *in v*]
apply − **apply** (*PLM-subst-method* $\neg((\neg\varphi)$ & $(\neg\psi))$ $\varphi \vee \psi$)
using *oth-class-taut-6-b*[*equiv-sym*] **by** *auto*
hence [$(\neg(\neg(\Diamond(\varphi \vee \psi)))) \equiv (\neg((\neg(\Diamond\varphi))\&(\neg(\Diamond\psi))))$ *in v*]
by (*rule oth-class-taut-5-d*[*equiv-lr*])
hence [$\Diamond(\varphi \vee \psi) \equiv (\neg((\neg(\Diamond\varphi))$ & $(\neg(\Diamond\psi))))$ *in v*]
apply − **apply** (*PLM-subst-method* $\neg(\neg(\Diamond(\varphi \vee \psi)))$ $\Diamond(\varphi \vee \psi)$)
using *oth-class-taut-4-b*[*equiv-sym*] **by** *auto*
thus *?thesis*
apply − **apply** (*PLM-subst-method* $\neg((\neg(\Diamond\varphi))$ & $(\neg(\Diamond\psi)))$ $(\Diamond\varphi) \vee (\Diamond\psi)$)
using *oth-class-taut-6-b*[*equiv-sym*] **by** *auto*
qed

lemma *KBasic2-7*[*PLM*]:

```
  [(□φ ∨ □ψ) → □(φ ∨ ψ) in v]
  proof −
    have ⋀ v . [φ → (φ ∨ ψ) in v]
      by (metis contraposition-1 contraposition-2 useful-tautologies-3 disj-def)
    hence [□φ → □(φ ∨ ψ) in v] using RM-1 by auto
    moreover {
       have ⋀ v . [ψ → (φ ∨ ψ) in v]
         by (simp only: pl-1[axiom-instance] disj-def)
       hence [□ψ → □(φ ∨ ψ) in v]
         using RM-1 by auto
    }
    ultimately show ?thesis
      using oth-class-taut-10-d vdash-properties-10 by blast
  qed

lemma KBasic2-8[PLM]:
  [◇(φ & ψ) → (◇φ & ◇ψ) in v]
  by (metis CP RM-2 &I oth-class-taut-9-a
          oth-class-taut-9-b vdash-properties-10)

lemma KBasic2-9[PLM]:
  [◇(φ → ψ) ≡ (□φ → ◇ψ) in v]
  apply (PLM-subst-method (¬(□φ)) ∨ (◇ψ) □φ → ◇ψ)
   using oth-class-taut-5-k[equiv-sym] apply simp
  apply (PLM-subst-method (¬φ) ∨ ψ φ → ψ)
   using oth-class-taut-5-k[equiv-sym] apply simp
  apply (PLM-subst-method ◇(¬φ) ¬(□φ))
   using KBasic2-2[equiv-sym] apply simp
  using KBasic2-6 .

lemma KBasic2-10[PLM]:
  [◇(□φ) ≡ (¬(□◇(¬φ))) in v]
  unfolding diamond-def apply (PLM-subst-method φ ¬¬φ)
  using oth-class-taut-4-b oth-class-taut-4-a by auto

lemma KBasic2-11[PLM]:
  [◇◇φ ≡ (¬(□□(¬φ))) in v]
  unfolding diamond-def
  apply (PLM-subst-method □(¬φ) ¬(¬(□(¬φ))))
  using oth-class-taut-4-b oth-class-taut-4-a by auto

lemma KBasic2-12[PLM]: [□(φ ∨ ψ) → (□φ ∨ ◇ψ) in v]
  proof −
    have [□(ψ ∨ φ) → (□(¬ψ) → □φ) in v]
      using CP RM-1-b ∨E(2) by blast
    hence [□(ψ ∨ φ) → (◇ψ ∨ □φ) in v]
      unfolding diamond-def disj-def
      by (meson CP ¬¬E vdash-properties-6)
    thus ?thesis apply −
      apply (PLM-subst-method (◇ψ ∨ □φ) (□φ ∨ ◇ψ))
       apply (simp add: PLM.oth-class-taut-3-e)
      apply (PLM-subst-method (ψ ∨ φ) (φ ∨ ψ))
       apply (simp add: PLM.oth-class-taut-3-e)
      by assumption
  qed

lemma TBasic[PLM]:
  [φ → ◇φ in v]
```

```
  unfolding diamond-def
  apply (subst contraposition-1)
  apply (PLM-subst-method □¬φ ¬¬□¬φ)
   apply (simp add: PLM.oth-class-taut-4-b)
  using qml-2[where φ=¬φ, axiom-instance]
  by simp
lemmas T◇ = TBasic

lemma S5Basic-1[PLM]:
  [◇□φ → □φ in v]
  proof (rule CP)
    assume [◇□φ in v]
    hence [¬□◇¬φ in v]
      using KBasic2-10[equiv-lr] by simp
    moreover have [◇(¬φ) → □◇(¬φ) in v]
      by (simp add: qml-3[axiom-instance])
    ultimately have [¬◇¬φ in v]
      by (simp add: PLM.modus-tollens-1)
    thus [□φ in v]
      unfolding diamond-def apply -
      apply (PLM-subst-method ¬¬φ φ)
       using oth-class-taut-4-b[equiv-sym] apply simp
      unfolding diamond-def using oth-class-taut-4-b[equiv-rl]
      by simp
  qed
lemmas 5◇ = S5Basic-1

lemma S5Basic-2[PLM]:
  [□φ ≡ ◇□φ in v]
  using 5◇ T◇ ≡I by blast

lemma S5Basic-3[PLM]:
  [◇φ ≡ □◇φ in v]
  using qml-3[axiom-instance] qml-2[axiom-instance] ≡I by blast

lemma S5Basic-4[PLM]:
  [φ → □◇φ in v]
  using T◇[deduction, THEN S5Basic-3[equiv-lr]]
  by (rule CP)

lemma S5Basic-5[PLM]:
  [◇□φ → φ in v]
  using S5Basic-2[equiv-rl, THEN qml-2[axiom-instance, deduction]]
  by (rule CP)
lemmas B◇ = S5Basic-5

lemma S5Basic-6[PLM]:
  [□φ → □□φ in v]
  using S5Basic-4[deduction] RM-1[OF S5Basic-1, deduction] CP by auto
lemmas 4□ = S5Basic-6

lemma S5Basic-7[PLM]:
  [□φ ≡ □□φ in v]
  using 4□ qml-2[axiom-instance] by (rule ≡I)

lemma S5Basic-8[PLM]:
  [◇◇φ → ◇φ in v]
  using S5Basic-6[where φ=¬φ, THEN contraposition-1[THEN iffD1], deduction]
```

KBasic2-11[*equiv-lr*] *CP* **unfolding** *diamond-def* **by** *auto*
lemmas *4◇ = S5Basic-8*

lemma *S5Basic-9*[*PLM*]:
[◇◇φ ≡ ◇φ *in v*]
using *4◇ T◇* **by** (*rule* ≡*I*)

lemma *S5Basic-10*[*PLM*]:
[□(φ ∨ □ψ) ≡ (□φ ∨ □ψ) *in v*]
apply (*rule* ≡*I*)
apply (*PLM-subst-goal-method* λ χ . □(φ ∨ □ψ) → (□φ ∨ χ) ◇□ψ)
using *S5Basic-2*[*equiv-sym*] **apply** *simp*
using *KBasic2-12* **apply** *assumption*
apply (*PLM-subst-goal-method* λ χ .(□φ ∨ χ) → □(φ ∨ □ψ) □□ψ)
using *S5Basic-7*[*equiv-sym*] **apply** *simp*
using *KBasic2-7* **by** *auto*

lemma *S5Basic-11*[*PLM*]:
[□(φ ∨ ◇ψ) ≡ (□φ ∨ ◇ψ) *in v*]
apply (*rule* ≡*I*)
apply (*PLM-subst-goal-method* λ χ . □(φ ∨ ◇ψ) → (□φ ∨ χ) ◇◇ψ)
using *S5Basic-9* **apply** *simp*
using *KBasic2-12* **apply** *assumption*
apply (*PLM-subst-goal-method* λ χ .(□φ ∨ χ) → □(φ ∨ ◇ψ) □◇ψ)
using *S5Basic-3*[*equiv-sym*] **apply** *simp*
using *KBasic2-7* **by** *assumption*

lemma *S5Basic-12*[*PLM*]:
[◇(φ & ◇ψ) ≡ (◇φ & ◇ψ) *in v*]
proof −
have [□((¬φ) ∨ □(¬ψ)) ≡ (□(¬φ) ∨ □(¬ψ)) *in v*]
using *S5Basic-10* **by** *auto*
hence *1*: [(¬□((¬φ) ∨ □(¬ψ))) ≡ ¬(□(¬φ) ∨ □(¬ψ)) *in v*]
using *oth-class-taut-5-d*[*equiv-lr*] **by** *auto*
have *2*: [(◇(¬((¬φ) ∨ (¬(◇ψ))))) ≡ (¬((¬(◇φ)) ∨ (¬(◇ψ)))) *in v*]
apply (*PLM-subst-method* □¬ψ ¬◇ψ)
using *KBasic2-4* **apply** *simp*
apply (*PLM-subst-method* □¬φ ¬◇φ)
using *KBasic2-4* **apply** *simp*
apply (*PLM-subst-method* (¬□((¬φ) ∨ □(¬ψ))) (◇(¬((¬φ) ∨ (□(¬ψ))))))
unfolding *diamond-def*
apply (*simp add*: *RN oth-class-taut-4-b rule-sub-lem-1-a rule-sub-lem-1-f*)
using *1* **by** *assumption*
show *?thesis*
apply (*PLM-subst-method* ¬((¬φ) ∨ (¬◇ψ)) φ & ◇ψ)
using *oth-class-taut-6-a*[*equiv-sym*] **apply** *simp*
apply (*PLM-subst-method* ¬((¬(◇φ)) ∨ (¬◇ψ)) ◇φ & ◇ψ)
using *oth-class-taut-6-a*[*equiv-sym*] **apply** *simp*
using *2* **by** *assumption*
qed

lemma *S5Basic-13*[*PLM*]:
[◇(φ & (□ψ)) ≡ (◇φ & (□ψ)) *in v*]
apply (*PLM-subst-method* ◇□ψ □ψ)
using *S5Basic-2*[*equiv-sym*] **apply** *simp*
using *S5Basic-12* **by** *simp*

lemma *S5Basic-14*[*PLM*]:

```
  [□(φ → (□ψ)) ≡ □(◇φ → ψ) in v]
  proof (rule ≡I; rule CP)
    assume [□(φ → □ψ) in v]
    moreover {
      have ⋀v.[□(φ → □ψ) → (◇φ → ψ) in v]
        proof (rule CP)
          fix v
          assume [□(φ → □ψ) in v]
          hence [◇φ → ◇□ψ in v]
            using K◇[deduction] by auto
          thus [◇φ → ψ in v]
            using B◇ ded-thm-cor-3 by blast
        qed
      hence [□(□(φ → □ψ) → (◇φ → ψ)) in v]
        by (rule RN)
      hence [□(□(φ → □ψ)) → □((◇φ → ψ)) in v]
        using qml-1[axiom-instance, deduction] by auto
    }
    ultimately show [□(◇φ → ψ) in v]
      using S5Basic-6 CP vdash-properties-10 by meson
  next
    assume [□(◇φ → ψ) in v]
    moreover {
      fix v
      {
        assume [□(◇φ → ψ) in v]
        hence 1: [□◇φ → □ψ in v]
          using qml-1[axiom-instance, deduction] by auto
        assume [φ in v]
        hence [□◇φ in v]
          using S5Basic-4[deduction] by auto
        hence [□ψ in v]
          using 1[deduction] by auto
      }
      hence [□(◇φ → ψ) in v] ⟹ [φ → □ψ in v]
        using CP by auto
    }
    ultimately show [□(φ → □ψ) in v]
      using S5Basic-6 RN-2 vdash-properties-10 by blast
  qed

lemma sc-eq-box-box-1[PLM]:
  [□(φ → □φ) → (◇φ ≡ □φ) in v]
  proof(rule CP)
    assume 1: [□(φ → □φ) in v]
    hence [□(◇φ → φ) in v]
      using S5Basic-14[equiv-lr] by auto
    hence [◇φ → φ in v]
      using qml-2[axiom-instance, deduction] by auto
    moreover from 1 have [φ → □φ in v]
      using qml-2[axiom-instance, deduction] by auto
    ultimately have [◇φ → □φ in v]
      using ded-thm-cor-3 by auto
    moreover have [□φ → ◇φ in v]
      using qml-2[axiom-instance] T◇
      by (rule ded-thm-cor-3)
    ultimately show [◇φ ≡ □φ in v]
      by (rule ≡I)
```

qed

lemma *sc-eq-box-box-2*[*PLM*]:
[□(φ → □φ) → ((¬□φ) ≡ (□(¬φ))) *in v*]
proof (*rule CP*)
assume [□(φ → □φ) *in v*]
hence [(¬□(¬φ)) ≡ □φ *in v*]
using *sc-eq-box-box-1*[*deduction*] **unfolding** *diamond-def* **by** *auto*
thus [((¬□φ) ≡ (□(¬φ))) *in v*]
by (*meson CP* ≡*I* ≡*E*(*3*)
≡*E*(*4*) ¬¬*I* ¬¬*E*)
qed

lemma *sc-eq-box-box-3*[*PLM*]:
[(□(φ → □φ) & □(ψ → □ψ)) → ((□φ ≡ □ψ) → □(φ ≡ ψ)) *in v*]
proof (*rule CP*)
assume *1*: [(□(φ → □φ) & □(ψ → □ψ)) *in v*]
{
assume [□φ ≡ □ψ *in v*]
hence [(□φ & □ψ) ∨ ((¬(□φ)) & (¬(□ψ))) *in v*]
using *oth-class-taut-5-i*[*equiv-lr*] **by** *auto*
moreover {
assume [□φ & □ψ *in v*]
hence [□(φ ≡ ψ) *in v*]
using *KBasic-7*[*deduction*] **by** *auto*
}
moreover {
assume [(¬(□φ)) & (¬(□ψ)) *in v*]
hence [□(¬φ) & □(¬ψ) *in v*]
using *1* &*E* &*I* *sc-eq-box-box-2*[*deduction*, *equiv-lr*]
by *metis*
hence [□((¬φ) & (¬ψ)) *in v*]
using *KBasic-3*[*equiv-rl*] **by** *auto*
hence [□(φ ≡ ψ) *in v*]
using *KBasic-9*[*deduction*] **by** *auto*
}
ultimately have [□(φ ≡ ψ) *in v*]
using *CP* ∨*E*(*1*) **by** *blast*
}
thus [□φ ≡ □ψ → □(φ ≡ ψ) *in v*]
using *CP* **by** *auto*
qed

lemma *derived-S5-rules-1-a*[*PLM*]:
assumes ⋀*v*. [χ *in v*] ⟹ [◇φ → ψ *in v*]
shows [□χ *in v*] ⟹ [φ → □ψ *in v*]
proof −
have [□χ *in v*] ⟹ [□◇φ → □ψ *in v*]
using *assms RM-1-b* **by** *metis*
thus [□χ *in v*] ⟹ [φ → □ψ *in v*]
using *S5Basic-4 vdash-properties-10 CP* **by** *metis*
qed

lemma *derived-S5-rules-1-b*[*PLM*]:
assumes ⋀*v*. [◇φ → ψ *in v*]
shows [φ → □ψ *in v*]
using *derived-S5-rules-1-a all-self-eq-1 assms* **by** *blast*

lemma *derived-S5-rules-2-a*[*PLM*]:
assumes $\bigwedge v.\ [\chi\ in\ v] \Longrightarrow [\varphi \rightarrow \Box\psi\ in\ v]$
shows $[\Box\chi\ in\ v] \Longrightarrow [\Diamond\varphi \rightarrow \psi\ in\ v]$
proof −
have $[\Box\chi\ in\ v] \Longrightarrow [\Diamond\varphi \rightarrow \Diamond\Box\psi\ in\ v]$
using *RM-2-b assms* **by** *metis*
thus $[\Box\chi\ in\ v] \Longrightarrow [\Diamond\varphi \rightarrow \psi\ in\ v]$
using *B*$\Diamond$ *vdash-properties-10 CP* **by** *metis*
qed

lemma *derived-S5-rules-2-b*[*PLM*]:
assumes $\bigwedge v.\ [\varphi \rightarrow \Box\psi\ in\ v]$
shows $[\Diamond\varphi \rightarrow \psi\ in\ v]$
using *assms derived-S5-rules-2-a all-self-eq-1* **by** *blast*

lemma *BFs-1*[*PLM*]: $[(\forall\,\alpha.\ \Box(\varphi\ \alpha)) \rightarrow \Box(\forall\,\alpha.\ \varphi\ \alpha)\ in\ v]$
proof (*rule derived-S5-rules-1-b*)
fix v
{
fix α
have $\bigwedge v.[(\forall\,\alpha\ .\ \Box(\varphi\ \alpha)) \rightarrow \Box(\varphi\ \alpha)\ in\ v]$
using *cqt-orig-1* **by** *metis*
hence $[\Diamond(\forall\,\alpha.\ \Box(\varphi\ \alpha)) \rightarrow \Diamond\Box(\varphi\ \alpha)\ in\ v]$
using *RM-2* **by** *metis*
moreover have $[\Diamond\Box(\varphi\ \alpha) \rightarrow (\varphi\ \alpha)\ in\ v]$
using *B*$\Diamond$ **by** *auto*
ultimately have $[\Diamond(\forall\,\alpha.\ \Box(\varphi\ \alpha)) \rightarrow (\varphi\ \alpha)\ in\ v]$
using *ded-thm-cor-3* **by** *auto*
}
hence $[\forall\ \alpha\ .\ \Diamond(\forall\,\alpha.\ \Box(\varphi\ \alpha)) \rightarrow (\varphi\ \alpha)\ in\ v]$
using $\forall I$ **by** *metis*
thus $[\Diamond(\forall\,\alpha.\ \Box(\varphi\ \alpha)) \rightarrow (\forall\,\alpha.\ \varphi\ \alpha)\ in\ v]$
using *cqt-orig-2*[*deduction*] **by** *auto*
qed
lemmas *BF* = *BFs-1*

lemma *BFs-2*[*PLM*]:
$[\Box(\forall\,\alpha.\ \varphi\ \alpha) \rightarrow (\forall\,\alpha.\ \Box(\varphi\ \alpha))\ in\ v]$
proof −
{
fix α
{
fix v
have $[(\forall\,\alpha.\ \varphi\ \alpha) \rightarrow \varphi\ \alpha\ in\ v]$ **using** *cqt-orig-1* **by** *metis*
}
hence $[\Box(\forall\,\alpha\ .\ \varphi\ \alpha) \rightarrow \Box(\varphi\ \alpha)\ in\ v]$ **using** *RM-1* **by** *auto*
}
hence $[\forall\,\alpha\ .\ \Box(\forall\,\alpha\ .\ \varphi\ \alpha) \rightarrow \Box(\varphi\ \alpha)\ in\ v]$ **using** $\forall I$ **by** *metis*
thus *?thesis* **using** *cqt-orig-2*[*deduction*] **by** *metis*
qed
lemmas *CBF* = *BFs-2*

lemma *BFs-3*[*PLM*]:
$[\Diamond(\exists\ \alpha.\ \varphi\ \alpha) \rightarrow (\exists\ \alpha\ .\ \Diamond(\varphi\ \alpha))\ in\ v]$
proof −
have $[(\forall\,\alpha.\ \Box(\neg(\varphi\ \alpha))) \rightarrow \Box(\forall\,\alpha.\ \neg(\varphi\ \alpha))\ in\ v]$
using *BF* **by** *metis*
hence *1*: $[(\neg(\Box(\forall\,\alpha.\ \neg(\varphi\ \alpha)))) \rightarrow (\neg(\forall\,\alpha.\ \Box(\neg(\varphi\ \alpha))))\ in\ v]$

```
    using contraposition-1 by simp
  have 2: [◇(¬(∀ α. ¬(φ α))) → (¬(∀ α. □(¬(φ α)))) in v]
    apply (PLM-subst-method ¬□(∀ α . ¬(φ α)) ◇(¬(∀ α. ¬(φ α))))
    using KBasic2-2 1 by simp+
  have [◇(¬(∀ α. ¬(φ α))) → (∃ α . ¬(□(¬(φ α)))) in v]
    apply (PLM-subst-method ¬(∀ α. □(¬(φ α))) ∃ α. ¬(□(¬(φ α))))
     using cqt-further-2 apply metis
    using 2 by metis
  thus ?thesis
    unfolding exists-def diamond-def by auto
qed
lemmas BF◇ = BFs-3

lemma BFs-4[PLM]:
[(∃ α . ◇(φ α)) → ◇(∃ α. φ α) in v]
proof −
  have 1: [□(∀ α . ¬(φ α)) → (∀ α. □(¬(φ α))) in v]
    using CBF by auto
  have 2: [(∃ α . (¬(□(¬(φ α))))) → (¬(□(∀ α. ¬(φ α)))) in v]
    apply (PLM-subst-method ¬(∀ α. □(¬(φ α))) (∃ α . (¬(□(¬(φ α))))))
     using cqt-further-2 apply blast
    using 1 using contraposition-1 by metis
  have [(∃ α . (¬(□(¬(φ α))))) → ◇(¬(∀ α . ¬(φ α))) in v]
    apply (PLM-subst-method ¬(□(∀ α. ¬(φ α))) ◇(¬(∀ α. ¬(φ α))))
     using KBasic2-2 apply blast
    using 2 by assumption
  thus ?thesis
    unfolding diamond-def exists-def by auto
qed
lemmas CBF◇ = BFs-4

lemma sign-S5-thm-1[PLM]:
[(∃ α. □(φ α)) → □(∃ α. φ α) in v]
proof (rule CP)
  assume [∃ α . □(φ α) in v]
  then obtain τ where [□(φ τ) in v]
    by (rule ∃E)
  moreover {
    fix v
    assume [φ τ in v]
    hence [∃ α . φ α in v]
      by (rule ∃I)
  }
  ultimately show [□(∃ α . φ α) in v]
    using RN-2 by blast
qed
lemmas Buridan = sign-S5-thm-1

lemma sign-S5-thm-2[PLM]:
[◇(∀ α . φ α) → (∀ α . ◇(φ α)) in v]
proof −
  {
    fix α
    {
      fix v
      have [(∀ α . φ α) → φ α in v]
        using cqt-orig-1 by metis
    }
```

```
    hence [◇(∀ α . φ α) → ◇(φ α) in v]
      using RM-2 by metis
  }
  hence [∀ α . ◇(∀ α . φ α) → ◇(φ α) in v]
    using ∀ I by metis
  thus ?thesis
    using cqt-orig-2[deduction] by metis
qed
lemmas Buridan◇ = sign-S5-thm-2

lemma sign-S5-thm-3[PLM]:
[◇(∃ α . φ α & ψ α) → ◇((∃ α . φ α) & (∃ α . ψ α)) in v]
by (simp only: RM-2 cqt-further-5)

lemma sign-S5-thm-4[PLM]:
[((□(∀ α. φ α → ψ α)) & (□(∀ α . ψ α → χ α))) → □(∀α. φ α → χ α) in v]
proof (rule CP)
  assume [□(∀α. φ α → ψ α) & □(∀α. ψ α → χ α) in v]
  hence [□((∀α. φ α → ψ α) & (∀α. ψ α → χ α)) in v]
    using KBasic-3[equiv-rl] by blast
  moreover {
    fix v
    assume [((∀α. φ α → ψ α) & (∀α. ψ α → χ α)) in v]
    hence [(∀ α . φ α → χ α) in v]
      using cqt-basic-9[deduction] by blast
  }
  ultimately show [□(∀α. φ α → χ α) in v]
    using RN-2 by blast
qed

lemma sign-S5-thm-5[PLM]:
[((□(∀α. φ α ≡ ψ α)) & (□(∀α. ψ α ≡ χ α))) → (□(∀α. φ α ≡ χ α)) in v]
proof (rule CP)
  assume [□(∀α. φ α ≡ ψ α) & □(∀α. ψ α ≡ χ α) in v]
  hence [□((∀α. φ α ≡ ψ α) & (∀α. ψ α ≡ χ α)) in v]
    using KBasic-3[equiv-rl] by blast
  moreover {
    fix v
    assume [((∀α. φ α ≡ ψ α) & (∀α. ψ α ≡ χ α)) in v]
    hence [(∀ α . φ α ≡ χ α) in v]
      using cqt-basic-10[deduction] by blast
  }
  ultimately show [□(∀α. φ α ≡ χ α) in v]
    using RN-2 by blast
qed

lemma id-nec2-1[PLM]:
[◇((α::'a::id-eq) = β) ≡ (α = β) in v]
apply (rule ≡I; rule CP)
 using id-nec[equiv-lr] derived-S5-rules-2-b CP modus-ponens apply blast
using T◇[deduction] by auto

lemma id-nec2-2-Aux:
[(◇φ) ≡ ψ in v] ⟹ [(¬ψ) ≡ □(¬φ) in v]
proof −
  assume [(◇φ) ≡ ψ in v]
  moreover have ⋀φ ψ. [(¬φ) ≡ ψ in v] ⟹ [(¬ψ) ≡ φ in v]
    by PLM-solver
```

ultimately show *?thesis*
unfolding *diamond-def* **by** *blast*
qed

lemma *id-nec2-2*[*PLM*]:
[((α::'a::*id-eq*) ≠ β) ≡ □(α ≠ β) *in v*]
using *id-nec2-1*[*THEN id-nec2-2-Aux*] **by** *auto*

lemma *id-nec2-3*[*PLM*]:
[(◇((α::'a::*id-eq*) ≠ β)) ≡ (α ≠ β) *in v*]
using *T◇* ≡*I id-nec2-2*[*equiv-lr*]
CP derived-S5-rules-2-b **by** *metis*

lemma *exists-desc-box-1*[*PLM*]:
[(∃ *y* . (y^P) = (ι*x*. φ *x*)) → (∃ *y* . □((y^P) = (ι*x*. φ *x*))) *in v*]
proof (*rule CP*)
assume [∃ *y*. (y^P) = (ι*x*. φ *x*) *in v*]
then obtain *y* **where** [(y^P) = (ι*x*. φ *x*) *in v*]
by (*rule* ∃ *E*)
hence [□(y^P = (ι*x*. φ *x*)) *in v*]
using *l-identity*[*axiom-instance*, *deduction*, *deduction*]
cqt-1[*axiom-instance*] *all-self-eq-2*[**where** *'a*=ν]
modus-ponens **unfolding** *identity-ν-def* **by** *fast*
thus [∃ *y*. □((y^P) = (ι*x*. φ *x*)) *in v*]
by (*rule* ∃ *I*)
qed

lemma *exists-desc-box-2*[*PLM*]:
[(∃ *y* . (y^P) = (ι*x*. φ *x*)) → □(∃ *y* .((y^P) = (ι*x*. φ *x*))) *in v*]
using *exists-desc-box-1 Buridan ded-thm-cor-3* **by** *fast*

lemma *en-eq-1*[*PLM*]:
[◇{|*x*,*F*|} ≡ □{|*x*,*F*|} *in v*]
using *encoding*[*axiom-instance*] *RN*
sc-eq-box-box-1 modus-ponens **by** *blast*
lemma *en-eq-2*[*PLM*]:
[{|*x*,*F*|} ≡ □{|*x*,*F*|} *in v*]
using *encoding*[*axiom-instance*] *qml-2*[*axiom-instance*] **by** (*rule* ≡*I*)
lemma *en-eq-3*[*PLM*]:
[◇{|*x*,*F*|} ≡ {|*x*,*F*|} *in v*]
using *encoding*[*axiom-instance*] *derived-S5-rules-2-b* ≡*I T◇* **by** *auto*
lemma *en-eq-4*[*PLM*]:
[({|*x*,*F*|} ≡ {|*y*,*G*|}) ≡ (□{|*x*,*F*|} ≡ □{|*y*,*G*|}) *in v*]
by (*metis CP en-eq-2* ≡*I* ≡*E*(*1*) ≡*E*(*2*))
lemma *en-eq-5*[*PLM*]:
[□({|*x*,*F*|} ≡ {|*y*,*G*|}) ≡ (□{|*x*,*F*|} ≡ □{|*y*,*G*|}) *in v*]
using ≡*I KBasic-6 encoding*[*axiom-necessitation*, *axiom-instance*]
sc-eq-box-box-3[*deduction*] &*I* **by** *simp*
lemma *en-eq-6*[*PLM*]:
[({|*x*,*F*|} ≡ {|*y*,*G*|}) ≡ □({|*x*,*F*|} ≡ {|*y*,*G*|}) *in v*]
using *en-eq-4 en-eq-5 oth-class-taut-4-a* ≡*E*(*6*) **by** *meson*
lemma *en-eq-7*[*PLM*]:
[(¬{|*x*,*F*|}) ≡ □(¬{|*x*,*F*|}) *in v*]
using *en-eq-3*[*THEN id-nec2-2-Aux*] **by** *blast*
lemma *en-eq-8*[*PLM*]:
[◇(¬{|*x*,*F*|}) ≡ (¬{|*x*,*F*|}) *in v*]
unfolding *diamond-def* **apply** (*PLM-subst-method* {|*x*,*F*|} ¬¬{|*x*,*F*|})
using *oth-class-taut-4-b* **apply** *simp*

apply (*PLM-subst-method* ⦃*x*,*F*⦄ □⦃*x*,*F*⦄)
using *en-eq-2* **apply** *simp*
using *oth-class-taut-4-a* **by** *assumption*
lemma *en-eq-9*[*PLM*]:
[◇(¬⦃*x*,*F*⦄) ≡ □(¬⦃*x*,*F*⦄) *in v*]
using *en-eq-8 en-eq-7* ≡*E*(*5*) **by** *blast*
lemma *en-eq-10*[*PLM*]:
[𝓐⦃*x*,*F*⦄ ≡ ⦃*x*,*F*⦄ *in v*]
apply (*rule* ≡*I*)
using *encoding*[*axiom-actualization*, *axiom-instance*,
THEN logic-actual-nec-2[*axiom-instance*, *equiv-lr*],
deduction, *THEN qml-act-2*[*axiom-instance*, *equiv-rl*],
THEN en-eq-2[*equiv-rl*]] *CP*
apply *simp*
using *encoding*[*axiom-instance*] *nec-imp-act ded-thm-cor-3* **by** *blast*

