

# Algorithms

WS 2012/13

## Exercises 2

1. **Network Flow (Niveau II)** Prove the Theorem:

For a network  $(V, E, s, t)$  with capacities  $\text{cap} : E \rightarrow \mathbb{R}_+$  the maximum value of a flow is equal to the minimum capacity of an  $(s, t)$ -cut:

$$\max\{\text{val}(f) \mid f \text{ is a flow}\} = \min\{\text{cap}(S, T) \mid (S, T) \text{ is an } (s, t)\text{-cut}\}$$

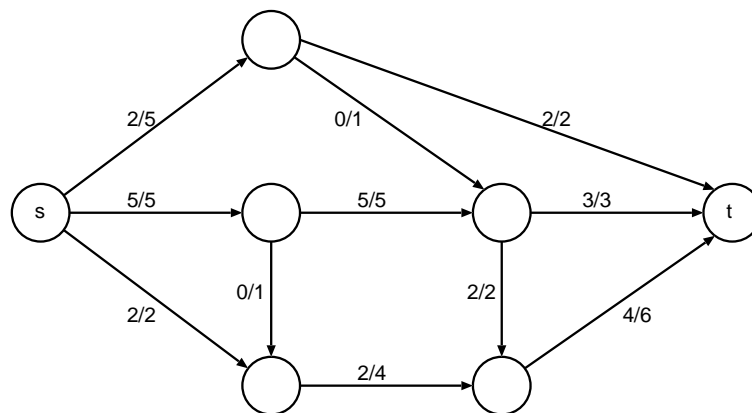
Hint: Show that the following conditions are equivalent:

- (a)  $f$  is a maximum flow.
- (b) The residual network  $G_f$  contains no augmenting path.
- (c)  $\text{val}(f) = \text{cap}(S, T)$  for some cut  $(S, T)$  of  $G$

2. **Network Flow (Niveau I)** Assume a flow network with edge and additional vertex capacities. Each vertex  $v$  has a limit on the flow that can pass through it. Explain how to transform this flow network into an equivalent flow network without vertex capacities.

3. **Ford-Fulkerson (Niveau I)**

- (a) Use the Ford-Fulkerson algorithm to find a maximum flow in the network



Start with the initial flow  $f$ . An edge label  $f/c$  means initial flow  $f$  and capacity  $c$ .

- (b) Find a minimum cut proving the maximality of the flow.

#### 4. Matching and Bipartite Graphs (Niveau I)

- (a) Apply the matching augmenting algorithm for bipartite graphs to the graph below and compute a maximum cardinality matching from the initial matching.