A.9.11. The Theory of Relations

lemma *beta-equiv-eq-1-1*[*PLM*]:
assumes *IsProperInX* φ
and *IsProperInX* ψ
and ⋀*x*.[φ (*x*^{P}) ≡ ψ (*x*^{P}) *in v*]
shows [(|**λ** *y*. φ (*y*^{P}), *x*^{P}|) ≡ (|**λ** *y*. ψ (*y*^{P}), *x*^{P}|) *in v*]
using *lambda-predicates-2-1*[*OF assms*(*1*), *axiom-instance*]
using *lambda-predicates-2-1*[*OF assms*(*2*), *axiom-instance*]
using *assms*(*3*) **by** (*meson* ≡*E*(*6*) *oth-class-taut-4-a*)

lemma *beta-equiv-eq-1-2*[*PLM*]:
assumes *IsProperInXY* φ
and *IsProperInXY* ψ
and ⋀*x y*.[φ (*x*^{P}) (*y*^{P}) ≡ ψ (*x*^{P}) (*y*^{P}) *in v*]
shows [(|**λ**^{2} (λ *x y*. φ (*x*^{P}) (*y*^{P})), *x*^{P}, *y*^{P}|)
≡ (|**λ**^{2} (λ *x y*. ψ (*x*^{P}) (*y*^{P})), *x*^{P}, *y*^{P}|) *in v*]
using *lambda-predicates-2-2*[*OF assms*(*1*), *axiom-instance*]
using *lambda-predicates-2-2*[*OF assms*(*2*), *axiom-instance*]
using *assms*(*3*) **by** (*meson* ≡*E*(*6*) *oth-class-taut-4-a*)

lemma *beta-equiv-eq-1-3*[*PLM*]:
assumes *IsProperInXYZ* φ
and *IsProperInXYZ* ψ
and ⋀*x y z*.[φ (*x*^{P}) (*y*^{P}) (*z*^{P}) ≡ ψ (*x*^{P}) (*y*^{P}) (*z*^{P}) *in v*]
shows [(|**λ**^{3} (λ *x y z*. φ (*x*^{P}) (*y*^{P}) (*z*^{P})), *x*^{P}, *y*^{P}, *z*^{P}|)
≡ (|**λ**^{3} (λ *x y z*. ψ (*x*^{P}) (*y*^{P}) (*z*^{P})), *x*^{P}, *y*^{P}, *z*^{P}|) *in v*]
using *lambda-predicates-2-3*[*OF assms*(*1*), *axiom-instance*]
using *lambda-predicates-2-3*[*OF assms*(*2*), *axiom-instance*]
using *assms*(*3*) **by** (*meson* ≡*E*(*6*) *oth-class-taut-4-a*)

lemma *beta-equiv-eq-2-1*[*PLM*]:
assumes *IsProperInX* φ
and *IsProperInX* ψ
shows [(□(∀ *x* . φ (*x*^{P}) ≡ ψ (*x*^{P}))) →
(□(∀ *x* . (|**λ** *y*. φ (*y*^{P}), *x*^{P}|) ≡ (|**λ** *y*. ψ (*y*^{P}), *x*^{P}|))) *in v*]
apply (*rule qml-1*[*axiom-instance*, *deduction*])
apply (*rule RN*)
proof (*rule CP*, *rule* ∀*I*)
fix *v x*
assume [∀ *x*. φ (*x*^{P}) ≡ ψ (*x*^{P}) *in v*]
hence ⋀*x*.[φ (*x*^{P}) ≡ ψ (*x*^{P}) *in v*]

by *PLM-solver*
thus $[(\!|\boldsymbol{\lambda}\ y.\ \varphi\ (y^P),\ x^P|\!) \equiv (\!|\boldsymbol{\lambda}\ y.\ \psi\ (y^P),\ x^P|\!)\ in\ v]$
using *assms beta-equiv-eq-1-1* **by** *auto*
qed

lemma *beta-equiv-eq-2-2*[*PLM*]:
assumes *IsProperInXY* φ
and *IsProperInXY* ψ
shows $[(\Box(\forall\ x\ y\ .\ \varphi\ (x^P)\ (y^P) \equiv \psi\ (x^P)\ (y^P))) \rightarrow$
$(\Box(\forall\ x\ y\ .\ (\!|\boldsymbol{\lambda}^2\ (\lambda\ x\ y.\ \varphi\ (x^P)\ (y^P)),\ x^P,\ y^P|\!)$
$\equiv (\!|\boldsymbol{\lambda}^2\ (\lambda\ x\ y.\ \psi\ (x^P)\ (y^P)),\ x^P,\ y^P|\!)))\ in\ v]$
apply (*rule qml-1*[*axiom-instance*, *deduction*])
apply (*rule RN*)
proof (*rule CP*, *rule* $\forall I$, *rule* $\forall I$)
fix *v x y*
assume $[\forall x\ y.\ \varphi\ (x^P)\ (y^P) \equiv \psi\ (x^P)\ (y^P)\ in\ v]$
hence $(\bigwedge x\ y.[\varphi\ (x^P)\ (y^P) \equiv \psi\ (x^P)\ (y^P)\ in\ v])$
by (*meson* $\forall E$)
thus $[(\!|\boldsymbol{\lambda}^2\ (\lambda\ x\ y.\ \varphi\ (x^P)\ (y^P)),\ x^P,\ y^P|\!)$
$\equiv (\!|\boldsymbol{\lambda}^2\ (\lambda\ x\ y.\ \psi\ (x^P)\ (y^P)),\ x^P,\ y^P|\!)\ in\ v]$
using *assms beta-equiv-eq-1-2* **by** *auto*
qed

lemma *beta-equiv-eq-2-3*[*PLM*]:
assumes *IsProperInXYZ* φ
and *IsProperInXYZ* ψ
shows $[(\Box(\forall\ x\ y\ z\ .\ \varphi\ (x^P)\ (y^P)\ (z^P) \equiv \psi\ (x^P)\ (y^P)\ (z^P))) \rightarrow$
$(\Box(\forall\ x\ y\ z\ .\ (\!|\boldsymbol{\lambda}^3\ (\lambda\ x\ y\ z.\ \varphi\ (x^P)\ (y^P)\ (z^P)),\ x^P,\ y^P,\ z^P|\!)$
$\equiv (\!|\boldsymbol{\lambda}^3\ (\lambda\ x\ y\ z.\ \psi\ (x^P)\ (y^P)\ (z^P)),\ x^P,\ y^P,\ z^P|\!)))\ in\ v]$
apply (*rule qml-1*[*axiom-instance*, *deduction*])
apply (*rule RN*)
proof (*rule CP*, *rule* $\forall I$, *rule* $\forall I$, *rule* $\forall I$)
fix *v x y z*
assume $[\forall x\ y\ z.\ \varphi\ (x^P)\ (y^P)\ (z^P) \equiv \psi\ (x^P)\ (y^P)\ (z^P)\ in\ v]$
hence $(\bigwedge x\ y\ z.[\varphi\ (x^P)\ (y^P)\ (z^P) \equiv \psi\ (x^P)\ (y^P)\ (z^P)\ in\ v])$
by (*meson* $\forall E$)
thus $[(\!|\boldsymbol{\lambda}^3\ (\lambda\ x\ y\ z.\ \varphi\ (x^P)\ (y^P)\ (z^P)),\ x^P,\ y^P,\ z^P|\!)$
$\equiv (\!|\boldsymbol{\lambda}^3\ (\lambda\ x\ y\ z.\ \psi\ (x^P)\ (y^P)\ (z^P)),\ x^P,\ y^P,\ z^P|\!)\ in\ v]$
using *assms beta-equiv-eq-1-3* **by** *auto*
qed

lemma *beta-C-meta-1*[*PLM*]:
assumes *IsProperInX* φ
shows $[(\!|\boldsymbol{\lambda}\ y.\ \varphi\ (y^P),\ x^P|\!) \equiv \varphi\ (x^P)\ in\ v]$
using *lambda-predicates-2-1*[*OF assms*, *axiom-instance*] **by** *auto*

lemma *beta-C-meta-2*[*PLM*]:
assumes *IsProperInXY* φ
shows $[(\!|\boldsymbol{\lambda}^2\ (\lambda\ x\ y.\ \varphi\ (x^P)\ (y^P)),\ x^P,\ y^P|\!) \equiv \varphi\ (x^P)\ (y^P)\ in\ v]$
using *lambda-predicates-2-2*[*OF assms*, *axiom-instance*] **by** *auto*

lemma *beta-C-meta-3*[*PLM*]:
assumes *IsProperInXYZ* φ
shows $[(\!|\boldsymbol{\lambda}^3\ (\lambda\ x\ y\ z.\ \varphi\ (x^P)\ (y^P)\ (z^P)),\ x^P,\ y^P,\ z^P|\!) \equiv \varphi\ (x^P)\ (y^P)\ (z^P)\ in\ v]$
using *lambda-predicates-2-3*[*OF assms*, *axiom-instance*] **by** *auto*

lemma *relations-1*[*PLM*]:
assumes *IsProperInX* φ

shows $[\exists\ F.\ \Box(\forall\ x.\ (\!|F,x^P|\!) \equiv \varphi\ (x^P))\ in\ v]$
using *assms* **apply** $-$ **by** *PLM-solver*

lemma *relations-2*[*PLM*]:
assumes *IsProperInXY* φ
shows $[\exists\ F.\ \Box(\forall\ x\ y.\ (\!|F,x^P,y^P|\!) \equiv \varphi\ (x^P)\ (y^P))\ in\ v]$
using *assms* **apply** $-$ **by** *PLM-solver*

lemma *relations-3*[*PLM*]:
assumes *IsProperInXYZ* φ
shows $[\exists\ F.\ \Box(\forall\ x\ y\ z.\ (\!|F,x^P,y^P,z^P|\!) \equiv \varphi\ (x^P)\ (y^P)\ (z^P))\ in\ v]$
using *assms* **apply** $-$ **by** *PLM-solver*

lemma *prop-equiv*[*PLM*]:
shows $[(\forall\ x\ .\ (\{\!|x^P,F|\!\} \equiv \{\!|x^P,G|\!\})) \rightarrow F = G\ in\ v]$
proof (*rule CP*)
assume *1*: $[\forall x.\ \{\!|x^P,F|\!\} \equiv \{\!|x^P,G|\!\}\ in\ v]$
{
fix *x*
have $[\{\!|x^P,F|\!\} \equiv \{\!|x^P,G|\!\}\ in\ v]$
using *1* **by** (*rule* $\forall E$)
hence $[\Box(\{\!|x^P,F|\!\} \equiv \{\!|x^P,G|\!\})\ in\ v]$
using *PLM.en-eq-6* $\equiv E(1)$ **by** *blast*
}
hence $[\forall x.\ \Box(\{\!|x^P,F|\!\} \equiv \{\!|x^P,G|\!\})\ in\ v]$
by (*rule* $\forall I$)
thus $[F = G\ in\ v]$
unfolding *identity-defs*
by (*rule BF*[*deduction*])
qed

lemma *propositions-lemma-1*[*PLM*]:
$[\boldsymbol{\lambda}^0\ \varphi = \varphi\ in\ v]$
using *lambda-predicates-3-0*[*axiom-instance*] **.**

lemma *propositions-lemma-2*[*PLM*]:
$[\boldsymbol{\lambda}^0\ \varphi \equiv \varphi\ in\ v]$
using *lambda-predicates-3-0*[*axiom-instance*, *THEN id-eq-prop-prop-8-b*[*deduction*]]
apply (*rule l-identity*[*axiom-instance*, *deduction*, *deduction*])
by *PLM-solver*

lemma *propositions-lemma-4*[*PLM*]:
assumes $\bigwedge x.[\boldsymbol{\mathcal{A}}(\varphi\ x \equiv \psi\ x)\ in\ v]$
shows $[(\chi{::}\kappa{\Rightarrow}o)\ (\boldsymbol{\iota}x.\ \varphi\ x) = \chi\ (\boldsymbol{\iota}x.\ \psi\ x)\ in\ v]$
proof $-$
have $[\boldsymbol{\lambda}^0\ (\chi\ (\boldsymbol{\iota}x.\ \varphi\ x)) = \boldsymbol{\lambda}^0\ (\chi\ (\boldsymbol{\iota}x.\ \psi\ x))\ in\ v]$
using *assms lambda-predicates-4-0*[*axiom-instance*]
by *blast*
hence $[(\chi\ (\boldsymbol{\iota}x.\ \varphi\ x)) = \boldsymbol{\lambda}^0\ (\chi\ (\boldsymbol{\iota}x.\ \psi\ x))\ in\ v]$
using *propositions-lemma-1*[*THEN id-eq-prop-prop-8-b*[*deduction*]]
id-eq-prop-prop-9-b[*deduction*] **&***I*
by *blast*
thus *?thesis*
using *propositions-lemma-1 id-eq-prop-prop-9-b*[*deduction*] **&***I*
by *blast*
qed

lemma *propositions*[*PLM*]:

$[\exists\ p\ .\ \Box(p \equiv p')\ in\ v]$
by *PLM-solver*

lemma *pos-not-equiv-then-not-eq*[*PLM*]:
$[\Diamond(\neg(\forall\, x.\ (\!|F,x^P|\!) \equiv (\!|G,x^P|\!))) \rightarrow F \neq G\ in\ v]$
unfolding *diamond-def*
proof (*subst contraposition-1*[*symmetric*], *rule CP*)
assume $[F = G\ in\ v]$
thus $[\Box(\neg(\neg(\forall\, x.\ (\!|F,x^P|\!) \equiv (\!|G,x^P|\!))))\ in\ v]$
apply (*rule l-identity*[*axiom-instance*, *deduction*, *deduction*])
by *PLM-solver*
qed

lemma *thm-relation-negation-1-1*[*PLM*]:
$[(\!|F^-,\ x^P|\!) \equiv \neg(\!|F,\ x^P|\!)\ in\ v]$
unfolding *propnot-defs*
apply (*rule lambda-predicates-2-1*[*axiom-instance*])
by *show-proper*

lemma *thm-relation-negation-1-2*[*PLM*]:
$[(\!|F^-,\ x^P,\ y^P|\!) \equiv \neg(\!|F,\ x^P,\ y^P|\!)\ in\ v]$
unfolding *propnot-defs*
apply (*rule lambda-predicates-2-2*[*axiom-instance*])
by *show-proper*

lemma *thm-relation-negation-1-3*[*PLM*]:
$[(\!|F^-,\ x^P,\ y^P,\ z^P|\!) \equiv \neg(\!|F,\ x^P,\ y^P,\ z^P|\!)\ in\ v]$
unfolding *propnot-defs*
apply (*rule lambda-predicates-2-3*[*axiom-instance*])
by *show-proper*

lemma *thm-relation-negation-2-1*[*PLM*]:
$[(\neg(\!|F^-,\ x^P|\!)) \equiv (\!|F,\ x^P|\!)\ in\ v]$
using *thm-relation-negation-1-1*[*THEN oth-class-taut-5-d*[*equiv-lr*]]
apply − **by** *PLM-solver*

lemma *thm-relation-negation-2-2*[*PLM*]:
$[(\neg(\!|F^-,\ x^P,\ y^P|\!)) \equiv (\!|F,\ x^P,\ y^P|\!)\ in\ v]$
using *thm-relation-negation-1-2*[*THEN oth-class-taut-5-d*[*equiv-lr*]]
apply − **by** *PLM-solver*

lemma *thm-relation-negation-2-3*[*PLM*]:
$[(\neg(\!|F^-,\ x^P,\ y^P,\ z^P|\!)) \equiv (\!|F,\ x^P,\ y^P,\ z^P|\!)\ in\ v]$
using *thm-relation-negation-1-3*[*THEN oth-class-taut-5-d*[*equiv-lr*]]
apply − **by** *PLM-solver*

lemma *thm-relation-negation-3*[*PLM*]:
$[(p)^- \equiv \neg p\ in\ v]$
unfolding *propnot-defs*
using *propositions-lemma-2* **by** *simp*

lemma *thm-relation-negation-4*[*PLM*]:
$[(\neg((p{::}\mathrm{o})^-)) \equiv p\ in\ v]$
using *thm-relation-negation-3*[*THEN oth-class-taut-5-d*[*equiv-lr*]]
apply − **by** *PLM-solver*

lemma *thm-relation-negation-5-1*[*PLM*]:
$[(F{::}\Pi_1) \neq (F^-)\ in\ v]$

using *id-eq-prop-prop-2*[*deduction*]
l-identity[**where** $\varphi{=}\lambda\ G\ .\ (\!|G,x^P|\!) \equiv (\!|F^-,x^P|\!)$, *axiom-instance*, *deduction*, *deduction*]
oth-class-taut-4-a thm-relation-negation-1-1 $\equiv E(5)$
oth-class-taut-1-b modus-tollens-1 CP
by *meson*

lemma *thm-relation-negation-5-2*[*PLM*]:
$[(F::\Pi_2) \neq (F^-)\ in\ v]$
using *id-eq-prop-prop-5-a*[*deduction*]
l-identity[**where** $\varphi{=}\lambda\ G\ .\ (\!|G,x^P,y^P|\!) \equiv (\!|F^-,x^P,y^P|\!)$, *axiom-instance*, *deduction*, *deduction*]
oth-class-taut-4-a thm-relation-negation-1-2 $\equiv E(5)$
oth-class-taut-1-b modus-tollens-1 CP
by *meson*

lemma *thm-relation-negation-5-3*[*PLM*]:
$[(F::\Pi_3) \neq (F^-)\ in\ v]$
using *id-eq-prop-prop-5-b*[*deduction*]
l-identity[**where** $\varphi{=}\lambda\ G\ .\ (\!|G,x^P,y^P,z^P|\!) \equiv (\!|F^-,x^P,y^P,z^P|\!)$, *axiom-instance*, *deduction*, *deduction*]
oth-class-taut-4-a thm-relation-negation-1-3 $\equiv E(5)$
oth-class-taut-1-b modus-tollens-1 CP
by *meson*

lemma *thm-relation-negation-6*[*PLM*]:
$[(p::o) \neq (p^-)\ in\ v]$
using *id-eq-prop-prop-8-b*[*deduction*]
l-identity[**where** $\varphi{=}\lambda\ G\ .\ G \equiv (p^-)$, *axiom-instance*, *deduction*, *deduction*]
oth-class-taut-4-a thm-relation-negation-3 $\equiv E(5)$
oth-class-taut-1-b modus-tollens-1 CP
by *meson*

lemma *thm-relation-negation-7*[*PLM*]:
$[((p::o)^-) = \neg p\ in\ v]$
unfolding *propnot-defs* **using** *propositions-lemma-1* **by** *simp*

lemma *thm-relation-negation-8*[*PLM*]:
$[(p::o) \neq \neg p\ in\ v]$
unfolding *propnot-defs*
using *id-eq-prop-prop-8-b*[*deduction*]
l-identity[**where** $\varphi{=}\lambda\ G\ .\ G \equiv \neg(p)$, *axiom-instance*, *deduction*, *deduction*]
oth-class-taut-4-a oth-class-taut-1-b
modus-tollens-1 CP
by *meson*

lemma *thm-relation-negation-9*[*PLM*]:
$[((p::o) = q) \to ((\neg p) = (\neg q))\ in\ v]$
using *l-identity*[**where** $\alpha{=}p$ **and** $\beta{=}q$ **and** $\varphi{=}\lambda\ x\ .\ (\neg p) = (\neg x)$, *axiom-instance*, *deduction*]
id-eq-prop-prop-7-b **using** *CP modus-ponens* **by** *blast*

lemma *thm-relation-negation-10*[*PLM*]:
$[((p::o) = q) \to ((p^-) = (q^-))\ in\ v]$
using *l-identity*[**where** $\alpha{=}p$ **and** $\beta{=}q$ **and** $\varphi{=}\lambda\ x\ .\ (p^-) = (x^-)$, *axiom-instance*, *deduction*]

id-eq-prop-prop-7-b **using** *CP modus-ponens* **by** *blast*

lemma *thm-cont-prop-1*[*PLM*]:
[*NonContingent* (F::Π_1) $\equiv$ *NonContingent* (F^-) *in* v]
proof (*rule* $\equiv I$; *rule CP*)
assume [*NonContingent F in* v]
hence [$\Box(\forall x.(\!|F,x^P|\!)) \vee \Box(\forall x.\neg(\!|F,x^P|\!))$ *in* v]
unfolding *NonContingent-def Necessary-defs Impossible-defs* **.**
hence [$\Box(\forall x.\ \neg(\!|F^-,x^P|\!)) \vee \Box(\forall x.\ \neg(\!|F,x^P|\!))$ *in* v]
apply $-$
apply (*PLM-subst-method* $\lambda\ x\ .\ (\!|F,x^P|\!)\ \lambda\ x\ .\ \neg(\!|F^-,x^P|\!)$)
using *thm-relation-negation-2-1*[*equiv-sym*] **by** *auto*
hence [$\Box(\forall x.\ \neg(\!|F^-,x^P|\!)) \vee \Box(\forall x.\ (\!|F^-,x^P|\!))$ *in* v]
apply $-$
apply (*PLM-subst-goal-method*
$\lambda\ \varphi\ .\ \Box(\forall x.\ \neg(\!|F^-,x^P|\!)) \vee \Box(\forall x.\ \varphi\ x)\ \lambda\ x\ .\ \neg(\!|F,x^P|\!)$)
using *thm-relation-negation-1-1*[*equiv-sym*] **by** *auto*
hence [$\Box(\forall x.\ (\!|F^-,x^P|\!)) \vee \Box(\forall x.\ \neg(\!|F^-,x^P|\!))$ *in* v]
by (*rule oth-class-taut-3-e*[*equiv-lr*])
thus [*NonContingent* (F^-) *in* v]
unfolding *NonContingent-def Necessary-defs Impossible-defs* **.**
next
assume [*NonContingent* (F^-) *in* v]
hence [$\Box(\forall x.\ \neg(\!|F^-,x^P|\!)) \vee \Box(\forall x.\ (\!|F^-,x^P|\!))$ *in* v]
unfolding *NonContingent-def Necessary-defs Impossible-defs*
by (*rule oth-class-taut-3-e*[*equiv-lr*])
hence [$\Box(\forall x.(\!|F,x^P|\!)) \vee \Box(\forall x.(\!|F^-,x^P|\!))$ *in* v]
apply $-$
apply (*PLM-subst-method* $\lambda\ x\ .\ \neg(\!|F^-,x^P|\!)\ \lambda\ x\ .\ (\!|F,x^P|\!)$)
using *thm-relation-negation-2-1* **by** *auto*
hence [$\Box(\forall x.\ (\!|F,x^P|\!)) \vee \Box(\forall x.\ \neg(\!|F,x^P|\!))$ *in* v]
apply $-$
apply (*PLM-subst-method* $\lambda\ x\ .\ (\!|F^-,x^P|\!)\ \lambda\ x\ .\ \neg(\!|F,x^P|\!)$)
using *thm-relation-negation-1-1* **by** *auto*
thus [*NonContingent F in* v]
unfolding *NonContingent-def Necessary-defs Impossible-defs* **.**
qed

lemma *thm-cont-prop-2*[*PLM*]:
[*Contingent* $F \equiv \Diamond(\exists\ x\ .\ (\!|F,x^P|\!))$ **&** $\Diamond(\exists\ x\ .\ \neg(\!|F,x^P|\!))$ *in* v]
proof (*rule* $\equiv I$; *rule CP*)
assume [*Contingent F in* v]
hence [$\neg(\Box(\forall x.(\!|F,x^P|\!)) \vee \Box(\forall x.\neg(\!|F,x^P|\!)))$ *in* v]
unfolding *Contingent-def Necessary-defs Impossible-defs* **.**
hence [$(\neg\Box(\forall x.(\!|F,x^P|\!)))$ **&** $(\neg\Box(\forall x.\neg(\!|F,x^P|\!)))$ *in* v]
by (*rule oth-class-taut-6-d*[*equiv-lr*])
hence [$(\Diamond\neg(\forall x.\neg(\!|F,x^P|\!)))$ **&** $(\Diamond\neg(\forall x.(\!|F,x^P|\!)))$ *in* v]
using *KBasic2-2*[*equiv-lr*] **&***I* **&***E* **by** *meson*
thus [$(\Diamond(\exists\ x.(\!|F,x^P|\!)))$ **&** $(\Diamond(\exists x.\ \neg(\!|F,x^P|\!)))$ *in* v]
unfolding *exists-def* **apply** $-$
apply (*PLM-subst-method* $\lambda\ x\ .\ (\!|F,x^P|\!)\ \lambda\ x\ .\ \neg\neg(\!|F,x^P|\!)$)
using *oth-class-taut-4-b* **by** *auto*
next
assume [$(\Diamond(\exists\ x.(\!|F,x^P|\!)))$ **&** $(\Diamond(\exists x.\ \neg(\!|F,x^P|\!)))$ *in* v]
hence [$(\Diamond\neg(\forall x.\neg(\!|F,x^P|\!)))$ **&** $(\Diamond\neg(\forall x.(\!|F,x^P|\!)))$ *in* v]
unfolding *exists-def* **apply** $-$
apply (*PLM-subst-goal-method*
$\lambda\ \varphi\ .\ (\Diamond\neg(\forall x.\neg(\!|F,x^P|\!)))$ **&** $(\Diamond\neg(\forall x.\ \varphi\ x))\ \lambda\ x\ .\ \neg\neg(\!|F,x^P|\!)$)

using *oth-class-taut-4-b*[*equiv-sym*] **by** *auto*
hence [(¬□(∀ *x*.(|*F*,*x*^{P}|)) & (¬□(∀ *x*.¬(|*F*,*x*^{P}|))) *in v*]
using *KBasic2-2*[*equiv-rl*] &*I* &*E* **by** *meson*
hence [¬(□(∀ *x*.(|*F*,*x*^{P}|)) ∨ □(∀ *x*.¬(|*F*,*x*^{P}|))) *in v*]
by (*rule oth-class-taut-6-d*[*equiv-rl*])
thus [*Contingent F in v*]
unfolding *Contingent-def Necessary-defs Impossible-defs* .
qed

lemma *thm-cont-prop-3*[*PLM*]:
[*Contingent* (*F*::Π$_1$) ≡ *Contingent* (*F*$^-$) *in v*]
using *thm-cont-prop-1*
unfolding *NonContingent-def Contingent-def*
by (*rule oth-class-taut-5-d*[*equiv-lr*])

lemma *lem-cont-e*[*PLM*]:
[◇(∃ *x* . (|*F*,*x*^{P}|) & (◇(¬(|*F*,*x*^{P}|)))) ≡ ◇(∃ *x* . ((¬(|*F*,*x*^{P}|)) & ◇(|*F*,*x*^{P}|)))*in v*]
proof −
have [◇(∃ *x* . (|*F*,*x*^{P}|) & (◇(¬(|*F*,*x*^{P}|)))) *in v*]
= [(∃ *x* . ◇((|*F*,*x*^{P}|) & ◇(¬(|*F*,*x*^{P}|)))) *in v*]
using *BF*◇[*deduction*] *CBF*◇[*deduction*] **by** *fast*
also have ... = [∃ *x* . (◇(|*F*,*x*^{P}|) & ◇(¬(|*F*,*x*^{P}|))) *in v*]
apply (*PLM-subst-method*
λ *x* . ◇((|*F*,*x*^{P}|) & ◇(¬(|*F*,*x*^{P}|)))
λ *x* . ◇(|*F*,*x*^{P}|) & ◇(¬(|*F*,*x*^{P}|)))
using *S5Basic-12* **by** *auto*
also have ... = [∃ *x* . ◇(¬(|*F*,*x*^{P}|)) & ◇(|*F*,*x*^{P}|) *in v*]
apply (*PLM-subst-method*
λ *x* . ◇(|*F*,*x*^{P}|) & ◇(¬(|*F*,*x*^{P}|))
λ *x* . ◇(¬(|*F*,*x*^{P}|)) & ◇(|*F*,*x*^{P}|))
using *oth-class-taut-3-b* **by** *auto*
also have ... = [∃ *x* . ◇((¬(|*F*,*x*^{P}|)) & ◇(|*F*,*x*^{P}|)) *in v*]
apply (*PLM-subst-method*
λ *x* . ◇(¬(|*F*,*x*^{P}|)) & ◇(|*F*,*x*^{P}|)
λ *x* . ◇((¬(|*F*,*x*^{P}|)) & ◇(|*F*,*x*^{P}|)))
using *S5Basic-12*[*equiv-sym*] **by** *auto*
also have ... = [◇ (∃ *x* . ((¬(|*F*,*x*^{P}|)) & ◇(|*F*,*x*^{P}|))) *in v*]
using *CBF*◇[*deduction*] *BF*◇[*deduction*] **by** *fast*
finally show *?thesis* **using** ≡*I CP* **by** *blast*
qed

lemma *lem-cont-e-2*[*PLM*]:
[◇(∃ *x* . (|*F*,*x*^{P}|) & ◇(¬(|*F*,*x*^{P}|))) ≡ ◇(∃ *x* . (|*F*$^-$,*x*^{P}|) & ◇(¬(|*F*$^-$,*x*^{P}|))) *in v*]
apply (*PLM-subst-method* λ *x* . (|*F*,*x*^{P}|) λ *x* . ¬(|*F*$^-$,*x*^{P}|))
using *thm-relation-negation-2-1*[*equiv-sym*] **apply** *simp*
apply (*PLM-subst-method* λ *x* . ¬(|*F*,*x*^{P}|) λ *x* . (|*F*$^-$,*x*^{P}|))
using *thm-relation-negation-1-1*[*equiv-sym*] **apply** *simp*
using *lem-cont-e* **by** *simp*

lemma *thm-cont-e-1*[*PLM*]:
[◇(∃ *x* . ((¬(|*E*!,*x*^{P}|)) & (◇(|*E*!,*x*^{P}|)))) *in v*]
using *lem-cont-e*[**where** *F*=*E*!, *equiv-lr*] *qml-4*[*axiom-instance*,*conj1*]
by *blast*

lemma *thm-cont-e-2*[*PLM*]:
[*Contingent* (*E*!) *in v*]
using *thm-cont-prop-2*[*equiv-rl*] &*I qml-4*[*axiom-instance*, *conj1*]
KBasic2-8[*deduction*, *OF sign-S5-thm-3*[*deduction*], *conj1*]

KBasic2-8[*deduction*, *OF sign-S5-thm-3*[*deduction*, *OF thm-cont-e-1*], *conj1*]
by *fast*

lemma *thm-cont-e-3*[*PLM*]:
[*Contingent* ($E!^-$) *in v*]
using *thm-cont-e-2 thm-cont-prop-3*[*equiv-lr*] **by** *blast*

lemma *thm-cont-e-4*[*PLM*]:
[$\exists$ (F::Π_1) G . ($F \neq G$ & *Contingent F* & *Contingent G*) *in v*]
apply (*rule-tac* α=$E!$ **in** $\exists\, I$, *rule-tac* α=$E!^-$ **in** $\exists\, I$)
using *thm-cont-e-2 thm-cont-e-3 thm-relation-negation-5-1* &*I* **by** *auto*

context
begin
qualified definition L **where** $L \equiv (\boldsymbol{\lambda}\ x\ .\ (\!|E!, x^P|\!) \rightarrow (\!|E!, x^P|\!))$

lemma *thm-noncont-e-e-1*[*PLM*]:
[*Necessary L in v*]
unfolding *Necessary-defs L-def* **apply** (*rule RN*, *rule* $\forall\, I$)
apply (*rule lambda-predicates-2-1*[*axiom-instance*, *equiv-rl*])
apply *show-proper*
using *if-p-then-p* **.**

lemma *thm-noncont-e-e-2*[*PLM*]:
[*Impossible* (L^-) *in v*]
unfolding *Impossible-defs L-def* **apply** (*rule RN*, *rule* $\forall\, I$)
apply (*rule thm-relation-negation-2-1*[*equiv-rl*])
apply (*rule lambda-predicates-2-1*[*axiom-instance*, *equiv-rl*])
apply *show-proper*
using *if-p-then-p* **.**

lemma *thm-noncont-e-e-3*[*PLM*]:
[*NonContingent* (*L*) *in v*]
unfolding *NonContingent-def* **using** *thm-noncont-e-e-1*
by (*rule* $\vee I(1)$)

lemma *thm-noncont-e-e-4*[*PLM*]:
[*NonContingent* (L^-) *in v*]
unfolding *NonContingent-def* **using** *thm-noncont-e-e-2*
by (*rule* $\vee I(2)$)

lemma *thm-noncont-e-e-5*[*PLM*]:
[$\exists$ (F::Π_1) G . $F \neq G$ & *NonContingent F* & *NonContingent G in v*]
apply (*rule-tac* α=L **in** $\exists\, I$, *rule-tac* α=L^- **in** $\exists\, I$)
using $\exists\, I$ *thm-relation-negation-5-1 thm-noncont-e-e-3*
thm-noncont-e-e-4 &*I*
by *simp*

lemma *four-distinct-1*[*PLM*]:
[*NonContingent* (F::Π_1) $\rightarrow$ $\neg$($\exists$ G . (*Contingent G* & $G = F$)) *in v*]
proof (*rule CP*)
assume [*NonContingent F in v*]
hence [$\neg$(*Contingent F*) *in v*]
unfolding *NonContingent-def Contingent-def*
apply $-$ **by** *PLM-solver*
moreover {
assume [$\exists$ G . *Contingent G* & $G = F$ *in v*]

```
    then obtain P where [Contingent P & P = F in v]
     by (rule ∃ E)
    hence [Contingent F in v]
      using &E l-identity[axiom-instance, deduction, deduction]
      by blast
  }
  ultimately show [¬(∃ G. Contingent G & G = F) in v]
    using modus-tollens-1 CP by blast
qed

lemma four-distinct-2[PLM]:
 [Contingent (F::Π1) → ¬(∃ G . (NonContingent G & G = F)) in v]
 proof (rule CP)
  assume [Contingent F in v]
  hence [¬(NonContingent F) in v]
    unfolding NonContingent-def Contingent-def
    apply − by PLM-solver
  moreover {
    assume [∃ G . NonContingent G & G = F in v]
    then obtain P where [NonContingent P & P = F in v]
     by (rule ∃ E)
    hence [NonContingent F in v]
      using &E l-identity[axiom-instance, deduction, deduction]
      by blast
  }
  ultimately show [¬(∃ G. NonContingent G & G = F) in v]
    using modus-tollens-1 CP by blast
qed

 lemma four-distinct-3[PLM]:
  [L ≠ (L⁻) & L ≠ E! & L ≠ (E!⁻) & (L⁻) ≠ E!
    & (L⁻) ≠ (E!⁻) & E! ≠ (E!⁻) in v]
  proof (rule &I)+
    show [L ≠ (L⁻) in v]
    by (rule thm-relation-negation-5-1)
  next
    {
      assume [L = E! in v]
      hence [NonContingent L & L = E! in v]
        using thm-noncont-e-e-3 &I by auto
      hence [∃ G . NonContingent G & G = E! in v]
        using thm-noncont-e-e-3 &I ∃ I by fast
    }
    thus [L ≠ E! in v]
      using four-distinct-2[deduction, OF thm-cont-e-2]
            modus-tollens-1 CP
      by blast
  next
    {
      assume [L = (E!⁻) in v]
      hence [NonContingent L & L = (E!⁻) in v]
        using thm-noncont-e-e-3 &I by auto
      hence [∃ G . NonContingent G & G = (E!⁻) in v]
        using thm-noncont-e-e-3 &I ∃ I by fast
    }
    thus [L ≠ (E!⁻) in v]
      using four-distinct-2[deduction, OF thm-cont-e-3]
            modus-tollens-1 CP
```

```
      by blast
    next
      {
        assume [(L⁻) = E! in v]
        hence [NonContingent (L⁻) & (L⁻) = E! in v]
          using thm-noncont-e-e-4 &I by auto
        hence [∃ G . NonContingent G & G = E! in v]
          using thm-noncont-e-e-3 &I ∃ I by fast
      }
      thus [(L⁻) ≠ E! in v]
        using four-distinct-2[deduction, OF thm-cont-e-2]
              modus-tollens-1 CP
        by blast
    next
      {
        assume [(L⁻) = (E!⁻) in v]
        hence [NonContingent (L⁻) & (L⁻) = (E!⁻) in v]
          using thm-noncont-e-e-4 &I by auto
        hence [∃ G . NonContingent G & G = (E!⁻) in v]
          using thm-noncont-e-e-3 &I ∃ I by fast
      }
      thus [(L⁻) ≠ (E!⁻) in v]
        using four-distinct-2[deduction, OF thm-cont-e-3]
              modus-tollens-1 CP
        by blast
    next
      show [E! ≠ (E!⁻) in v]
        by (rule thm-relation-negation-5-1)
    qed
end

lemma thm-cont-propos-1[PLM]:
  [NonContingent (p::o) ≡ NonContingent (p⁻) in v]
  proof (rule ≡I; rule CP)
    assume [NonContingent p in v]
    hence [□p ∨ □¬p in v]
      unfolding NonContingent-def Necessary-defs Impossible-defs .
    hence [□(¬(p⁻)) ∨ □(¬p) in v]
      apply −
      apply (PLM-subst-method p ¬(p⁻))
      using thm-relation-negation-4[equiv-sym] by auto
    hence [□(¬(p⁻)) ∨ □(p⁻) in v]
      apply −
      apply (PLM-subst-goal-method λφ . □(¬(p⁻)) ∨ □(φ) ¬p)
      using thm-relation-negation-3[equiv-sym] by auto
    hence [□(p⁻) ∨ □(¬(p⁻)) in v]
      by (rule oth-class-taut-3-e[equiv-lr])
    thus [NonContingent (p⁻) in v]
      unfolding NonContingent-def Necessary-defs Impossible-defs .
  next
    assume [NonContingent (p⁻) in v]
    hence [□(¬(p⁻)) ∨ □(p⁻) in v]
      unfolding NonContingent-def Necessary-defs Impossible-defs
      by (rule oth-class-taut-3-e[equiv-lr])
    hence [□(p) ∨ □(p⁻) in v]
      apply −
      apply (PLM-subst-goal-method λφ . □φ ∨ □(p⁻) ¬(p⁻))
      using thm-relation-negation-4 by auto
```

hence $[\Box(p) \vee \Box(\neg p)\ in\ v]$
apply $-$
apply $(PLM\text{-}subst\text{-}method\ p^{-}\ \neg p)$
using *thm-relation-negation-3* **by** *auto*
thus $[NonContingent\ p\ in\ v]$
unfolding *NonContingent-def Necessary-defs Impossible-defs* **.**
qed

lemma *thm-cont-propos-2*$[PLM]$:
$[Contingent\ p \equiv \Diamond p\ \&\ \Diamond(\neg p)\ in\ v]$
proof $(rule \equiv I;\ rule\ CP)$
assume $[Contingent\ p\ in\ v]$
hence $[\neg(\Box p \vee \Box(\neg p))\ in\ v]$
unfolding *Contingent-def Necessary-defs Impossible-defs* **.**
hence $[(\neg\Box p)\ \&\ (\neg\Box(\neg p))\ in\ v]$
by $(rule\ oth\text{-}class\text{-}taut\text{-}6\text{-}d[equiv\text{-}lr])$
hence $[(\Diamond\neg(\neg p))\ \&\ (\Diamond\neg p)\ in\ v]$
using *KBasic2-2*$[equiv\text{-}lr]$ $\&I$ $\&E$ **by** *meson*
thus $[(\Diamond p)\ \&\ (\Diamond(\neg p))\ in\ v]$
apply $-$ **apply** *PLM-solver*
apply $(PLM\text{-}subst\text{-}method\ \neg\neg p\ p)$
using *oth-class-taut-4-b*$[equiv\text{-}sym]$ **by** *auto*
next
assume $[(\Diamond p)\ \&\ (\Diamond\neg(p))\ in\ v]$
hence $[(\Diamond\neg(\neg p))\ \&\ (\Diamond\neg(p))\ in\ v]$
apply $-$ **apply** *PLM-solver*
apply $(PLM\text{-}subst\text{-}method\ p\ \neg\neg p)$
using *oth-class-taut-4-b* **by** *auto*
hence $[(\neg\Box p)\ \&\ (\neg\Box(\neg p))\ in\ v]$
using *KBasic2-2*$[equiv\text{-}rl]$ $\&I$ $\&E$ **by** *meson*
hence $[\neg(\Box(p) \vee \Box(\neg p))\ in\ v]$
by $(rule\ oth\text{-}class\text{-}taut\text{-}6\text{-}d[equiv\text{-}rl])$
thus $[Contingent\ p\ in\ v]$
unfolding *Contingent-def Necessary-defs Impossible-defs* **.**
qed

lemma *thm-cont-propos-3*$[PLM]$:
$[Contingent\ (p::o) \equiv Contingent\ (p^{-})\ in\ v]$
using *thm-cont-propos-1*
unfolding *NonContingent-def Contingent-def*
by $(rule\ oth\text{-}class\text{-}taut\text{-}5\text{-}d[equiv\text{-}lr])$

context
begin
private definition p_0 **where**
$p_0 \equiv \forall x.\ (\!|E!,x^P|\!) \rightarrow (\!|E!,x^P|\!)$

lemma *thm-noncont-propos-1*$[PLM]$:
$[Necessary\ p_0\ in\ v]$
unfolding *Necessary-defs* p_0*-def*
apply $(rule\ RN,\ rule \forall I)$
using *if-p-then-p* **.**

lemma *thm-noncont-propos-2*$[PLM]$:
$[Impossible\ (p_0^{-})\ in\ v]$
unfolding *Impossible-defs*
apply $(PLM\text{-}subst\text{-}method\ \neg p_0\ p_0^{-})$
using *thm-relation-negation-3*$[equiv\text{-}sym]$ **apply** *simp*

apply (*PLM-subst-method* p_0 $\neg\neg p_0$)
using *oth-class-taut-4-b* **apply** *simp*
using *thm-noncont-propos-1* **unfolding** *Necessary-defs*
by *simp*

lemma *thm-noncont-propos-3*[*PLM*]:
[*NonContingent* (p_0) *in v*]
unfolding *NonContingent-def* **using** *thm-noncont-propos-1*
by (*rule* $\vee I(1)$)

lemma *thm-noncont-propos-4*[*PLM*]:
[*NonContingent* ($p_0{}^-$) *in v*]
unfolding *NonContingent-def* **using** *thm-noncont-propos-2*
by (*rule* $\vee I(2)$)

lemma *thm-noncont-propos-5*[*PLM*]:
[$\exists$ (p::o) q . $p \neq q$ & *NonContingent p* & *NonContingent q in v*]
apply (*rule-tac* α=p_0 **in** $\exists I$, *rule-tac* α=$p_0{}^-$ **in** $\exists I$)
using $\exists I$ *thm-relation-negation-6 thm-noncont-propos-3*
thm-noncont-propos-4 &*I* **by** *simp*

private definition q_0 **where**
$q_0 \equiv \exists\ x\ .\ (\!|E!,x^P|\!)$ & $\Diamond(\neg(\!|E!,x^P|\!))$

lemma *basic-prop-1*[*PLM*]:
[$\exists\ p\ .\ \Diamond p$ & $\Diamond(\neg p)$ *in v*]
apply (*rule-tac* α=q_0 **in** $\exists I$) **unfolding** q_0-*def*
using *qml-4*[*axiom-instance*] **by** *simp*

lemma *basic-prop-2*[*PLM*]:
[*Contingent* q_0 *in v*]
unfolding *Contingent-def Necessary-defs Impossible-defs*
apply (*rule oth-class-taut-6-d*[*equiv-rl*])
apply (*PLM-subst-goal-method* $\lambda\ \varphi\ .\ (\neg\Box(\varphi))$ & $\neg\Box\neg q_0$ $\neg\neg q_0$)
using *oth-class-taut-4-b*[*equiv-sym*] **apply** *simp*
using *qml-4*[*axiom-instance*,*conj-sym*]
unfolding q_0-*def diamond-def* **by** *simp*

lemma *basic-prop-3*[*PLM*]:
[*Contingent* ($q_0{}^-$) *in v*]
apply (*rule thm-cont-propos-3*[*equiv-lr*])
using *basic-prop-2* **.**

lemma *basic-prop-4*[*PLM*]:
[$\exists$ (p::o) q . $p \neq q$ & *Contingent p* & *Contingent q in v*]
apply (*rule-tac* α=q_0 **in** $\exists I$, *rule-tac* α=$q_0{}^-$ **in** $\exists I$)
using *thm-relation-negation-6 basic-prop-2 basic-prop-3* &*I* **by** *simp*

lemma *four-distinct-props-1*[*PLM*]:
[*NonContingent* (p::Π_0) $\rightarrow$ ($\neg$($\exists\ q$. *Contingent q* & $q = p$)) *in v*]
proof (*rule CP*)
assume [*NonContingent p in v*]
hence [$\neg$(*Contingent p*) *in v*]
unfolding *NonContingent-def Contingent-def*
apply − **by** *PLM-solver*
moreover {
assume [$\exists\ q$. *Contingent q* & $q = p$ *in v*]
then obtain r **where** [*Contingent r* & $r = p$ *in v*]

by (*rule* ∃ *E*)
hence [*Contingent p in v*]
using &*E l-identity*[*axiom-instance*, *deduction*, *deduction*]
by *blast*
}
ultimately show [¬(∃ *q*. *Contingent q* & *q* = *p*) *in v*]
using *modus-tollens-1 CP* **by** *blast*
qed

lemma *four-distinct-props-2*[*PLM*]:
[*Contingent* (*p*::o) → ¬(∃ *q* . (*NonContingent q* & *q* = *p*)) *in v*]
proof (*rule CP*)
assume [*Contingent p in v*]
hence [¬(*NonContingent p*) *in v*]
unfolding *NonContingent-def Contingent-def*
apply − **by** *PLM-solver*
moreover {
assume [∃ *q* . *NonContingent q* & *q* = *p in v*]
then obtain *r* **where** [*NonContingent r* & *r* = *p in v*]
by (*rule* ∃ *E*)
hence [*NonContingent p in v*]
using &*E l-identity*[*axiom-instance*, *deduction*, *deduction*]
by *blast*
}
ultimately show [¬(∃ *q*. *NonContingent q* & *q* = *p*) *in v*]
using *modus-tollens-1 CP* **by** *blast*
qed

lemma *four-distinct-props-4*[*PLM*]:
[p_0 ≠ (p_0^-) & p_0 ≠ q_0 & p_0 ≠ (q_0^-) & (p_0^-) ≠ q_0
& (p_0^-) ≠ (q_0^-) & q_0 ≠ (q_0^-) *in v*]
proof (*rule* &*I*)+
show [p_0 ≠ (p_0^-) *in v*]
by (*rule thm-relation-negation-6*)
next
{
assume [p_0 = q_0 *in v*]
hence [∃ *q* . *NonContingent q* & *q* = q_0 *in v*]
using &*I thm-noncont-propos-3* ∃ *I*[**where** α=p_0]
by *simp*
}
thus [p_0 ≠ q_0 *in v*]
using *four-distinct-props-2*[*deduction*, *OF basic-prop-2*]
modus-tollens-1 CP
by *blast*
next
{
assume [p_0 = (q_0^-) *in v*]
hence [∃ *q* . *NonContingent q* & *q* = (q_0^-) *in v*]
using *thm-noncont-propos-3* &*I* ∃ *I*[**where** α=p_0] **by** *simp*
}
thus [p_0 ≠ (q_0^-) *in v*]
using *four-distinct-props-2*[*deduction*, *OF basic-prop-3*]
modus-tollens-1 CP
by *blast*
next
{
assume [(p_0^-) = q_0 *in v*]

```
    hence [∃ q . NonContingent q & q = q₀ in v]
      using thm-noncont-propos-4 &I ∃ I[where α=p₀⁻] by auto
  }
  thus [(p₀⁻) ≠ q₀ in v]
    using four-distinct-props-2[deduction, OF basic-prop-2]
          modus-tollens-1 CP
    by blast
next
  {
    assume [(p₀⁻) = (q₀⁻) in v]
    hence [∃ q . NonContingent q & q = (q₀⁻) in v]
      using thm-noncont-propos-4 &I ∃ I[where α=p₀⁻] by auto
  }
  thus [(p₀⁻) ≠ (q₀⁻) in v]
    using four-distinct-props-2[deduction, OF basic-prop-3]
          modus-tollens-1 CP
    by blast
next
  show [q₀ ≠ (q₀⁻) in v]
    by (rule thm-relation-negation-6)
qed

lemma cont-true-cont-1[PLM]:
  [ContingentlyTrue p → Contingent p in v]
  apply (rule CP, rule thm-cont-propos-2[equiv-rl])
  unfolding ContingentlyTrue-def
  apply (rule &I, drule &E(1))
   using T◊[deduction] apply simp
  by (rule &E(2))

lemma cont-true-cont-2[PLM]:
  [ContingentlyFalse p → Contingent p in v]
  apply (rule CP, rule thm-cont-propos-2[equiv-rl])
  unfolding ContingentlyFalse-def
  apply (rule &I, drule &E(2))
   apply simp
  apply (drule &E(1))
  using T◊[deduction] by simp

lemma cont-true-cont-3[PLM]:
  [ContingentlyTrue p ≡ ContingentlyFalse (p⁻) in v]
  unfolding ContingentlyTrue-def ContingentlyFalse-def
  apply (PLM-subst-method ¬p p⁻)
   using thm-relation-negation-3[equiv-sym] apply simp
  apply (PLM-subst-method p ¬¬p)
  by PLM-solver+

lemma cont-true-cont-4[PLM]:
  [ContingentlyFalse p ≡ ContingentlyTrue (p⁻) in v]
  unfolding ContingentlyTrue-def ContingentlyFalse-def
  apply (PLM-subst-method ¬p p⁻)
   using thm-relation-negation-3[equiv-sym] apply simp
  apply (PLM-subst-method p ¬¬p)
  by PLM-solver+

lemma cont-tf-thm-1[PLM]:
  [ContingentlyTrue q₀ ∨ ContingentlyFalse q₀ in v]
  proof −
```

```
    have [q_0 ∨ ¬q_0 in v]
      by PLM-solver
    moreover {
      assume [q_0 in v]
      hence [q_0 & ◇¬q_0 in v]
        unfolding q_0-def
        using qml-4[axiom-instance,conj2] &I
        by auto
    }
    moreover {
      assume [¬q_0 in v]
      hence [(¬q_0) & ◇q_0 in v]
        unfolding q_0-def
        using qml-4[axiom-instance,conj1] &I
        by auto
    }
    ultimately show ?thesis
      unfolding ContingentlyTrue-def ContingentlyFalse-def
      using ∨E(4) CP by auto
  qed

lemma cont-tf-thm-2[PLM]:
  [ContingentlyFalse q_0 ∨ ContingentlyFalse (q_0⁻) in v]
  using cont-tf-thm-1 cont-true-cont-3[where p=q_0]
        cont-true-cont-4[where p=q_0]
  apply − by PLM-solver

lemma cont-tf-thm-3[PLM]:
  [∃ p . ContingentlyTrue p in v]
  proof (rule ∨E(1); (rule CP)?)
    show [ContingentlyTrue q_0 ∨ ContingentlyFalse q_0 in v]
      using cont-tf-thm-1 .
  next
    assume [ContingentlyTrue q_0 in v]
    thus ?thesis
      using ∃I by metis
  next
    assume [ContingentlyFalse q_0 in v]
    hence [ContingentlyTrue (q_0⁻) in v]
      using cont-true-cont-4[equiv-lr] by simp
    thus ?thesis
      using ∃I by metis
  qed

lemma cont-tf-thm-4[PLM]:
  [∃ p . ContingentlyFalse p in v]
  proof (rule ∨E(1); (rule CP)?)
    show [ContingentlyTrue q_0 ∨ ContingentlyFalse q_0 in v]
      using cont-tf-thm-1 .
  next
    assume [ContingentlyTrue q_0 in v]
    hence [ContingentlyFalse (q_0⁻) in v]
      using cont-true-cont-3[equiv-lr] by simp
    thus ?thesis
      using ∃I by metis
  next
    assume [ContingentlyFalse q_0 in v]
    thus ?thesis
```

```
      using ∃ I by metis
    qed

  lemma cont-tf-thm-5[PLM]:
    [ContingentlyTrue p & Necessary q → p ≠ q in v]
    proof (rule CP)
      assume [ContingentlyTrue p & Necessary q in v]
      hence 1: [◇(¬p) & □ q in v]
        unfolding ContingentlyTrue-def Necessary-defs
        using &E &I by blast
      hence [¬□p in v]
        apply − apply (drule &E(1))
        unfolding diamond-def
        apply (PLM-subst-method ¬¬p p)
        using oth-class-taut-4-b[equiv-sym] by auto
      moreover {
        assume [p = q in v]
        hence [□p in v]
          using l-identity[where α=q and β=p and φ=λ x . □ x,
                           axiom-instance, deduction, deduction]
                1[conj2] id-eq-prop-prop-8-b[deduction]
          by blast
      }
      ultimately show [p ≠ q in v]
        using modus-tollens-1 CP by blast
    qed

  lemma cont-tf-thm-6[PLM]:
    [(ContingentlyFalse p & Impossible q) → p ≠ q in v]
    proof (rule CP)
      assume [ContingentlyFalse p & Impossible q in v]
      hence 1: [◇p & □(¬q) in v]
        unfolding ContingentlyFalse-def Impossible-defs
        using &E &I by blast
      hence [¬◇q in v]
        unfolding diamond-def apply − by PLM-solver
      moreover {
        assume [p = q in v]
        hence [◇q in v]
          using l-identity[axiom-instance, deduction, deduction] 1[conj1]
                id-eq-prop-prop-8-b[deduction]
          by blast
      }
      ultimately show [p ≠ q in v]
        using modus-tollens-1 CP by blast
    qed
end

lemma oa-contingent-1[PLM]:
  [O! ≠ A! in v]
  proof −
    {
      assume [O! = A! in v]
      hence [(λx. ◇(|E!,x^P|)) = (λx. ¬◇(|E!,x^P|)) in v]
        unfolding Ordinary-def Abstract-def .
      moreover have [(|(λx. ◇(|E!,x^P|)), x^P|) ≡ ◇(|E!,x^P|) in v]
        apply (rule beta-C-meta-1)
        by show-proper
```

ultimately have $[(\!|(\boldsymbol{\lambda}x.\ \neg\Diamond(\!|E!,x^P|\!)),\ x^P|\!) \equiv \Diamond(\!|E!,x^P|\!)\ in\ v]$
using *l-identity*[*axiom-instance*, *deduction*, *deduction*] **by** *fast*
moreover have $[(\!|(\boldsymbol{\lambda}x.\ \neg\Diamond(\!|E!,x^P|\!)),\ x^P|\!) \equiv \neg\Diamond(\!|E!,x^P|\!)\ in\ v]$
apply (*rule beta-C-meta-1*)
by *show-proper*
ultimately have $[\Diamond(\!|E!,x^P|\!) \equiv \neg\Diamond(\!|E!,x^P|\!)\ in\ v]$
apply − **by** *PLM-solver*
}
thus *?thesis*
using *oth-class-taut-1-b modus-tollens-1 CP*
by *blast*
qed

lemma *oa-contingent-2*[*PLM*]:
$[(\!|O!,x^P|\!) \equiv \neg(\!|A!,x^P|\!)\ in\ v]$
proof −
have $[(\!|(\boldsymbol{\lambda}x.\ \neg\Diamond(\!|E!,x^P|\!)),\ x^P|\!) \equiv \neg\Diamond(\!|E!,x^P|\!)\ in\ v]$
apply (*rule beta-C-meta-1*)
by *show-proper*
hence $[(\neg(\!|(\boldsymbol{\lambda}x.\ \neg\Diamond(\!|E!,x^P|\!)),\ x^P|\!)) \equiv \Diamond(\!|E!,x^P|\!)\ in\ v]$
using *oth-class-taut-5-d*[*equiv-lr*] *oth-class-taut-4-b*[*equiv-sym*]
$\equiv E(5)$ **by** *blast*
moreover have $[(\!|(\boldsymbol{\lambda}x.\ \Diamond(\!|E!,x^P|\!)),\ x^P|\!) \equiv \Diamond(\!|E!,x^P|\!)\ in\ v]$
apply (*rule beta-C-meta-1*)
by *show-proper*
ultimately show *?thesis*
unfolding *Ordinary-def Abstract-def*
apply − **by** *PLM-solver*
qed

lemma *oa-contingent-3*[*PLM*]:
$[(\!|A!,x^P|\!) \equiv \neg(\!|O!,x^P|\!)\ in\ v]$
using *oa-contingent-2*
apply − **by** *PLM-solver*

lemma *oa-contingent-4*[*PLM*]:
[*Contingent O*! *in v*]
apply (*rule thm-cont-prop-2*[*equiv-rl*], *rule* &*I*)
subgoal
unfolding *Ordinary-def*
apply (*PLM-subst-method* $\lambda\ x\ .\ \Diamond(\!|E!,x^P|\!)\ \lambda\ x\ .\ (\!|\boldsymbol{\lambda}x.\ \Diamond(\!|E!,x^P|\!),x^P|\!)$)
apply (*safe intro*!: *beta-C-meta-1*[*equiv-sym*])
apply *show-proper*
using *BF*$\Diamond$[*deduction*, *OF thm-cont-prop-2*[*equiv-lr*, *OF thm-cont-e-2*, *conj1*]]
by (*rule T*$\Diamond$[*deduction*])
subgoal
apply (*PLM-subst-method* $\lambda\ x\ .\ (\!|A!,x^P|\!)\ \lambda\ x\ .\ \neg(\!|O!,x^P|\!)$)
using *oa-contingent-3* **apply** *simp*
using *cqt-further-5*[*deduction*,*conj1*, *OF A-objects*[*axiom-instance*]]
by (*rule T*$\Diamond$[*deduction*])
done

lemma *oa-contingent-5*[*PLM*]:
[*Contingent A*! *in v*]
apply (*rule thm-cont-prop-2*[*equiv-rl*], *rule* &*I*)
subgoal
using *cqt-further-5*[*deduction*,*conj1*, *OF A-objects*[*axiom-instance*]]
by (*rule T*$\Diamond$[*deduction*])

subgoal
unfolding *Abstract-def*
apply (*PLM-subst-method* $\lambda\ x\ .\ \neg\Diamond(\!|E!,x^P|\!)\ \lambda\ x\ .\ (\!|\boldsymbol{\lambda}x.\ \neg\Diamond(\!|E!,x^P|\!),x^P|\!)$)
apply (*safe intro*!: *beta-C-meta-1*[*equiv-sym*])
apply *show-proper*
apply (*PLM-subst-method* $\lambda\ x\ .\ \Diamond(\!|E!,x^P|\!)\ \lambda\ x\ .\ \neg\neg\Diamond(\!|E!,x^P|\!)$)
using *oth-class-taut-4-b* **apply** *simp*
using *BF*$\Diamond$[*deduction*, *OF thm-cont-prop-2*[*equiv-lr*, *OF thm-cont-e-2*, *conj1*]]
by (*rule T*$\Diamond$[*deduction*])
done

lemma *oa-contingent-6*[*PLM*]:
[$(O!^-) \neq (A!^-)$ *in* v]
proof −
{
assume [$(O!^-) = (A!^-)$ *in* v]
hence [$(\boldsymbol{\lambda}x.\ \neg(\!|O!,x^P|\!)) = (\boldsymbol{\lambda}x.\ \neg(\!|A!,x^P|\!))$ *in* v]
unfolding *propnot-defs* **.**
moreover have [$(\!|(\boldsymbol{\lambda}x.\ \neg(\!|O!,x^P|\!)),\ x^P|\!) \equiv \neg(\!|O!,x^P|\!)$ *in* v]
apply (*rule beta-C-meta-1*)
by *show-proper*
ultimately have [$(\!|\boldsymbol{\lambda}x.\ \neg(\!|A!,x^P|\!),x^P|\!) \equiv\ \neg(\!|O!,x^P|\!)$ *in* v]
using *l-identity*[*axiom-instance*, *deduction*, *deduction*]
by *fast*
hence [$(\neg(\!|A!,x^P|\!)) \equiv \neg(\!|O!,x^P|\!)$ *in* v]
apply −
apply (*PLM-subst-method* $(\!|\boldsymbol{\lambda}x.\ \neg(\!|A!,x^P|\!),x^P|\!)\ (\neg(\!|A!,x^P|\!))$)
apply (*safe intro*!: *beta-C-meta-1*)
by *show-proper*
hence [$(\!|O!,x^P|\!) \equiv \neg(\!|O!,x^P|\!)$ *in* v]
using *oa-contingent-2* **apply** − **by** *PLM-solver*
}
thus *?thesis*
using *oth-class-taut-1-b modus-tollens-1 CP*
by *blast*
qed

lemma *oa-contingent-7*[*PLM*]:
[$(\!|O!^-,x^P|\!) \equiv \neg(\!|A!^-,x^P|\!)$ *in* v]
proof −
have [$(\neg(\!|\boldsymbol{\lambda}x.\ \neg(\!|A!,x^P|\!),x^P|\!)) \equiv (\!|A!,x^P|\!)$ *in* v]
apply (*PLM-subst-method* $(\neg(\!|A!,x^P|\!))\ (\!|\boldsymbol{\lambda}x.\ \neg(\!|A!,x^P|\!),x^P|\!)$)
apply (*safe intro*!: *beta-C-meta-1*[*equiv-sym*])
apply *show-proper*
using *oth-class-taut-4-b*[*equiv-sym*] **by** *auto*
moreover have [$(\!|\boldsymbol{\lambda}x.\ \neg(\!|O!,x^P|\!),x^P|\!) \equiv \neg(\!|O!,x^P|\!)$ *in* v]
apply (*rule beta-C-meta-1*)
by *show-proper*
ultimately show *?thesis*
unfolding *propnot-defs*
using *oa-contingent-3*
apply − **by** *PLM-solver*
qed

lemma *oa-contingent-8*[*PLM*]:
[*Contingent* $(O!^-)$ *in* v]
using *oa-contingent-4 thm-cont-prop-3*[*equiv-lr*] **by** *auto*

lemma *oa-contingent-9*[*PLM*]:
[*Contingent* (*A*!$^-$) *in v*]
using *oa-contingent-5 thm-cont-prop-3*[*equiv-lr*] **by** *auto*

lemma *oa-facts-1*[*PLM*]:
[(|*O*!,x^P|) → □(|*O*!,x^P|) *in v*]
proof (*rule CP*)
assume [(|*O*!,x^P|) *in v*]
hence [◇(|*E*!,x^P|) *in v*]
unfolding *Ordinary-def* **apply** −
apply (*rule beta-C-meta-1*[*equiv-lr*])
by *show-proper*
hence [□◇(|*E*!,x^P|) *in v*]
using *qml-3*[*axiom-instance*, *deduction*] **by** *auto*
thus [□(|*O*!,x^P|) *in v*]
unfolding *Ordinary-def*
apply −
apply (*PLM-subst-method* ◇(|*E*!,x^P|) (|**λ***x*. ◇(|*E*!,x^P|),x^P|))
apply (*safe intro*!: *beta-C-meta-1*[*equiv-sym*])
by *show-proper*
qed

lemma *oa-facts-2*[*PLM*]:
[(|*A*!,x^P|) → □(|*A*!,x^P|) *in v*]
proof (*rule CP*)
assume [(|*A*!,x^P|) *in v*]
hence [¬◇(|*E*!,x^P|) *in v*]
unfolding *Abstract-def* **apply** −
apply (*rule beta-C-meta-1*[*equiv-lr*])
by *show-proper*
hence [□□¬(|*E*!,x^P|) *in v*]
using *KBasic2-4*[*equiv-rl*] *4*□[*deduction*] **by** *auto*
hence [□¬◇(|*E*!,x^P|) *in v*]
apply −
apply (*PLM-subst-method* □¬(|*E*!,x^P|) ¬◇(|*E*!,x^P|))
using *KBasic2-4* **by** *auto*
thus [□(|*A*!,x^P|) *in v*]
unfolding *Abstract-def*
apply −
apply (*PLM-subst-method* ¬◇(|*E*!,x^P|) (|**λ***x*. ¬◇(|*E*!,x^P|),x^P|))
apply (*safe intro*!: *beta-C-meta-1*[*equiv-sym*])
by *show-proper*
qed

lemma *oa-facts-3*[*PLM*]:
[◇(|*O*!,x^P|) → (|*O*!,x^P|) *in v*]
using *oa-facts-1* **by** (*rule derived-S5-rules-2-b*)

lemma *oa-facts-4*[*PLM*]:
[◇(|*A*!,x^P|) → (|*A*!,x^P|) *in v*]
using *oa-facts-2* **by** (*rule derived-S5-rules-2-b*)

lemma *oa-facts-5*[*PLM*]:
[◇(|*O*!,x^P|) ≡ □(|*O*!,x^P|) *in v*]
using *oa-facts-1*[*deduction*, *OF oa-facts-3*[*deduction*]]
T◇[*deduction*, *OF qml-2*[*axiom-instance*, *deduction*]]
≡*I CP* **by** *blast*

lemma *oa-facts-6*[*PLM*]:
[$\Diamond(\!|A!,x^P|\!) \equiv \Box(\!|A!,x^P|\!)$ *in* v]
using *oa-facts-2*[*deduction*, *OF oa-facts-4*[*deduction*]]
T$\Diamond$[*deduction*, *OF qml-2*[*axiom-instance*, *deduction*]]
$\equiv I$ *CP* **by** *blast*

lemma *oa-facts-7*[*PLM*]:
[$(\!|O!,x^P|\!) \equiv \boldsymbol{\mathcal{A}}(\!|O!,x^P|\!)$ *in* v]
apply (*rule* $\equiv I$; *rule CP*)
apply (*rule nec-imp-act*[*deduction*, *OF oa-facts-1*[*deduction*]]; *assumption*)
proof −
assume [$\boldsymbol{\mathcal{A}}(\!|O!,x^P|\!)$ *in* v]
hence [$\boldsymbol{\mathcal{A}}(\Diamond(\!|E!,x^P|\!))$ *in* v]
unfolding *Ordinary-def* **apply** −
apply (*PLM-subst-method* $(\!|\boldsymbol{\lambda}x.\ \Diamond(\!|E!,x^P|\!),x^P|\!)$ $\Diamond(\!|E!,x^P|\!)$)
apply (*safe intro*!: *beta-C-meta-1*)
by *show-proper*
hence [$\Diamond(\!|E!,x^P|\!)$ *in* v]
using *Act-Basic-6*[*equiv-rl*] **by** *auto*
thus [$(\!|O!,x^P|\!)$ *in* v]
unfolding *Ordinary-def* **apply** −
apply (*PLM-subst-method* $\Diamond(\!|E!,x^P|\!)$ $(\!|\boldsymbol{\lambda}x.\ \Diamond(\!|E!,x^P|\!),x^P|\!)$)
apply (*safe intro*!: *beta-C-meta-1*[*equiv-sym*])
by *show-proper*
qed

lemma *oa-facts-8*[*PLM*]:
[$(\!|A!,x^P|\!) \equiv \boldsymbol{\mathcal{A}}(\!|A!,x^P|\!)$ *in* v]
apply (*rule* $\equiv I$; *rule CP*)
apply (*rule nec-imp-act*[*deduction*, *OF oa-facts-2*[*deduction*]]; *assumption*)
proof −
assume [$\boldsymbol{\mathcal{A}}(\!|A!,x^P|\!)$ *in* v]
hence [$\boldsymbol{\mathcal{A}}(\neg\Diamond(\!|E!,x^P|\!))$ *in* v]
unfolding *Abstract-def* **apply** −
apply (*PLM-subst-method* $(\!|\boldsymbol{\lambda}x.\ \neg\Diamond(\!|E!,x^P|\!),x^P|\!)$ $\neg\Diamond(\!|E!,x^P|\!)$)
apply (*safe intro*!: *beta-C-meta-1*)
by *show-proper*
hence [$\boldsymbol{\mathcal{A}}(\Box\neg(\!|E!,x^P|\!))$ *in* v]
apply −
apply (*PLM-subst-method* $(\neg\Diamond(\!|E!,x^P|\!))$ $(\Box\neg(\!|E!,x^P|\!))$)
using *KBasic2-4*[*equiv-sym*] **by** *auto*
hence [$\neg\Diamond(\!|E!,x^P|\!)$ *in* v]
using *qml-act-2*[*axiom-instance*, *equiv-rl*] *KBasic2-4*[*equiv-lr*] **by** *auto*
thus [$(\!|A!,x^P|\!)$ *in* v]
unfolding *Abstract-def* **apply** −
apply (*PLM-subst-method* $\neg\Diamond(\!|E!,x^P|\!)$ $(\!|\boldsymbol{\lambda}x.\ \neg\Diamond(\!|E!,x^P|\!),x^P|\!)$)
apply (*safe intro*!: *beta-C-meta-1*[*equiv-sym*])
by *show-proper*
qed

lemma *cont-nec-fact1-1*[*PLM*]:
[*WeaklyContingent* $F \equiv$ *WeaklyContingent* (F^-) *in* v]
proof (*rule* $\equiv I$; *rule CP*)
assume [*WeaklyContingent* F *in* v]
hence *wc-def*: [*Contingent* F & $(\forall\ x\ .\ (\Diamond(\!|F,x^P|\!) \rightarrow \Box(\!|F,x^P|\!)))$ *in* v]
unfolding *WeaklyContingent-def* **.**
have [*Contingent* (F^-) *in* v]
using *wc-def*[*conj1*] **by** (*rule thm-cont-prop-3*[*equiv-lr*])

```
  moreover {
   {
    fix x
    assume [◇(|F⁻,x^P|) in v]
    hence [¬□(|F,x^P|) in v]
     unfolding diamond-def apply −
     apply (PLM-subst-method ¬(|F⁻,x^P|) (|F,x^P|))
      using thm-relation-negation-2-1 by auto
    moreover {
     assume [¬□(|F⁻,x^P|) in v]
     hence [¬□(|λx. ¬(|F,x^P|),x^P|) in v]
      unfolding propnot-defs .
     hence [◇(|F,x^P|) in v]
      unfolding diamond-def
      apply − apply (PLM-subst-method (|λx. ¬(|F,x^P|),x^P|) ¬(|F,x^P|))
      apply (safe intro!: beta-C-meta-1)
      by show-proper
     hence [□(|F,x^P|) in v]
      using wc-def[conj2] cqt-1[axiom-instance, deduction]
            modus-ponens by fast
    }
    ultimately have [□(|F⁻, x^P|) in v]
     using ¬¬E modus-tollens-1 CP by blast
   }
   hence [∀ x . ◇(|F⁻,x^P|) → □(|F⁻, x^P|) in v]
    using ∀I CP by fast
  }
  ultimately show [WeaklyContingent (F⁻) in v]
   unfolding WeaklyContingent-def by (rule &I)
next
  assume [WeaklyContingent (F⁻) in v]
  hence wc-def: [Contingent (F⁻) & (∀ x . (◇(|F⁻,x^P|) → □(|F⁻,x^P|))) in v]
   unfolding WeaklyContingent-def .
  have [Contingent F in v]
   using wc-def[conj1] by (rule thm-cont-prop-3[equiv-rl])
  moreover {
   {
    fix x
    assume [◇(|F,x^P|) in v]
    hence [¬□(|F⁻,x^P|) in v]
     unfolding diamond-def apply −
     apply (PLM-subst-method ¬(|F,x^P|) (|F⁻,x^P|))
     using thm-relation-negation-1-1[equiv-sym] by auto
    moreover {
     assume [¬□(|F,x^P|) in v]
     hence [◇(|F⁻,x^P|) in v]
      unfolding diamond-def
      apply − apply (PLM-subst-method (|F,x^P|) ¬(|F⁻,x^P|))
      using thm-relation-negation-2-1[equiv-sym] by auto
     hence [□(|F⁻,x^P|) in v]
      using wc-def[conj2] cqt-1[axiom-instance, deduction]
            modus-ponens by fast
    }
    ultimately have [□(|F, x^P|) in v]
     using ¬¬E modus-tollens-1 CP by blast
   }
   hence [∀ x . ◇(|F,x^P|) → □(|F, x^P|) in v]
    using ∀I CP by fast
```

```
    }
    ultimately show [WeaklyContingent (F) in v]
      unfolding WeaklyContingent-def by (rule &I)
  qed

lemma cont-nec-fact1-2[PLM]:
  [(WeaklyContingent F & ¬(WeaklyContingent G)) → (F ≠ G) in v]
  using l-identity[axiom-instance,deduction,deduction] &E &I
        modus-tollens-1 CP by metis

lemma cont-nec-fact2-1[PLM]:
  [WeaklyContingent (O!) in v]
  unfolding WeaklyContingent-def
  apply (rule &I)
   using oa-contingent-4 apply simp
  using oa-facts-5 unfolding equiv-def
  using &E(1) ∀ I by fast

lemma cont-nec-fact2-2[PLM]:
  [WeaklyContingent (A!) in v]
  unfolding WeaklyContingent-def
  apply (rule &I)
   using oa-contingent-5 apply simp
  using oa-facts-6 unfolding equiv-def
  using &E(1) ∀ I by fast

lemma cont-nec-fact2-3[PLM]:
  [¬(WeaklyContingent (E!)) in v]
  proof (rule modus-tollens-1, rule CP)
    assume [WeaklyContingent E! in v]
    thus [∀ x . ◇(|E!,x^P|) → □(|E!,x^P|) in v]
    unfolding WeaklyContingent-def using &E(2) by fast
  next
    {
      assume 1: [∀ x . ◇(|E!,x^P|) → □(|E!,x^P|) in v]
      have [∃ x . ◇((|E!,x^P|) & ◇(¬(|E!,x^P|))) in v]
        using qml-4[axiom-instance,conj1, THEN BFs-3[deduction]] .
      then obtain x where [◇((|E!,x^P|) & ◇(¬(|E!,x^P|))) in v]
        by (rule ∃ E)
      hence [◇(|E!,x^P|) & ◇(¬(|E!,x^P|)) in v]
        using KBasic2-8[deduction] S5Basic-8[deduction]
              &I &E by blast
      hence [□(|E!,x^P|) & (¬□(|E!,x^P|)) in v]
        using 1[THEN ∀ E, deduction] &E &I
              KBasic2-2[equiv-rl] by blast
      hence [¬(∀ x . ◇(|E!,x^P|) → □(|E!,x^P|)) in v]
        using oth-class-taut-1-a modus-tollens-1 CP by blast
    }
    thus [¬(∀ x . ◇(|E!,x^P|) → □(|E!,x^P|)) in v]
      using reductio-aa-2 if-p-then-p CP by meson
  qed

lemma cont-nec-fact2-4[PLM]:
  [¬(WeaklyContingent (PLM.L)) in v]
  proof −
    {
      assume [WeaklyContingent PLM.L in v]
      hence [Contingent PLM.L in v]
```

unfolding *WeaklyContingent-def* **using** &$E(1)$ **by** *blast*
}
thus *?thesis*
using *thm-noncont-e-e-3*
unfolding *Contingent-def NonContingent-def*
using *modus-tollens-2 CP* **by** *blast*
qed

lemma *cont-nec-fact2-5*[*PLM*]:
[$O! \neq E!$ & $O! \neq (E!^{-})$ & $O! \neq PLM.L$ & $O! \neq (PLM.L^{-})$ *in* v]
proof ((*rule* &I)+)
show [$O! \neq E!$ *in* v]
using *cont-nec-fact2-1 cont-nec-fact2-3*
cont-nec-fact1-2[*deduction*] &I **by** *simp*
next
have [$\neg$(*WeaklyContingent* $(E!^{-})$) *in* v]
using *cont-nec-fact1-1*[*THEN oth-class-taut-5-d*[*equiv-lr*], *equiv-lr*]
cont-nec-fact2-3 **by** *auto*
thus [$O! \neq (E!^{-})$ *in* v]
using *cont-nec-fact2-1 cont-nec-fact1-2*[*deduction*] &I **by** *simp*
next
show [$O! \neq PLM.L$ *in* v]
using *cont-nec-fact2-1 cont-nec-fact2-4*
cont-nec-fact1-2[*deduction*] &I **by** *simp*
next
have [$\neg$(*WeaklyContingent* $(PLM.L^{-})$) *in* v]
using *cont-nec-fact1-1*[*THEN oth-class-taut-5-d*[*equiv-lr*], *equiv-lr*]
cont-nec-fact2-4 **by** *auto*
thus [$O! \neq (PLM.L^{-})$ *in* v]
using *cont-nec-fact2-1 cont-nec-fact1-2*[*deduction*] &I **by** *simp*
qed

lemma *cont-nec-fact2-6*[*PLM*]:
[$A! \neq E!$ & $A! \neq (E!^{-})$ & $A! \neq PLM.L$ & $A! \neq (PLM.L^{-})$ *in* v]
proof ((*rule* &I)+)
show [$A! \neq E!$ *in* v]
using *cont-nec-fact2-2 cont-nec-fact2-3*
cont-nec-fact1-2[*deduction*] &I **by** *simp*
next
have [$\neg$(*WeaklyContingent* $(E!^{-})$) *in* v]
using *cont-nec-fact1-1*[*THEN oth-class-taut-5-d*[*equiv-lr*], *equiv-lr*]
cont-nec-fact2-3 **by** *auto*
thus [$A! \neq (E!^{-})$ *in* v]
using *cont-nec-fact2-2 cont-nec-fact1-2*[*deduction*] &I **by** *simp*
next
show [$A! \neq PLM.L$ *in* v]
using *cont-nec-fact2-2 cont-nec-fact2-4*
cont-nec-fact1-2[*deduction*] &I **by** *simp*
next
have [$\neg$(*WeaklyContingent* $(PLM.L^{-})$) *in* v]
using *cont-nec-fact1-1*[*THEN oth-class-taut-5-d*[*equiv-lr*],
equiv-lr] *cont-nec-fact2-4* **by** *auto*
thus [$A! \neq (PLM.L^{-})$ *in* v]
using *cont-nec-fact2-2 cont-nec-fact1-2*[*deduction*] &I **by** *simp*
qed

lemma *id-nec3-1*[*PLM*]:
[$((x^{P}) =_{E} (y^{P})) \equiv (\Box((x^{P}) =_{E} (y^{P})))$ *in* v]

proof (*rule* ≡*I*; *rule CP*)
assume $[(x^P) =_E (y^P)\ in\ v]$
hence $[(\!|O!,x^P|\!)\ in\ v] \wedge [(\!|O!,y^P|\!)\ in\ v] \wedge [\Box(\forall\ F\ .\ (\!|F,x^P|\!) \equiv (\!|F,y^P|\!))\ in\ v]$
using *eq-E-simple-1*[*equiv-lr*] **using** &*E* **by** *blast*
hence $[\Box(\!|O!,x^P|\!)\ in\ v] \wedge [\Box(\!|O!,y^P|\!)\ in\ v]$
$\wedge\ [\Box\Box(\forall\ F\ .\ (\!|F,x^P|\!) \equiv (\!|F,y^P|\!))\ in\ v]$
using *oa-facts-1*[*deduction*] *S5Basic-6*[*deduction*] **by** *blast*
hence $[\Box((\!|O!,x^P|\!)\ \&\ (\!|O!,y^P|\!)\ \&\ \Box(\forall\ F.\ (\!|F,x^P|\!) \equiv (\!|F,y^P|\!)))\ in\ v]$
using &*I KBasic-3*[*equiv-rl*] **by** *presburger*
thus $[\Box((x^P) =_E (y^P))\ in\ v]$
apply −
apply (*PLM-subst-method*
$(\!|O!,x^P|\!)\ \&\ (\!|O!,y^P|\!)\ \&\ \Box(\forall\ F.\ (\!|F,x^P|\!) \equiv (\!|F,y^P|\!))$
$(x^P) =_E (y^P)$)
using *eq-E-simple-1*[*equiv-sym*] **by** *auto*
next
assume $[\Box((x^P) =_E (y^P))\ in\ v]$
thus $[((x^P) =_E (y^P))\ in\ v]$
using *qml-2*[*axiom-instance*,*deduction*] **by** *simp*
qed

lemma *id-nec3-2*[*PLM*]:
$[\Diamond((x^P) =_E (y^P)) \equiv ((x^P) =_E (y^P))\ in\ v]$
proof (*rule* ≡*I*; *rule CP*)
assume $[\Diamond((x^P) =_E (y^P))\ in\ v]$
thus $[(x^P) =_E (y^P)\ in\ v]$
using *derived-S5-rules-2-b*[*deduction*] *id-nec3-1*[*equiv-lr*]
CP modus-ponens **by** *blast*
next
assume $[(x^P) =_E (y^P)\ in\ v]$
thus $[\Diamond((x^P) =_E (y^P))\ in\ v]$
by (*rule TBasic*[*deduction*])
qed

lemma *thm-neg-eqE*[*PLM*]:
$[((x^P) \neq_E (y^P)) \equiv (\neg((x^P) =_E (y^P)))\ in\ v]$
proof −
have $[(x^P) \neq_E (y^P)\ in\ v] = [(\!|(\boldsymbol{\lambda}^2\ (\lambda\ x\ y\ .\ (x^P) =_E (y^P)))^-,\ x^P,\ y^P|\!)\ in\ v]$
unfolding *not-identical*$_E$*-def* **by** *simp*
also have ... $= [\neg(\!|(\boldsymbol{\lambda}^2\ (\lambda\ x\ y\ .\ (x^P) =_E (y^P))),\ x^P,\ y^P|\!)\ in\ v]$
unfolding *propnot-defs*
apply (*safe intro*!: *beta-C-meta-2*[*equiv-lr*] *beta-C-meta-2*[*equiv-rl*])
by *show-proper*+
also have ... $= [\neg((x^P) =_E (y^P))\ in\ v]$
apply (*PLM-subst-method*
$(\!|(\boldsymbol{\lambda}^2\ (\lambda\ x\ y\ .\ (x^P) =_E (y^P))),\ x^P,\ y^P|\!)$
$(x^P) =_E (y^P)$)
apply (*safe intro*!: *beta-C-meta-2*)
unfolding *identity-defs* **by** *show-proper*
finally show *?thesis*
using ≡*I CP* **by** *presburger*
qed

lemma *id-nec4-1*[*PLM*]:
$[((x^P) \neq_E (y^P)) \equiv \Box((x^P) \neq_E (y^P))\ in\ v]$
proof −
have $[(\neg((x^P) =_E (y^P))) \equiv \Box(\neg((x^P) =_E (y^P)))\ in\ v]$
using *id-nec3-2*[*equiv-sym*] *oth-class-taut-5-d*[*equiv-lr*]

KBasic2-4[*equiv-sym*] *intro-elim-6-e* **by** *fast*
thus *?thesis*
apply −
apply (*PLM-subst-method* (¬((x^P) =_E (y^P))) (x^P) ≠_E (y^P))
using *thm-neg-eqE*[*equiv-sym*] **by** *auto*
qed

lemma *id-nec4-2*[*PLM*]:
[◇((x^P) ≠_E (y^P)) ≡ ((x^P) ≠_E (y^P)) *in v*]
using ≡*I* *id-nec4-1*[*equiv-lr*] *derived-S5-rules-2-b CP T*◇ **by** *simp*

lemma *id-act-1*[*PLM*]:
[((x^P) =_E (y^P)) ≡ (𝒜((x^P) =_E (y^P))) *in v*]
proof (*rule* ≡*I*; *rule CP*)
assume [(x^P) =_E (y^P) *in v*]
hence [□((x^P) =_E (y^P)) *in v*]
using *id-nec3-1*[*equiv-lr*] **by** *auto*
thus [𝒜((x^P) =_E (y^P)) *in v*]
using *nec-imp-act*[*deduction*] **by** *fast*
next
assume [𝒜((x^P) =_E (y^P)) *in v*]
hence [𝒜(⌈O!,x^P⌋ & ⌈O!,y^P⌋ & □(∀ F . ⌈F,x^P⌋ ≡ ⌈F,y^P⌋)) *in v*]
apply −
apply (*PLM-subst-method*
(x^P) =_E (y^P)
(⌈O!,x^P⌋ & ⌈O!,y^P⌋ & □(∀ F . ⌈F,x^P⌋ ≡ ⌈F,y^P⌋)))
using *eq-E-simple-1* **by** *auto*
hence [𝒜⌈O!,x^P⌋ & 𝒜⌈O!,y^P⌋ & 𝒜(□(∀ F . ⌈F,x^P⌋ ≡ ⌈F,y^P⌋)) *in v*]
using *Act-Basic-2*[*equiv-lr*] &*I* &*E* **by** *meson*
thus [(x^P) =_E (y^P) *in v*]
apply − **apply** (*rule eq-E-simple-1*[*equiv-rl*])
using *oa-facts-7*[*equiv-rl*] *qml-act-2*[*axiom-instance*, *equiv-rl*]
&*I* &*E* **by** *meson*
qed

lemma *id-act-2*[*PLM*]:
[((x^P) ≠_E (y^P)) ≡ (𝒜((x^P) ≠_E (y^P))) *in v*]
apply (*PLM-subst-method* (¬((x^P) =_E (y^P))) ((x^P) ≠_E (y^P)))
using *thm-neg-eqE*[*equiv-sym*] **apply** *simp*
using *id-act-1 oth-class-taut-5-d*[*equiv-lr*] *thm-neg-eqE intro-elim-6-e*
logic-actual-nec-1[*axiom-instance*,*equiv-sym*] **by** *meson*

end

class *id-act* = *id-eq* +
assumes *id-act-prop*: [𝒜(α = β) *in v*] ⟹ [(α = β) *in v*]

instantiation ν :: *id-act*
begin
instance proof
interpret *PLM* .
fix *x*::ν **and** *y*::ν **and** *v*::*i*
assume [𝒜(x = y) *in v*]
hence [𝒜(((x^P) =_E (y^P)) ∨ (⌈A!,x^P⌋ & ⌈A!,y^P⌋
& □(∀ F . ⦃x^P,F⦄ ≡ ⦃y^P,F⦄))) *in v*]
unfolding *identity-defs* **by** *auto*
hence [𝒜(((x^P) =_E (y^P))) ∨ 𝒜((⌈A!,x^P⌋ & ⌈A!,y^P⌋
& □(∀ F . ⦃x^P,F⦄ ≡ ⦃y^P,F⦄))) *in v*]

```
    using Act-Basic-10[equiv-lr] by auto
  moreover {
    assume [𝓐(((x^P) =_E (y^P))) in v]
    hence [(x^P) = (y^P) in v]
     using id-act-1[equiv-rl] eq-E-simple-2[deduction] by auto
  }
  moreover {
    assume [𝓐((⌈A!,x^P⌉) & (⌈A!,y^P⌉) & □(∀ F . {|x^P,F|} ≡ {|y^P,F|})) in v]
    hence [𝓐(⌈A!,x^P⌉) & 𝓐(⌈A!,y^P⌉) & 𝓐(□(∀ F . {|x^P,F|} ≡ {|y^P,F|})) in v]
      using Act-Basic-2[equiv-lr] &I &E by meson
    hence [(⌈A!,x^P⌉) & (⌈A!,y^P⌉) & (□(∀ F . {|x^P,F|} ≡ {|y^P,F|})) in v]
      using oa-facts-8[equiv-rl] qml-act-2[axiom-instance,equiv-rl]
        &I &E by meson
    hence [(x^P) = (y^P) in v]
     unfolding identity-defs using ∨I by auto
  }
  ultimately have [(x^P) = (y^P) in v]
    using intro-elim-4-a CP by meson
  thus [x = y in v]
    unfolding identity-defs by auto
 qed
end

instantiation Π1 :: id-act
begin
 instance proof
  interpret PLM .
  fix F::Π1 and G::Π1 and v::i
  show [𝓐(F = G) in v] ⟹ [(F = G) in v]
    unfolding identity-defs
    using qml-act-2[axiom-instance,equiv-rl] by auto
 qed
end

instantiation o :: id-act
begin
 instance proof
  interpret PLM .
  fix p :: o and q :: o and v::i
  show [𝓐(p = q) in v] ⟹ [p = q in v]
    unfolding identity_o-def using id-act-prop by blast
 qed
end

instantiation Π2 :: id-act
begin
 instance proof
  interpret PLM .
  fix F::Π2 and G::Π2 and v::i
  assume a: [𝓐(F = G) in v]
  {
    fix x
    have [𝓐((λy. (⌈F,x^P,y^P⌉)) = (λy. (⌈G,x^P,y^P⌉))
          & (λy. (⌈F,y^P,x^P⌉)) = (λy. (⌈G,y^P,x^P⌉))) in v]
      using a logic-actual-nec-3[axiom-instance, equiv-lr] cqt-basic-4[equiv-lr] ∀E
      unfolding identity_2-def by fast
    hence [((λy. (⌈F,x^P,y^P⌉)) = (λy. (⌈G,x^P,y^P⌉)))
          & ((λy. (⌈F,y^P,x^P⌉)) = (λy. (⌈G,y^P,x^P⌉))) in v]
```

```
      using &I &E id-act-prop Act-Basic-2[equiv-lr] by metis
    }
    thus [F = G in v] unfolding identity-defs by (rule ∀ I)
  qed
end

instantiation Π3 :: id-act
begin
  instance proof
    interpret PLM .
    fix F::Π3 and G::Π3 and v::i
    assume a: [𝒜(F = G) in v]
    let ?p = λ x y . (λz. (|F,z^P,x^P,y^P|)) = (λz. (|G,z^P,x^P,y^P|))
                   & (λz. (|F,x^P,z^P,y^P|)) = (λz. (|G,x^P,z^P,y^P|))
                   & (λz. (|F,x^P,y^P,z^P|)) = (λz. (|G,x^P,y^P,z^P|))
    {
      fix x
      {
        fix y
        have [𝒜(?p x y) in v]
          using a logic-actual-nec-3[axiom-instance, equiv-lr]
                cqt-basic-4[equiv-lr] ∀E[where 'a=ν]
          unfolding identity3-def by blast
        hence [?p x y in v]
          using &I &E id-act-prop Act-Basic-2[equiv-lr] by metis
      }
      hence [∀ y . ?p x y in v]
        by (rule ∀ I)
    }
    thus [F = G in v]
      unfolding identity3-def by (rule ∀ I)
  qed
end

context PLM
begin
  lemma id-act-3[PLM]:
    [((α::('a::id-act)) = β) ≡ 𝒜(α = β) in v]
    using ≡I CP id-nec[equiv-lr, THEN nec-imp-act[deduction]]
          id-act-prop by metis

  lemma id-act-4[PLM]:
    [((α::('a::id-act)) ≠ β) ≡ 𝒜(α ≠ β) in v]
    using id-act-3[THEN oth-class-taut-5-d[equiv-lr]]
          logic-actual-nec-1[axiom-instance, equiv-sym]
          intro-elim-6-e by blast

  lemma id-act-desc[PLM]:
    [(y^P) = (ιx . x = y) in v]
    using descriptions[axiom-instance,equiv-rl]
          id-act-3[equiv-sym] ∀ I by fast

  lemma eta-conversion-lemma-1[PLM]:
    [(λ x . (|F,x^P|)) = F in v]
    using lambda-predicates-3-1[axiom-instance] .

  lemma eta-conversion-lemma-0[PLM]:
    [(λ^0 p) = p in v]
```

using *lambda-predicates-3-0*[*axiom-instance*] **.**

lemma *eta-conversion-lemma-2*[*PLM*]:
[($\boldsymbol{\lambda}^2$ (λ *x y* . (|*F*,x^P,y^P|))) = *F in v*]
using *lambda-predicates-3-2*[*axiom-instance*] **.**

lemma *eta-conversion-lemma-3*[*PLM*]:
[($\boldsymbol{\lambda}^3$ (λ *x y z* . (|*F*,x^P,y^P,z^P|))) = *F in v*]
using *lambda-predicates-3-3*[*axiom-instance*] **.**

lemma *lambda-p-q-p-eq-q*[*PLM*]:
[(($\boldsymbol{\lambda}^0$ *p*) = ($\boldsymbol{\lambda}^0$ *q*)) ≡ (*p* = *q*) *in v*]
using *eta-conversion-lemma-0*
l-identity[*axiom-instance*, *deduction*, *deduction*]
eta-conversion-lemma-0[*eq-sym*] ≡*I CP*
by *metis*

A.9.12. The Theory of Objects

lemma *partition-1*[*PLM*]:
[∀ *x* . (|*O*!,x^P|) ∨ (|*A*!,x^P|) *in v*]
proof (*rule* ∀*I*)
fix *x*
have [◇(|*E*!,x^P|) ∨ ¬◇(|*E*!,x^P|) *in v*]
by *PLM-solver*
moreover have [◇(|*E*!,x^P|) ≡ (|$\boldsymbol{\lambda}$ *y* . ◇(|*E*!,y^P|), x^P|) *in v*]
apply (*rule beta-C-meta-1*[*equiv-sym*])
by *show-proper*
moreover have [(¬◇(|*E*!,x^P|)) ≡ (|$\boldsymbol{\lambda}$ *y* . ¬◇(|*E*!,y^P|), x^P|) *in v*]
apply (*rule beta-C-meta-1*[*equiv-sym*])
by *show-proper*
ultimately show [(|*O*!, x^P|) ∨ (|*A*!, x^P|) *in v*]
unfolding *Ordinary-def Abstract-def* **by** *PLM-solver*
qed

lemma *partition-2*[*PLM*]:
[¬(∃ *x* . (|*O*!,x^P|) & (|*A*!,x^P|)) *in v*]
proof −
{
assume [∃ *x* . (|*O*!,x^P|) & (|*A*!,x^P|) *in v*]
then obtain *b* **where** [(|*O*!,b^P|) & (|*A*!,b^P|) *in v*]
by (*rule* ∃*E*)
hence *?thesis*
using &*E oa-contingent-2*[*equiv-lr*]
reductio-aa-2 **by** *fast*
}
thus *?thesis*
using *reductio-aa-2* **by** *blast*
qed

lemma *ord-eq-Eequiv-1*[*PLM*]:
[(|*O*!,*x*|) → (*x* =$_E$ *x*) *in v*]
proof (*rule CP*)
assume [(|*O*!,*x*|) *in v*]
moreover have [□(∀ *F* . (|*F*,*x*|) ≡ (|*F*,*x*|)) *in v*]
by *PLM-solver*
ultimately show [(*x*) =$_E$ (*x*) *in v*]
using &*I eq-E-simple-1*[*equiv-rl*] **by** *blast*

qed

lemma *ord-eq-Eequiv-2*[*PLM*]:
[$(x =_E y) \rightarrow (y =_E x)$ *in* v]
proof (*rule CP*)
assume [$x =_E y$ *in* v]
hence *1*: [$(\!|O!,x|\!)$ & $(\!|O!,y|\!)$ & $\Box(\forall\ F\ .\ (\!|F,x|\!) \equiv (\!|F,y|\!))$ *in* v]
using *eq-E-simple-1*[*equiv-lr*] **by** *simp*
have [$\Box(\forall\ F\ .\ (\!|F,y|\!) \equiv (\!|F,x|\!))$ *in* v]
apply (*PLM-subst-method*
$\lambda\ F\ .\ (\!|F,x|\!) \equiv (\!|F,y|\!)$
$\lambda\ F\ .\ (\!|F,y|\!) \equiv (\!|F,x|\!)$)
using *oth-class-taut-3-g 1*[*conj2*] **by** *auto*
thus [$y =_E x$ *in* v]
using *eq-E-simple-1*[*equiv-rl*] *1*[*conj1*]
&*E* &*I* **by** *meson*
qed

lemma *ord-eq-Eequiv-3*[*PLM*]:
[$((x =_E y)$ & $(y =_E z)) \rightarrow (x =_E z)$ *in* v]
proof (*rule CP*)
assume *a*: [$(x =_E y)$ & $(y =_E z)$ *in* v]
have [$\Box((\forall\ F\ .\ (\!|F,x|\!) \equiv (\!|F,y|\!))$ & $(\forall\ F\ .\ (\!|F,y|\!) \equiv (\!|F,z|\!)))$ *in* v]
using *KBasic-3*[*equiv-rl*] *a*[*conj1*, *THEN eq-E-simple-1*[*equiv-lr*,*conj2*]]
a[*conj2*, *THEN eq-E-simple-1*[*equiv-lr*,*conj2*]] &*I* **by** *blast*
moreover {
{
fix *w*
have [$((\forall\ F\ .\ (\!|F,x|\!) \equiv (\!|F,y|\!))$ & $(\forall\ F\ .\ (\!|F,y|\!) \equiv (\!|F,z|\!)))$
$\rightarrow (\forall\ F\ .\ (\!|F,x|\!) \equiv (\!|F,z|\!))$ *in* w]
by *PLM-solver*
}
hence [$\Box(((\forall\ F\ .\ (\!|F,x|\!) \equiv (\!|F,y|\!))$ & $(\forall\ F\ .\ (\!|F,y|\!) \equiv (\!|F,z|\!)))$
$\rightarrow (\forall\ F\ .\ (\!|F,x|\!) \equiv (\!|F,z|\!)))$ *in* v]
by (*rule RN*)
}
ultimately have [$\Box(\forall\ F\ .\ (\!|F,x|\!) \equiv (\!|F,z|\!))$ *in* v]
using *qml-1*[*axiom-instance*,*deduction*,*deduction*] **by** *blast*
thus [$x =_E z$ *in* v]
using *a*[*conj1*, *THEN eq-E-simple-1*[*equiv-lr*,*conj1*,*conj1*]]
using *a*[*conj2*, *THEN eq-E-simple-1*[*equiv-lr*,*conj1*,*conj2*]]
eq-E-simple-1[*equiv-rl*] &*I*
by *presburger*
qed

lemma *ord-eq-E-eq*[*PLM*]:
[$((\!|O!,x^P|\!) \vee (\!|O!,y^P|\!)) \rightarrow ((x^P = y^P) \equiv (x^P =_E y^P))$ *in* v]
proof (*rule CP*)
assume [$(\!|O!,x^P|\!) \vee (\!|O!,y^P|\!)$ *in* v]
moreover {
assume [$(\!|O!,x^P|\!)$ *in* v]
hence [$(x^P = y^P) \equiv (x^P =_E y^P)$ *in* v]
using $\equiv$*I CP l-identity*[*axiom-instance*, *deduction*, *deduction*]
ord-eq-Eequiv-1[*deduction*] *eq-E-simple-2*[*deduction*] **by** *metis*
}
moreover {
assume [$(\!|O!,y^P|\!)$ *in* v]
hence [$(x^P = y^P) \equiv (x^P =_E y^P)$ *in* v]

using $\equiv I$ *CP l-identity*[*axiom-instance*, *deduction*, *deduction*]
ord-eq-Eequiv-1[*deduction*] *eq-E-simple-2*[*deduction*] *id-eq-2*[*deduction*]
ord-eq-Eequiv-2[*deduction*] *identity-ν-def* **by** *metis*
}
ultimately show $[(x^P = y^P) \equiv (x^P =_E y^P)\ in\ v]$
using *intro-elim-4-a CP* **by** *blast*
qed

lemma *ord-eq-E*[*PLM*]:
$[((|O!,x^P|)\ \&\ (|O!,y^P|)) \rightarrow ((\forall\ F\ .\ (|F,x^P|) \equiv (|F,y^P|)) \rightarrow x^P =_E y^P)\ in\ v]$
proof (*rule CP*; *rule CP*)
assume *ord-xy*: $[(|O!,x^P|)\ \&\ (|O!,y^P|)\ in\ v]$
assume $[\forall\ F\ .\ (|F,x^P|) \equiv (|F,y^P|)\ in\ v]$
hence $[(|\boldsymbol{\lambda}\ z\ .\ z^P =_E x^P,\ x^P|) \equiv (|\boldsymbol{\lambda}\ z\ .\ z^P =_E x^P,\ y^P|)\ in\ v]$
by (*rule* $\forall E$)
moreover have $[(|\boldsymbol{\lambda}\ z\ .\ z^P =_E x^P,\ x^P|)\ in\ v]$
apply (*rule beta-C-meta-1*[*equiv-rl*])
unfolding $identity_E$*-infix-def*
apply *show-proper*
using *ord-eq-Eequiv-1*[*deduction*] *ord-xy*[*conj1*]
unfolding $identity_E$*-infix-def* **by** *simp*
ultimately have $[(|\boldsymbol{\lambda}\ z\ .\ z^P =_E x^P,\ y^P|)\ in\ v]$
using $\equiv E$ **by** *blast*
hence $[y^P =_E x^P\ in\ v]$
unfolding $identity_E$*-infix-def*
apply (*safe intro*!:
beta-C-meta-1[**where** $\varphi = \lambda\ z\ .\ (|basic\text{-}identity_E,z,x^P|)$, *equiv-lr*])
by *show-proper*
thus $[x^P =_E y^P\ in\ v]$
by (*rule ord-eq-Eequiv-2*[*deduction*])
qed

lemma *ord-eq-E2*[*PLM*]:
$[((|O!,x^P|)\ \&\ (|O!,y^P|)) \rightarrow$
$((x^P \neq y^P) \equiv (\boldsymbol{\lambda} z\ .\ z^P =_E x^P) \neq (\boldsymbol{\lambda} z\ .\ z^P =_E y^P))\ in\ v]$
proof (*rule CP*; *rule* $\equiv I$; *rule CP*)
assume *ord-xy*: $[(|O!,x^P|)\ \&\ (|O!,y^P|)\ in\ v]$
assume $[x^P \neq y^P\ in\ v]$
hence $[\neg(x^P =_E y^P)\ in\ v]$
using *eq-E-simple-2 modus-tollens-1* **by** *fast*
moreover {
assume $[(\boldsymbol{\lambda} z\ .\ z^P =_E x^P) = (\boldsymbol{\lambda} z\ .\ z^P =_E y^P)\ in\ v]$
moreover have $[(|\boldsymbol{\lambda} z\ .\ z^P =_E x^P,\ x^P|)\ in\ v]$
apply (*rule beta-C-meta-1*[*equiv-rl*])
unfolding $identity_E$*-infix-def*
apply *show-proper*
using *ord-eq-Eequiv-1*[*deduction*] *ord-xy*[*conj1*]
unfolding $identity_E$*-infix-def* **by** *presburger*
ultimately have $[(|\boldsymbol{\lambda} z\ .\ z^P =_E y^P,\ x^P|)\ in\ v]$
using *l-identity*[*axiom-instance*, *deduction*, *deduction*] **by** *fast*
hence $[x^P =_E y^P\ in\ v]$
unfolding $identity_E$*-infix-def*
apply (*safe intro*!:
beta-C-meta-1[**where** $\varphi = \lambda\ z\ .\ (|basic\text{-}identity_E,z,y^P|)$, *equiv-lr*])
by *show-proper*
}
ultimately show $[(\boldsymbol{\lambda} z\ .\ z^P =_E x^P) \neq (\boldsymbol{\lambda} z\ .\ z^P =_E y^P)\ in\ v]$
using *modus-tollens-1 CP* **by** *blast*

next
assume *ord-xy*: $[(\!|O!,x^P|\!)$ **&** $(\!|O!,y^P|\!)$ *in* $v]$
assume $[(\boldsymbol{\lambda} z\ .\ z^P =_E x^P) \neq (\boldsymbol{\lambda} z\ .\ z^P =_E y^P)$ *in* $v]$
moreover {
assume $[x^P = y^P$ *in* $v]$
hence $[(\boldsymbol{\lambda} z\ .\ z^P =_E x^P) = (\boldsymbol{\lambda} z\ .\ z^P =_E y^P)$ *in* $v]$
using *id-eq-1 l-identity*[*axiom-instance*, *deduction*, *deduction*]
by *fast*
}
ultimately show $[x^P \neq y^P$ *in* $v]$
using *modus-tollens-1 CP* **by** *blast*
qed

lemma *ab-obey-1*[*PLM*]:
$[((\!|A!,x^P|\!)$ **&** $(\!|A!,y^P|\!)) \rightarrow ((\forall\ F\ .\ \{\!|x^P, F|\!\} \equiv \{\!|y^P, F|\!\}) \rightarrow x^P = y^P)$ *in* $v]$
proof(*rule CP*; *rule CP*)
assume *abs-xy*: $[(\!|A!,x^P|\!)$ **&** $(\!|A!,y^P|\!)$ *in* $v]$
assume *enc-equiv*: $[\forall\ F\ .\ \{\!|x^P, F|\!\} \equiv \{\!|y^P, F|\!\}$ *in* $v]$
{
fix P
have $[\{\!|x^P, P|\!\} \equiv \{\!|y^P, P|\!\}$ *in* $v]$
using *enc-equiv* **by** (*rule* $\forall E$)
hence $[\Box(\{\!|x^P, P|\!\} \equiv \{\!|y^P, P|\!\})$ *in* $v]$
using *en-eq-2 intro-elim-6-e intro-elim-6-f*
en-eq-5[*equiv-rl*] **by** *meson*
}
hence $[\Box(\forall\ F\ .\ \{\!|x^P, F|\!\} \equiv \{\!|y^P, F|\!\})$ *in* $v]$
using *BF*[*deduction*] $\forall I$ **by** *fast*
thus $[x^P = y^P$ *in* $v]$
unfolding *identity-defs*
using $\vee I(2)$ *abs-xy* $\&I$ **by** *presburger*
qed

lemma *ab-obey-2*[*PLM*]:
$[((\!|A!,x^P|\!)$ **&** $(\!|A!,y^P|\!)) \rightarrow ((\exists\ F\ .\ \{\!|x^P, F|\!\}$ **&** $\neg\{\!|y^P, F|\!\}) \rightarrow x^P \neq y^P)$ *in* $v]$
proof(*rule CP*; *rule CP*)
assume *abs-xy*: $[(\!|A!,x^P|\!)$ **&** $(\!|A!,y^P|\!)$ *in* $v]$
assume $[\exists\ F\ .\ \{\!|x^P, F|\!\}$ **&** $\neg\{\!|y^P, F|\!\}$ *in* $v]$
then obtain P **where** *P-prop*:
$[\{\!|x^P, P|\!\}$ **&** $\neg\{\!|y^P, P|\!\}$ *in* $v]$
by (*rule* $\exists E$)
{
assume $[x^P = y^P$ *in* $v]$
hence $[\{\!|x^P, P|\!\} \equiv \{\!|y^P, P|\!\}$ *in* $v]$
using *l-identity*[*axiom-instance*, *deduction*, *deduction*]
oth-class-taut-4-a **by** *fast*
hence $[\{\!|y^P, P|\!\}$ *in* $v]$
using *P-prop*[*conj1*] **by** (*rule* $\equiv E$)
}
thus $[x^P \neq y^P$ *in* $v]$
using *P-prop*[*conj2*] *modus-tollens-1 CP* **by** *blast*
qed

lemma *ordnecfail*[*PLM*]:
$[(\!|O!,x^P|\!) \rightarrow \Box(\neg(\exists\ F\ .\ \{\!|x^P, F|\!\}))$ *in* $v]$
proof (*rule CP*)
assume $[(\!|O!,x^P|\!)$ *in* $v]$
hence $[\Box(\!|O!,x^P|\!)$ *in* $v]$

```
    using oa-facts-1[deduction] by simp
  moreover hence [□(⦇O!,x^P⦈ → (¬(∃ F . ⦃x^P, F⦄))) in v]
    using nocoder[axiom-necessitation, axiom-instance] by simp
  ultimately show [□(¬(∃ F . ⦃x^P, F⦄)) in v]
    using qml-1[axiom-instance, deduction, deduction] by fast
 qed

lemma o-objects-exist-1[PLM]:
 [◇(∃ x . ⦇E!,x^P⦈) in v]
 proof −
  have [◇(∃ x . ⦇E!,x^P⦈ & ◇(¬⦇E!,x^P⦈)) in v]
    using qml-4[axiom-instance, conj1] .
  hence [◇((∃ x . ⦇E!,x^P⦈) & (∃ x . ◇(¬⦇E!,x^P⦈))) in v]
    using sign-S5-thm-3[deduction] by fast
  hence [◇(∃ x . ⦇E!,x^P⦈) & ◇(∃ x . ◇(¬⦇E!,x^P⦈)) in v]
    using KBasic2-8[deduction] by blast
  thus ?thesis using &E by blast
 qed

lemma o-objects-exist-2[PLM]:
 [□(∃ x . ⦇O!,x^P⦈) in v]
 apply (rule RN) unfolding Ordinary-def
 apply (PLM-subst-method λ x . ◇⦇E!,x^P⦈ λ x . (⦇λy. ◇⦇E!,y^P⦈, x^P⦈)
  apply (safe intro!: beta-C-meta-1[equiv-sym])
  apply show-proper
 using o-objects-exist-1 BF◇[deduction] by blast

lemma o-objects-exist-3[PLM]:
 [□(¬(∀ x . ⦇A!,x^P⦈)) in v]
 apply (PLM-subst-method (∃ x. ¬⦇A!,x^P⦈) ¬(∀ x. ⦇A!,x^P⦈))
  using cqt-further-2[equiv-sym] apply fast
 apply (PLM-subst-method λ x . ⦇O!,x^P⦈ λ x . ¬⦇A!,x^P⦈)
 using oa-contingent-2 o-objects-exist-2 by auto

lemma a-objects-exist-1[PLM]:
 [□(∃ x . ⦇A!,x^P⦈) in v]
 proof −
  {
   fix v
   have [∃ x . ⦇A!,x^P⦈ & (∀ F . ⦃x^P, F⦄ ≡ (F = F)) in v]
    using A-objects[axiom-instance] by simp
   hence [∃ x . ⦇A!,x^P⦈ in v]
    using cqt-further-5[deduction,conj1] by fast
  }
  thus ?thesis by (rule RN)
 qed

lemma a-objects-exist-2[PLM]:
 [□(¬(∀ x . ⦇O!,x^P⦈)) in v]
 apply (PLM-subst-method (∃ x. ¬⦇O!,x^P⦈) ¬(∀ x. ⦇O!,x^P⦈))
  using cqt-further-2[equiv-sym] apply fast
 apply (PLM-subst-method λ x . ⦇A!,x^P⦈ λ x . ¬⦇O!,x^P⦈)
  using oa-contingent-3 a-objects-exist-1 by auto

lemma a-objects-exist-3[PLM]:
 [□(¬(∀ x . ⦇E!,x^P⦈)) in v]
 proof −
  {
```

```
    fix v
    have [∃ x . (|A!,x^P|) & (∀ F . {|x^P, F|} ≡ (F = F)) in v]
      using A-objects[axiom-instance] by simp
    hence [∃ x . (|A!,x^P|) in v]
      using cqt-further-5[deduction,conj1] by fast
    then obtain a where
      [(|A!,a^P|) in v]
      by (rule ∃E)
    hence [¬(◇(|E!,a^P|)) in v]
      unfolding Abstract-def
      apply (safe intro!: beta-C-meta-1[equiv-lr])
      by show-proper
    hence [(¬(|E!,a^P|)) in v]
      using KBasic2-4[equiv-rl] qml-2[axiom-instance,deduction]
      by simp
    hence [¬(∀ x . (|E!,x^P|)) in v]
      using ∃I cqt-further-2[equiv-rl]
      by fast
  }
  thus ?thesis
    by (rule RN)
qed

lemma encoders-are-abstract[PLM]:
  [(∃ F . {|x^P, F|}) → (|A!,x^P|) in v]
  using nocoder[axiom-instance] contraposition-2
        oa-contingent-2[THEN oth-class-taut-5-d[equiv-lr], equiv-lr]
        useful-tautologies-1[deduction]
        vdash-properties-10 CP by metis

lemma A-objects-unique[PLM]:
  [∃! x . (|A!,x^P|) & (∀ F . {|x^P, F|} ≡ φ F) in v]
  proof −
    have [∃ x . (|A!,x^P|) & (∀ F . {|x^P, F|} ≡ φ F) in v]
      using A-objects[axiom-instance] by simp
    then obtain a where a-prop:
      [(|A!,a^P|) & (∀ F . {|a^P, F|} ≡ φ F) in v] by (rule ∃E)
    moreover have [∀ y . (|A!,y^P|) & (∀ F . {|y^P, F|} ≡ φ F) → (y = a) in v]
      proof (rule ∀I; rule CP)
        fix b
        assume b-prop: [(|A!,b^P|) & (∀ F . {|b^P, F|} ≡ φ F) in v]
        {
          fix P
          have [{|b^P,P|} ≡ {|a^P, P|} in v]
            using a-prop[conj2] b-prop[conj2] ≡I ≡E(1) ≡E(2)
                  CP vdash-properties-10 ∀E by metis
        }
        hence [∀ F . {|b^P,F|} ≡ {|a^P, F|} in v]
          using ∀I by fast
        thus [b = a in v]
          unfolding identity-ν-def
          using ab-obey-1[deduction, deduction]
                a-prop[conj1] b-prop[conj1] &I by blast
      qed
    ultimately show ?thesis
      unfolding exists-unique-def
      using &I ∃I by fast
  qed
```

lemma *obj-oth-1*[*PLM*]:
[∃! x . (|*A*!,x^P|) & (∀ F . {|x^P, F|} ≡ (|F, y^P|)) *in* v]
using *A-objects-unique* .

lemma *obj-oth-2*[*PLM*]:
[∃! x . (|*A*!,x^P|) & (∀ F . {|x^P, F|} ≡ ((|F, y^P|) & (|F, z^P|))) *in* v]
using *A-objects-unique* .

lemma *obj-oth-3*[*PLM*]:
[∃! x . (|*A*!,x^P|) & (∀ F . {|x^P, F|} ≡ ((|F, y^P|) ∨ (|F, z^P|))) *in* v]
using *A-objects-unique* .

lemma *obj-oth-4*[*PLM*]:
[∃! x . (|*A*!,x^P|) & (∀ F . {|x^P, F|} ≡ (□(|F, y^P|))) *in* v]
using *A-objects-unique* .

lemma *obj-oth-5*[*PLM*]:
[∃! x . (|*A*!,x^P|) & (∀ F . {|x^P, F|} ≡ ($F = G$)) *in* v]
using *A-objects-unique* .

lemma *obj-oth-6*[*PLM*]:
[∃! x . (|*A*!,x^P|) & (∀ F . {|x^P, F|} ≡ □(∀ y . (|G, y^P|) → (|F, y^P|))) *in* v]
using *A-objects-unique* .

lemma *A-Exists-1*[*PLM*]:
[**A**(∃! x :: ('a :: *id-act*) . φ x) ≡ (∃! x . **A**(φ x)) *in* v]
unfolding *exists-unique-def*
proof (*rule* ≡*I*; *rule CP*)
assume [**A**(∃ α. φ α & (∀ β. φ β → $\beta = \alpha$)) *in* v]
hence [∃ α. **A**(φ α & (∀ β. φ β → $\beta = \alpha$)) *in* v]
using *Act-Basic-11*[*equiv-lr*] **by** *blast*
then obtain α **where**
[**A**(φ α & (∀ β. φ β → $\beta = \alpha$)) *in* v]
by (*rule* ∃ *E*)
hence *1*: [**A**(φ α) & **A**(∀ β. φ β → $\beta = \alpha$) *in* v]
using *Act-Basic-2*[*equiv-lr*] **by** *blast*
have *2*: [∀ β. **A**(φ β → $\beta = \alpha$) *in* v]
using *1*[*conj2*] *logic-actual-nec-3*[*axiom-instance*, *equiv-lr*] **by** *blast*
{
fix β
have [**A**(φ β → $\beta = \alpha$) *in* v]
using *2* **by** (*rule* ∀ *E*)
hence [**A**(φ β) → ($\beta = \alpha$) *in* v]
using *logic-actual-nec-2*[*axiom-instance*, *equiv-lr*, *deduction*]
id-act-3[*equiv-rl*] *CP* **by** *blast*
}
hence [∀ β . **A**(φ β) → ($\beta = \alpha$) *in* v]
by (*rule* ∀ *I*)
thus [∃ α. **A**φ α & (∀ β. **A**φ β → $\beta = \alpha$) *in* v]
using *1*[*conj1*] &*I* ∃ *I* **by** *fast*
next
assume [∃ α. **A**φ α & (∀ β. **A**φ β → $\beta = \alpha$) *in* v]
then obtain α **where** *1*:
[**A**φ α & (∀ β. **A**φ β → $\beta = \alpha$) *in* v]
by (*rule* ∃ *E*)
{
fix β

have $[\mathcal{A}(\varphi\ \beta) \rightarrow \beta = \alpha\ \textit{in}\ v]$
using *1*[*conj2*] **by** (*rule* $\forall E$)
hence $[\mathcal{A}(\varphi\ \beta \rightarrow \beta = \alpha)\ \textit{in}\ v]$
using *logic-actual-nec-2*[*axiom-instance*, *equiv-rl*] *id-act-3*[*equiv-lr*]
vdash-properties-10 CP **by** *blast*
}
hence $[\forall\ \beta\ .\ \mathcal{A}(\varphi\ \beta \rightarrow \beta = \alpha)\ \textit{in}\ v]$
by (*rule* $\forall I$)
hence $[\mathcal{A}(\forall\ \beta\ .\ \varphi\ \beta \rightarrow \beta = \alpha)\ \textit{in}\ v]$
using *logic-actual-nec-3*[*axiom-instance*, *equiv-rl*] **by** *fast*
hence $[\mathcal{A}(\varphi\ \alpha\ \&\ (\forall\ \beta\ .\ \varphi\ \beta \rightarrow \beta = \alpha))\ \textit{in}\ v]$
using *1*[*conj1*] *Act-Basic-2*[*equiv-rl*] $\&I$ **by** *blast*
hence $[\exists\,\alpha.\ \mathcal{A}(\varphi\ \alpha\ \&\ (\forall\,\beta.\ \varphi\ \beta \rightarrow \beta = \alpha))\ \textit{in}\ v]$
using $\exists I$ **by** *fast*
thus $[\mathcal{A}(\exists\,\alpha.\ \varphi\ \alpha\ \&\ (\forall\,\beta.\ \varphi\ \beta \rightarrow \beta = \alpha))\ \textit{in}\ v]$
using *Act-Basic-11*[*equiv-rl*] **by** *fast*
qed

lemma *A-Exists-2*[*PLM*]:
$[(\exists\ y\ .\ y^P = (\iota x\ .\ \varphi\ x)) \equiv \mathcal{A}(\exists\,!x\ .\ \varphi\ x)\ \textit{in}\ v]$
using *actual-desc-1 A-Exists-1*[*equiv-sym*]
intro-elim-6-e **by** *blast*

lemma *A-descriptions*[*PLM*]:
$[\exists\ y\ .\ y^P = (\iota x\ .\ (\!|A!,x^P|\!)\ \&\ (\forall\ F\ .\ \{\!|x^P,F|\!\} \equiv \varphi\ F))\ \textit{in}\ v]$
using *A-objects-unique*[*THEN RN*, *THEN nec-imp-act*[*deduction*]]
A-Exists-2[*equiv-rl*] **by** *auto*

lemma *thm-can-terms2*[*PLM*]:
$[(y^P = (\iota x\ .\ (\!|A!,x^P|\!)\ \&\ (\forall\ F\ .\ \{\!|x^P,F|\!\} \equiv \varphi\ F)))$
$\rightarrow ((\!|A!,y^P|\!)\ \&\ (\forall\ F\ .\ \{\!|y^P,F|\!\} \equiv \varphi\ F))\ \textit{in}\ dw]$
using *y-in-2* **by** *auto*

lemma *can-ab2*[*PLM*]:
$[(y^P = (\iota x\ .\ (\!|A!,x^P|\!)\ \&\ (\forall\ F\ .\ \{\!|x^P,F|\!\} \equiv \varphi\ F))) \rightarrow (\!|A!,y^P|\!)\ \textit{in}\ v]$
proof (*rule CP*)
assume $[y^P = (\iota x\ .\ (\!|A!,x^P|\!)\ \&\ (\forall\ F\ .\ \{\!|x^P,F|\!\} \equiv \varphi\ F))\ \textit{in}\ v]$
hence $[\mathcal{A}(\!|A!,y^P|\!)\ \&\ \mathcal{A}(\forall\ F\ .\ \{\!|y^P,F|\!\} \equiv \varphi\ F)\ \textit{in}\ v]$
using *nec-hintikka-scheme*[*equiv-lr*, *conj1*]
Act-Basic-2[*equiv-lr*] **by** *blast*
thus $[(\!|A!,y^P|\!)\ \textit{in}\ v]$
using *oa-facts-8*[*equiv-rl*] $\&E$ **by** *blast*
qed

lemma *desc-encode*[*PLM*]:
$[\{\!|\iota x\ .\ (\!|A!,x^P|\!)\ \&\ (\forall\ F\ .\ \{\!|x^P,F|\!\} \equiv \varphi\ F),\ G|\!\} \equiv \varphi\ G\ \textit{in}\ dw]$
proof −
obtain *a* **where**
$[a^P = (\iota x\ .\ (\!|A!,x^P|\!)\ \&\ (\forall\ F\ .\ \{\!|x^P,F|\!\} \equiv \varphi\ F))\ \textit{in}\ dw]$
using *A-descriptions* **by** (*rule* $\exists E$)
moreover hence $[\{\!|a^P,\ G|\!\} \equiv \varphi\ G\ \textit{in}\ dw]$
using *hintikka*[*equiv-lr*, *conj1*] $\&E\ \forall E$ **by** *fast*
ultimately show *?thesis*
using *l-identity*[*axiom-instance*, *deduction*, *deduction*] **by** *fast*
qed

lemma *desc-nec-encode*[*PLM*]:
$[\{\!|\iota x\ .\ (\!|A!,x^P|\!)\ \&\ (\forall\ F\ .\ \{\!|x^P,F|\!\} \equiv \varphi\ F),\ G|\!\} \equiv \mathcal{A}(\varphi\ G)\ \textit{in}\ v]$

```
  proof −
    obtain a where
      [a^P = (ιx . (|A!,x^P|) & (∀ F . {|x^P,F|} ≡ φ F)) in v]
      using A-descriptions by (rule ∃E)
    moreover {
      hence [A((|A!,a^P|) & (∀ F . {|a^P,F|} ≡ φ F)) in v]
        using nec-hintikka-scheme[equiv-lr, conj1] by fast
      hence [A(∀ F . {|a^P,F|} ≡ φ F) in v]
        using Act-Basic-2[equiv-lr,conj2] by blast
      hence [∀ F . A( {|a^P,F|} ≡ φ F) in v]
        using logic-actual-nec-3[axiom-instance, equiv-lr] by blast
      hence [A({|a^P, G|} ≡ φ G) in v]
        using ∀E by fast
      hence [A{|a^P, G|} ≡ A(φ G) in v]
        using Act-Basic-5[equiv-lr] by fast
      hence [{|a^P, G|} ≡ A(φ G) in v]
        using en-eq-10[equiv-sym] intro-elim-6-e by blast
    }
    ultimately show ?thesis
      using l-identity[axiom-instance, deduction, deduction] by fast
  qed

notepad
begin
    fix v
    let ?x = ιx . (|A!,x^P|) & (∀ F . {|x^P, F|} ≡ (∃ q . q & F = (λ y . q)))
    have [□(∃ p . ContingentlyTrue p) in v]
      using cont-tf-thm-3 RN by auto
    hence [A(∃ p . ContingentlyTrue p) in v]
      using nec-imp-act[deduction] by simp
    hence [∃ p . A(ContingentlyTrue p) in v]
      using Act-Basic-11[equiv-lr] by auto
    then obtain p_1 where
      [A(ContingentlyTrue p_1) in v]
      by (rule ∃E)
    hence [Ap_1 in v]
      unfolding ContingentlyTrue-def
      using Act-Basic-2[equiv-lr] &E by fast
    hence [Ap_1 & A((λ y . p_1) = (λ y . p_1)) in v]
      using &I id-eq-1[THEN RN, THEN nec-imp-act[deduction]] by fast
    hence [A(p_1 & (λ y . p_1) = (λ y . p_1)) in v]
      using Act-Basic-2[equiv-rl] by fast
    hence [∃ q . A( q & (λ y . p_1) = (λ y . q)) in v]
      using ∃I by fast
    hence [A(∃ q . q & (λ y . p_1) = (λ y . q)) in v]
      using Act-Basic-11[equiv-rl] by fast
    moreover have [{|?x, λ y . p_1|} ≡ A(∃ q . q & (λ y . p_1) = (λ y . q)) in v]
      using desc-nec-encode by fast
    ultimately have [{|?x, λ y . p_1|} in v]
      using ≡E by blast
end

lemma Box-desc-encode-1[PLM]:
  [□(φ G) → {|(ιx . (|A!,x^P|) & (∀ F . {|x^P, F|} ≡ φ F)), G|} in v]
  proof (rule CP)
    assume [□(φ G) in v]
    hence [A(φ G) in v]
      using nec-imp-act[deduction] by auto
```

```
    thus [⦃ιx . (|A!,x^P|) & (∀ F . ⦃x^P,F⦄ ≡ φ F), G⦄ in v]
      using desc-nec-encode[equiv-rl] by simp
  qed

lemma Box-desc-encode-2[PLM]:
  [□(φ G) → □(⦃(ιx . (|A!,x^P|) & (∀ F . ⦃x^P, F⦄ ≡ φ F)), G⦄ ≡ φ G) in v]
  proof (rule CP)
    assume a: [□(φ G) in v]
    hence [□(⦃(ιx . (|A!,x^P|) & (∀ F . ⦃x^P, F⦄ ≡ φ F)), G⦄ → φ G) in v]
      using KBasic-1[deduction] by simp
    moreover {
      have [⦃(ιx . (|A!,x^P|) & (∀ F . ⦃x^P, F⦄ ≡ φ F)), G⦄ in v]
        using a Box-desc-encode-1[deduction] by auto
      hence [□⦃(ιx . (|A!,x^P|) & (∀ F . ⦃x^P, F⦄ ≡ φ F)), G⦄ in v]
        using encoding[axiom-instance,deduction] by blast
      hence [□(φ G → ⦃(ιx . (|A!,x^P|) & (∀ F . ⦃x^P, F⦄ ≡ φ F)), G⦄) in v]
        using KBasic-1[deduction] by simp
    }
    ultimately show [□(⦃(ιx . (|A!,x^P|) & (∀ F . ⦃x^P, F⦄ ≡ φ F)), G⦄
                    ≡ φ G) in v]
      using &I KBasic-4[equiv-rl] by blast
  qed

lemma box-phi-a-1[PLM]:
  assumes [□(∀ F . φ F → □(φ F)) in v]
  shows [((|A!,x^P|) & (∀ F . ⦃x^P, F⦄ ≡ φ F)) → □((|A!,x^P|)
        & (∀ F . ⦃x^P, F⦄ ≡ φ F)) in v]
  proof (rule CP)
    assume a: [((|A!,x^P|) & (∀ F . ⦃x^P, F⦄ ≡ φ F)) in v]
    have [□(|A!,x^P|) in v]
      using oa-facts-2[deduction] a[conj1] by auto
    moreover have [□(∀ F . ⦃x^P, F⦄ ≡ φ F) in v]
      proof (rule BF[deduction]; rule ∀I)
        fix F
        have ϑ: [□(φ F → □(φ F)) in v]
          using assms[THEN CBF[deduction]] by (rule ∀E)
        moreover have [□(⦃x^P, F⦄ → □⦃x^P, F⦄) in v]
          using encoding[axiom-necessitation, axiom-instance] by simp
        moreover have [□⦃x^P, F⦄ ≡ □(φ F) in v]
          proof (rule ≡I; rule CP)
            assume [□⦃x^P, F⦄ in v]
            hence [⦃x^P, F⦄ in v]
              using qml-2[axiom-instance, deduction] by blast
            hence [φ F in v]
              using a[conj2] ∀E[where 'a=Π1] ≡E by blast
            thus [□(φ F) in v]
              using ϑ[THEN qml-2[axiom-instance, deduction], deduction] by simp
          next
            assume [□(φ F) in v]
            hence [φ F in v]
              using qml-2[axiom-instance, deduction] by blast
            hence [⦃x^P, F⦄ in v]
              using a[conj2] ∀E[where 'a=Π1] ≡E by blast
            thus [□⦃x^P, F⦄ in v]
              using encoding[axiom-instance, deduction] by simp
          qed
        ultimately show [□(⦃x^P,F⦄ ≡ φ F) in v]
          using sc-eq-box-box-3[deduction, deduction] &I by blast
```

```
    qed
  ultimately show [□(⦇A!,x^P⦈ & (∀ F. ⦃x^P,F⦄ ≡ φ F)) in v]
   using &I KBasic-3[equiv-rl] by blast
 qed

lemma box-phi-a-2[PLM]:
 assumes [□(∀ F . φ F → □(φ F)) in v]
 shows [y^P = (ιx . ⦇A!,x^P⦈ & (∀ F. ⦃x^P, F⦄ ≡ φ F))
       → (⦇A!,y^P⦈ & (∀ F . ⦃y^P, F⦄ ≡ φ F)) in v]
 proof −
  let ?ψ = λ x . ⦇A!,x^P⦈ & (∀ F . ⦃x^P, F⦄ ≡ φ F)
  have [∀ x . ?ψ x → □(?ψ x) in v]
   using box-phi-a-1[OF assms] ∀I by fast
  hence [(∃! x . ?ψ x) → (∀ y . y^P = (ιx . ?ψ x) → ?ψ y) in v]
   using unique-box-desc[deduction] by fast
  hence [(∀ y . y^P = (ιx . ?ψ x) → ?ψ y) in v]
   using A-objects-unique modus-ponens by blast
  thus ?thesis by (rule ∀E)
 qed

lemma box-phi-a-3[PLM]:
 assumes [□(∀ F . φ F → □(φ F)) in v]
 shows [⦃ιx . ⦇A!,x^P⦈ & (∀ F . ⦃x^P, F⦄ ≡ φ F), G⦄ ≡ φ G in v]
 proof −
  obtain a where
   [a^P = (ιx . ⦇A!,x^P⦈ & (∀ F . ⦃x^P, F⦄ ≡ φ F)) in v]
   using A-descriptions by (rule ∃E)
  moreover {
   hence [(∀ F . ⦃a^P, F⦄ ≡ φ F) in v]
    using box-phi-a-2[OF assms, deduction, conj2] by blast
   hence [⦃a^P, G⦄ ≡ φ G in v] by (rule ∀E)
  }
  ultimately show ?thesis
   using l-identity[axiom-instance, deduction, deduction] by fast
 qed

lemma null-uni-uniq-1[PLM]:
 [∃! x . Null (x^P) in v]
 proof −
  have [∃ x . ⦇A!,x^P⦈ & (∀ F . ⦃x^P, F⦄ ≡ (F ≠ F)) in v]
   using A-objects[axiom-instance] by simp
  then obtain a where a-prop:
   [⦇A!,a^P⦈ & (∀ F . ⦃a^P, F⦄ ≡ (F ≠ F)) in v]
   by (rule ∃E)
  have 1: [⦇A!,a^P⦈ & (¬(∃ F . ⦃a^P, F⦄)) in v]
   using a-prop[conj1] apply (rule &I)
   proof −
    {
     assume [∃ F . ⦃a^P, F⦄ in v]
     then obtain P where
      [⦃a^P, P⦄ in v] by (rule ∃E)
     hence [P ≠ P in v]
      using a-prop[conj2, THEN ∀E, equiv-lr] by simp
     hence [¬(∃ F . ⦃a^P, F⦄) in v]
      using id-eq-1 reductio-aa-1 by fast
    }
    thus [¬(∃ F . ⦃a^P, F⦄) in v]
     using reductio-aa-1 by blast
```

```
    qed
  moreover have [∀ y . ((|A!,y^P|) & (¬(∃ F . {|y^P, F|}))) → y = a in v]
    proof (rule ∀I; rule CP)
      fix y
      assume 2: [(|A!,y^P|) & (¬(∃ F . {|y^P, F|})) in v]
      have [∀ F . {|y^P, F|} ≡ {|a^P, F|} in v]
        using cqt-further-12[deduction] 1[conj2] 2[conj2] &I by blast
      thus [y = a in v]
        using ab-obey-1[deduction, deduction]
        &I[OF 2[conj1] 1[conj1]] identity-ν-def by presburger
    qed
  ultimately show ?thesis
    using &I ∃I
    unfolding Null-def exists-unique-def by fast
qed

lemma null-uni-uniq-2[PLM]:
  [∃! x . Universal (x^P) in v]
  proof −
    have [∃ x . (|A!,x^P|) & (∀ F . {|x^P, F|} ≡ (F = F)) in v]
      using A-objects[axiom-instance] by simp
    then obtain a where a-prop:
      [(|A!,a^P|) & (∀ F . {|a^P, F|} ≡ (F = F)) in v]
      by (rule ∃E)
    have 1: [(|A!,a^P|) & (∀ F . {|a^P, F|}) in v]
      using a-prop[conj1] apply (rule &I)
      using ∀I a-prop[conj2, THEN ∀E, equiv-rl] id-eq-1 by fast
    moreover have [∀ y . ((|A!,y^P|) & (∀ F . {|y^P, F|})) → y = a in v]
      proof (rule ∀I; rule CP)
        fix y
        assume 2: [(|A!,y^P|) & (∀ F . {|y^P, F|}) in v]
        have [∀ F . {|y^P, F|} ≡ {|a^P, F|} in v]
          using cqt-further-11[deduction] 1[conj2] 2[conj2] &I by blast
        thus [y = a in v]
          using ab-obey-1[deduction, deduction]
            &I[OF 2[conj1] 1[conj1]] identity-ν-def
          by presburger
      qed
    ultimately show ?thesis
      using &I ∃I
      unfolding Universal-def exists-unique-def by fast
  qed

lemma null-uni-uniq-3[PLM]:
  [∃ y . y^P = (ιx . Null (x^P)) in v]
  using null-uni-uniq-1[THEN RN, THEN nec-imp-act[deduction]]
        A-Exists-2[equiv-rl] by auto

lemma null-uni-uniq-4[PLM]:
  [∃ y . y^P = (ιx . Universal (x^P)) in v]
  using null-uni-uniq-2[THEN RN, THEN nec-imp-act[deduction]]
        A-Exists-2[equiv-rl] by auto

lemma null-uni-facts-1[PLM]:
  [Null (x^P) → □(Null (x^P)) in v]
  proof (rule CP)
    assume [Null (x^P) in v]
    hence 1: [(|A!,x^P|) & (¬(∃ F . {|x^P,F|})) in v]
```

```
    unfolding Null-def .
  have [□(|A!,x^P|) in v]
    using 1[conj1] oa-facts-2[deduction] by simp
  moreover have [□(¬(∃ F . {|x^P,F|})) in v]
    proof −
      {
        assume [¬□(¬(∃ F . {|x^P,F|})) in v]
        hence [◇(∃ F . {|x^P,F|}) in v]
          unfolding diamond-def .
        hence [∃ F . ◇{|x^P,F|} in v]
          using BF◇[deduction] by blast
        then obtain P where [◇{|x^P,P|} in v]
          by (rule ∃E)
        hence [{|x^P, P|} in v]
          using en-eq-3[equiv-lr] by simp
        hence [∃ F . {|x^P, F|} in v]
          using ∃I by fast
      }
      thus ?thesis
        using 1[conj2] modus-tollens-1 CP
              useful-tautologies-1[deduction] by metis
    qed
  ultimately show [□Null (x^P) in v]
    unfolding Null-def
    using &I KBasic-3[equiv-rl] by blast
qed

lemma null-uni-facts-2[PLM]:
  [Universal (x^P) → □(Universal (x^P)) in v]
  proof (rule CP)
    assume [Universal (x^P) in v]
    hence 1: [(|A!,x^P|) & (∀ F . {|x^P,F|}) in v]
      unfolding Universal-def .
    have [□(|A!,x^P|) in v]
      using 1[conj1] oa-facts-2[deduction] by simp
    moreover have [□(∀ F . {|x^P,F|}) in v]
      proof (rule BF[deduction]; rule ∀I)
        fix F
        have [{|x^P, F|} in v]
          using 1[conj2] by (rule ∀E)
        thus [□{|x^P, F|} in v]
          using encoding[axiom-instance, deduction] by auto
      qed
    ultimately show [□Universal (x^P) in v]
      unfolding Universal-def
      using &I KBasic-3[equiv-rl] by blast
  qed

lemma null-uni-facts-3[PLM]:
  [Null (a_∅) in v]
  proof −
    let ?ψ = λ x . Null x
    have [((∃! x . ?ψ (x^P)) → (∀ y . y^P = (ιx . ?ψ (x^P)) → ?ψ (y^P))) in v]
      using unique-box-desc[deduction] null-uni-facts-1[THEN ∀I] by fast
    have 1: [(∀ y . y^P = (ιx . ?ψ (x^P)) → ?ψ (y^P)) in v]
      using unique-box-desc[deduction, deduction] null-uni-uniq-1
            null-uni-facts-1[THEN ∀I] by fast
    have [∃ y . y^P = (a_∅) in v]
```

unfolding *NullObject-def* **using** *null-uni-uniq-3* **.**
then obtain *y* **where** $[y^P = (\mathbf{a}_\emptyset)\ in\ v]$
by (*rule* $\exists\, E$)
moreover hence $[?\psi\ (y^P)\ in\ v]$
using *1*[*THEN* $\forall\, E$, *deduction*] **unfolding** *NullObject-def* **by** *simp*
ultimately show $[?\psi\ (\mathbf{a}_\emptyset)\ in\ v]$
using *l-identity*[*axiom-instance*, *deduction*, *deduction*] **by** *blast*
qed

lemma *null-uni-facts-4*[*PLM*]:
$[Universal\ (\mathbf{a}_V)\ in\ v]$
proof −
let $?\psi = \lambda\ x\ .\ Universal\ x$
have $[((\exists !\ x\ .\ ?\psi\ (x^P)) \rightarrow (\forall\ y\ .\ y^P = (\iota x\ .\ ?\psi\ (x^P)) \rightarrow\ ?\psi\ (y^P)))\ in\ v]$
using *unique-box-desc*[*deduction*] *null-uni-facts-2*[*THEN* $\forall\, I$] **by** *fast*
have *1*: $[(\forall\ y\ .\ y^P = (\iota x\ .\ ?\psi\ (x^P)) \rightarrow\ ?\psi\ (y^P))\ in\ v]$
using *unique-box-desc*[*deduction*, *deduction*] *null-uni-uniq-2*
null-uni-facts-2[*THEN* $\forall\, I$] **by** *fast*
have $[\exists\ y\ .\ y^P = (\mathbf{a}_V)\ in\ v]$
unfolding *UniversalObject-def* **using** *null-uni-uniq-4* **.**
then obtain *y* **where** $[y^P = (\mathbf{a}_V)\ in\ v]$
by (*rule* $\exists\, E$)
moreover hence $[?\psi\ (y^P)\ in\ v]$
using *1*[*THEN* $\forall\, E$, *deduction*]
unfolding *UniversalObject-def* **by** *simp*
ultimately show $[?\psi\ (\mathbf{a}_V)\ in\ v]$
using *l-identity*[*axiom-instance*, *deduction*, *deduction*] **by** *blast*
qed

lemma *aclassical-1*[*PLM*]:
$[\forall\ R\ .\ \exists\ x\ y\ .\ (\!|A!,x^P|\!)\ \&\ (\!|A!,y^P|\!)\ \&\ (x \neq y)$
$\&\ (\boldsymbol{\lambda}\ z\ .\ (\!|R,z^P,x^P|\!)) = (\boldsymbol{\lambda}\ z\ .\ (\!|R,z^P,y^P|\!))\ in\ v]$
proof (*rule* $\forall\, I$)
fix *R*
obtain *a* **where** ϑ:
$[(\!|A!,a^P|\!)\ \&\ (\forall\ F\ .\ \{\!|a^P, F|\!\} \equiv (\exists\ y\ .\ (\!|A!,y^P|\!)$
$\&\ F = (\boldsymbol{\lambda}\ z\ .\ (\!|R,z^P,y^P|\!))\ \&\ \neg\{\!|y^P, F|\!\}))\ in\ v]$
using *A-objects*[*axiom-instance*] **by** (*rule* $\exists\, E$)
{
assume $[\neg\{\!|a^P, (\boldsymbol{\lambda}\ z\ .\ (\!|R,z^P,a^P|\!))|\!\}\ in\ v]$
hence $[\neg((\!|A!,a^P|\!)\ \&\ (\boldsymbol{\lambda}\ z\ .\ (\!|R,z^P,a^P|\!)) = (\boldsymbol{\lambda}\ z\ .\ (\!|R,z^P,a^P|\!))$
$\&\ \neg\{\!|a^P, (\boldsymbol{\lambda}\ z\ .\ (\!|R,z^P,a^P|\!))|\!\})\ in\ v]$
using ϑ[*conj2*, *THEN* $\forall\, E$, *THEN oth-class-taut-5-d*[*equiv-lr*], *equiv-lr*]
cqt-further-4[*equiv-lr*] $\forall\, E$ **by** *fast*
hence $[(\!|A!,a^P|\!)\ \&\ (\boldsymbol{\lambda}\ z\ .\ (\!|R,z^P,a^P|\!)) = (\boldsymbol{\lambda}\ z\ .\ (\!|R,z^P,a^P|\!))$
$\rightarrow \{\!|a^P, (\boldsymbol{\lambda}\ z\ .\ (\!|R,z^P,a^P|\!))|\!\}\ in\ v]$
apply − **by** *PLM-solver*
hence $[\{\!|a^P, (\boldsymbol{\lambda}\ z\ .\ (\!|R,z^P,a^P|\!))|\!\}\ in\ v]$
using ϑ[*conj1*] *id-eq-1* &*I vdash-properties-10* **by** *fast*
}
hence *1*: $[\{\!|a^P, (\boldsymbol{\lambda}\ z\ .\ (\!|R,z^P,a^P|\!))|\!\}\ in\ v]$
using *reductio-aa-1 CP if-p-then-p* **by** *blast*
then obtain *b* **where** ξ:
$[(\!|A!,b^P|\!)\ \&\ (\boldsymbol{\lambda}\ z\ .\ (\!|R,z^P,a^P|\!)) = (\boldsymbol{\lambda}\ z\ .\ (\!|R,z^P,b^P|\!))$
$\&\ \neg\{\!|b^P, (\boldsymbol{\lambda}\ z\ .\ (\!|R,z^P,a^P|\!))|\!\}\ in\ v]$
using ϑ[*conj2*, *THEN* $\forall\, E$, *equiv-lr*] $\exists\, E$ **by** *blast*
have $[a \neq b\ in\ v]$
proof −

```
    {
      assume [a = b in v]
      hence [⦃b^P, (λ z . (|R,z^P,a^P|))⦄ in v]
        using 1 l-identity[axiom-instance, deduction, deduction] by fast
      hence ?thesis
        using ξ[conj2] reductio-aa-1 by blast
    }
    thus ?thesis using reductio-aa-1 by blast
  qed
 hence [(|A!,a^P|) & (|A!,b^P|) & a ≠ b
       & (λ z . (|R,z^P,a^P|)) = (λ z . (|R,z^P,b^P|)) in v]
  using ϑ[conj1] ξ[conj1, conj1] ξ[conj1, conj2] &I by presburger
 hence [∃ y . (|A!,a^P|) & (|A!,y^P|) & a ≠ y
       & (λz. (|R,z^P,a^P|)) = (λz. (|R,z^P,y^P|)) in v]
  using ∃ I by fast
 thus [∃ x y . (|A!,x^P|) & (|A!,y^P|) & x ≠ y
       & (λz. (|R,z^P,x^P|)) = (λz. (|R,z^P,y^P|)) in v]
  using ∃ I by fast
qed

lemma aclassical-2[PLM]:
 [∀ R . ∃ x y . (|A!,x^P|) & (|A!,y^P|) & (x ≠ y)
  & (λ z . (|R,x^P,z^P|)) = (λ z . (|R,y^P,z^P|)) in v]
 proof (rule ∀ I)
  fix R
  obtain a where ϑ:
   [(|A!,a^P|) & (∀ F . ⦃a^P, F⦄ ≡ (∃ y . (|A!,y^P|)
    & F = (λ z . (|R,y^P,z^P|)) & ¬⦃y^P, F⦄)) in v]
   using A-objects[axiom-instance] by (rule ∃ E)
  {
   assume [¬⦃a^P, (λ z . (|R,a^P,z^P|))⦄ in v]
   hence [¬((|A!,a^P|) & (λ z . (|R,a^P,z^P|)) = (λ z . (|R,a^P,z^P|))
         & ¬⦃a^P, (λ z . (|R,a^P,z^P|))⦄) in v]
    using ϑ[conj2, THEN ∀ E, THEN oth-class-taut-5-d[equiv-lr], equiv-lr]
          cqt-further-4[equiv-lr] ∀ E by fast
   hence [(|A!,a^P|) & (λ z . (|R,a^P,z^P|)) = (λ z . (|R,a^P,z^P|))
         → ⦃a^P, (λ z . (|R,a^P,z^P|))⦄ in v]
    apply − by PLM-solver
   hence [⦃a^P, (λ z . (|R,a^P,z^P|))⦄ in v]
    using ϑ[conj1] id-eq-1 &I vdash-properties-10 by fast
  }
  hence 1: [⦃a^P, (λ z . (|R,a^P,z^P|))⦄ in v]
   using reductio-aa-1 CP if-p-then-p by blast
  then obtain b where ξ:
   [(|A!,b^P|) & (λ z . (|R,a^P,z^P|)) = (λ z . (|R,b^P,z^P|))
    & ¬⦃b^P, (λ z . (|R,a^P,z^P|))⦄ in v]
   using ϑ[conj2, THEN ∀ E, equiv-lr] ∃ E by blast
  have [a ≠ b in v]
   proof −
    {
     assume [a = b in v]
     hence [⦃b^P, (λ z . (|R,a^P,z^P|))⦄ in v]
      using 1 l-identity[axiom-instance, deduction, deduction] by fast
     hence ?thesis using ξ[conj2] reductio-aa-1 by blast
    }
    thus ?thesis using ξ[conj2] reductio-aa-1 by blast
   qed
  hence [(|A!,a^P|) & (|A!,b^P|) & a ≠ b
```

```
        & (λ z . (|R,a^P,z^P|)) = (λ z . (|R,b^P,z^P|)) in v]
    using ϑ[conj1] ξ[conj1, conj1] ξ[conj1, conj2] &I by presburger
  hence [∃ y . (|A!,a^P|) & (|A!,y^P|) & a ≠ y
        & (λz. (|R,a^P,z^P|)) = (λz. (|R,y^P,z^P|)) in v]
    using ∃ I by fast
  thus [∃ x y . (|A!,x^P|) & (|A!,y^P|) & x ≠ y
        & (λz. (|R,x^P,z^P|)) = (λz. (|R,y^P,z^P|)) in v]
    using ∃ I by fast
qed

lemma aclassical-3[PLM]:
 [∀ F . ∃ x y . (|A!,x^P|) & (|A!,y^P|) & (x ≠ y)
  & ((λ^0 (|F,x^P|)) = (λ^0 (|F,y^P|))) in v]
 proof (rule ∀ I)
  fix R
  obtain a where ϑ:
   [(|A!,a^P|) & (∀ F . {|a^P, F|} ≡ (∃ y . (|A!,y^P|)
    & F = (λ z . (|R,y^P|)) & ¬{|y^P, F|})) in v]
   using A-objects[axiom-instance] by (rule ∃ E)
  {
   assume [¬{|a^P, (λ z . (|R,a^P|))|} in v]
   hence [¬((|A!,a^P|) & (λ z . (|R,a^P|)) = (λ z . (|R,a^P|))
        & ¬{|a^P, (λ z . (|R,a^P|))|}) in v]
    using ϑ[conj2, THEN ∀ E, THEN oth-class-taut-5-d[equiv-lr], equiv-lr]
        cqt-further-4[equiv-lr] ∀ E by fast
   hence [(|A!,a^P|) & (λ z . (|R,a^P|)) = (λ z . (|R,a^P|))
        → {|a^P, (λ z . (|R,a^P|))|} in v]
    apply − by PLM-solver
   hence [{|a^P, (λ z . (|R,a^P|))|} in v]
    using ϑ[conj1] id-eq-1 &I vdash-properties-10 by fast
  }
  hence 1: [{|a^P, (λ z . (|R,a^P|))|} in v]
   using reductio-aa-1 CP if-p-then-p by blast
  then obtain b where ξ:
   [(|A!,b^P|) & (λ z . (|R,a^P|)) = (λ z . (|R,b^P|))
    & ¬{|b^P, (λ z . (|R,a^P|))|} in v]
   using ϑ[conj2, THEN ∀ E, equiv-lr] ∃ E by blast
  have [a ≠ b in v]
   proof −
    {
     assume [a = b in v]
     hence [{|b^P, (λ z . (|R,a^P|))|} in v]
      using 1 l-identity[axiom-instance, deduction, deduction] by fast
     hence ?thesis
      using ξ[conj2] reductio-aa-1 by blast
    }
    thus ?thesis using reductio-aa-1 by blast
   qed
  moreover {
   have [(|R,a^P|) = (|R,b^P|) in v]
    unfolding identity_o-def
    using ξ[conj1, conj2] by auto
   hence [(λ^0 (|R,a^P|)) = (λ^0 (|R,b^P|)) in v]
    using lambda-p-q-p-eq-q[equiv-rl] by simp
  }
  ultimately have [(|A!,a^P|) & (|A!,b^P|) & a ≠ b
        & ((λ^0 (|R,a^P|)) =(λ^0 (|R,b^P|))) in v]
   using ϑ[conj1] ξ[conj1, conj1] ξ[conj1, conj2] &I
```

by *presburger*
hence $[\exists\ y\ .\ (\!|A!,a^P|\!)\ \&\ (\!|A!,y^P|\!)\ \&\ a \neq y$
$\&\ (\boldsymbol{\lambda}^0\ (\!|R,a^P|\!)) = (\boldsymbol{\lambda}^0\ (\!|R,y^P|\!))\ in\ v]$
using $\exists\, I$ **by** *fast*
thus $[\exists\ x\ y\ .\ (\!|A!,x^P|\!)\ \&\ (\!|A!,y^P|\!)\ \&\ x \neq y$
$\&\ (\boldsymbol{\lambda}^0\ (\!|R,x^P|\!)) = (\boldsymbol{\lambda}^0\ (\!|R,y^P|\!))\ in\ v]$
using $\exists\, I$ **by** *fast*
qed

lemma *aclassical2*[*PLM*]:
$[\exists\ x\ y\ .\ (\!|A!,x^P|\!)\ \&\ (\!|A!,y^P|\!)\ \&\ x \neq y\ \&\ (\forall\ F\ .\ (\!|F,x^P|\!) \equiv (\!|F,y^P|\!))\ in\ v]$
proof −
let $?R_1 = \boldsymbol{\lambda}^2\ (\lambda\ x\ y\ .\ \forall\ F\ .\ (\!|F,x^P|\!) \equiv (\!|F,y^P|\!))$
have $[\exists\ x\ y\ .\ (\!|A!,x^P|\!)\ \&\ (\!|A!,y^P|\!)\ \&\ x \neq y$
$\&\ (\boldsymbol{\lambda}z.\ (\!|?R_1,z^P,x^P|\!)) = (\boldsymbol{\lambda}z.\ (\!|?R_1,z^P,y^P|\!))\ in\ v]$
using *aclassical-1* **by** (*rule* $\forall\, E$)
then obtain a **where**
$[\exists\ y\ .\ (\!|A!,a^P|\!)\ \&\ (\!|A!,y^P|\!)\ \&\ a \neq y$
$\&\ (\boldsymbol{\lambda}z.\ (\!|?R_1,z^P,a^P|\!)) = (\boldsymbol{\lambda}z.\ (\!|?R_1,z^P,y^P|\!))\ in\ v]$
by (*rule* $\exists\, E$)
then obtain b **where** *ab-prop*:
$[(\!|A!,a^P|\!)\ \&\ (\!|A!,b^P|\!)\ \&\ a \neq b$
$\&\ (\boldsymbol{\lambda}z.\ (\!|?R_1,z^P,a^P|\!)) = (\boldsymbol{\lambda}z.\ (\!|?R_1,z^P,b^P|\!))\ in\ v]$
by (*rule* $\exists\, E$)
have $[(\!|?R_1,\ a^P,\ a^P|\!)\ in\ v]$
apply (*rule beta-C-meta-2*[*equiv-rl*])
apply *show-proper*
using *oth-class-taut-4-a*[*THEN* $\forall\, I$] **by** *fast*
hence $[(\!|\boldsymbol{\lambda}\ z\ .\ (\!|?R_1,\ z^P,\ a^P|\!),\ a^P|\!)\ in\ v]$
apply − **apply** (*rule beta-C-meta-1*[*equiv-rl*])
apply *show-proper*
by *auto*
hence $[(\!|\boldsymbol{\lambda}\ z\ .\ (\!|?R_1,\ z^P,\ b^P|\!),\ a^P|\!)\ in\ v]$
using *ab-prop*[*conj2*] *l-identity*[*axiom-instance*, *deduction*, *deduction*]
by *fast*
hence $[(\!|?R_1,\ a^P,\ b^P|\!)\ in\ v]$
apply (*safe intro*!: *beta-C-meta-1*[**where** $\varphi =$
$\lambda z\ .\ (\!|\boldsymbol{\lambda}^2\ (\lambda x\ y.\ \forall F.\ (\!|F,x^P|\!) \equiv (\!|F,y^P|\!)),z,b^P|\!),\ equiv\text{-}lr$])
by *show-proper*
moreover have *IsProperInXY* $(\lambda x\ y.\ \forall F.\ (\!|F,x|\!) \equiv (\!|F,y|\!))$
by *show-proper*
ultimately have $[\forall F.\ (\!|F,a^P|\!) \equiv (\!|F,b^P|\!)\ in\ v]$
using *beta-C-meta-2*[*equiv-lr*] **by** *blast*
hence $[(\!|A!,a^P|\!)\ \&\ (\!|A!,b^P|\!)\ \&\ a \neq b\ \&\ (\forall F.\ (\!|F,a^P|\!) \equiv (\!|F,b^P|\!))\ in\ v]$
using *ab-prop*[*conj1*] **&***I* **by** *presburger*
hence $[\exists\ y\ .\ (\!|A!,a^P|\!)\ \&\ (\!|A!,y^P|\!)\ \&\ a \neq y\ \&\ (\forall F.\ (\!|F,a^P|\!) \equiv (\!|F,y^P|\!))\ in\ v]$
using $\exists\, I$ **by** *fast*
thus *?thesis* **using** $\exists\, I$ **by** *fast*
qed

A.9.13. Propositional Properties

lemma *prop-prop2-1*:
$[\forall\ p\ .\ \exists\ F\ .\ F = (\boldsymbol{\lambda}\ x\ .\ p)\ in\ v]$
proof (*rule* $\forall\, I$)
fix p
have $[(\boldsymbol{\lambda}\ x\ .\ p) = (\boldsymbol{\lambda}\ x\ .\ p)\ in\ v]$
using *id-eq-prop-prop-1* **by** *auto*

thus $[\exists\ \ F\ .\ F = (\boldsymbol{\lambda}\ x\ .\ p)\ \mathit{in}\ v]$
by *PLM-solver*
qed

lemma *prop-prop2-2*:
$[F = (\boldsymbol{\lambda}\ x\ .\ p) \rightarrow \Box(\forall\ x\ .\ (\!|F,x^P|\!) \equiv p)\ \mathit{in}\ v]$
proof (*rule CP*)
assume *1*: $[F = (\boldsymbol{\lambda}\ x\ .\ p)\ \mathit{in}\ v]$
{
fix v
{
fix x
have $[(\!|(\boldsymbol{\lambda}\ x\ .\ p),\ x^P|\!) \equiv p\ \mathit{in}\ v]$
apply (*rule beta-C-meta-1*)
by *show-proper*
}
hence $[\forall\ x\ .\ (\!|(\boldsymbol{\lambda}\ x\ .\ p),\ x^P|\!) \equiv p\ \mathit{in}\ v]$
by (*rule* $\forall I$)
}
hence $[\Box(\forall\ x\ .\ (\!|(\boldsymbol{\lambda}\ x\ .\ p),\ x^P|\!) \equiv p)\ \mathit{in}\ v]$
by (*rule RN*)
thus $[\Box(\forall x.\ (\!|F,x^P|\!) \equiv p)\ \mathit{in}\ v]$
using *l-identity*[*axiom-instance*,*deduction*,*deduction*,
OF 1[*THEN id-eq-prop-prop-2*[*deduction*]]] **by** *fast*
qed

lemma *prop-prop2-3*:
$[\mathit{Propositional}\ F \rightarrow \Box(\mathit{Propositional}\ F)\ \mathit{in}\ v]$
proof (*rule CP*)
assume $[\mathit{Propositional}\ F\ \mathit{in}\ v]$
hence $[\exists\ p\ .\ F = (\boldsymbol{\lambda}\ x\ .\ p)\ \mathit{in}\ v]$
unfolding *Propositional-def* **.**
then obtain q **where** $[F = (\boldsymbol{\lambda}\ x\ .\ q)\ \mathit{in}\ v]$
by (*rule* $\exists E$)
hence $[\Box(F = (\boldsymbol{\lambda}\ x\ .\ q))\ \mathit{in}\ v]$
using *id-nec*[*equiv-lr*] **by** *auto*
hence $[\exists\ p\ .\ \Box(F = (\boldsymbol{\lambda}\ x\ .\ p))\ \mathit{in}\ v]$
using $\exists I$ **by** *fast*
thus $[\Box(\mathit{Propositional}\ F)\ \mathit{in}\ v]$
unfolding *Propositional-def*
using *sign-S5-thm-1*[*deduction*] **by** *fast*
qed

lemma *prop-indis*:
$[\mathit{Indiscriminate}\ F \rightarrow (\neg(\exists\ x\ y\ .\ (\!|F,x^P|\!)\ \&\ (\neg(\!|F,y^P|\!))))\ \mathit{in}\ v]$
proof (*rule CP*)
assume $[\mathit{Indiscriminate}\ F\ \mathit{in}\ v]$
hence *1*: $[\Box((\exists x.\ (\!|F,x^P|\!)) \rightarrow (\forall x.\ (\!|F,x^P|\!)))\ \mathit{in}\ v]$
unfolding *Indiscriminate-def* **.**
{
assume $[\exists\ x\ y\ .\ (\!|F,x^P|\!)\ \&\ \neg(\!|F,y^P|\!)\ \mathit{in}\ v]$
then obtain x **where** $[\exists\ y\ .\ (\!|F,x^P|\!)\ \&\ \neg(\!|F,y^P|\!)\ \mathit{in}\ v]$
by (*rule* $\exists E$)
then obtain y **where** *2*: $[(\!|F,x^P|\!)\ \&\ \neg(\!|F,y^P|\!)\ \mathit{in}\ v]$
by (*rule* $\exists E$)
hence $[\exists\ x\ .\ (\!|F,\ x^P|\!)\ \mathit{in}\ v]$
using $\&E(1)$ $\exists I$ **by** *fast*

```
    hence [∀ x . (|F,x^P|) in v]
      using 1[THEN qml-2[axiom-instance, deduction], deduction] by fast
    hence [(|F,y^P|) in v]
      using cqt-orig-1[deduction] by fast
    hence [(|F,y^P|) & (¬(|F,y^P|)) in v]
      using 2 &I &E by fast
    hence [¬(∃ x y . (|F,x^P|) & ¬(|F,y^P|)) in v]
      using pl-1[axiom-instance, deduction, THEN modus-tollens-1]
            oth-class-taut-1-a by blast
  }
  thus [¬(∃ x y . (|F,x^P|) & ¬(|F,y^P|)) in v]
    using reductio-aa-2 if-p-then-p deduction-theorem by blast
qed


lemma prop-in-thm:
  [Propositional F → Indiscriminate F in v]
  proof (rule CP)
    assume [Propositional F in v]
    hence [□(Propositional F) in v]
      using prop-prop2-3[deduction] by auto
    moreover {
      fix w
      assume [∃ p . (F = (λ y . p)) in w]
      then obtain q where q-prop: [F = (λ y . q) in w]
        by (rule ∃E)
      {
        assume [∃ x . (|F,x^P|) in w]
        then obtain a where [(|F,a^P|) in w]
          by (rule ∃E)
        hence [(|λ y . q, a^P|) in w]
          using q-prop l-identity[axiom-instance,deduction,deduction] by fast
        hence q: [q in w]
          apply (safe intro!: beta-C-meta-1[where φ=λy. q, equiv-lr])
           apply show-proper
          by simp
        {
          fix x
          have [(|λ y . q, x^P|) in w]
            apply (safe intro!: q beta-C-meta-1[equiv-rl])
            by show-proper
          hence [(|F,x^P|) in w]
            using q-prop[eq-sym] l-identity[axiom-instance, deduction, deduction]
            by fast
        }
        hence [∀ x . (|F,x^P|) in w]
          by (rule ∀I)
      }
      hence [(∃ x . (|F,x^P|)) → (∀ x . (|F, x^P|)) in w]
        by (rule CP)
    }
    ultimately show [Indiscriminate F in v]
      unfolding Propositional-def Indiscriminate-def
      using RM-1[deduction] deduction-theorem by blast
  qed

lemma prop-in-f-1:
  [Necessary F → Indiscriminate F in v]
```

unfolding *Necessary-defs Indiscriminate-def*
using *pl-1*[*axiom-instance*, *THEN RM-1*] **by** *simp*

lemma *prop-in-f-2*:
[*Impossible F* $\rightarrow$ *Indiscriminate F in v*]
proof −
{
fix *w*
have [(¬(∃ *x* . (|*F*,x^P|))) → ((∃ *x* . (|*F*,x^P|)) → (∀ *x* . (|*F*,x^P|))) *in w*]
using *useful-tautologies-3* **by** *auto*
hence [(∀ *x* . ¬(|*F*,x^P|)) → ((∃ *x* . (|*F*,x^P|)) → (∀ *x* . (|*F*,x^P|))) *in w*]
apply − **apply** (*PLM-subst-method* ¬(∃ *x*. (|*F*,x^P|)) (∀ *x*. ¬(|*F*,x^P|)))
using *cqt-further-4* **unfolding** *exists-def* **by** *fast+*
}
thus *?thesis*
unfolding *Impossible-defs Indiscriminate-def* **using** *RM-1 CP* **by** *blast*
qed

lemma *prop-in-f-3-a*:
[¬(*Indiscriminate* (*E*!)) *in v*]
proof (*rule reductio-aa-2*)
show [□¬(∀ *x*. (|*E*!,x^P|)) *in v*]
using *a-objects-exist-3* **.**
next
assume [*Indiscriminate E*! *in v*]
thus [¬□¬(∀ *x* . (|*E*!,x^P|)) *in v*]
unfolding *Indiscriminate-def*
using *o-objects-exist-1 KBasic2-5*[*deduction*,*deduction*]
unfolding *diamond-def* **by** *blast*
qed

lemma *prop-in-f-3-b*:
[¬(*Indiscriminate* (*E*!$^-$)) *in v*]
proof (*rule reductio-aa-2*)
assume [*Indiscriminate* (*E*!$^-$) *in v*]
moreover have [□(∃ *x* . (|*E*!$^-$, x^P|)) *in v*]
apply (*PLM-subst-method* λ *x* . ¬(|*E*!, x^P|) λ *x* . (|*E*!$^-$, x^P|))
using *thm-relation-negation-1-1*[*equiv-sym*] **apply** *simp*
unfolding *exists-def*
apply (*PLM-subst-method* λ *x* . (|*E*!, x^P|) λ *x* . ¬¬(|*E*!, x^P|))
using *oth-class-taut-4-b* **apply** *simp*
using *a-objects-exist-3* **by** *auto*
ultimately have [□(∀ *x*. (|*E*!$^-$,x^P|)) *in v*]
unfolding *Indiscriminate-def*
using *qml-1*[*axiom-instance*, *deduction*, *deduction*] **by** *blast*
thus [□(∀ *x*. ¬(|*E*!,x^P|)) *in v*]
apply −
apply (*PLM-subst-method* λ *x* . (|*E*!$^-$, x^P|) λ *x* . ¬(|*E*!, x^P|))
using *thm-relation-negation-1-1* **by** *auto*
next
show [¬□(∀ *x* . ¬(|*E*!, x^P|)) *in v*]
using *o-objects-exist-1*
unfolding *diamond-def exists-def*
apply −
apply (*PLM-subst-method* ¬¬(∀ *x*. ¬(|*E*!,x^P|)) ∀ *x*. ¬(|*E*!,x^P|))
using *oth-class-taut-4-b*[*equiv-sym*] **by** *auto*
qed

lemma *prop-in-f-3-c*:
[¬(*Indiscriminate* (*O*!)) *in v*]
proof (*rule reductio-aa-2*)
show [¬(∀ *x* . (|*O*!,x^P|)) *in v*]
using *a-objects-exist-2*[*THEN qml-2*[*axiom-instance*, *deduction*]]
by *blast*
next
assume [*Indiscriminate O*! *in v*]
thus [(∀ *x* . (|*O*!,x^P|)) *in v*]
unfolding *Indiscriminate-def*
using *o-objects-exist-2 qml-1*[*axiom-instance*, *deduction*, *deduction*]
qml-2[*axiom-instance*, *deduction*] **by** *blast*
qed

lemma *prop-in-f-3-d*:
[¬(*Indiscriminate* (*A*!)) *in v*]
proof (*rule reductio-aa-2*)
show [¬(∀ *x* . (|*A*!,x^P|)) *in v*]
using *o-objects-exist-3*[*THEN qml-2*[*axiom-instance*, *deduction*]]
by *blast*
next
assume [*Indiscriminate A*! *in v*]
thus [(∀ *x* . (|*A*!,x^P|)) *in v*]
unfolding *Indiscriminate-def*
using *a-objects-exist-1 qml-1*[*axiom-instance*, *deduction*, *deduction*]
qml-2[*axiom-instance*, *deduction*] **by** *blast*
qed

lemma *prop-in-f-4-a*:
[¬(*Propositional E*!) *in v*]
using *prop-in-thm*[*deduction*] *prop-in-f-3-a modus-tollens-1 CP*
by *meson*

lemma *prop-in-f-4-b*:
[¬(*Propositional* (*E*!$^-$)) *in v*]
using *prop-in-thm*[*deduction*] *prop-in-f-3-b modus-tollens-1 CP*
by *meson*

lemma *prop-in-f-4-c*:
[¬(*Propositional* (*O*!)) *in v*]
using *prop-in-thm*[*deduction*] *prop-in-f-3-c modus-tollens-1 CP*
by *meson*

lemma *prop-in-f-4-d*:
[¬(*Propositional* (*A*!)) *in v*]
using *prop-in-thm*[*deduction*] *prop-in-f-3-d modus-tollens-1 CP*
by *meson*

lemma *prop-prop-nec-1*:
[◇(∃ *p* . *F* = (**λ** *x* . *p*)) → (∃ *p* . *F* = (**λ** *x* . *p*)) *in v*]
proof (*rule CP*)
assume [◇(∃ *p* . *F* = (**λ** *x* . *p*)) *in v*]
hence [∃ *p* . ◇(*F* = (**λ** *x* . *p*)) *in v*]
using *BF*◇[*deduction*] **by** *auto*
then obtain *p* **where** [◇(*F* = (**λ** *x* . *p*)) *in v*]
by (*rule* ∃ *E*)
hence [◇□(∀ *x*. {|x^P,*F*|} ≡ {|x^P,**λ***x*. *p*|}) *in v*]
unfolding *identity-defs* **.**

hence $[\Box(\forall\, x.\ \{\!|x^P,F|\!\} \equiv \{\!|x^P,\boldsymbol{\lambda} x.\ p|\!\})\ \textit{in}\ v]$
using *5◇[deduction]* **by** *auto*
hence $[(F = (\boldsymbol{\lambda}\ x\ .\ p))\ \textit{in}\ v]$
unfolding *identity-defs* **.**
thus $[\exists\ p\ .\ (F = (\boldsymbol{\lambda}\ x\ .\ p))\ \textit{in}\ v]$
by *PLM-solver*
qed

lemma *prop-prop-nec-2*:
$[(\forall\ p\ .\ F \neq (\boldsymbol{\lambda}\ x\ .\ p)) \rightarrow \Box(\forall\ p\ .\ F \neq (\boldsymbol{\lambda}\ x\ .\ p))\ \textit{in}\ v]$
apply (*PLM-subst-method*
$\neg(\exists\ p\ .\ (F = (\boldsymbol{\lambda}\ x\ .\ p)))$
$(\forall\ p\ .\ \neg(F = (\boldsymbol{\lambda}\ x\ .\ p))))$
using *cqt-further-4* **apply** *blast*
apply (*PLM-subst-method*
$\neg\Diamond(\exists\, p.\ F = (\boldsymbol{\lambda}x.\ p))$
$\Box\neg(\exists\, p.\ F = (\boldsymbol{\lambda}x.\ p)))$
using *KBasic2-4*[*equiv-sym*] *prop-prop-nec-1*
contraposition-1 **by** *auto*

lemma *prop-prop-nec-3*:
$[(\exists\ p\ .\ F = (\boldsymbol{\lambda}\ x\ .\ p)) \rightarrow \Box(\exists\ p\ .\ F = (\boldsymbol{\lambda}\ x\ .\ p))\ \textit{in}\ v]$
using *prop-prop-nec-1 derived-S5-rules-1-b* **by** *simp*

lemma *prop-prop-nec-4*:
$[\Diamond(\forall\ p\ .\ F \neq (\boldsymbol{\lambda}\ x\ .\ p)) \rightarrow (\forall\ p\ .\ F \neq (\boldsymbol{\lambda}\ x\ .\ p))\ \textit{in}\ v]$
using *prop-prop-nec-2 derived-S5-rules-2-b* **by** *simp*

lemma *enc-prop-nec-1*:
$[\Diamond(\forall\ F\ .\ \{\!|x^P,\ F|\!\} \rightarrow (\exists\ p\ .\ F = (\boldsymbol{\lambda}\ x\ .\ p)))$
$\rightarrow (\forall\ F\ .\ \{\!|x^P,\ F|\!\} \rightarrow (\exists\ p\ .\ F = (\boldsymbol{\lambda}\ x\ .\ p)))\ \textit{in}\ v]$
proof (*rule CP*)
assume $[\Diamond(\forall\, F.\ \{\!|x^P,F|\!\} \rightarrow (\exists\, p.\ F = (\boldsymbol{\lambda}x.\ p)))\ \textit{in}\ v]$
hence *1*: $[(\forall\, F.\ \Diamond(\{\!|x^P,F|\!\} \rightarrow (\exists\, p.\ F = (\boldsymbol{\lambda}x.\ p))))\ \textit{in}\ v]$
using *Buridan◇*[*deduction*] **by** *auto*
{
fix *Q*
assume $[\{\!|x^P,Q|\!\}\ \textit{in}\ v]$
hence $[\Box\{\!|x^P,Q|\!\}\ \textit{in}\ v]$
using *encoding*[*axiom-instance*, *deduction*] **by** *auto*
moreover have $[\Diamond(\{\!|x^P,Q|\!\} \rightarrow (\exists\, p.\ Q = (\boldsymbol{\lambda}x.\ p)))\ \textit{in}\ v]$
using *cqt-1*[*axiom-instance*, *deduction*] *1* **by** *fast*
ultimately have $[\Diamond(\exists\, p.\ Q = (\boldsymbol{\lambda}x.\ p))\ \textit{in}\ v]$
using *KBasic2-9*[*equiv-lr*,*deduction*] **by** *auto*
hence $[(\exists\, p.\ Q = (\boldsymbol{\lambda}x.\ p))\ \textit{in}\ v]$
using *prop-prop-nec-1*[*deduction*] **by** *auto*
}
thus $[(\forall\ F\ .\ \{\!|x^P,\ F|\!\} \rightarrow (\exists\ p\ .\ F = (\boldsymbol{\lambda}\ x\ .\ p)))\ \textit{in}\ v]$
apply − **by** *PLM-solver*
qed

lemma *enc-prop-nec-2*:
$[(\forall\ F\ .\ \{\!|x^P,\ F|\!\} \rightarrow (\exists\ p\ .\ F = (\boldsymbol{\lambda}\ x\ .\ p))) \rightarrow \Box(\forall\ F\ .\ \{\!|x^P,\ F|\!\}$
$\rightarrow (\exists\ p\ .\ F = (\boldsymbol{\lambda}\ x\ .\ p)))\ \textit{in}\ v]$
using *derived-S5-rules-1-b enc-prop-nec-1* **by** *blast*
end
end

A.10. Possible Worlds

locale *PossibleWorlds* = *PLM*
begin

A.10.1. Definitions

definition *Situation* **where**
Situation $x \equiv (\!|A!,x|\!)$ **&** $(\forall\ F.\ \{\!|x,F|\!\} \rightarrow \textit{Propositional}\ F)$

definition *EncodeProposition* (**infixl** $\mathbf{\Sigma}$ *70*) **where**
$x\mathbf{\Sigma}p \equiv (\!|A!,x|\!)$ **&** $\{\!|x, \boldsymbol{\lambda}\ x\ .\ p|\!\}$
definition *TrueInSituation* (**infixl** $\models$ *10*) **where**
$x \models p \equiv \textit{Situation}\ x$ **&** $x\mathbf{\Sigma}p$
definition *PossibleWorld* **where**
PossibleWorld $x \equiv \textit{Situation}\ x$ **&** $\Diamond(\forall\ p\ .\ x\mathbf{\Sigma}p \equiv p)$

A.10.2. Auxiliary Lemmas

lemma *possit-sit-1*:
$[\textit{Situation}\ (x^P) \equiv \Box(\textit{Situation}\ (x^P))\ \textit{in}\ v]$
proof (*rule* $\equiv I$; *rule CP*)
assume $[\textit{Situation}\ (x^P)\ \textit{in}\ v]$
hence *1*: $[(\!|A!,x^P|\!)$ **&** $(\forall\ F.\ \{\!|x^P,F|\!\} \rightarrow \textit{Propositional}\ F)\ \textit{in}\ v]$
unfolding *Situation-def* **by** *auto*
have $[\Box(\!|A!,x^P|\!)\ \textit{in}\ v]$
using *1*[*conj1*, *THEN oa-facts-2*[*deduction*]] **.**
moreover have $[\Box(\forall\ F.\ \{\!|x^P,F|\!\} \rightarrow \textit{Propositional}\ F)\ \textit{in}\ v]$
using *1*[*conj2*] **unfolding** *Propositional-def*
by (*rule enc-prop-nec-2*[*deduction*])
ultimately show $[\Box\textit{Situation}\ (x^P)\ \textit{in}\ v]$
unfolding *Situation-def*
apply *cut-tac* **apply** (*rule KBasic-3*[*equiv-rl*])
by (*rule intro-elim-1*)
next
assume $[\Box\textit{Situation}\ (x^P)\ \textit{in}\ v]$
thus $[\textit{Situation}\ (x^P)\ \textit{in}\ v]$
using *qml-2*[*axiom-instance*, *deduction*] **by** *auto*
qed

lemma *possworld-nec*:
$[\textit{PossibleWorld}\ (x^P) \equiv \Box(\textit{PossibleWorld}\ (x^P))\ \textit{in}\ v]$
apply (*rule* $\equiv I$; *rule CP*)
subgoal unfolding *PossibleWorld-def*
apply (*rule KBasic-3*[*equiv-rl*])
apply (*rule intro-elim-1*)
using *possit-sit-1*[*equiv-lr*] **&**$E(1)$ **apply** *blast*
using *qml-3*[*axiom-instance*, *deduction*] **&**$E(2)$ **by** *blast*
using *qml-2*[*axiom-instance*,*deduction*] **by** *auto*

lemma *TrueInWorldNecc*:
$[((x^P) \models p) \equiv \Box((x^P) \models p)\ \textit{in}\ v]$
proof (*rule* $\equiv I$; *rule CP*)
assume $[x^P \models p\ \textit{in}\ v]$
hence $[\textit{Situation}\ (x^P)$ **&** $((\!|A!,x^P|\!)$ **&** $\{\!|x^P,\boldsymbol{\lambda}x.\ p|\!\})\ \textit{in}\ v]$
unfolding *TrueInSituation-def EncodeProposition-def* **.**
hence $[(\Box\textit{Situation}\ (x^P)$ **&** $\Box(\!|A!,x^P|\!))$ **&** $\Box\{\!|x^P, \boldsymbol{\lambda}x.\ p|\!\}\ \textit{in}\ v]$

using &I &E *possit-sit-1*[*equiv-lr*] *oa-facts-2*[*deduction*]
encoding[*axiom-instance*,*deduction*] **by** *metis*
thus [$\Box((x^P) \models p)$ *in* v]
unfolding *TrueInSituation-def EncodeProposition-def*
using *KBasic-3*[*equiv-rl*] &I &E **by** *metis*
next
assume [$\Box(x^P \models p)$ *in* v]
thus [$x^P \models p$ *in* v]
using *qml-2*[*axiom-instance*,*deduction*] **by** *auto*
qed

lemma *PossWorldAux*:
[($(\!|A!,x^P|\!)$ & ($\forall$ F . ($\{\!|x^P,F|\!\} \equiv (\exists$ p . p & $(F = (\boldsymbol{\lambda}$ x . $p))))))$
$\rightarrow$ (*PossibleWorld* (x^P)) *in* v]
proof (*rule CP*)
assume *DefX*: [$(\!|A!,x^P|\!)$ & ($\forall$ F . ($\{\!|x^P,F|\!\} \equiv$
$(\exists$ p . p & $(F = (\boldsymbol{\lambda}$ x . $p)))))$ *in* v]

have [*Situation* (x^P) *in* v]
proof −
have [$(\!|A!,x^P|\!)$ *in* v]
using *DefX*[*conj1*] **.**
moreover have [($\forall$ F. $\{\!|x^P,F|\!\} \rightarrow$ *Propositional* F) *in* v]
proof (*rule* $\forall I$; *rule CP*)
fix F
assume [$\{\!|x^P,F|\!\}$ *in* v]
moreover have [$\{\!|x^P,F|\!\} \equiv (\exists$ p . p & $(F = (\boldsymbol{\lambda}$ x . $p)))$ *in* v]
using *DefX*[*conj2*] *cqt-1*[*axiom-instance*, *deduction*] **by** *auto*
ultimately have [$(\exists$ p . p & $(F = (\boldsymbol{\lambda}$ x . $p)))$ *in* v]
using $\equiv E(1)$ **by** *blast*
then obtain p **where** [p & $(F = (\boldsymbol{\lambda}$ x . $p))$ *in* v]
by (*rule* $\exists E$)
hence [$(F = (\boldsymbol{\lambda}$ x . $p))$ *in* v]
by (*rule* &$E(2)$)
hence [$(\exists$ p . $(F = (\boldsymbol{\lambda}$ x . $p)))$ *in* v]
by *PLM-solver*
thus [*Propositional* F *in* v]
unfolding *Propositional-def* **.**
qed
ultimately show [*Situation* (x^P) *in* v]
unfolding *Situation-def* **by** (*rule* &I)
qed
moreover have [$\Diamond(\forall p.\ x^P\ \boldsymbol{\Sigma}\ p \equiv p)$ *in* v]
unfolding *EncodeProposition-def*
proof (*rule TBasic*[*deduction*]; *rule* $\forall I$)
fix q
have *EncodeLambda*:
[$\{\!|x^P, \boldsymbol{\lambda}x.\ q|\!\} \equiv (\exists$ p . p & $((\boldsymbol{\lambda}x.\ q) = (\boldsymbol{\lambda}$ x . $p)))$ *in* v]
using *DefX*[*conj2*] **by** (*rule cqt-1*[*axiom-instance*, *deduction*])
moreover {
assume [q *in* v]
moreover have [$(\boldsymbol{\lambda}x.\ q) = (\boldsymbol{\lambda}$ x . $q)$ *in* v]
using *id-eq-prop-prop-1* **by** *auto*
ultimately have [q & $((\boldsymbol{\lambda}x.\ q) = (\boldsymbol{\lambda}$ x . $q))$ *in* v]
by (*rule* &I)
hence [$\exists$ p . p & $((\boldsymbol{\lambda}x.\ q) = (\boldsymbol{\lambda}$ x . $p))$ *in* v]
by *PLM-solver*

```
      moreover have [(|A!,x^P|) in v]
        using DefX[conj1] .
      ultimately have [(|A!,x^P|) & {|x^P, λx. q|} in v]
        using EncodeLambda[equiv-rl] &I by auto
    }
    moreover {
      assume [(|A!,x^P|) & {|x^P, λx. q|} in v]
      hence [{|x^P, λx. q|} in v]
        using &E(2) by auto
      hence [∃ p . p & ((λx. q) = (λ x . p)) in v]
        using EncodeLambda[equiv-lr] by auto
      then obtain p where p-and-lambda-q-is-lambda-p:
        [p & ((λx. q) = (λ x . p)) in v]
        by (rule ∃E)
      have [(|(λ x . p), x^P|) ≡ p in v]
        apply (rule beta-C-meta-1)
        by show-proper
      hence [(|(λ x . p), x^P|) in v]
        using p-and-lambda-q-is-lambda-p[conj1] ≡E(2) by auto
      hence [(|(λ x . q), x^P|) in v]
        using p-and-lambda-q-is-lambda-p[conj2, THEN id-eq-prop-prop-2[deduction]]
          l-identity[axiom-instance, deduction, deduction] by fast
      moreover have [(|(λ x . q), x^P|) ≡ q in v]
        apply (rule beta-C-meta-1) by show-proper
      ultimately have [q in v]
        using ≡E(1) by blast
    }
    ultimately show [(|A!,x^P|) & {|x^P,λx. q|} ≡ q in v]
      using &I ≡I CP by auto
  qed

  ultimately show [PossibleWorld (x^P) in v]
    unfolding PossibleWorld-def by (rule &I)
qed
```

A.10.3. For every syntactic Possible World there is a semantic Possible World

```
theorem SemanticPossibleWorldForSyntacticPossibleWorlds:
 ∀ x . [PossibleWorld (x^P) in w] ⟶
  (∃ v . ∀ p . [(x^P |= p) in w] ⟷ [p in v])
 proof
  fix x
  {
   assume PossWorldX: [PossibleWorld (x^P) in w]
   hence SituationX: [Situation (x^P) in w]
     unfolding PossibleWorld-def apply cut-tac by PLM-solver
   have PossWorldExpanded:
    [(|A!,x^P|) & (∀ F. {|x^P,F|} → (∃ p. F = (λx. p)))
      & ◇(∀ p. (|A!,x^P|) & {|x^P,λx. p|} ≡ p) in w]
     using PossWorldX
     unfolding PossibleWorld-def Situation-def
           Propositional-def EncodeProposition-def .
   have AbstractX: [(|A!,x^P|) in w]
     using PossWorldExpanded[conj1,conj1] .

   have [◇(∀ p. {|x^P,λx. p|} ≡ p) in w]
     apply (PLM-subst-method
          λp. (|A!,x^P|) & {|x^P,λx. p|}
```

```
        λ p . ⦇x^P,λx. p⦈)
    subgoal using PossWorldExpanded[conj1,conj1,THEN oa-facts-2[deduction]]
          using Semantics.T6 apply cut-tac by PLM-solver
    using PossWorldExpanded[conj2] .

  hence ∃ v. ∀ p. ([⦇x^P,λx. p⦈ in v])
              = [p in v]
   unfolding diamond-def equiv-def conj-def
   apply (simp add: Semantics.T4 Semantics.T6 Semantics.T5
                Semantics.T8)
   by auto

  then obtain v where PropsTrueInSemWorld:
   ∀ p. ([⦇x^P,λx. p⦈ in v]) = [p in v]
   by auto
  {
   fix p
   {
    assume [((x^P) ⊨ p) in w]
    hence [((x^P) ⊨ p) in v]
     using TrueInWorldNecc[equiv-lr] Semantics.T6 by simp
    hence [Situation (x^P) & (⦇A!,x^P⦈) & ⦇x^P,λx. p⦈) in v]
     unfolding TrueInSituation-def EncodeProposition-def .
    hence [⦇x^P,λx. p⦈ in v]
     using &E(2) by blast
    hence [p in v]
     using PropsTrueInSemWorld by blast
   }
   moreover {
    assume [p in v]
    hence [⦇x^P,λx. p⦈ in v]
     using PropsTrueInSemWorld by blast
    hence [(x^P) ⊨ p in v]
     apply cut-tac unfolding TrueInSituation-def EncodeProposition-def
     apply (rule &I) using SituationX[THEN possit-sit-1[equiv-lr]]
     subgoal using Semantics.T6 by auto
     apply (rule &I)
     subgoal using AbstractX[THEN oa-facts-2[deduction]]
       using Semantics.T6 by auto
     by assumption
    hence [□((x^P) ⊨ p) in v]
     using TrueInWorldNecc[equiv-lr] by simp
    hence [(x^P) ⊨ p in w]
     using Semantics.T6 by simp
   }
   ultimately have [p in v] ⟷ [(x^P) ⊨ p in w]
    by auto
  }
  hence (∃ v . ∀ p . [p in v] ⟷ [(x^P) ⊨ p in w])
   by blast
 }
 thus [PossibleWorld (x^P) in w] ⟶
     (∃ v. ∀ p . [(x^P) ⊨ p in w] ⟷ [p in v])
  by blast
qed
```

A.10.4. For every semantic Possible World there is a syntactic Possible World

theorem *SyntacticPossibleWorldForSemanticPossibleWorlds*:
$\forall\ v\ .\ \exists\ x\ .\ [PossibleWorld\ (x^P)\ in\ w] \wedge$
$(\forall\ p\ .\ [p\ in\ v] \longleftrightarrow [((x^P) \models p)\ in\ w])$
proof
fix v
have $[\exists\, x.\ (\!|A!,x^P|\!)\ \&\ (\forall\ F\ .\ (\{\!|x^P,F|\!\} \equiv$
$(\exists\ p\ .\ p\ \&\ (F = (\boldsymbol{\lambda}\ x\ .\ p)))))\ in\ v]$
using *A-objects*[*axiom-instance*] **by** *fast*
then obtain x **where** *DefX*:
$[(\!|A!,x^P|\!)\ \&\ (\forall\ F\ .\ (\{\!|x^P,F|\!\} \equiv (\exists\ p\ .\ p\ \&\ (F = (\boldsymbol{\lambda}\ x\ .\ p)))))\ in\ v]$
by (*rule* $\exists\, E$)
hence *PossWorldX*: $[PossibleWorld\ (x^P)\ in\ v]$
using *PossWorldAux*[*deduction*] **by** *blast*
hence $[PossibleWorld\ (x^P)\ in\ w]$
using *possworld-nec*[*equiv-lr*] *Semantics.T6* **by** *auto*
moreover have $(\forall\ p\ .\ [p\ in\ v] \longleftrightarrow [(x^P) \models p\ in\ w])$
proof
fix q
{
assume $[q\ in\ v]$
moreover have $[(\boldsymbol{\lambda}\ x\ .\ q) = (\boldsymbol{\lambda}\ x\ .\ q)\ in\ v]$
using *id-eq-prop-prop-1* **by** *auto*
ultimately have $[q\ \&\ (\boldsymbol{\lambda}\ x\ .\ q) = (\boldsymbol{\lambda}\ x\ .\ q)\ in\ v]$
using $\&I$ **by** *auto*
hence $[(\exists\ p\ .\ p\ \&\ ((\boldsymbol{\lambda}\ x\ .\ q) = (\boldsymbol{\lambda}\ x\ .\ p)))\ in\ v]$
by *PLM-solver*
hence *4*: $[\{\!|x^P,\ (\boldsymbol{\lambda}\ x\ .\ q)|\!\}\ in\ v]$
using *cqt-1*[*axiom-instance*,*deduction*, *OF DefX*[*conj2*], *equiv-rl*]
by *blast*
have $[(x^P \models q)\ in\ v]$
unfolding *TrueInSituation-def* **apply** (*rule* $\&I$)
using *PossWorldX* **unfolding** *PossibleWorld-def*
using $\&E(1)$ **apply** *blast*
unfolding *EncodeProposition-def* **apply** (*rule* $\&I$)
using *DefX*[*conj1*] **apply** *simp*
using *4* **.**
hence $[(x^P \models q)\ in\ w]$
using *TrueInWorldNecc*[*equiv-lr*] *Semantics.T6* **by** *auto*
}
moreover {
assume $[(x^P \models q)\ in\ w]$
hence $[(x^P \models q)\ in\ v]$
using *TrueInWorldNecc*[*equiv-lr*] *Semantics.T6*
by *auto*
hence $[\{\!|x^P,\ (\boldsymbol{\lambda}\ x\ .\ q)|\!\}\ in\ v]$
unfolding *TrueInSituation-def EncodeProposition-def*
using $\&E(2)$ **by** *blast*
hence $[(\exists\ p\ .\ p\ \&\ ((\boldsymbol{\lambda}\ x\ .\ q) = (\boldsymbol{\lambda}\ x\ .\ p)))\ in\ v]$
using *cqt-1*[*axiom-instance*,*deduction*, *OF DefX*[*conj2*], *equiv-lr*]
by *blast*
then obtain p **where** *4*:
$[(p\ \&\ ((\boldsymbol{\lambda}\ x\ .\ q) = (\boldsymbol{\lambda}\ x\ .\ p)))\ in\ v]$
by (*rule* $\exists\, E$)
have $[(\!|(\boldsymbol{\lambda}\ x\ .\ p),x^P|\!) \equiv p\ in\ v]$
apply (*rule beta-C-meta-1*)
by *show-proper*

hence $[(\!|(\boldsymbol{\lambda}\ x\ .\ q), x^P|\!) \equiv p\ in\ v]$
using *l-identity*[**where** $\beta{=}(\boldsymbol{\lambda}\ x\ .\ q)$ **and** $\alpha{=}(\boldsymbol{\lambda}\ x\ .\ p)$,
axiom-instance, *deduction*, *deduction*]
using *4*[*conj2*,*THEN id-eq-prop-prop-2*[*deduction*]] **by** *meson*
hence $[(\!|(\boldsymbol{\lambda}\ x\ .\ q), x^P|\!)\ in\ v]$ **using** *4*[*conj1*] $\equiv E(2)$ **by** *blast*
moreover have $[(\!|(\boldsymbol{\lambda}\ x\ .\ q), x^P|\!) \equiv q\ in\ v]$
apply (*rule beta-C-meta-1*)
by *show-proper*
ultimately have $[q\ in\ v]$
using $\equiv E(1)$ **by** *blast*
}
ultimately show $[q\ in\ v] \longleftrightarrow [(x^P) \models q\ in\ w]$
by *blast*
qed
ultimately show $\exists\ x\ .\ [PossibleWorld\ (x^P)\ in\ w]$
$\wedge\ (\forall\ p\ .\ [p\ in\ v] \longleftrightarrow [(x^P) \models p\ in\ w])$
by *auto*
qed
end

A.11. Artificial Theorems

Remark. *Some examples of theorems that can be derived from the model structure, but which are not derivable from the deductive system PLM itself.*

locale *ArtificialTheorems*
begin

lemma *lambda-enc-1*:
$[(\!|\boldsymbol{\lambda} x\ .\ \{\!|x^P, F|\!\} \equiv \{\!|x^P, F|\!\}, y^P|\!)\ in\ v]$
by (*auto simp*: *meta-defs meta-aux conn-defs forall-*Π_1*-def*)

lemma *lambda-enc-2*:
$[(\!|\boldsymbol{\lambda}\ x\ .\ \{\!|y^P, G|\!\}, x^P|\!) \equiv \{\!|y^P, G|\!\}\ in\ v]$
by (*auto simp*: *meta-defs meta-aux conn-defs forall-*Π_1*-def*)

Remark. *The following is* not *a theorem and nitpick can find a countermodel. This is expected and important. If this were a theorem, the theory would become inconsistent.*

lemma *lambda-enc-3*:
$[((\!|\boldsymbol{\lambda}\ x\ .\ \{\!|x^P, F|\!\}, x^P|\!) \rightarrow \{\!|x^P, F|\!\})\ in\ v]$
apply (*simp add*: *meta-defs meta-aux conn-defs forall-*Π_1*-def*)
nitpick[*user-axioms*, *expect=genuine*]
oops — countermodel by nitpick

Remark. *Instead the following two statements hold.*

lemma *lambda-enc-4*:
$[(\!|(\boldsymbol{\lambda}\ x\ .\ \{\!|x^P, F|\!\}), x^P|\!)\ in\ v] = (\exists\ y\ .\ \nu\upsilon\ y = \nu\upsilon\ x \wedge [\{\!|y^P, F|\!\}\ in\ v])$
by (*simp add*: *meta-defs meta-aux*)

lemma *lambda-ex*:
$[(\!|(\boldsymbol{\lambda}\ x\ .\ \varphi\ (x^P)), x^P|\!)\ in\ v] = (\exists\ y\ .\ \nu\upsilon\ y = \nu\upsilon\ x \wedge [\varphi\ (y^P)\ in\ v])$
by (*simp add*: *meta-defs meta-aux*)

Remark. *These statements can be translated to statements in the embedded logic.*

lemma *lambda-ex-emb*:
[(|(**λ** x . φ (x^P)), x^P|) ≡ (∃ y . (∀ F . (|F,x^P|) ≡ (|F,y^P|)) & φ (y^P)) *in* v]
proof(*rule MetaSolver.EquivI*)
interpret *MetaSolver* **.**
{
assume [(|(**λ** x . φ (x^P)), x^P|) *in* v]
then obtain y **where** $\nu\upsilon$ $y = \nu\upsilon$ $x \wedge$ [φ (y^P) *in* v]
using *lambda-ex* **by** *blast*
moreover hence [(∀ F . (|F,x^P|) ≡ (|F,y^P|)) *in* v]
apply − **apply** *meta-solver*
by (*simp add*: *Semantics.*d_κ*-proper Semantics.ex1-def*)
ultimately have [∃ y . (∀ F . (|F,x^P|) ≡ (|F,y^P|)) & φ (y^P) *in* v]
using *ExIRule ConjI* **by** *fast*
}
moreover {
assume [∃ y . (∀ F . (|F,x^P|) ≡ (|F,y^P|)) & φ (y^P) *in* v]
then obtain y **where** *y-def*: [(∀ F . (|F,x^P|) ≡ (|F,y^P|)) & φ (y^P) *in* v]
by (*rule ExERule*)
hence ⋀ F . [(|F,x^P|) *in* v] = [(|F,y^P|) *in* v]
apply − **apply** (*drule ConjE*) **apply** (*drule conjunct1*)
apply (*drule AllE*) **apply** (*drule EquivE*) **by** *simp*
hence [(|*make*Π_1 (λ u s w . $\nu\upsilon$ $y = u$),x^P|) *in* v]
= [(|*make*Π_1 (λ u s w . $\nu\upsilon$ $y = u$),y^P|) *in* v] **by** *auto*
hence $\nu\upsilon$ $y = \nu\upsilon$ x **by** (*simp add*: *meta-defs meta-aux*)
moreover have [φ (y^P) *in* v] **using** *y-def ConjE* **by** *blast*
ultimately have [(|(**λ** x . φ (x^P)), x^P|) *in* v]
using *lambda-ex* **by** *blast*
}
ultimately show [(|**λ**x. φ (x^P),x^P|) *in* v]
= [∃ y. (∀ F. (|F,x^P|) ≡ (|F,y^P|)) & φ (y^P) *in* v]
by *auto*
qed

lemma *lambda-enc-emb*:
[(|(**λ** x . {|x^P, F|}), x^P|) ≡ (∃ y . (∀ F . (|F,x^P|) ≡ (|F,y^P|)) & {|y^P, F|}) *in* v]
using *lambda-ex-emb* **by** *fast*

Remark. *In the case of proper maps, the generalized* β*-conversion reduces to classical* β*-conversion.*

lemma *proper-beta*:
assumes *IsProperInX* φ
shows [(∃ y . (∀ F . (|F,x^P|) ≡ (|F,y^P|)) & φ (y^P)) ≡ φ (x^P) *in* v]
proof (*rule MetaSolver.EquivI*; *rule*)
interpret *MetaSolver* **.**
assume [∃ y. (∀ F. (|F,x^P|) ≡ (|F,y^P|)) & φ (y^P) *in* v]
then obtain y **where** *y-def*: [(∀ F. (|F,x^P|) ≡ (|F,y^P|)) & φ (y^P) *in* v] **by** (*rule ExERule*)
hence [(|*make*Π_1 (λ u s w . $\nu\upsilon$ $y = u$), x^P|) *in* v] = [(|*make*Π_1 (λ u s w . $\nu\upsilon$ $y = u$), y^P|) *in* v]
using *EquivS AllE ConjE* **by** *blast*
hence $\nu\upsilon$ $y = \nu\upsilon$ x **by** (*simp add*: *meta-defs meta-aux*)
thus [φ (x^P) *in* v]
using *y-def*[*THEN ConjE*[*THEN conjunct2*]]
assms IsProperInX.rep-eq valid-in.rep-eq
by *blast*
next
interpret *MetaSolver* **.**

assume [$\varphi\ (x^P)$ *in* v]
moreover have [$\forall F.\ (\!|F,x^P|\!) \equiv (\!|F,x^P|\!)$ *in* v] **apply** *meta-solver* **by** *blast*
ultimately show [$\exists y.\ (\forall F.\ (\!|F,x^P|\!) \equiv (\!|F,y^P|\!))$ **&** $\varphi\ (y^P)$ *in* v]
by (*meson ConjI ExI*)
qed

Remark. *The following theorem is a consequence of the constructed Aczel-model, but not part of PLM. Separate research on possible modifications of the embedding suggest that this artificial theorem can be avoided by introducing a dependency on states for the mapping from abstract objects to special urelements.*

lemma *lambda-rel-extensional*:
assumes [$\forall F\ .\ (\!|F,a^P|\!) \equiv (\!|F,b^P|\!)$ *in* v]
shows ($\boldsymbol{\lambda} x.\ (\!|R,x^P,a^P|\!)) = (\boldsymbol{\lambda} x\ .\ (\!|R,x^P,\ b^P|\!))$
proof −
interpret *MetaSolver* **.**
obtain F **where** *F-def*: $F = make\Pi_1\ (\lambda\ u\ s\ w\ .\ u = \nu v\ a)$ **by** *auto*
have [$(\!|F,\ a^P|\!) \equiv (\!|F,\ b^P|\!)$ *in* v] **using** *assms* **by** (*rule AllE*)
moreover have [$(\!|F,\ a^P|\!)$ *in* v]
unfolding *F-def* **by** (*simp add*: *meta-defs meta-aux*)
ultimately have [$(\!|F,\ b^P|\!)$ *in* v] **using** *EquivE* **by** *auto*
hence $\nu v\ a = \nu v\ b$ **using** *F-def* **by** (*simp add*: *meta-defs meta-aux*)
thus *?thesis* **by** (*simp add*: *meta-defs meta-aux*)
qed

end

A.12. Sanity Tests

locale *SanityTests*
begin
interpretation *MetaSolver*.
interpretation *Semantics*.

A.12.1. Consistency

lemma *True*
nitpick[*expect*=*genuine*, *user-axioms*, *satisfy*]
by *auto*

A.12.2. Intensionality

lemma [$(\boldsymbol{\lambda} y.\ (q \vee \neg q)) = (\boldsymbol{\lambda} y.\ (p \vee \neg p))$ *in* v]
unfolding *identity*-Π_1-*def conn-defs*
apply (*rule* $Eq_1 I$) **apply** (*simp add*: *meta-defs*)
nitpick[*expect* = *genuine*, *user-axioms*=*true*, *card i* = *2*,
card j = *2*, *card* ω = *1*, *card* σ = *1*,
sat-solver = *MiniSat-JNI*, *verbose*, *show-all*]
oops — Countermodel by Nitpick
lemma [$(\boldsymbol{\lambda} y.\ (p \vee q)) = (\boldsymbol{\lambda} y.\ (q \vee p))$ *in* v]
unfolding *identity*-Π_1-*def*
apply (*rule* $Eq_1 I$) **apply** (*simp add*: *meta-defs*)
nitpick[*expect* = *genuine*, *user-axioms*=*true*,
sat-solver = *MiniSat-JNI*, *card i* = *2*,

card j = 2, *card σ = 1*, *card ω = 1*,
card υ = 2, *verbose*, *show-all*]
oops — Countermodel by Nitpick

A.12.3. Concreteness coindices with Object Domains

lemma *OrdCheck*:
[(|**λ** *x* . ¬□(¬(|*E*!, x^P|)), *x*|) *in v*] ⟷
(*proper x*) ∧ (*case* (*rep x*) *of* *ων y* ⇒ *True* | - ⇒ *False*)
using *OrdinaryObjectsPossiblyConcreteAxiom*
apply (*simp add*: *meta-defs meta-aux split*: *ν.split υ.split*)
using *νυ-ων-is-ωυ* **by** *fastforce*
lemma *AbsCheck*:
[(|**λ** *x* . □(¬(|*E*!, x^P|)), *x*|) *in v*] ⟷
(*proper x*) ∧ (*case* (*rep x*) *of* *αν y* ⇒ *True* | - ⇒ *False*)
using *OrdinaryObjectsPossiblyConcreteAxiom*
apply (*simp add*: *meta-defs meta-aux split*: *ν.split υ.split*)
using *no-αω* **by** *blast*

A.12.4. Justification for Meta-Logical Axioms

Remark. *OrdinaryObjectsPossiblyConcreteAxiom is equivalent to "all ordinary objects are possibly concrete".*

lemma *OrdAxiomCheck*:
OrdinaryObjectsPossiblyConcrete ⟷
(∀ *x*. ([(|**λ** *x* . ¬□(¬(|*E*!, x^P|)), x^P|) *in v*]
⟷ (*case x of* *ων y* ⇒ *True* | - ⇒ *False*)))
unfolding *Concrete-def*
apply (*simp add*: *meta-defs meta-aux split*: *ν.split υ.split*)
using *νυ-ων-is-ωυ* **by** *fastforce*

Remark. *OrdinaryObjectsPossiblyConcreteAxiom is equivalent to "all abstract objects are necessarily not concrete".*

lemma *AbsAxiomCheck*:
OrdinaryObjectsPossiblyConcrete ⟷
(∀ *x*. ([(|**λ** *x* . □(¬(|*E*!, x^P|)), x^P|) *in v*]
⟷ (*case x of* *αν y* ⇒ *True* | - ⇒ *False*)))
apply (*simp add*: *meta-defs meta-aux split*: *ν.split υ.split*)
using *νυ-ων-is-ωυ no-αω* **by** *fastforce*

Remark. *PossiblyContingentObjectExistsAxiom is equivalent to the corresponding statement in the embedded logic.*

lemma *PossiblyContingentObjectExistsCheck*:
PossiblyContingentObjectExists ⟷ [¬(□(∀ *x*. (|*E*!,x^P|) → □(|*E*!,x^P|))) *in v*]
apply (*simp add*: *meta-defs forall-ν-def meta-aux split*: *ν.split υ.split*)
by (*metis ν.simps*(*5*) *νυ-def υ.simps*(*1*) *no-σω υ.exhaust*)

Remark. *PossiblyNoContingentObjectExistsAxiom is equivalent to the corresponding statement in the embedded logic.*

lemma *PossiblyNoContingentObjectExistsCheck*:
PossiblyNoContingentObjectExists ⟷ [¬(□(¬(∀ *x*. (|*E*!,x^P|) → □(|*E*!,x^P|)))) *in v*]
apply (*simp add*: *meta-defs forall-ν-def meta-aux split*: *ν.split υ.split*)
using *νυ-ων-is-ωυ* **by** *blast*

A.12.5. Relations in the Meta-Logic

Remark. *Material equality in the embedded logic corresponds to equality in the actual state in the meta-logic.*

lemma *mat-eq-is-eq-dj*:
$[\forall\ x\ .\ \Box((\!|F,x^P|\!) \equiv (\!|G,x^P|\!))\ in\ v] \longleftrightarrow$
$((\lambda\ x\ .\ (eval\Pi_1\ F)\ x\ dj) = (\lambda\ x\ .\ (eval\Pi_1\ G)\ x\ dj))$
proof
assume *1*: $[\forall x.\ \Box((\!|F,x^P|\!) \equiv (\!|G,x^P|\!))\ in\ v]$
{
fix v
fix y
obtain x **where** *y-def*: $y = \nu v\ x$
by (*meson* νv*-surj surj-def*)
have $(\exists r\ o_1.\ Some\ r = d_1\ F \wedge Some\ o_1 = d_\kappa\ (x^P) \wedge o_1 \in ex1\ r\ v) =$
$(\exists r\ o_1.\ Some\ r = d_1\ G \wedge Some\ o_1 = d_\kappa\ (x^P) \wedge o_1 \in ex1\ r\ v)$
using *1* **apply** − **by** *meta-solver*
moreover obtain r **where** *r-def*: $Some\ r = d_1\ F$
unfolding d_1*-def* **by** *auto*
moreover obtain s **where** *s-def*: $Some\ s = d_1\ G$
unfolding d_1*-def* **by** *auto*
moreover have $Some\ x = d_\kappa\ (x^P)$
using d_κ*-proper* **by** *simp*
ultimately have $(x \in ex1\ r\ v) = (x \in ex1\ s\ v)$
by (*metis option.inject*)
hence $(eval\Pi_1\ F)\ y\ dj\ v = (eval\Pi_1\ G)\ y\ dj\ v$
using *r-def s-def y-def* **by** (*simp add*: d_1*.rep-eq ex1-def*)
}
thus $(\lambda x.\ eval\Pi_1\ F\ x\ dj) = (\lambda x.\ eval\Pi_1\ G\ x\ dj)$
by *auto*
next
assume *1*: $(\lambda x.\ eval\Pi_1\ F\ x\ dj) = (\lambda x.\ eval\Pi_1\ G\ x\ dj)$
{
fix $y\ v$
obtain x **where** *x-def*: $x = \nu v\ y$
by *simp*
hence $eval\Pi_1\ F\ x\ dj = eval\Pi_1\ G\ x\ dj$
using *1* **by** *metis*
moreover obtain r **where** *r-def*: $Some\ r = d_1\ F$
unfolding d_1*-def* **by** *auto*
moreover obtain s **where** *s-def*: $Some\ s = d_1\ G$
unfolding d_1*-def* **by** *auto*
ultimately have $(y \in ex1\ r\ v) = (y \in ex1\ s\ v)$
by (*simp add*: d_1*.rep-eq ex1-def* νv*-surj x-def*)
hence $[(\!|F,y^P|\!) \equiv (\!|G,y^P|\!)\ in\ v]$
apply − **apply** *meta-solver*
using *r-def s-def* **by** (*metis Semantics.*d_κ*-proper option.inject*)
}
thus $[\forall x.\ \Box((\!|F,x^P|\!) \equiv (\!|G,x^P|\!))\ in\ v]$
using *T6 T8* **by** *fast*
qed

Remark. *Materially equivalent relations are equal in the embedded logic if and only if they also coincide in all other states.*

lemma *mat-eq-is-eq-if-eq-forall-j*:

assumes $[\forall\ x\ .\ \Box((\!|F,x^P|\!) \equiv (\!|G,x^P|\!))\ in\ v]$
shows $[F = G\ in\ v] \longleftrightarrow$
$(\forall\ s\ .\ s \neq dj \longrightarrow (\forall\ x\ .\ (eval\Pi_1\ F)\ x\ s = (eval\Pi_1\ G)\ x\ s))$
proof
interpret *MetaSolver* **.**
assume $[F = G\ in\ v]$
hence $F = G$
apply − **unfolding** *identity-*Π_1*-def* **by** *meta-solver*
thus $\forall s.\ s \neq dj \longrightarrow (\forall x.\ eval\Pi_1\ F\ x\ s = eval\Pi_1\ G\ x\ s)$
by *auto*
next
interpret *MetaSolver* **.**
assume $\forall s.\ s \neq dj \longrightarrow (\forall x.\ eval\Pi_1\ F\ x\ s = eval\Pi_1\ G\ x\ s)$
moreover have $((\lambda\ x\ .\ (eval\Pi_1\ F)\ x\ dj) = (\lambda\ x\ .\ (eval\Pi_1\ G)\ x\ dj))$
using *assms mat-eq-is-eq-dj* **by** *auto*
ultimately have $\forall s\ x.\ eval\Pi_1\ F\ x\ s = eval\Pi_1\ G\ x\ s$
by *metis*
hence $eval\Pi_1\ F = eval\Pi_1\ G$
by *blast*
hence $F = G$
by (*metis eval*Π_1*-inverse*)
thus $[F = G\ in\ v]$
unfolding *identity-*Π_1*-def* **using** Eq_1I **by** *auto*
qed

Remark. *Under the assumption that all properties behave in all states like in the actual state the defined equality degenerates to material equality.*

lemma assumes $\forall\ F\ x\ s\ .\ (eval\Pi_1\ F)\ x\ s = (eval\Pi_1\ F)\ x\ dj$
shows $[\forall\ x\ .\ \Box((\!|F,x^P|\!) \equiv (\!|G,x^P|\!))\ in\ v] \longleftrightarrow [F = G\ in\ v]$
by (*metis* (*no-types*) *MetaSolver.*Eq_1S *assms identity-*Π_1*-def*
mat-eq-is-eq-dj mat-eq-is-eq-if-eq-forall-j)

A.12.6. Lambda Expressions

lemma *lambda-interpret-1*:
assumes $[a = b\ in\ v]$
shows $(\boldsymbol{\lambda}x.\ (\!|R,x^P,a|\!)) = (\boldsymbol{\lambda}x\ .\ (\!|R,x^P,\ b|\!))$
proof −
have $a = b$
using *MetaSolver.Eq*κ*S Semantics.*d_κ*-inject assms*
identity-κ*-def* **by** *auto*
thus *?thesis* **by** *simp*
qed

lemma *lambda-interpret-2*:
assumes $[a = (\boldsymbol{\iota}y.\ (\!|G,y^P|\!))\ in\ v]$
shows $(\boldsymbol{\lambda}x.\ (\!|R,x^P,a|\!)) = (\boldsymbol{\lambda}x\ .\ (\!|R,x^P,\ \boldsymbol{\iota}y.\ (\!|G,y^P|\!)|\!))$
proof −
have $a = (\boldsymbol{\iota}y.\ (\!|G,y^P|\!))$
using *MetaSolver.Eq*κ*S Semantics.*d_κ*-inject assms*
identity-κ*-def* **by** *auto*
thus *?thesis* **by** *simp*
qed

end

theory *TAO-99-Paradox*
imports *TAO-9-PLM TAO-98-ArtificialTheorems*
begin

A.13. Paradox

Under the additional assumption that expressions of the form $\lambda x.\ (\!|G,\iota y.\ \varphi\ y\ x|\!)$ for arbitrary φ are *proper maps*, for which β-conversion holds, the theory becomes inconsistent.

A.13.1. Auxiliary Lemmas

lemma *exe-impl-exists*:
$[(\!|(\lambda x\ .\ \forall\ p\ .\ p \to p),\ \iota y\ .\ \varphi\ y\ x|\!) \equiv (\exists !y\ .\ \mathcal{A}\varphi\ y\ x)\ in\ v]$
proof (*rule* $\equiv I$; *rule CP*)
fix $\varphi :: \nu \Rightarrow \nu \Rightarrow o$ **and** $x :: \nu$ **and** $v :: i$
assume $[(\!|(\lambda x\ .\ \forall\ p\ .\ p \to p),\iota y\ .\ \varphi\ y\ x|\!)\ in\ v]$
hence $[\exists\, y.\ \mathcal{A}\varphi\ y\ x\ \&\ (\forall\, z.\ \mathcal{A}\varphi\ z\ x \to z = y)$
$\&\ (\!|(\lambda x\ .\ \forall\ p\ .\ p \to p),\ y^P|\!)\ in\ v]$
using *nec-russell-axiom*[*equiv-lr*] *SimpleExOrEnc.intros* **by** *auto*
then obtain y **where**
$[\mathcal{A}\varphi\ y\ x\ \&\ (\forall\, z.\ \mathcal{A}\varphi\ z\ x \to z = y)$
$\&\ (\!|(\lambda x\ .\ \forall\ p\ .\ p \to p),\ y^P|\!)\ in\ v]$
by (*rule Instantiate*)
hence $[\mathcal{A}\varphi\ y\ x\ \&\ (\forall\, z.\ \mathcal{A}\varphi\ z\ x \to z = y)\ in\ v]$
using $\&E$ **by** *blast*
hence $[\exists\, y\ .\ \mathcal{A}\varphi\ y\ x\ \&\ (\forall\, z.\ \mathcal{A}\varphi\ z\ x \to z = y)\ in\ v]$
by (*rule existential*)
thus $[\exists !y.\ \mathcal{A}\varphi\ y\ x\ in\ v]$
unfolding *exists-unique-def* **by** *simp*
next
fix $\varphi :: \nu \Rightarrow \nu \Rightarrow o$ **and** $x :: \nu$ **and** $v :: i$
assume $[\exists !y.\ \mathcal{A}\varphi\ y\ x\ in\ v]$
hence $[\exists\, y.\ \mathcal{A}\varphi\ y\ x\ \&\ (\forall\, z.\ \mathcal{A}\varphi\ z\ x \to z = y)\ in\ v]$
unfolding *exists-unique-def* **by** *simp*
then obtain y **where**
$[\mathcal{A}\varphi\ y\ x\ \&\ (\forall\, z.\ \mathcal{A}\varphi\ z\ x \to z = y)\ in\ v]$
by (*rule Instantiate*)
moreover have $[(\!|(\lambda x\ .\ \forall\ p\ .\ p \to p),y^P|\!)\ in\ v]$
apply (*rule beta-C-meta-1*[*equiv-rl*])
apply *show-proper*
by *PLM-solver*
ultimately have $[\mathcal{A}\varphi\ y\ x\ \&\ (\forall\, z.\ \mathcal{A}\varphi\ z\ x \to z = y)$
$\&\ (\!|(\lambda x\ .\ \forall\ p\ .\ p \to p),y^P|\!)\ in\ v]$
using $\&I$ **by** *blast*
hence $[\exists\ \ y\ .\ \mathcal{A}\varphi\ y\ x\ \&\ (\forall\, z.\ \mathcal{A}\varphi\ z\ x \to z = y)$
$\&\ (\!|(\lambda x\ .\ \forall\ p\ .\ p \to p),y^P|\!)\ in\ v]$
by (*rule existential*)
thus $[(\!|(\lambda x\ .\ \forall\ p\ .\ p \to p),\ \iota y.\ \varphi\ y\ x|\!)\ in\ v]$
using *nec-russell-axiom*[*equiv-rl*]
SimpleExOrEnc.intros **by** *auto*
qed

lemma *exists-unique-actual-equiv*:
$[(\exists !y\ .\ \mathcal{A}(y = x\ \&\ \psi\ (x^P))) \equiv \mathcal{A}\psi\ (x^P)\ in\ v]$
proof (*rule* $\equiv I$; *rule CP*)
fix $x\ v$
let $?\varphi = \lambda\ y\ x.\ y = x\ \&\ \psi\ (x^P)$

assume [$\exists !y.\ \mathcal{A}?\varphi\ y\ x\ in\ v$]
hence [$\exists\,\alpha.\ \mathcal{A}?\varphi\ \alpha\ x\ \&\ (\forall\,\beta.\ \mathcal{A}?\varphi\ \beta\ x \rightarrow \beta = \alpha)\ in\ v$]
unfolding *exists-unique-def* **by** *simp*
then obtain α **where**
[$\mathcal{A}?\varphi\ \alpha\ x\ \&\ (\forall\,\beta.\ \mathcal{A}?\varphi\ \beta\ x \rightarrow \beta = \alpha)\ in\ v$]
by (*rule Instantiate*)
hence [$\mathcal{A}(\alpha = x\ \&\ \psi\ (x^P))\ in\ v$]
using &E **by** *blast*
thus [$\mathcal{A}(\psi\ (x^P))\ in\ v$]
using *Act-Basic-2*[*equiv-lr*] &E **by** *blast*
next
fix $x\ v$
let $?\varphi = \lambda\ y\ x.\ y = x\ \&\ \psi\ (x^P)$
assume *1*: [$\mathcal{A}\psi\ (x^P)\ in\ v$]
have [$x = x\ in\ v$]
using *id-eq-1*[**where** $'a{=}\nu$] **by** *simp*
hence [$\mathcal{A}(x = x)\ in\ v$]
using *id-act-3*[*equiv-lr*] **by** *fast*
hence [$\mathcal{A}(x = x\ \&\ \psi\ (x^P))\ in\ v$]
using *1 Act-Basic-2*[*equiv-rl*] &I **by** *blast*
hence [$\mathcal{A}?\varphi\ x\ x\ in\ v$]
by *simp*
moreover have [$\forall\,\beta.\ \mathcal{A}?\varphi\ \beta\ x \rightarrow \beta = x\ in\ v$]
proof (*rule* $\forall I$; *rule CP*)
fix β
assume [$\mathcal{A}?\varphi\ \beta\ x\ in\ v$]
hence [$\mathcal{A}(\beta = x)\ in\ v$]
using *Act-Basic-2*[*equiv-lr*] &E **by** *fast*
thus [$\beta = x\ in\ v$] **using** *id-act-3*[*equiv-rl*] **by** *fast*
qed
ultimately have [$\mathcal{A}?\varphi\ x\ x\ \&\ (\forall\,\beta.\ \mathcal{A}?\varphi\ \beta\ x \rightarrow \beta = x)\ in\ v$]
using &I **by** *fast*
hence [$\exists\,\alpha.\ \mathcal{A}?\varphi\ \alpha\ x\ \&\ (\forall\,\beta.\ \mathcal{A}?\varphi\ \beta\ x \rightarrow \beta = \alpha)\ in\ v$]
by (*rule existential*)
thus [$\exists !y.\ \mathcal{A}?\varphi\ y\ x\ in\ v$]
unfolding *exists-unique-def* **by** *simp*
qed

A.13.2. Fake β-Conversion using Description Backdoor

definition *backdoor* **where**
$backdoor \equiv \lambda\ \psi\ .\ \boldsymbol{\lambda}x\ .\ (\!|(\boldsymbol{\lambda}x\ .\ \forall\ \ p\ .\ p \rightarrow p),\ \boldsymbol{\iota}y\ .\ y = x\ \&\ \psi\ (x^P)|\!)$

lemma *fake-beta*:
assumes $\bigwedge\ G\ \varphi.\ IsProperInX\ (\lambda x\ .\ (\!|G,\boldsymbol{\iota}y\ .\ \varphi\ y\ x|\!))$
shows [$(\!|backdoor\ (\lambda\ x\ .\ \psi\ x),\ x^P|\!) \equiv \mathcal{A}\psi\ (x^P)\ in\ v$]
proof (*rule* $\equiv I$; *rule CP*)
assume [$(\!|backdoor\ \psi,x^P|\!)\ in\ v$]
hence [$(\!|\boldsymbol{\lambda}x.\ \forall\,p.\ p \rightarrow p,\boldsymbol{\iota}y.\ y = x\ \&\ \psi\ (x^P)|\!)\ in\ v$]
using *beta-C-meta-1*[*equiv-lr*, *OF assms*]
unfolding *backdoor-def identity-*ν*-def* **by** *fast*
hence [$\exists !y.\ \mathcal{A}\ (y = x\ \&\ \psi\ (x^P))\ in\ v$]
using *exe-impl-exists*[*equiv-lr*] **by** *fast*
thus [$\mathcal{A}\psi\ (x^P)\ in\ v$]
using *exists-unique-actual-equiv*[*equiv-lr*] **by** *blast*
next
assume [$\mathcal{A}\psi\ (x^P)\ in\ v$]
hence [$\exists !y.\ \mathcal{A}\ (y = x\ \&\ \psi\ (x^P))\ in\ v$]

using *exists-unique-actual-equiv*[*equiv-rl*] **by** *blast*
hence [(|**λ***x*. ∀ *p*. *p* → *p*,**ι***y*. *y* = *x* & *ψ* (*x*^{P})|) *in v*]
using *exe-impl-exists*[*equiv-rl*] **by** *fast*
thus [(|*backdoor ψ*,*x*^{P}|) *in v*]
using *beta-C-meta-1*[*equiv-rl*, *OF assms*]
unfolding *backdoor-def* **unfolding** *identity-ν-def* **by** *fast*
qed

lemma *fake-beta-act*:
assumes ⋀ *G φ*. *IsProperInX* (λ*x* . (| *G*,**ι***y* . *φ y x*|))
shows [(|*backdoor* (λ *x* . *ψ x*), *x*^{P}|) ≡ *ψ* (*x*^{P}) *in dw*]
using *fake-beta*[*OF assms*]
logic-actual[*necessitation-averse-axiom-instance*]
intro-elim-6-e **by** *blast*

A.13.3. Resulting Paradox

lemma *paradox*:
assumes ⋀ *G φ*. *IsProperInX* (λ*x* . (| *G*,**ι***y* . *φ y x*|))
shows *False*
proof −
obtain *K* **where** *K-def*:
K = *backdoor* (λ *x* . ∃ *F* . {|*x*,*F*|} & ¬(|*F*,*x*|)) **by** *auto*
have [∃ *x*. (|*A*!,*x*^{P}|) & (∀ *F*. {|*x*^{P},*F*|} ≡ (*F* = *K*)) *in dw*]
using *A-objects*[*axiom-instance*] **by** *fast*
then obtain *x* **where** *x-prop*:
[(|*A*!,*x*^{P}|) & (∀ *F*. {|*x*^{P},*F*|} ≡ (*F* = *K*)) *in dw*]
by (*rule Instantiate*)
{
assume [(|*K*,*x*^{P}|) *in dw*]
hence [∃ *F* . {|*x*^{P},*F*|} & ¬(|*F*,*x*^{P}|) *in dw*]
unfolding *K-def* **using** *fake-beta-act*[*OF assms*, *equiv-lr*]
by *blast*
then obtain *F* **where** *F-def*:
[{|*x*^{P},*F*|} & ¬(|*F*,*x*^{P}|) *in dw*] **by** (*rule Instantiate*)
hence [*F* = *K in dw*]
using *x-prop*[*conj2*, *THEN* ∀ *E*[**where** *β*=*F*], *equiv-lr*]
&*E* **unfolding** *K-def* **by** *blast*
hence [¬(|*K*,*x*^{P}|) *in dw*]
using *l-identity*[*axiom-instance*,*deduction*,*deduction*]
F-def[*conj2*] **by** *fast*
}
hence *1*: [¬(|*K*,*x*^{P}|) *in dw*]
using *reductio-aa-1* **by** *blast*
hence [¬(∃ *F* . {|*x*^{P},*F*|} & ¬(|*F*,*x*^{P}|)) *in dw*]
using *fake-beta-act*[*OF assms*,
THEN oth-class-taut-5-d[*equiv-lr*],
equiv-lr]
unfolding *K-def* **by** *blast*
hence [∀ *F* . {|*x*^{P},*F*|} → (|*F*,*x*^{P}|) *in dw*]
apply − **unfolding** *exists-def* **by** *PLM-solver*
moreover have [{|*x*^{P},*K*|} *in dw*]
using *x-prop*[*conj2*, *THEN* ∀ *E*[**where** *β*=*K*], *equiv-rl*]
id-eq-1 **by** *blast*
ultimately have [(|*K*,*x*^{P}|) *in dw*]
using ∀ *E vdash-properties-10* **by** *blast*
hence ⋀*φ*. [*φ in dw*]
using *raa-cor-2 1* **by** *blast*

thus *False* **using** *Semantics.T4* **by** *auto*
qed

A.13.4. Original Version of the Paradox

Originally the paradox was discovered using the following construction based on the comprehension theorem for relations without the explicit construction of the description backdoor and the resulting fake-β-conversion.

lemma assumes $\bigwedge$ *G* φ. *IsProperInX* (λx . ($|G,\iota y$. φ *y x*$|$))
shows *Fx-equiv-xH*: [$\forall$ *H* . $\exists$ *F* . $\Box(\forall x.$ ($|F,x^P|$) $\equiv$ $\{\!|x^P,H|\!\}$) *in v*]
proof (*rule* $\forall I$)
fix *H*
let *?G* = ($\boldsymbol{\lambda} x$. $\forall$ *p* . *p* $\rightarrow$ *p*)
obtain φ **where** φ*-def*: φ = (λ *y x* . (y^P) = *x* & $\{\!|x,H|\!\}$) **by** *auto*
have [$\exists$ *F*. $\Box(\forall x.$ ($|F,x^P|$) $\equiv$ ($|$*?G*,ιy . φ *y* (x^P)$|$)) *in v*]
using *relations-1*[*OF assms*] **by** *simp*
hence *1*: [$\exists$ *F*. $\Box(\forall x.$ ($|F,x^P|$) $\equiv$ ($\exists !y$. $\boldsymbol{\mathcal{A}}\varphi$ *y* (x^P))) *in v*]
apply − **apply** (*PLM-subst-method*
λ *x* . ($|$*?G*,ιy . φ *y* (x^P)$|$) λ *x* . ($\exists !y.$ $\boldsymbol{\mathcal{A}}\varphi$ *y* (x^P)))
using *exe-impl-exists* **by** *auto*
then obtain *F* **where** *F-def*: [$\Box(\forall x.$ ($|F,x^P|$) $\equiv$ ($\exists !y$. $\boldsymbol{\mathcal{A}}\varphi$ *y* (x^P))) *in v*]
by (*rule Instantiate*)
moreover have *2*: $\bigwedge$ *v x* . [($\exists !y$. $\boldsymbol{\mathcal{A}}\varphi$ *y* (x^P)) $\equiv$ $\{\!|x^P,H|\!\}$ *in v*]
proof (*rule* $\equiv I$; *rule CP*)
fix *x v*
assume [$\exists !y.$ $\boldsymbol{\mathcal{A}}\varphi$ *y* (x^P) *in v*]
hence [$\exists\,\alpha.$ $\boldsymbol{\mathcal{A}}\varphi$ α (x^P) & ($\forall\,\beta.$ $\boldsymbol{\mathcal{A}}\varphi$ β (x^P) $\rightarrow$ β = α) *in v*]
unfolding *exists-unique-def* **by** *simp*
then obtain α **where** [$\boldsymbol{\mathcal{A}}\varphi$ α (x^P) & ($\forall\,\beta.$ $\boldsymbol{\mathcal{A}}\varphi$ β (x^P) $\rightarrow$ β = α) *in v*]
by (*rule Instantiate*)
hence [$\boldsymbol{\mathcal{A}}(\alpha^P = x^P$ & $\{\!|x^P,H|\!\}$) *in v*]
unfolding φ*-def* **using** &*E* **by** *blast*
hence [$\boldsymbol{\mathcal{A}}(\{\!|x^P,H|\!\})$ *in v*]
using *Act-Basic-2*[*equiv-lr*] &*E* **by** *blast*
thus [$\{\!|x^P,H|\!\}$ *in v*]
using *en-eq-10*[*equiv-lr*] **by** *simp*
next
fix *x v*
assume [$\{\!|x^P,H|\!\}$ *in v*]
hence *1*: [$\boldsymbol{\mathcal{A}}(\{\!|x^P,H|\!\})$ *in v*]
using *en-eq-10*[*equiv-rl*] **by** *blast*
have [*x* = *x in v*]
using *id-eq-1*[**where** *'a*=ν] **by** *simp*
hence [$\boldsymbol{\mathcal{A}}(x = x)$ *in v*]
using *id-act-3*[*equiv-lr*] **by** *fast*
hence [$\boldsymbol{\mathcal{A}}(x^P = x^P$ & $\{\!|x^P,H|\!\}$) *in v*]
unfolding *identity-*ν*-def* **using** *1 Act-Basic-2*[*equiv-rl*] &*I* **by** *blast*
hence [$\boldsymbol{\mathcal{A}}\varphi$ *x* (x^P) *in v*]
unfolding φ*-def* **by** *simp*
moreover have [$\forall\,\beta.$ $\boldsymbol{\mathcal{A}}\varphi$ β (x^P) $\rightarrow$ β = *x in v*]
proof (*rule* $\forall I$; *rule CP*)
fix β
assume [$\boldsymbol{\mathcal{A}}\varphi$ β (x^P) *in v*]
hence [$\boldsymbol{\mathcal{A}}(\beta = x)$ *in v*]
unfolding φ*-def identity-*ν*-def*
using *Act-Basic-2*[*equiv-lr*] &*E* **by** *fast*
thus [β = *x in v*] **using** *id-act-3*[*equiv-rl*] **by** *fast*

```
    qed
    ultimately have [𝒜φ x (x^P) & (∀β. 𝒜φ β (x^P) → β = x) in v]
      using &I by fast
    hence [∃α. 𝒜φ α (x^P) & (∀β. 𝒜φ β (x^P) → β = α) in v]
      by (rule existential)
    thus [∃!y. 𝒜φ y (x^P) in v]
      unfolding exists-unique-def by simp
  qed
  have [□(∀x. (|F,x^P|) ≡ {|x^P,H|}) in v]
    apply (PLM-subst-goal-method
        λφ . □(∀x. (|F,x^P|) ≡ φ x)
        λ x . (∃!y . 𝒜φ y (x^P)))
    using 2 F-def by auto
  thus [∃ F . □(∀x. (|F,x^P|) ≡ {|x^P,H|}) in v]
    by (rule existential)
qed


lemma
  assumes is-propositional: (⋀G φ. IsProperInX (λx. (|G,ιy. φ y x|)))
      and Abs-x: [(|A!,x^P|) in v]
      and Abs-y: [(|A!,y^P|) in v]
      and noteq: [x ≠ y in v]
shows diffprop: [∃ F . ¬((|F,x^P|) ≡ (|F,y^P|)) in v]
proof −
  have [∃ F . ¬({|x^P, F|} ≡ {|y^P, F|}) in v]
    using noteq unfolding exists-def
  proof (rule reductio-aa-2)
    assume 1: [∀F. ¬¬({|x^P,F|} ≡ {|y^P,F|}) in v]
    {
      fix F
      have [({|x^P,F|} ≡ {|y^P,F|}) in v]
        using 1[THEN ∀E] useful-tautologies-1[deduction] by blast
    }
    hence [∀F. {|x^P,F|} ≡ {|y^P,F|} in v] by (rule ∀I)
    thus [x = y in v]
      unfolding identity-ν-def
      using ab-obey-1[deduction, deduction]
            Abs-x Abs-y &I by blast
  qed
  then obtain H where H-def: [¬({|x^P, H|} ≡ {|y^P, H|}) in v]
    by (rule Instantiate)
  hence 2: [({|x^P, H|} & ¬{|y^P, H|}) ∨ (¬{|x^P, H|} & {|y^P, H|}) in v]
    apply − by PLM-solver
  have [∃F. □(∀x. (|F,x^P|) ≡ {|x^P,H|}) in v]
    using Fx-equiv-xH[OF is-propositional, THEN ∀E] by simp
  then obtain F where [□(∀x. (|F,x^P|) ≡ {|x^P,H|}) in v]
    by (rule Instantiate)
  hence F-prop: [∀x. (|F,x^P|) ≡ {|x^P,H|} in v]
    using qml-2[axiom-instance, deduction] by blast
  hence a: [(|F,x^P|) ≡ {|x^P,H|} in v]
    using ∀E by blast
  have b: [(|F,y^P|) ≡ {|y^P,H|} in v]
    using F-prop ∀E by blast
  {
    assume 1: [{|x^P, H|} & ¬{|y^P, H|} in v]
    hence [(|F,x^P|) in v]
      using a[equiv-rl] &E by blast
```

```
    moreover have [¬(|F,y^P|) in v]
      using b[THEN oth-class-taut-5-d[equiv-lr], equiv-rl] 1[conj2] by auto
    ultimately have [(|F,x^P|) & (¬(|F,y^P|)) in v]
      by (rule &I)
    hence [((|F,x^P|) & ¬(|F,y^P|)) ∨ (¬(|F,x^P|) & (|F,y^P|)) in v]
      using ∨I by blast
    hence [¬((|F,x^P|) ≡ (|F,y^P|)) in v]
      using oth-class-taut-5-j[equiv-rl] by blast
  }
  moreover {
    assume 1: [¬{|x^P, H|} & {|y^P, H|} in v]
    hence [(|F,y^P|) in v]
      using b[equiv-rl] &E by blast
    moreover have [¬(|F,x^P|) in v]
      using a[THEN oth-class-taut-5-d[equiv-lr], equiv-rl] 1[conj1] by auto
    ultimately have [¬(|F,x^P|) & (|F,y^P|) in v]
      using &I by blast
    hence [((|F,x^P|) & ¬(|F,y^P|)) ∨ (¬(|F,x^P|) & (|F,y^P|)) in v]
      using ∨I by blast
    hence [¬((|F,x^P|) ≡ (|F,y^P|)) in v]
      using oth-class-taut-5-j[equiv-rl] by blast
  }
  ultimately have [¬((|F,x^P|) ≡ (|F,y^P|)) in v]
    using 2 intro-elim-4-b reductio-aa-1 by blast
  thus [∃ F . ¬((|F,x^P|) ≡ (|F,y^P|)) in v]
    by (rule existential)
qed

lemma original-paradox:
  assumes is-propositional: (⋀G φ. IsProperInX (λx. (|G,ιy. φ y x|)))
  shows False
proof −
  fix v
  have [∃ x y. (|A!,x^P|) & (|A!,y^P|) & x ≠ y & (∀ F. (|F,x^P|) ≡ (|F,y^P|)) in v]
    using aclassical2 by auto
  then obtain x where
    [∃ y. (|A!,x^P|) & (|A!,y^P|) & x ≠ y & (∀ F. (|F,x^P|) ≡ (|F,y^P|)) in v]
    by (rule Instantiate)
  then obtain y where xy-def:
    [(|A!,x^P|) & (|A!,y^P|) & x ≠ y & (∀ F. (|F,x^P|) ≡ (|F,y^P|)) in v]
    by (rule Instantiate)
  have [∃ F . ¬((|F,x^P|) ≡ (|F,y^P|)) in v]
    using diffprop[OF assms, OF xy-def[conj1,conj1,conj1],
                   OF xy-def[conj1,conj1,conj2],
                   OF xy-def[conj1,conj2]]
    by auto
  then obtain F where [¬((|F,x^P|) ≡ (|F,y^P|)) in v]
    by (rule Instantiate)
  moreover have [(|F,x^P|) ≡ (|F,y^P|) in v]
    using xy-def[conj2] by (rule ∀E)
  ultimately have ⋀φ.[φ in v]
    using PLM.raa-cor-2 by blast
  thus False
    using Semantics.T4 by auto
qed

end
```

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Selbstständigkeitserklärung

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